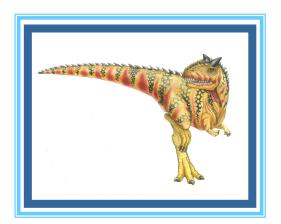
Chapter 3: Processes





Process Concept

- An operating system executes a variety of programs that run as a process.
- Process a program in execution; process execution must progress in sequential fashion. No parallel execution of instructions of a single process
- Multiple parts
 - The program code, also called text section
 - Current activity including program counter, processor registers
 - Stack containing temporary data
 - 4 Function parameters, return addresses, local variables
 - Data section containing global variables
 - Heap containing memory dynamically allocated during run time





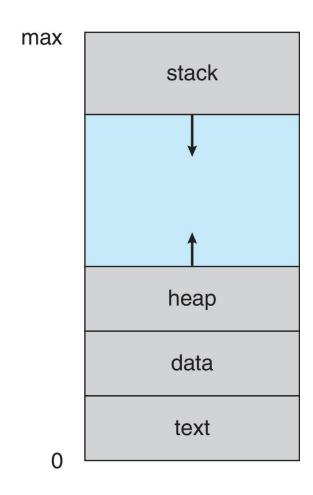
Process Concept (Cont.)

- Program is passive entity stored on disk (executable file);
 process is active
 - Program becomes process when an executable file is loaded into memory
- Execution of program started via GUI mouse clicks, command line entry of its name, etc.
- One program can be several processes
 - Consider multiple users executing the same program





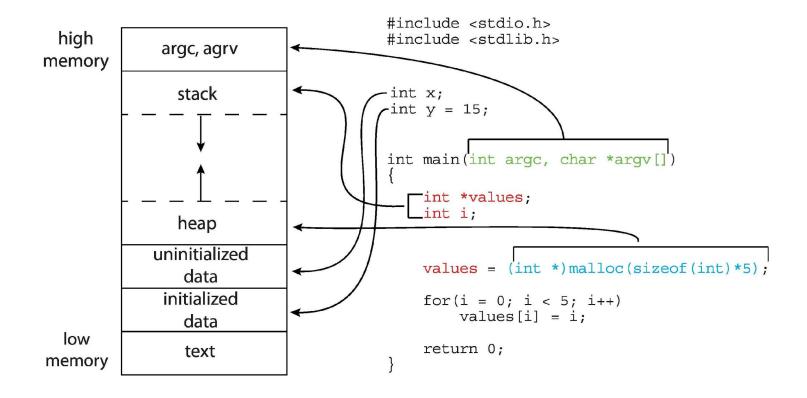
Process in Memory







Memory Layout of a C Program







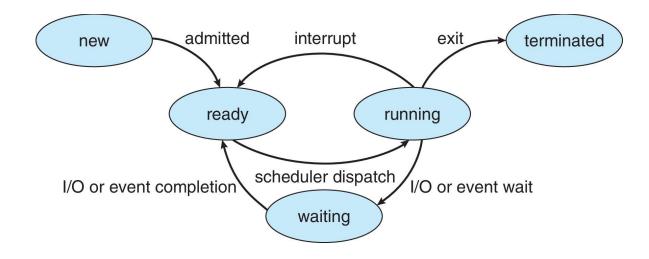
Process State

- As a process executes, it changes state
 - New: The process is being created
 - Running: Instructions are being executed
 - Waiting: The process is waiting for some event to occur
 - **Ready**: The process is waiting to be assigned to a processor
 - Terminated: The process has finished execution





Diagram of Process State







Process Control Block (PCB)

Information associated with each process(also called **task control block**)

- Process state running, waiting, etc.
- Program counter location of instruction to next execute
- CPU registers contents of all process-centric registers
- CPU scheduling information- priorities, scheduling queue pointers
- Memory-management information memory allocated to the process
- Accounting information CPU used, clock time elapsed since start, time limits
- I/O status information I/O devices allocated to process, list of open files

process state
process number
program counter
registers
memory limits
list of open files





