

QUANTUM ALGORITHMS
HOMEWORK 11 SELECTED SOLUTIONS

PROF. MATTHEW MOORE

1. Let $t \in \mathbb{N}$.

(i) Prove that

$$x^t - 1 = (x - 1) \sum_{k=0}^{t-1} x^k.$$

(ii) Prove that $x = e^{2\pi i(m/t)}$ is a solution to $x^t - 1$ for $m \in \mathbb{Z}$.

(iii) Let $m \in \mathbb{Z}$ with $0 \leq m < t$. Use the previous parts to prove that

$$\sum_{k=0}^{t-1} e^{2\pi i(km/t)} = \begin{cases} t & \text{if } m = 0, \\ 0 & \text{otherwise.} \end{cases}$$

Solution:

(i): *Proof.* We have

$$\begin{aligned} x^t - 1 &= x^t + (x^{t-1} - x^{t-1}) + (x^{t-2} - x^{t-2}) + \cdots + (x - x) - 1 \\ &= (x^t - x^{t-1}) + (x^{t-1} - x^{t-2}) + \cdots + (x^2 - x) + (x - 1) \\ &= (x - 1)x^{t-1} + (x - 1)x^{t-2} + \cdots + (x - 1)x + (x - 1) \\ &= (x - 1)(x^{t-1} + x^{t-2} + \cdots + x + 1) = (x - 1) \sum_{k=0}^{t-1} x^k, \end{aligned}$$

as claimed. □

(ii): *Proof.* Recall that $e^{2\pi i} = 1$. We have

$$\begin{aligned} (e^{2\pi i(m/t)})^t - 1 &= e^{2\pi i(mt/t)} - 1 = e^{2\pi i m} - 1 = (e^{2\pi i})^m - 1 = (1)^m - 1 \\ &= 1 - 1 = 0. \end{aligned}$$

Therefore $e^{2\pi i(m/t)}$ is a root of $x^t - 1$. □

(iii): *Proof.* Let $z = e^{2\pi i(m/t)}$. We have

$$0 = z^t - 1 = (z - 1) \sum_{k=0}^{t-1} z^k,$$

so $z = 1$ or $\sum z^k = 0$. If $m \neq 0$ then $z = e^{2\pi i(m/t)} \neq 1$, so it must be that $\sum z^k = 0$. If $m = 0$ then $z = e^{2\pi i(m/t)} = 1$, so $\sum z^k = t$. Hence

$$\sum_{k=0}^{t-1} e^{2\pi i(km/t)} = \sum_{k=0}^{t-1} z^k = \begin{cases} t & \text{if } m = 0, \\ 0 & \text{otherwise.} \end{cases} \quad \square$$

Date: April 29, 2021.

2. (i) Explicitly calculate $\mathcal{QFT}_n |0^n\rangle$.

(ii) Explicitly calculate $\mathcal{QFT}_n |1^n\rangle$.

Solution: We have

$$\begin{aligned}
 \mathcal{QFT}_n |0^n\rangle &= 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + e^{2\pi i(0)/2^k} |1\rangle = 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + |1\rangle = 2^{-n/2} (|0\rangle + |1\rangle)^{\otimes n} \\
 &= 2^{-n/2} \sum_{x \in \{0,1\}^n} |x\rangle, \\
 \mathcal{QFT}_n |1^n\rangle &= 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + e^{2\pi i[1^n]/2^k} |1\rangle = 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + e^{2\pi i(2^n-1)/2^k} |1\rangle \\
 &= 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + e^{-2\pi i/2^k} |1\rangle = 2^{-n/2} \sum_{x_1, \dots, x_n \in \{0,1\}} \left(\prod_{k=1}^n e^{-2\pi i x_k / 2^k} \right) |x_1 \dots x_n\rangle \\
 &= 2^{-n/2} \sum_{x_1, \dots, x_n \in \{0,1\}} e^{-2\pi i \sum x_k / 2^k} |x_1 \dots x_n\rangle \\
 &= 2^{-n/2} \sum_{x \in \{0,1\}^n} e^{-2\pi i [x]/2^n} |x\rangle.
 \end{aligned}$$

3. Show that

$$\mathcal{QFT}_n |x\rangle = 2^{-n/2} \sum_{y \in \{0,1\}^n} e^{2\pi i [x][y]/2^n} |y\rangle,$$

where $[x]$ represents the number with binary representation $x \in \{0,1\}^n$ (and so $[x][y]$ is the product of x and y , regarded as binary numbers).

