

Speed sub little low.

→ what barriers?

The Simply Typed Lambda Calc (STLC)

STLC = constants + sums + products + exponentials

↓
current
lang

↓
Gher
in Haskell

↓
Corresponds
to tuples

↓
functions

Products

$(\text{"Hello"}, 2) :: (\text{String}, \text{Int})$

$() :: ()$

syntax chart

types $T ::= \dots$
 $T_1 \times T_2$ -- product
 \perp -- unit

pre-term $e ::= \dots$
 $\langle e_1, e_2 \rangle$ -- pair
 $\pi_1(e)$ -- first proj.
 $\pi_2(e)$ -- second proj.
 $\langle \rangle$

- introduction rules create instances of a type

- elimination rules destruct / project from instances of a type

Statics:

UNIT $\frac{}{\Gamma \vdash \langle \rangle : \perp}$

PAIR (PROD) $\frac{\Gamma \vdash e_1 : T_1 \quad \Gamma \vdash e_2 : T_2}{\Gamma \vdash \langle e_1, e_2 \rangle : T_1 \times T_2}$

PROJ-1 $\frac{\Gamma \vdash e : T_1 \times T_2}{\Gamma \vdash \pi_1(e) : T_1}$

PROJ-2 $\frac{\Gamma \vdash e : T_1 \times T_2}{\Gamma \vdash \pi_2(e) : T_2}$

Dynamics:

VAL-UNIT $\frac{}{\langle \rangle \text{ val}}$

VAL-PROD $\frac{\langle e_1 \text{ val } e_2 \text{ val} \rangle}{\langle e_1, e_2 \rangle \text{ val}}$

$\pi_1(\langle e_1, e_2 \rangle)$
 \downarrow
 e_1

This is an axiom cos
our language is lazy.

With the dynamics we
will specify that we must
fully eval a pair before
we can proj from it

$\frac{}{\pi_1(\langle e_1, e_2 \rangle) \mapsto e_1}$ D-PROD-TUPLE-1

$\frac{}{\pi_2(\langle e_1, e_2 \rangle) \mapsto e_2}$ D-PROD-TUPLE-2

$\frac{e \mapsto e'}{\pi_1(e) \mapsto \pi_1(e')}$ D-PROD-1

$\frac{e \mapsto e'}{\pi_2(e) \mapsto \pi_2(e')}$ D-PROD-2

Sums

> data Either a b = Left a
| Right b

> data Void

> abort :: Void \rightarrow ()
> abort = ()

> abort :: Void \rightarrow Bool
> abort = True

types $T ::= \dots$
 $T_1 + T_2$

pre-terms $e ::=$
 $\text{abort}(e)$
 $\text{inl}(e)$
 $\text{inr}(e)$
 $\text{case}(e; \underline{x}.e_1; \underline{y}.e_2)$

$\text{fromEither} :: \text{Either } a \ b \rightarrow \text{Maybe } a$
 $\text{fromEither}(\text{Left } a) = \text{Just } a$
 $\text{fromEither}(\text{Right } b) = \text{Nothing}$

$\text{fromEither } e = \text{case } e \text{ of}$

$(\text{Left } x) \rightarrow e_1$

$(\text{Right } y) \rightarrow e_2$