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1 Given:

 V = Personalized pagerank vectors of a set of users.

Now, as defined in lecture (6th April, 29:45 timestamp)
 PPR vector for a user having teleport set S denotes "how close the user is from the teleportation set S ", and gives us vector s_{x_3} representing it.

Suppose pages $p_1, p_2 \& p_3$ have their teleport set vectors $s_1, s_2 \& s_3$ and PPR vectors $x_1, x_2 \& x_3$ respectively

such that $s_3 = \alpha s_1 + \beta s_2$
 Let M be $\alpha AD^{-1} + (1-\alpha) J$ as defined in class

$$\therefore M' x_3 = \lambda M x_3 + (1-\lambda) s_3 = x_3 \quad \text{--- (1)}$$

is the eq for x_3

Now,

$$\begin{aligned} & \alpha x_1 + \beta x_2 \\ &= \alpha (M x_1 + (1-\lambda) s_1) + \beta (M x_2 + (1-\lambda) s_2) \\ &= M(\alpha x_1 + \beta x_2) + (1-\lambda) (\underbrace{\alpha s_1 + \beta s_2}_{= s_3}) \\ &= M(\alpha x_1 + \beta x_2) + (1-\lambda) s_3 \quad \text{--- (2)} \end{aligned}$$

From (1) & (2), putting $\alpha x_1 + \beta x_2 = x_3$ satisfies.

∴ We proved that we can find the PPR of any user, whose teleport set is the span of the teleport set of vectors in V , by using linear combination of PPR vectors in V .

$$2. p_0 = \alpha \lambda + \alpha \sum_{i=1}^k \frac{p_i}{d_i} + \frac{(1-\alpha)}{N} \quad \text{--- (1)}$$

From the figure, we can see that all the k span pages are identical,
 $\therefore p_1 = p_2 = \dots p_k = (\text{say}) p_{\text{span}}$

Also, we can see that d_i (the outdegree) of all the span pages is 1

$$\therefore \sum_{i=1}^k \frac{p_i}{d_i} = \sum_{i=1}^k \frac{p_{\text{span}}}{1} = k p_{\text{span}}$$

$$\begin{aligned} \text{Now, } p_{\text{span}} &= \alpha \frac{p_0}{d_0} + \frac{1-\alpha}{N} \\ &= \frac{\alpha p_0}{k} + \frac{1-\alpha}{N} \end{aligned}$$

$$\therefore \sum_{i=1}^k \frac{p_i}{d_i} = k \left[\frac{\alpha p_0}{k} + \frac{1-\alpha}{N} \right] \quad \text{--- (2)}$$

From (1) and (2),

$$p_0 = \alpha \lambda + \alpha k \left(\frac{\alpha p_0}{k} + \frac{1-\alpha}{N} \right) + \frac{1-\alpha}{N}$$

$$p_0 = \alpha \lambda + \alpha^2 p_0 + \frac{\alpha k(1-\alpha)}{N} + \frac{1-\alpha}{N}$$

$$p_0(1-\alpha^2) = \alpha \lambda + \frac{1-\alpha}{N}(1+\alpha k)$$

$$\therefore p_0 = \frac{\alpha \lambda}{1-\alpha^2} + \left(\frac{1+\alpha k}{1+\alpha} \right) \frac{1}{N}$$

3. Given:
 • Turnstile stream of ' n ' distinct items.
 • Power law exp. = 3 distribution.
 i.e. no. of items with freq k is $\frac{C}{k^3}$

$$\therefore \sum_{k=1}^{\infty} \frac{C}{k^3} = n$$

$$\Rightarrow C \sum_{k=1}^{\infty} \frac{1}{k^3} = n \quad \text{constant (say) } \lambda$$

$$\Rightarrow C \lambda = n \quad \text{Apéry Constant, } \approx 1.202$$

$$\therefore C = O(n)$$

Now, as derived in class, the bounds for CM sketch and CS are:

$$f_x \leq \hat{f}_x \leq f_x + \varepsilon m$$

$$|f_x - \varepsilon|f_x|_2 \leq |\hat{f}_x|_2 \leq |f_x + \varepsilon|f_x|_2$$

$$\therefore \text{For CM, } |\hat{f}_x - f_x|_2 \leq \varepsilon |f_x|_2 \quad \left\{ \begin{array}{l} \text{For CS, } |\hat{f}_x - f_x|_2 \leq \varepsilon |f_x|_2 \\ \text{Both w.p. } \frac{1-\delta}{\delta} \end{array} \right.$$

$$\text{when we set } w = \log \left(\frac{1}{\delta} \right)$$

$$\text{But in CM, } d = O\left(\frac{1}{\varepsilon}\right)$$

$$\text{in CS, } d = O\left(\frac{1}{\varepsilon^2}\right)$$

$$\therefore \varepsilon = O(\varepsilon^2) \Rightarrow \varepsilon < \varepsilon' \quad [\because \text{both } < 1]$$

Let m = length of stream

$$\therefore \sum_{k=1}^{\infty} k \frac{C}{k^3} = m$$

$$\therefore m = \sum_{k=1}^{\infty} \frac{C}{k^2}$$

$$|\hat{f}_x|_2 = \sqrt{\sum_{i=1}^m f_i^2}$$

where f_i is the frequency of an item.

$$\text{No. of items with freq. } k = \frac{C}{k^3}$$

If $f_i = k$, no. of items with $f_i \geq k$ is $\frac{C}{k^3}$

$$\therefore \sum_{i=1}^{\infty} f_i^2 = \sum_{k=1}^{\infty} \sum_{i=1}^k \frac{C}{k^3} * k^2$$

$$= \sum_{k=1}^{\infty} \frac{C}{k^3} * k^2$$

$$= C \sum_{k=1}^{\infty} \frac{1}{k}$$

$$\therefore |\hat{f}_x|_2 = \sqrt{C \sum_{k=1}^{\infty} \frac{1}{k}}$$

$$\text{For CM, bound range} = \frac{\varepsilon m}{\varepsilon} = \frac{C}{k^2}$$

$$\text{For CS, bound range} = 2\varepsilon |\hat{f}_x|_2 = 2\varepsilon \sqrt{C \sum_{k=1}^{\infty} \frac{1}{k}}$$

$$\approx \frac{1}{k^2} = \frac{\pi^2}{6} \quad (\text{converges})$$

However, $\sum_{k=1}^{\infty} \frac{1}{k}$ does not converge for

∴ CM sketch bound is tighter and appropriate