On the Performance Prediction of BLAS-based Tensor Contractions

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November 16th 2014 — PMBS14





Tensor Contractions?

$$C := -A_i - B_i$$

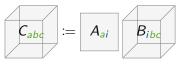
$$C \coloneqq \sum_{\mathbf{i}} A[\mathbf{i}]B[\mathbf{i}]$$

$$C_a := A_{ai} B_i$$

$$\forall \mathtt{a}.C[\mathtt{a}] \coloneqq \sum_{\mathtt{i}} A[\mathtt{a},\mathtt{i}]B[\mathtt{i}]$$

$$C_{ab} := A_{ai} B_{ib}$$

$$C_{ab} := A_{ai} \mid B_{ib} \quad \forall a, b, C[a,b] := \sum_{i} A[a,i]B[i,b]$$



$$C_{abc} := A_{ai}$$
 B_{ibc} $\forall a, b, c. C[a,b,c] := \sum_{i} A[a,i]B[i,b,c]$

$$C_{abc} := A_{ija}B_{jbic}$$

 $\forall a, b, c. C[a, b, c]$

$$:= \sum_{i,j} A[i,j,a]B[j,b,i,c]$$

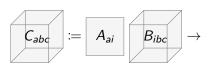
free indices

contracted indices

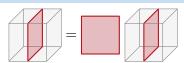
Use BLAS!

$$C_{ab}$$
 := A_{ai} B_{ib} \rightarrow

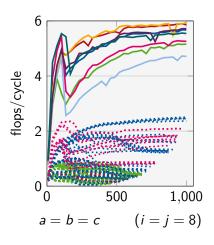
$$C_{ab} := A_{ai}B_{ib}$$
 b-gemv



$C_{abc} := A_{ai}B_{ibc}$ b-gemm



$C_{abc} := A_{aij}B_{jbic}$



Total: 176 Algorithms!

Goals

- Generate algorithms
- Predict their performance

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Outline

- Algorithm Generation
- 2 Performance Prediction
 - 1. Repeated Execution
 - 2. Cache Setup
 - 3. Prefetching
 - 4. Prefetching Failures
 - 5. First Iterations
- Results

$$C_a := A_{iaj}B_{ji}$$
 $C_{abc} := A_{aij}B_{jbic}$
Multithreading

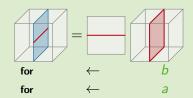
Algorithm Generation

- Select kernel
- Match kernel indices to tensor indices

- Cast remaining tensor indices as for-loops
- Assemble algorithm (AST, C-code)

Example

$$\begin{array}{cccc}
C_{\alpha} := A_{\alpha \iota} B_{\iota} & \to & C_{abc} = A_{ai} B_{ibc} \\
\iota & \to & i \\
\alpha & \to & c
\end{array}$$



```
C_{abc} := A_{ai}B_{ibc} ba-gemv

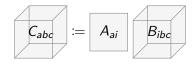
for b = 1:b

for a = 1:a

C[a,b,:] = A[a,:] B[:,b,:]
```

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Algorithms for $C_{abc} := A_{ai}B_{ibc}$



- BLAS-1
 - 6 dot-based: $(C := A_t B_t)$ abc-dot acb-dot bac-dot bca-dot cab-dot cba-dot • 18 axpy-based: $(C_{\alpha} := AB_{\alpha})$ ibc-axpy icb-axpy bic-axpy bci-axpy cib-axpy cbi-axpy iac-axpy ica-axpy aci-axpy cia-axpy cai-axpy iab-axpy iba-axpy aib-axpy abi-axpy bia-axpv bai-axpv
- BLAS-2
 - 6 gemv-based: $(C_{\alpha} := A_{\alpha \iota} B_{\iota})$ bc-gemv cb-gemv ac-gemv ca-gemv ab-gemv ba-gemv • 4 ger-based: $(C_{\alpha\beta} := A_{\alpha}B_{\beta})$ ic-ger ci-ger ib-ger bi-ger
- BLAS-3
 - 2 gemm-based: $(C_{\alpha\beta} := A_{\alpha \iota} B_{\iota\beta})$ c-gemm b-gemm

