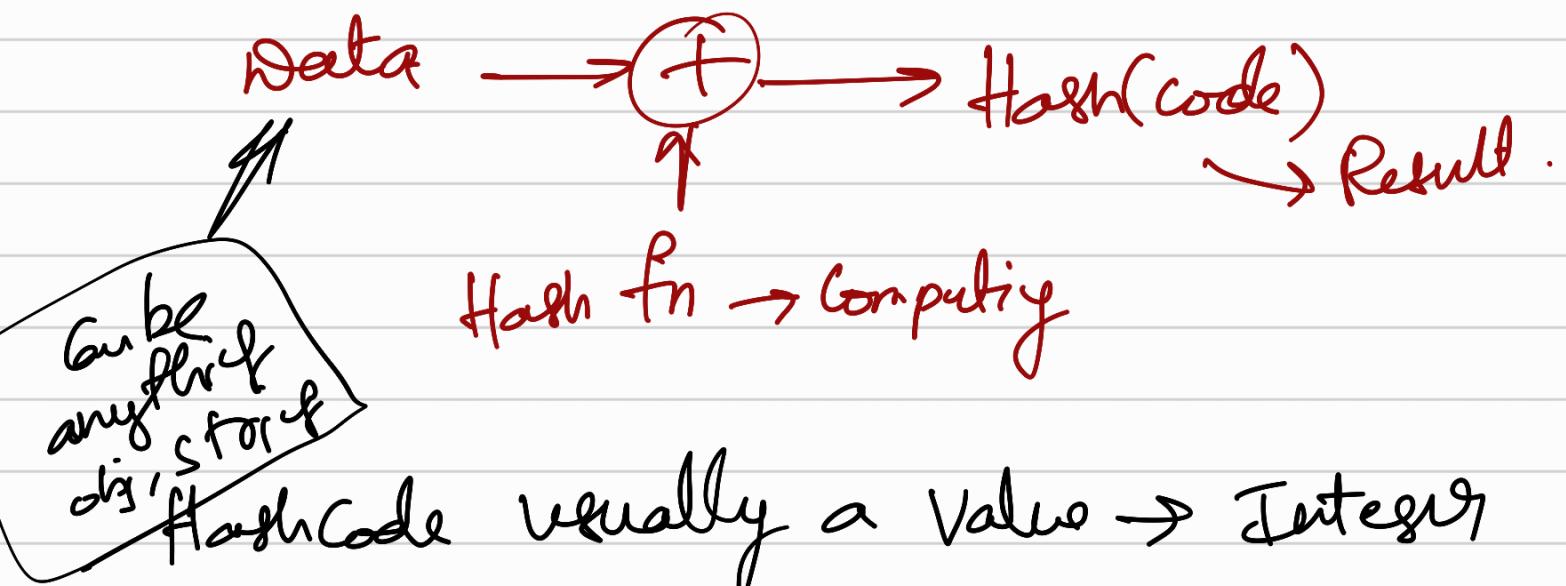
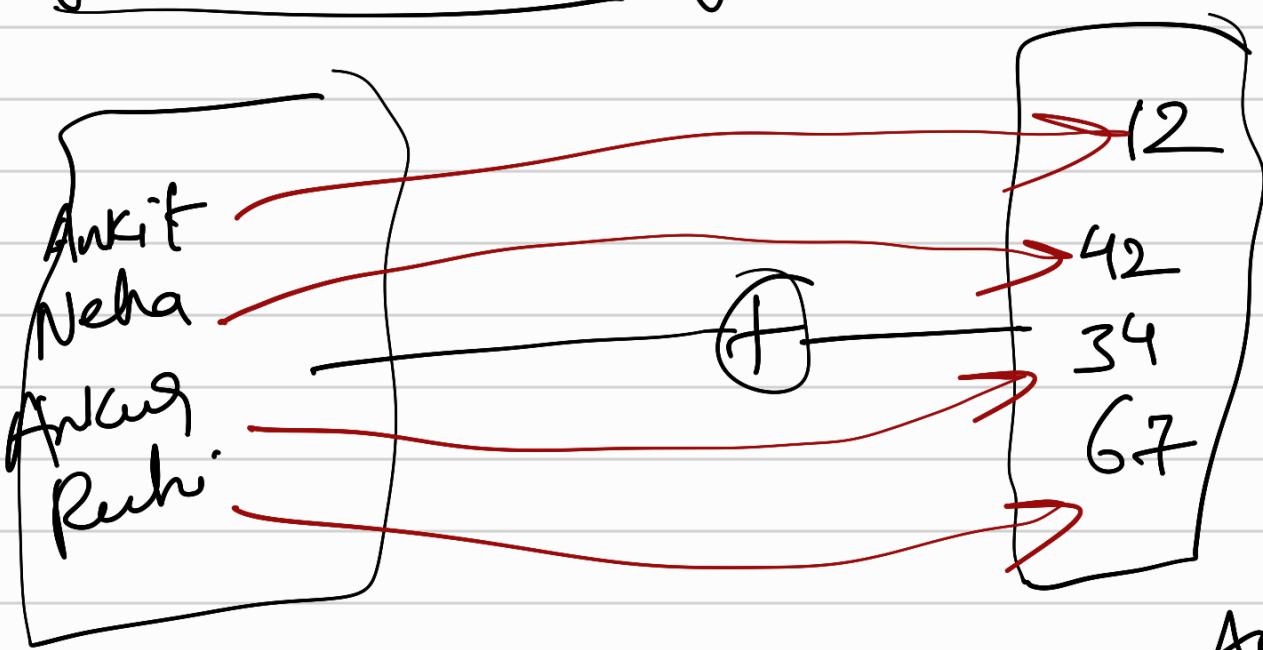


Hashing ?



HashCode usually a Value \rightarrow Integer
it can be other but
mostly int gr.

Purpose of Hashing :-



Search is easy
→ Array index

12

42

34

67

Ankit

Neha

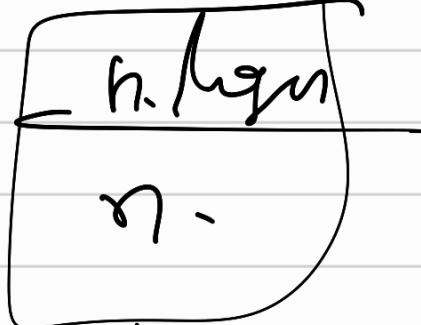
Ankur

Ruchi

✓ Accessing time becomes constant

but Hash fn ofp Constant

(Q) if binary Search used Complexity
becomes
not const



more time to search.

