

# Curry - Howard



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# Curry - Howard

Haskell



Brook



Curry



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Correspondence



# Curry - Howard - Lambek

The holy trinity



STLC

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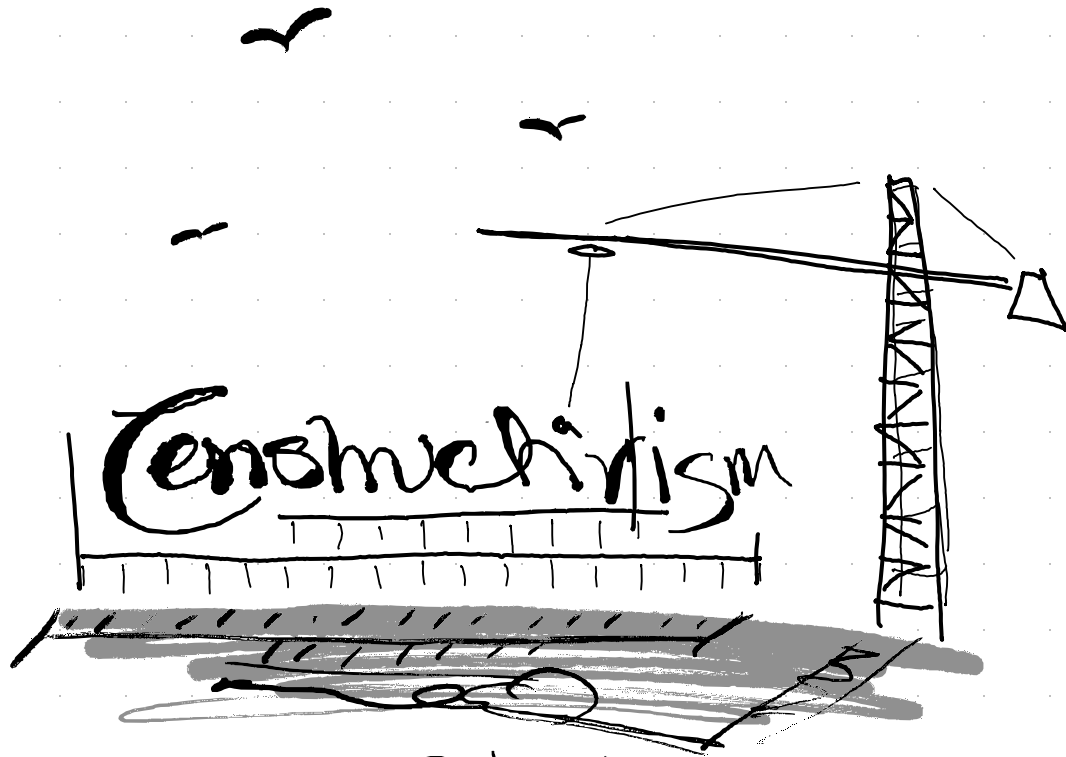


Intuitionistic  
Logic

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CCCs



Sam Frohlich

# Agenda

- What - is constructive
- Dependent types
- Curry-Howard
- Lean (using Curry Howard)
- HoTT



connecting STLC to Constructive logic

Construct	Evidence
false	<u>          </u>
true	<u>          </u>
$A \wedge B$	$eA$ and $eB$
$A \vee B$	$eA$ or $eB$
$A \Rightarrow B$	$eA \rightarrow eB$
$\neg A$	$eA \rightarrow \text{false}$

Not constructive

LEM :  $\neg A \vee A$   
 DNE :  $\neg \neg A \rightarrow A$   
 Peirce :  $((A \Rightarrow B) \Rightarrow A) \Rightarrow A$

$((A \Rightarrow B) \Rightarrow A) \Rightarrow A$        $((a \rightarrow b) \rightarrow a)$

1. Assume  $((A \Rightarrow B) \Rightarrow A)$
2. Proof by contra, ASS  $\neg A$
3. Subproof showing  $A \Rightarrow B$
4.     | ASS A
5.     |  $\perp$       -- 4 contradicts 2
6.     | B.
7. A      -- 1 @ 3
8.  $\perp$

## Dependent types

values at the type level

```
> head :: [a] -> a
> head (x:xs) = x
```

```
ghci> head []
```

```
> data Nat = Z | SNat
```

```
> two = S(S Z)
```

```
> data Vec (n:Nat) a where
```

```
>   VNil :: Vec 'Z a
```

```
>   Vcons :: a
```

```
>         -> Vec n a
```

```
>         -> Vec ('S n) a
```

+  
Plus  
:

```
> safeHead :: Vec ('S n) a -> a
```

```
> safeHead (Vcons x xs) = x
```

Construct	Evidence Meaning
$\exists x \in X. A$	$e \in A \xrightarrow{y \in X} e \in A[y/x]$
$\forall x \in X. A$	$(y, e \xrightarrow{y \in X} e \in A)$

Exists - Dependent pair

$:: (n:Nat, Vec n Bool)$

$(0, VNil)$

$(1, Vcons \perp VNil)$

$\sum_{(n:Nat)} Vec n Bool$

forall - Dependant Function

$:: n:Nat \rightarrow Vec n Bool$

```
> data SNat :: Nat -> * where
```

```
>   SZ :: SNat 'Z
```

```
>   SS :: SNat n -> SNat ('S n)
```

```
> replicate :: SNat n -> Vec n Bool
```

```
> replicate SZ = VNil
```

```
> replicate (SS n) = Vcons \ (replicate n)
```