Threads

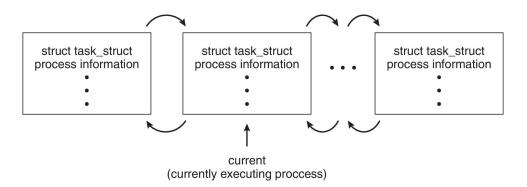
- So far, process has a single thread of execution
- Consider having multiple program counters per process
 - Multiple locations can execute at once
 - 4 Multiple threads of control -> threads
- Must then have storage for thread details, multiple program counters in PCB
- Explore in detail in Chapter 4



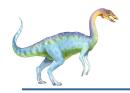


Process Representation in Linux

Represented by the C structure task struct







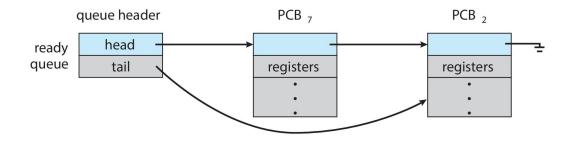
Process Scheduling

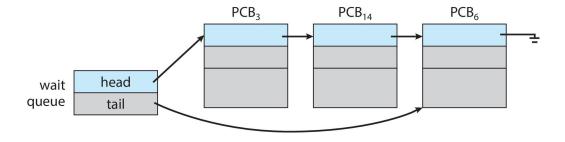
- Process scheduler selects among available processes for next execution on CPU core
- Goal -- Maximize CPU use, quickly switch processes onto CPU core
- Maintains scheduling queues of processes
 - Ready queue set of all processes residing in main memory, ready and waiting to execute
 - Wait queues set of processes waiting for an event (i.e., I/O)
 - Processes migrate among the various queues





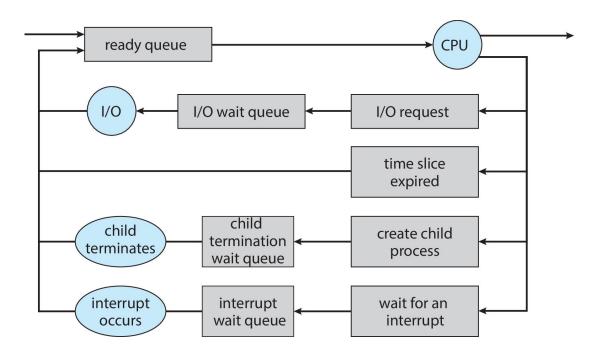
Ready and Wait Queues







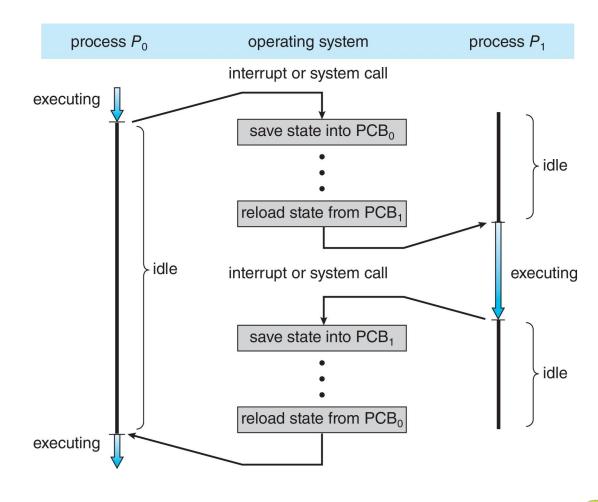
Representation of Process Scheduling





CPU Switch From Process to Process

A **context switch** occurs when the CPU switches from one process to another.





Context Switch

- When CPU switches to another process, the system must save the state of the old process and load the saved state for the new process via a context switch
- Context of a process represented in the PCB
- Context-switch time is pure overhead; the system does no useful work while switching
 - The more complex the OS and the PCB

 □ the longer the context switch
- Time dependent on hardware support
 - Some hardware provides multiple sets of registers per CPU

 multiple contexts loaded at once