✓ Hence if we have sorted q
(faster op's) accessing will be
easier.

Range of Hash Values

0-2³²

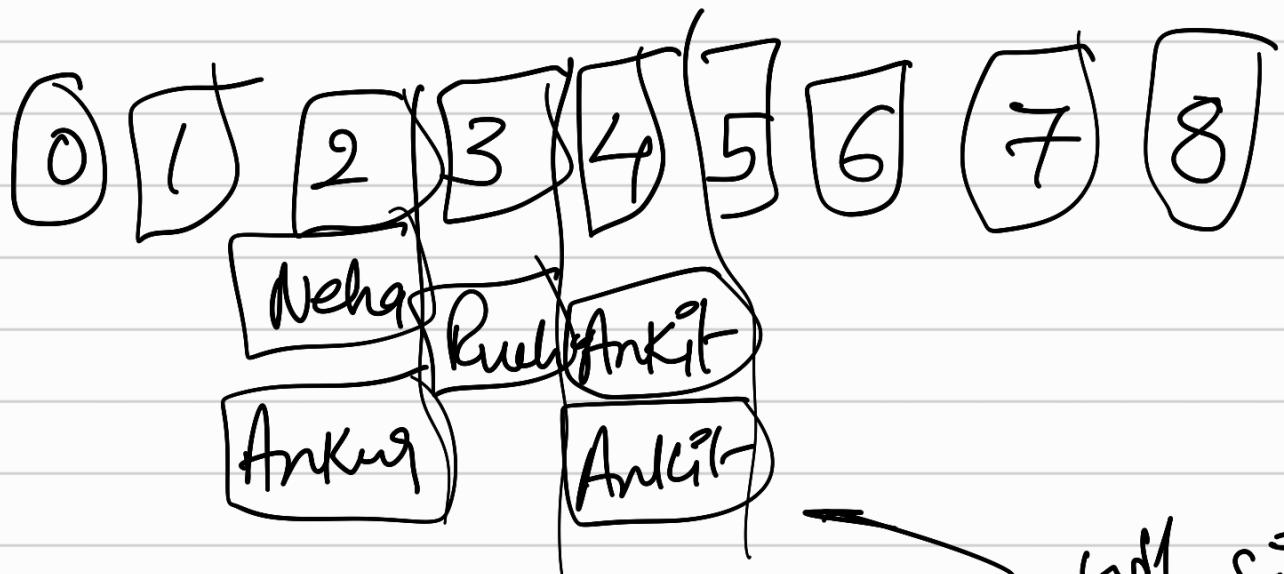
Many values have
we call fake modulus.

Ankit, Neha, Ankus, Ruchi

1

{2, 2, 13}

%8.

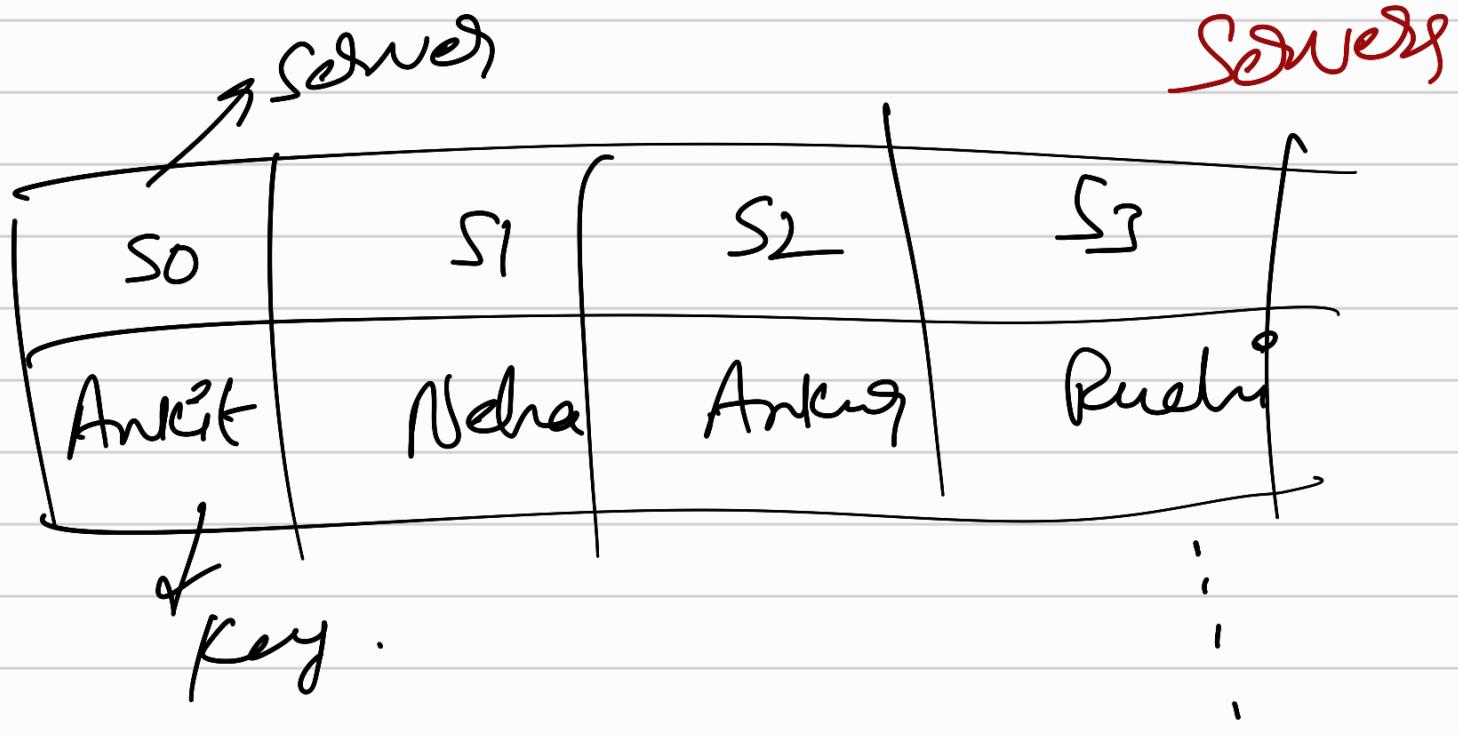
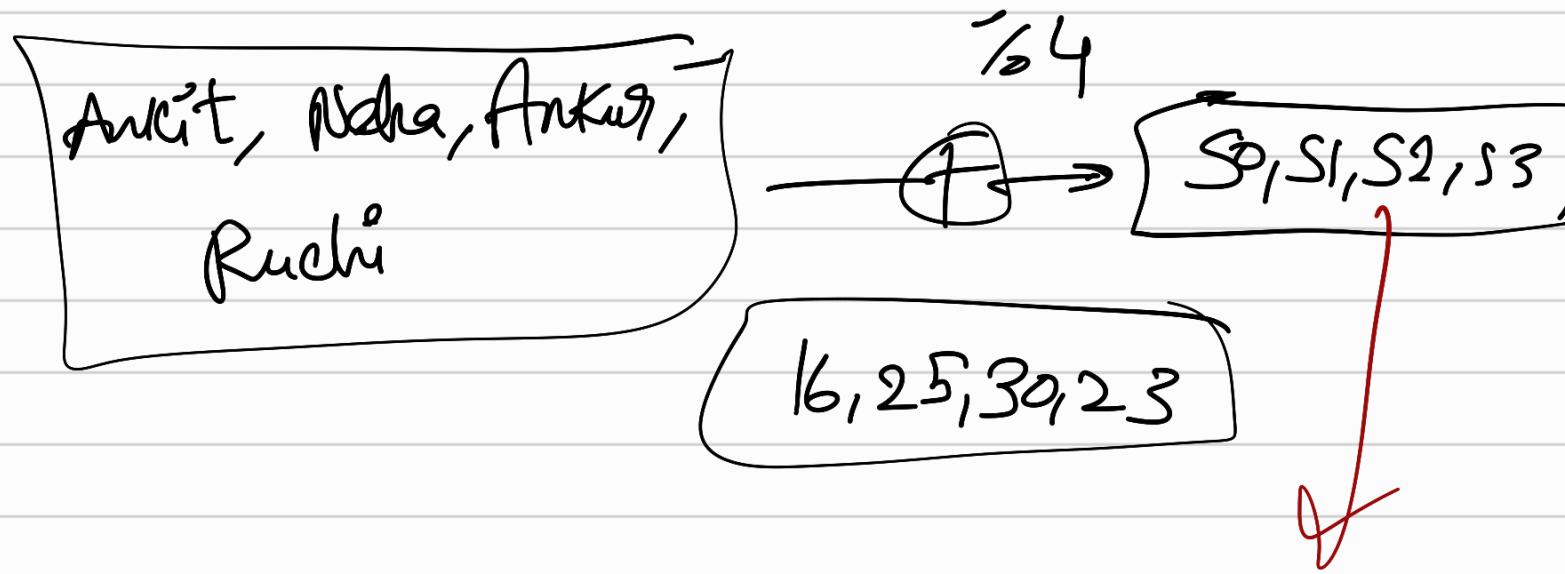