Solution:

Proof. We have

$$\begin{aligned}
 \mathcal{QFT}_n |x\rangle &= 2^{-n/2} \bigotimes_{k=1}^n |0\rangle + e^{2\pi i [x]/2^k} |1\rangle = 2^{-n/2} \bigotimes_{k=1}^n \sum_{y_k \in \{0,1\}} e^{2\pi i y_k [x]/2^k} |y_k\rangle \\
 &= 2^{-n/2} \sum_{y_1, \dots, y_n \in \{0,1\}} e^{2\pi i [x] \sum y_k / 2^k} |y_1 \dots y_n\rangle = 2^{-n/2} \sum_{y \in \{0,1\}^n} e^{2\pi i [x][y]/2^n} |y\rangle. \square
 \end{aligned}$$

4. Use the previous problem to prove that

$$\mathcal{QFT}_n^\dagger |x\rangle = 2^{-n/2} \sum_{y \in \{0,1\}^n} e^{-2\pi i [x][y]/2^n} |y\rangle$$

for basis vector $|x\rangle \in \{0,1\}^n$ defines the inverse of \mathcal{QFT}_n .

Hint 1: Show that $\mathcal{QFT}_n \circ \mathcal{QFT}_n^\dagger |x\rangle = \mathcal{QFT}_n^\dagger \circ \mathcal{QFT}_n |x\rangle = |x\rangle$.

Hint 2: You may find this identity useful

$$\sum_{k=0}^{2^n-1} e^{2\pi i k \ell / 2^n} = 0 \quad \text{if} \quad \ell \neq 0.$$

Solution:

Proof. We have

$$\begin{aligned}
\mathcal{QFT}_n \circ \mathcal{QFT}_n^\dagger |x\rangle &= \mathcal{QFT}_n \left(2^{-n/2} \sum_{y \in \{0,1\}^n} e^{-2\pi i [x][y]/2^n} |y\rangle \right) \\
&= 2^{-n/2} \sum_{y \in \{0,1\}^n} e^{-2\pi i [x][y]/2^n} \mathcal{QFT}_n |y\rangle \\
&= 2^{-n/2} \sum_{y \in \{0,1\}^n} e^{-2\pi i [x][y]/2^n} \left(2^{-n/2} \sum_{z \in \{0,1\}^n} e^{2\pi i [y][z]/2^n} |z\rangle \right) \\
&= 2^{-n} \sum_{z \in \{0,1\}^n} \left(\sum_{y \in \{0,1\}^n} e^{2\pi i [y]([x]-[z])/2^n} \right) |z\rangle.
\end{aligned}$$

Using the identity in Hint 2, the inner sum simplifies,

$$\sum_{y \in \{0,1\}^n} e^{2\pi i [y]([x]-[z])/2^n} = \begin{cases} 0 & \text{if } x \neq z \\ \sum_{y \in \{0,1\}^n} 1 & \text{if } x = z \end{cases} = \begin{cases} 0 & \text{if } x \neq z, \\ 2^n & \text{if } x = z. \end{cases}$$

Therefore the summand is nonzero only when $z = x$, in which case it is $2^n |x\rangle$. Thus,

$$2^{-n} \sum_{z \in \{0,1\}^n} \left(\sum_{y \in \{0,1\}^n} e^{2\pi i [y]([x]-[z])/2^n} \right) |z\rangle = 2^{-n} 2^n |x\rangle = |x\rangle,$$

as desired. The calculation showing $\mathcal{QFT}_n^\dagger \circ \mathcal{QFT}_n |x\rangle = |x\rangle$ is quite similar. \square