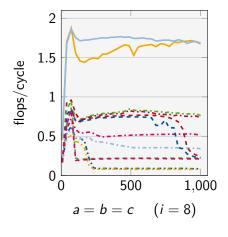
aic-axpy

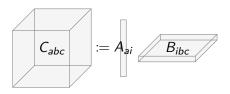
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```
C_a := A_{iaj}B_{ji}
C_{abc} := A_{aij}B_{jbic}
Multithreading
```

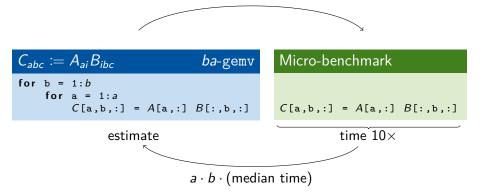
What are we predicting?





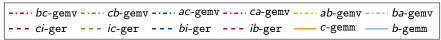
- Intel Penryn E5450 (Harpertown)
- Single-threaded OPENBLAS

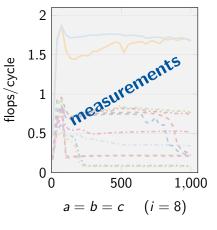
Micro-Benchmarks

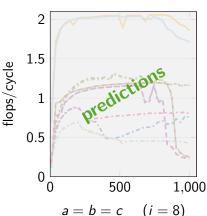


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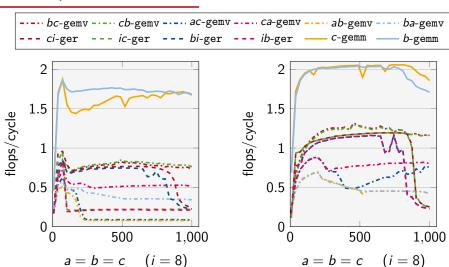
1. Repeated Execution







1. Repeated Execution



Problem: general overestimation

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2. Caching!

Problem: general overestimation

Cause: Cache locality not accounted for (micro-benchmark works in-cache)

Solution Approach

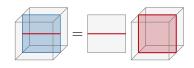
Recreate cache precondition within micro-benchmark

Assumption: Fully associative Least Recently Used (LRU) cache replacement policy

 \Rightarrow Cache state is defined by the order of memory accesses.

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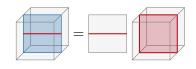
$C_{abc} := A_{ai}B_{ibc}$ ca-gemv for c = 1:cfor a = 1:aC[a,:,c] = A[a,:] B[:,:,c]



Access Distance d(M) = size(all data used since the last access to M)

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$C_{abc} := A_{ai}B_{ibc}$ ca-gemv for c = 1:cfor a = 1:aC[a,:,c] = A[a,:] B[:,:,c]

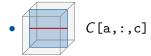


Access Distance d(M) = size(all data used since the last access to M)

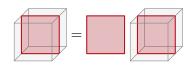


$$B[:,:,c]$$
 doesn't vary in for a $d(B[:,:,c]) = 0$ doubles





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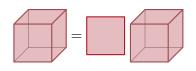
$$B[:,:,c]$$
 doesn't vary in for a $d(B[:,:,c]) = 0$ doubles



$$A[a,:]$$
 varies in for a; doesn't vary in for c
 $d(A[a,:]) = \text{size}(all \ operands \ in for a)$
 $= \text{size}(C[:,:,c]) + \text{size}(A[:,:]) + \text{size}(B[:,:,c])$
 $= a \cdot b + a \cdot i + i \cdot b$



$C_{abc} := A_{ai}B_{ibc}$ ca-gemv for c = 1:c for a = 1:a C[a,:,c] = A[a,:] B[:,:,c]



Access Distance d(M) = size(all data used since the last access to M)



$$B[:,:,c]$$
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$$C[a,:,c]$$
 varies in for a; varies in for c
 $d(C[a,:,c]) = size(all operands in for c)$
 $= size(C[:,:,:]) + size(A[:,:]) + size(B[:,:,:])$
 $= a \cdot b \cdot c + a \cdot i + i \cdot b \cdot c$

Setup

Access Distances d(B[:,:,c]) = 0 = 0 doubles $d(A[a,:]) = a \cdot b + a \cdot i + i \cdot b$ = 166,400 doubles $d(C[a,:,c]) = a \cdot b \cdot c + a \cdot i + i \cdot b \cdot c$ = 65,283,200 doubles

Sizes: a = b = c = 400, i = 8.