will sit in
buckets.

So Hashing helps us in storing

key value pairs & accessing key
value pairs very very efficiently

Hashing Use Cases :-

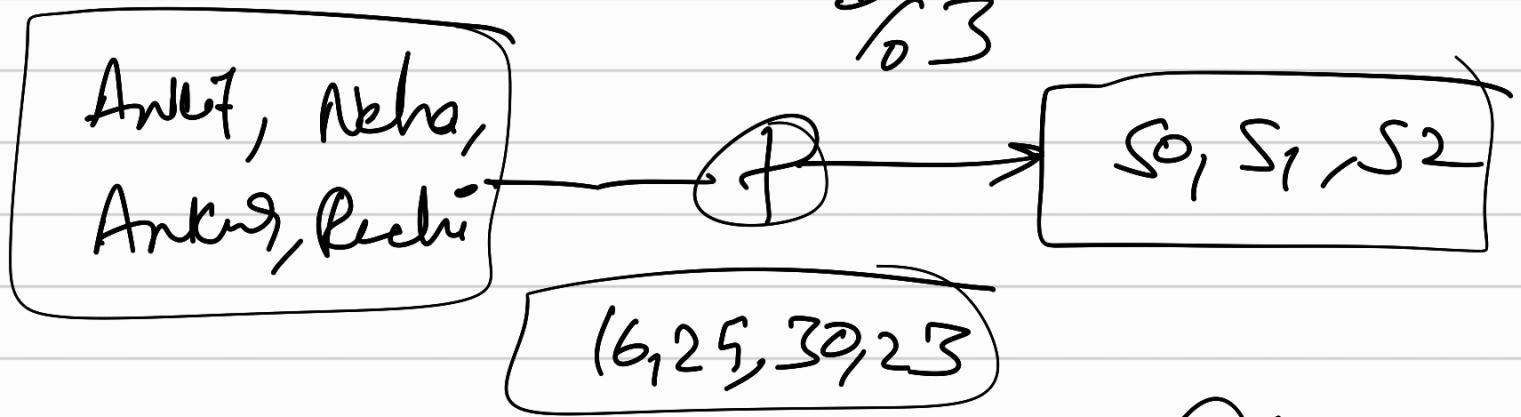


even increased key distribution
would be in similar manner.

problem with above approach

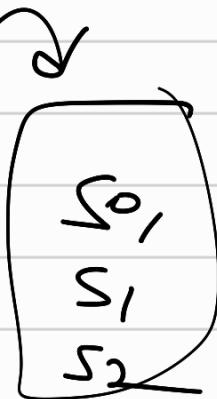
if Servers interested

/ Desired

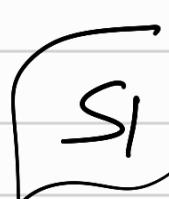


Rechi has to be assigned

but S₃ removed.



$$16 \% \cdot 3 = 1$$



$$25 \% \cdot 3 = 1$$



$$30 \% \cdot 3 = 0$$



$$23 \% \cdot 3 = 2$$



All keys are pre-assigned

Computation becomes more

if several added / removed.

This is the problem with
normal Hashif

↓
scale in } } → Computation
scale out } becomes more

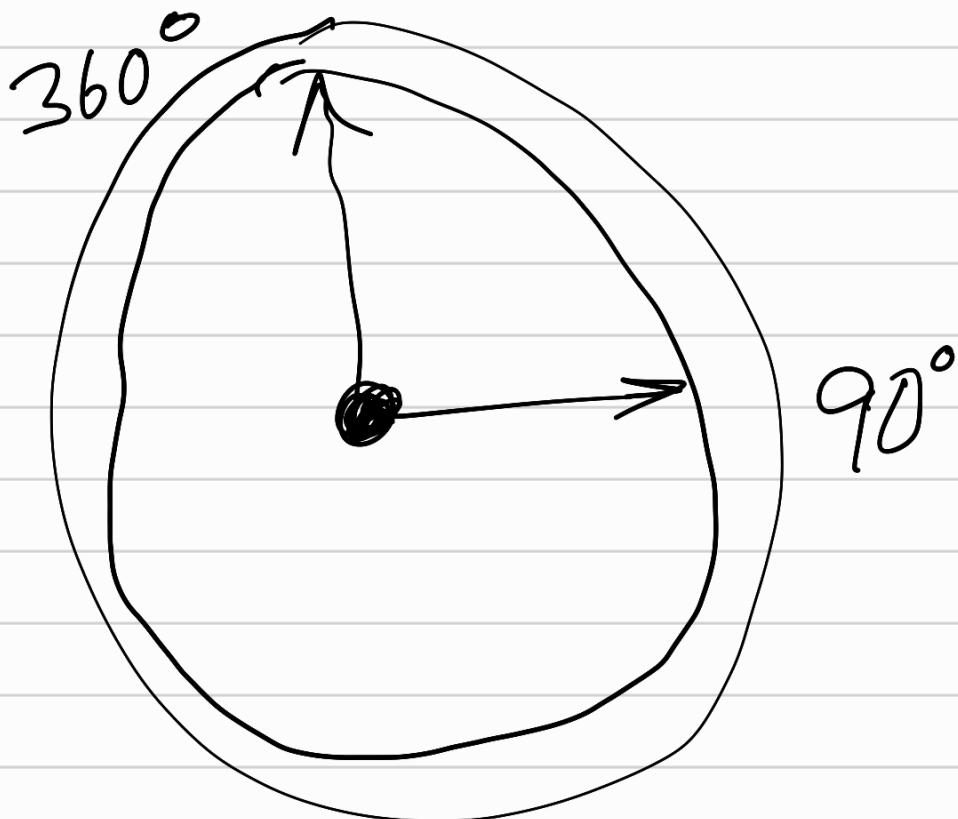
to minimize

↓
Concurrent Hashif.

Remapping of the
existing keys or
increasing/decreasing
of servers

?

Constant Hashing



Ex:-

Hash obj range

[0 - 100]

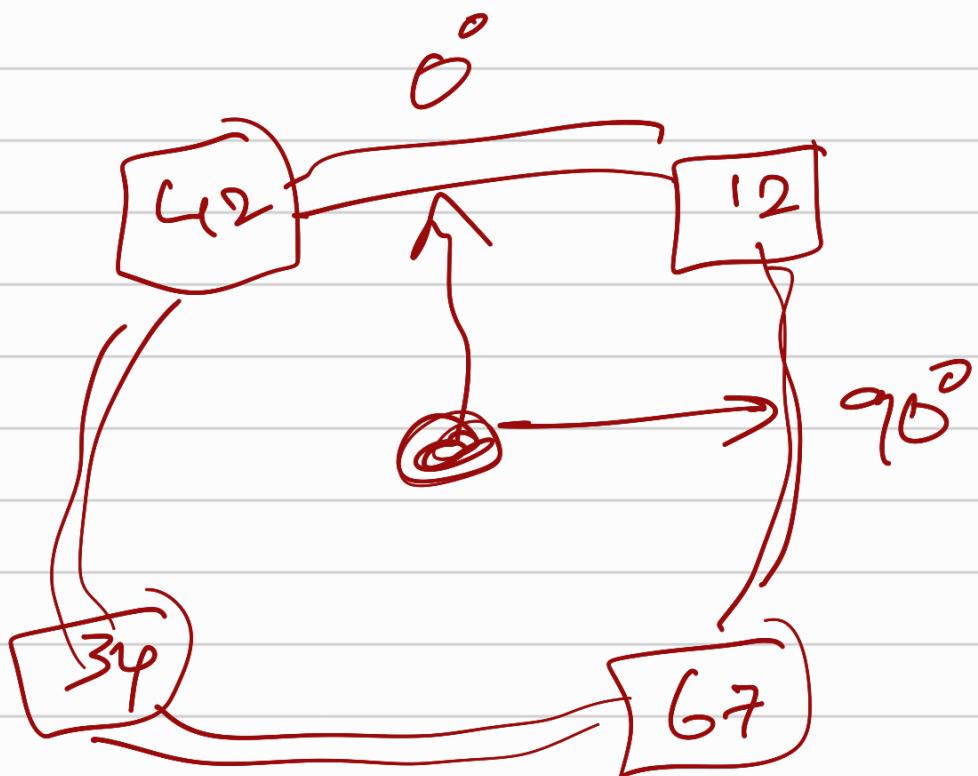
Angles on Circle

[0 - 360]

Ankit, Neha, Ankur, Ruhi

12, 42, 34, 67

Keys.