```
C_{abc} := A_{ai}B_{ibc} ca-gemv Micro-benchmark

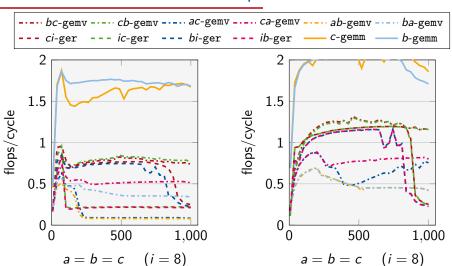
for c = 1:c
for a = 1:a
C[a,:,c] = A[a,:] B[:,:,c]
flush 816,632
touch A[a,:]
flush 163,200
touch B[:,:,c]
```

Limit at $\frac{5}{4}$ size(cache) = $\frac{5}{4} \cdot 6MB = 983,040$ doubles

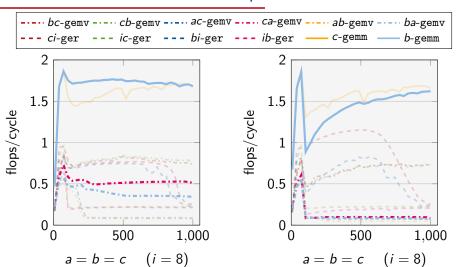
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C[a,:,c] = A[a,:] B[:,:,c]

2. Estimates with Cache Setup



2. Estimates with Cache Setup



Problem: underestimation

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3. Prefetching!

Problem: selective underestimation

Cause: Prefetching not accounted for

Solution Approach

Mimick prefetching

Prefetching:

access

prefetch





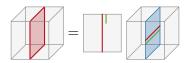






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Mimicking the Prefetching



$C_{abc} := A_{ai}B_{ibc}$

$A_{ai}B_{ibc}$ bi-ger

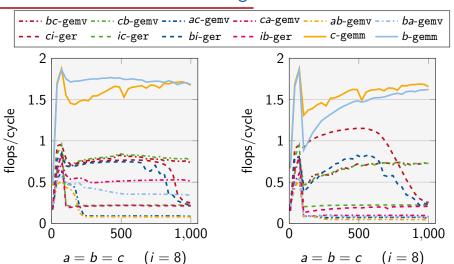
```
for b = 1:b
  for i = 1:i
        C[:,b,:] = A[:,i] B[i,b,:]
```

Micro-benchmark

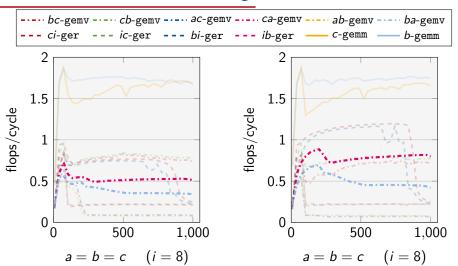
```
touch A[:,i]
flush 5,992
touch A[:8,i]
touch B[i,b,:]
touch C[:,b,:]
C[:,b,:] = A[:,i] B[i,b,:]
```

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3. Estimates with Prefetching



3. Estimates with Prefetching



Problem: some incorrect prefetching

Elmar Peise 18b

4. Prefetching Failures

Problem: selective prefetching failure

Cause: No Prefetching along 1st dimension across cache-lines. (every 8th iteration is not prefetched)

Solution Approach

Separate micro-benchmarks with and without prefetching

Micro-benchmark (pre)

```
touch A[:,i]
flush 5,992
touch A[:8,i]
touch B[i,b,:]
touch C[:,b,:]
C[:,b,:] = A[:,i] B[i,b,:]
```

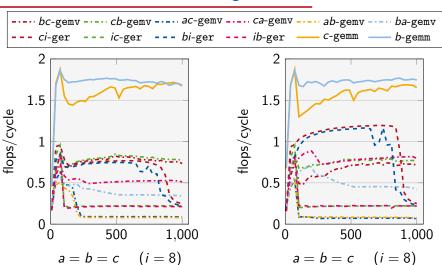
Micro-benchmark (no pre)

```
flush 816,240
touch A[:,i]
flush 5,992
touch A[:8,i]
touch C[:,b,:]
C[:,b,:] = A[:,i] B[i,b,:]
```

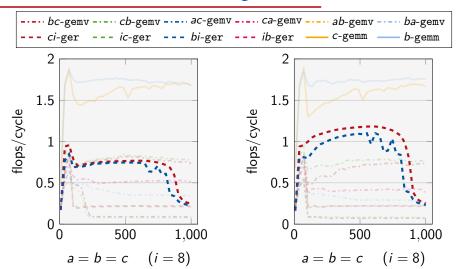
time $10\times$ time $10\times$

 $estimate = \frac{1}{8}(7median + 1median)$

4. Estimates with Prefetching Failures



4. Estimates with Prefetching Failures



Problem: selective overestimation

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5. Small Loops' First Iterations

Problem: selective overestimation

Cause: Innermost loop dimension too small (first iteration of innermost loop differs)

5. Small Loops' First Iterations

Problem: selective overestimation

Cause: Innermost loop dimension too small (first iteration of innermost loop differs)

```
bi-ger
C_{abc} := A_{ai}B_{ibc}
for b = 1:b
    for i = 1:i
         C[:,b,:] = A[:,i] B[i,b,:]
```



Solution Approach

Separate micro-benchmarks for first iteration of small loops

First Iteration Benchmark

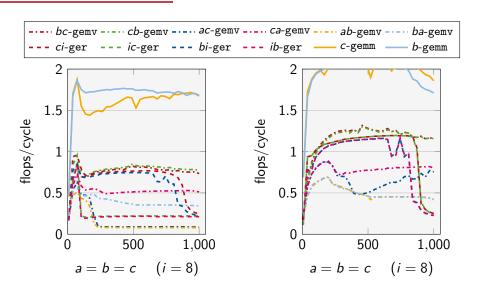
Access Distance:

Find last access to A[:,i], B[i,b,:], and C[:,b,:] within for if Find last access to A[:,:], B[:,b,:], and C[:,b,:] within for c

• Prefetching:



5. Final Estimates



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```
C_a \coloneqq A_{iaj}B_{ji}
C_{abc} \coloneqq A_{aij}B_{jbic}
Multithreading
Efficiency
```