S0

S1

S2

S3

Ankit

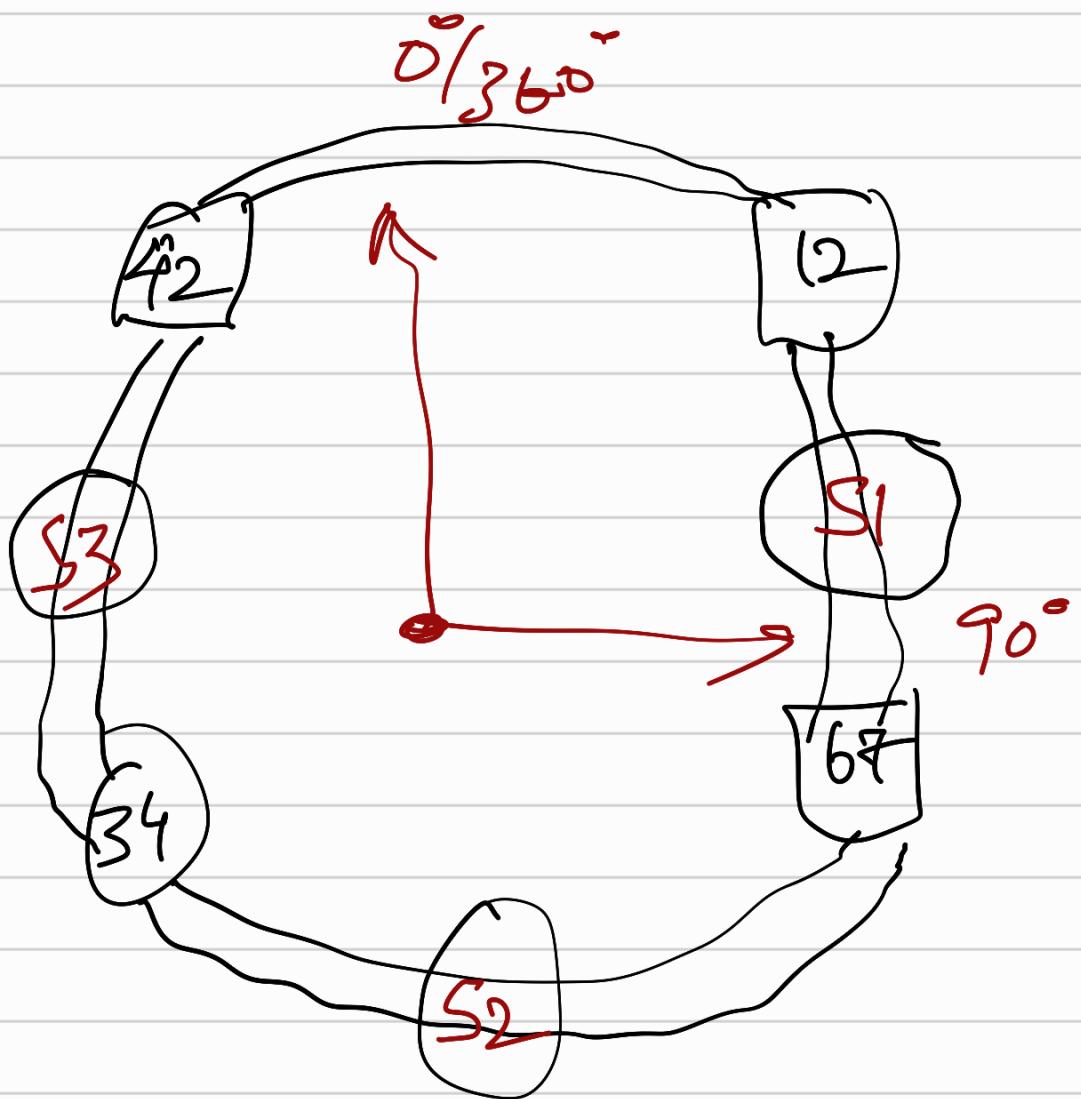
Neha

Ankur

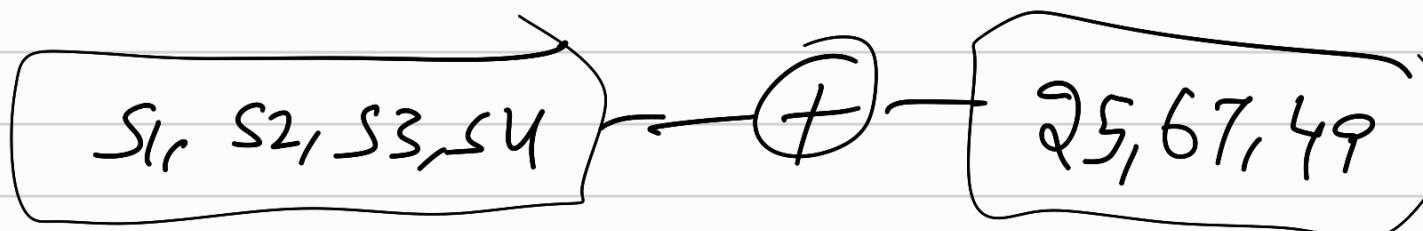
Ruhi

Similar hash for Server
each

Ex: based on IP add
of machine



Hash values of servers will also
going to sit on circle .



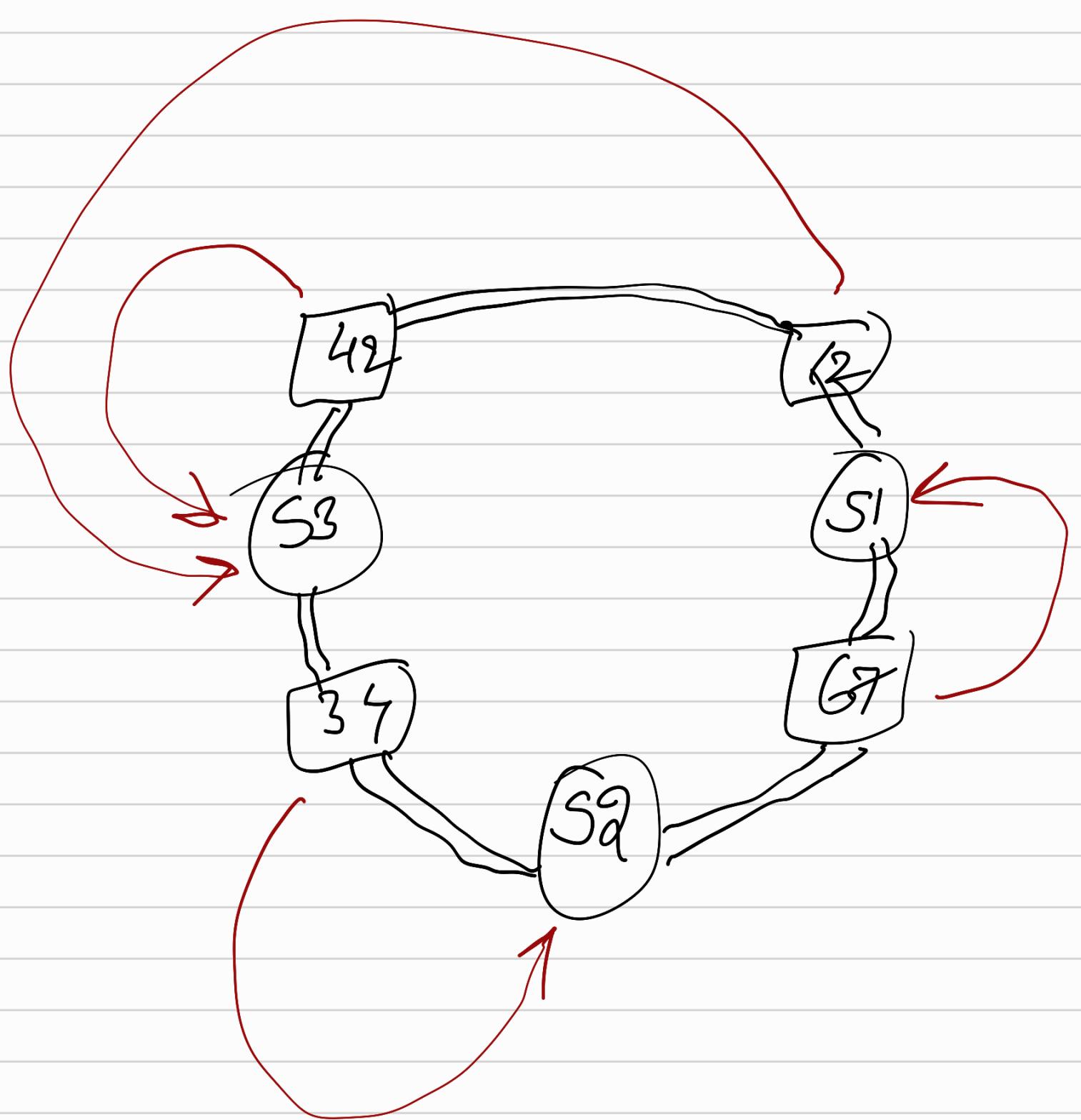
Here the servers and IP value
of one share the same IP range.
on circle.

Server Hash } share
Data Hash } some of range

Mapping of values to servers

clock
anti clock.

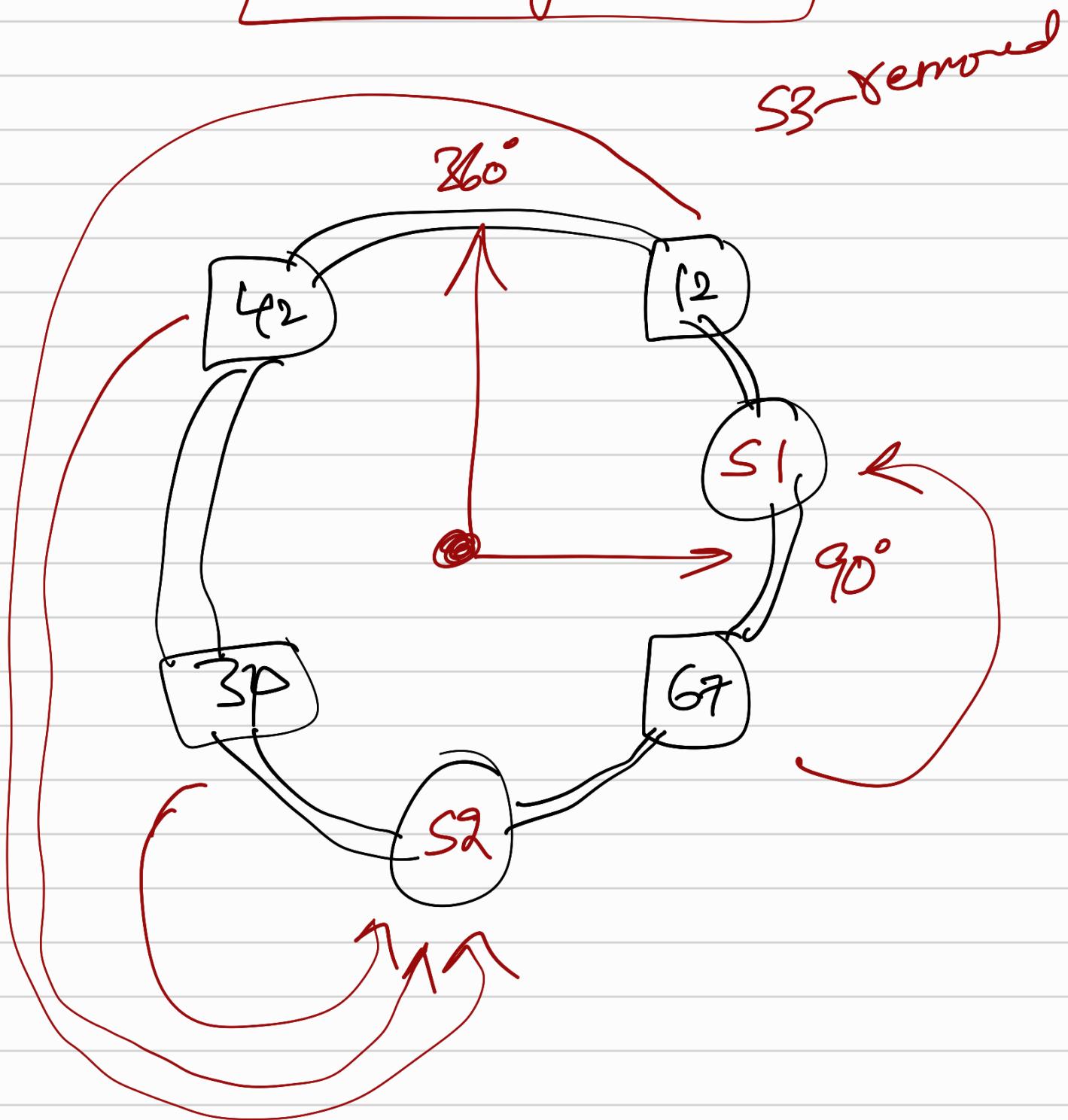
Every value will sit on next server
in clockwise direction
can be anti clockwise direction.