$C_a := A_{iaj}B_{ji}$ — Only BLAS-1 and BLAS-2



8 algorithms:

- 4 dot-based:
 aj-dot ja-dot ai-dot ia-dot
- 2 gemv-based:

$$C_a := A_{iaj}B_{ji} \qquad \qquad j\text{-gemv}$$

$$for j = 1:j \qquad \qquad C[:] += A[:,:,j] B[j,:]$$

$$C_a := A_{iaj}B_{ji} \qquad i'\text{-gemv}$$

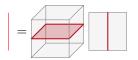
$$for i = 1:i$$

$$\widetilde{A}[:,:] = A[i,:,:]$$

$$C[:] += \widetilde{A}[:,:] B[:,i]$$

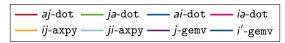
• $a = i = j = 8 \dots 1,000$

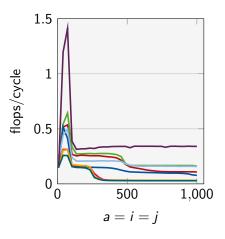
2 axpy-based:ij-axpy ji-axpy

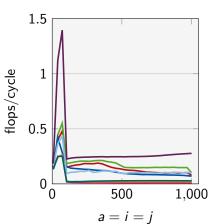


$C_a := A_{iaj}B_{ji}$ — Only BLAS-1 and BLAS-2

Results







$C_{abc} := A_{aij}B_{jbic}$ — Challenging Contraction

- $C_{abc} := A_{aij} B_{jbic}$ 176 Algorithms:
 - 48 dot-based
 - 72 axpy-based
 - 36 gemv-based
 - 12 ger-based
 - 8 gemm-based:

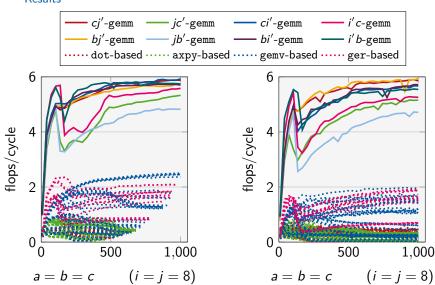
```
cj'-gemm jc'-gemm ci'-gemm i'c-gemm bj'-gemm jb'-gemm bi'-gemm i'b-gemm
```

- i = j = 8, $a = b = c = 8 \dots 1{,}000$
- Intel Ivy Bridge E5-2680 v2
- Single-threaded OPENBLAS

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$C_{abc} := A_{aij}B_{jbic}$ — Challenging Contraction

Results

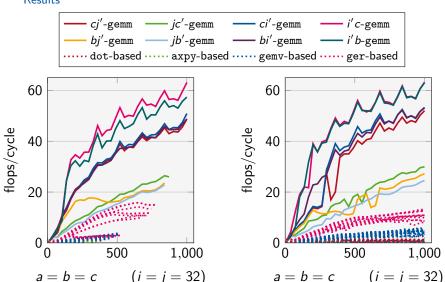


$C_{abc} := A_{aij}B_{jbic} - 10$ Threads

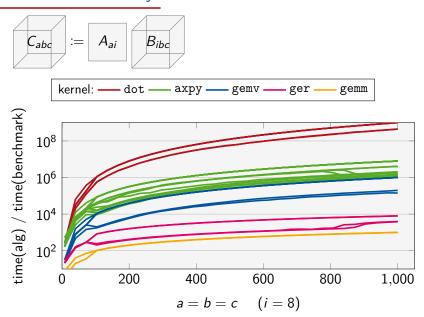
- $C_{abc} := A_{aij}B_{jbic}$
- i = j = 32, $a = b = c = 8 \dots 1,000$
- Intel Ivy Bridge E5-2680 v2
- OPENBLAS, 10 threads

$C_{abc} := A_{aij}B_{jbic} - 10$ Threads

Results



Prediction Efficiency



On the Performance Prediction of BLAS-based Tensor Contractions

- BLAS-based algorithm generation
- Micro-benchmarks with careful cache setup
- Applicable to wide range of challenging scenarios

Funding from Deutsche Forschungsgemeinschaft and Deutsche Telekom Stiftung is gratefully acknowledged.

Deutsche Forschungsgemeinschaft **DFG** Deutsche Telekom **Stiftung**

 \mathbf{T} ..

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