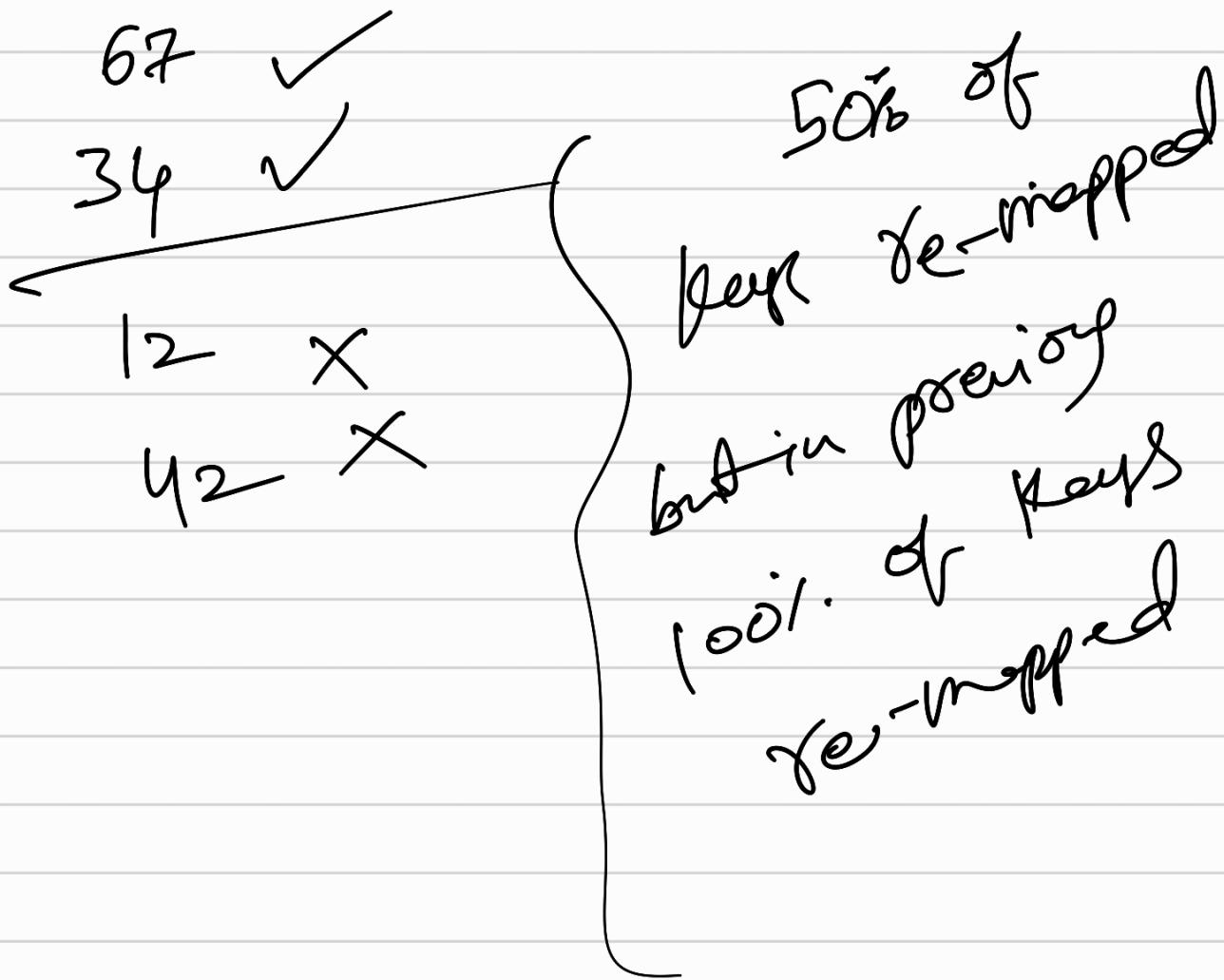


for every value we are trying the following
 the closest value is clockwise /
 Anticlockwise
 direction -

how to map key values to drift
scores.

Removing Scores





✓ less no of keys remapped

✓ less computation.

Once S3 removed → suddenly

12, 42, 34 mapped to S2

S2 → overloaded with reg's.

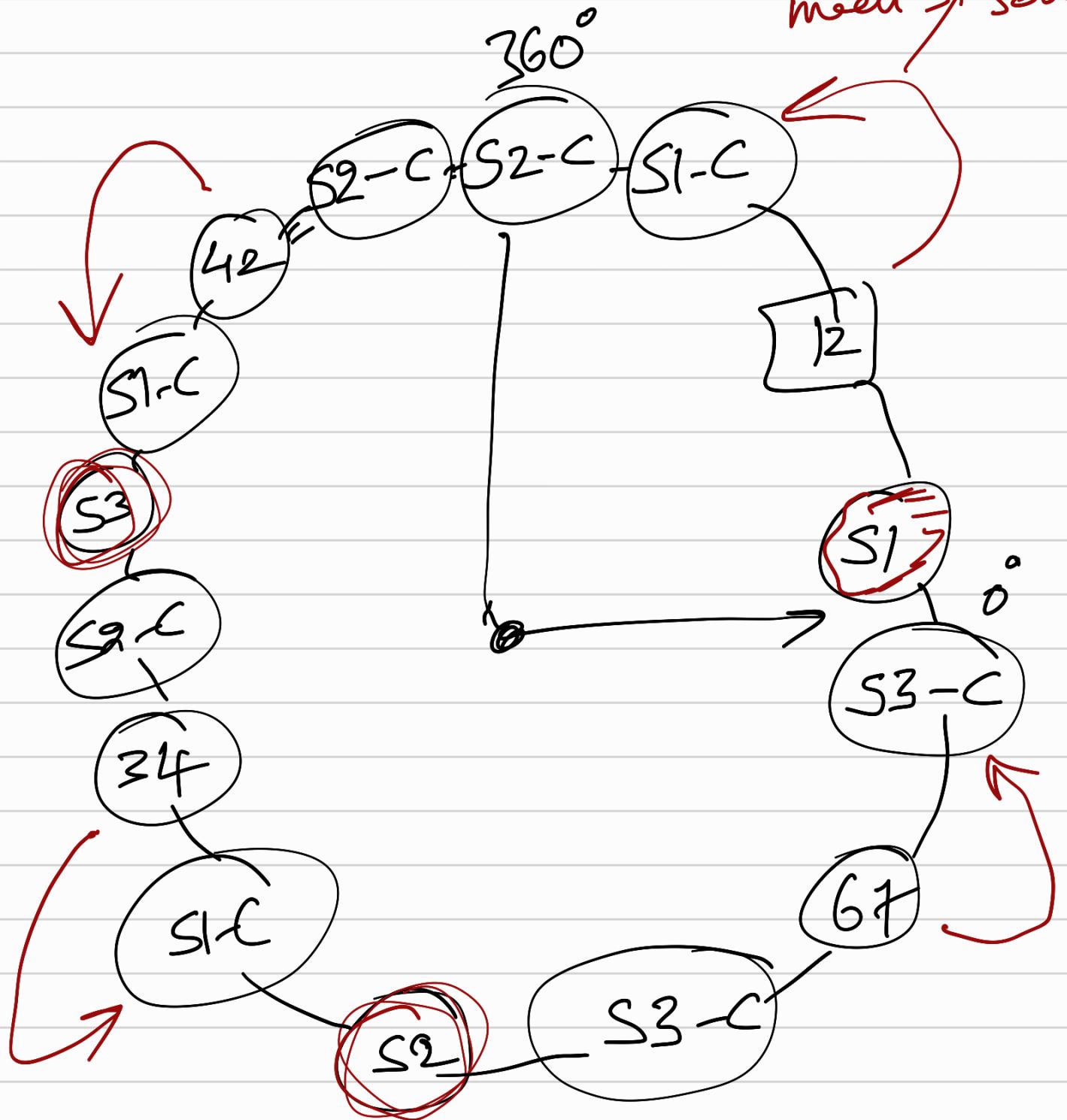
S1 - has less keys

S2 - has more more keys



create

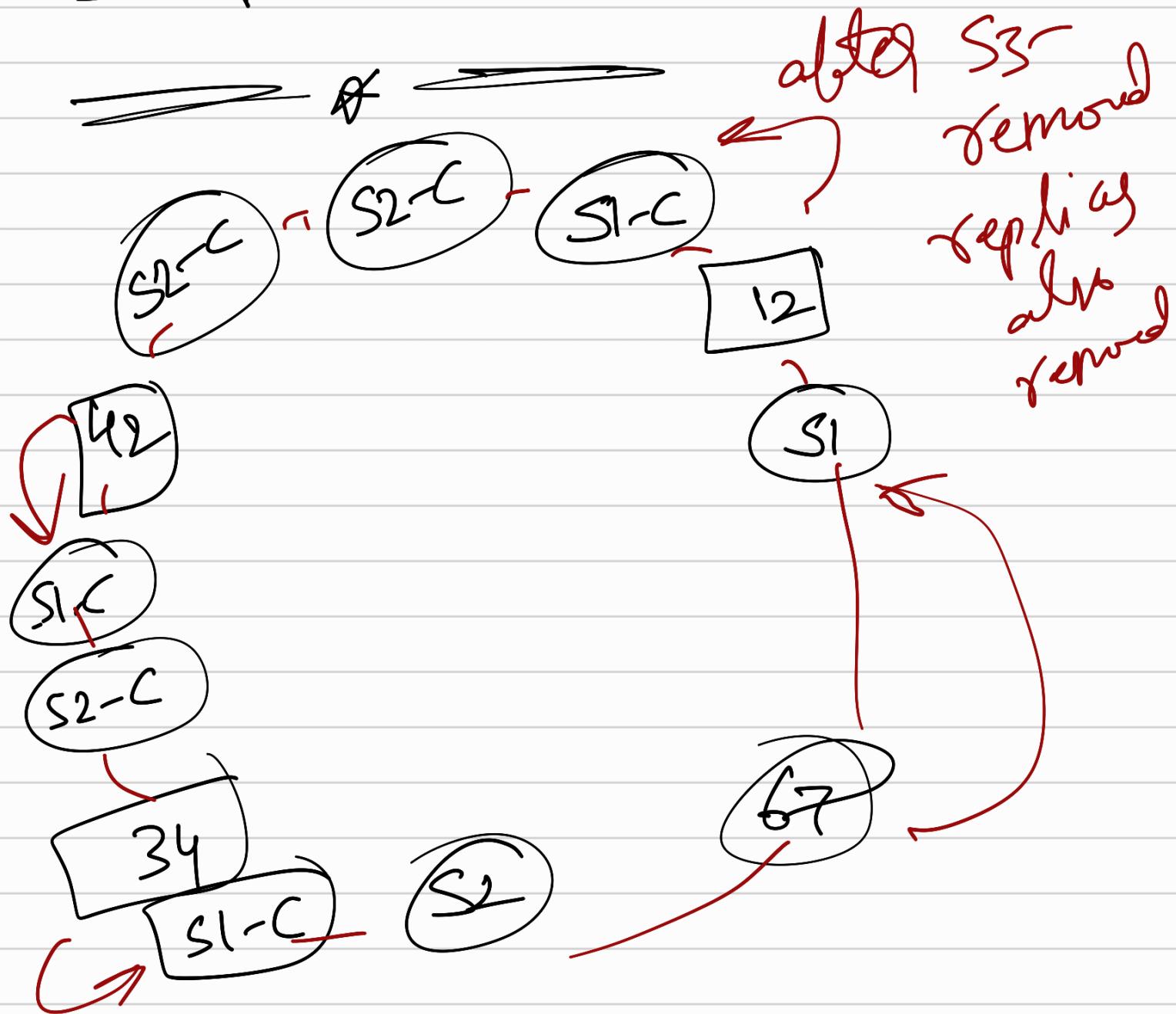
To overcome → Replica of
meeting servers



S_1 replicates $\rightarrow S_1-C$

S_2 replicates $\rightarrow S_2-C$

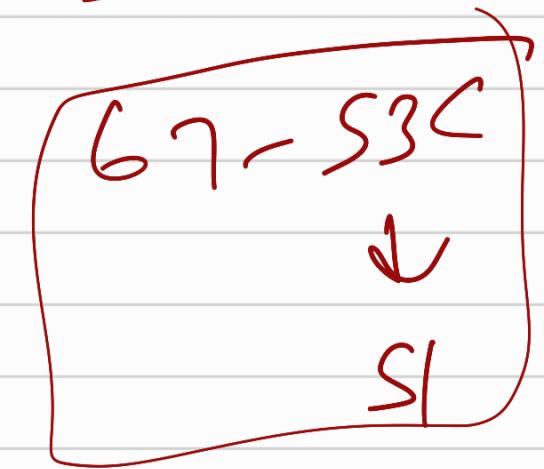
S_3 replicates $\rightarrow S_3-C$



12 - S_1-C ✓

42 - S_1-C ✓

34 - SC



only 1-Key have to be
re-mapped.



With replicas — 50% keys
re-mapped

With replicas — 25% keys
re-mapped.

