

EE2016 Midsem

Question 1)

In the first question, we are required to minimize the total number of LDR/STR instructions. Therefore, we require 6 registers to store values of $A[i]$ to $A[i+5]$. In addition, we require a register to store the minimum value of S obtained, one to store the corresponding index and one to store i to keep track of the number of iterations of the loop. Therefore in total, we require overall 11 registers in this code. Of these 11 registers, 4 registers need not be saved. Therefore, a total of 8 registers are to be saved and restored.

Initially, $r0$ is stored with address of $A[0]$ and $r1$ is stored with “ n ”. Before the loop, $A[5]$ to $A[0]$ are stored in $r1 - r6$.

Registers:

$r0$: To store $A[i+5]$

$r7$: To store S minimum

$r8$: To store index starting from $n-7$ which decrements by 1 to track no of iterations

$r9$: To store $n-6-i_0$. Here i_0 is the index which corresponds to S minimum

$r10$: To store present value of S

$r11$: To store $n-6$

Just before the loop starts, $r10$ is stored with S_0 . Since it is the first S calculated, it is stored as the minimum S and $n-6$ is stored in $r9$.

When loop starts, $i = 1$ and $A[i+4]$ is stored in $r1$, $A[i+3]$ is stored in $r2$, $A[i+2]$ is stored in $r3$, $A[i+1]$ is stored in $r4$ and $A[i]$ is stored in $r5$. Now $A[i+5]$ is stored in $r6$. And using these registers S_1 is calculated and stored in $r10$. Now we check whether S_1 is less than S minimum or not. If it is less, then we store S_1 in $r7$ and store $r8$ as index corresponding to S_1 in $r9$ and decrements $r8$. If it is not less, then it just decrements $r8$.

After incrementing $r8$, we now have $r8 = n-8$ which is same as $i = 2$. Now, $A[i+4]$ is stored in $r6$, $A[i+3]$ is stored in $r1$, $A[i+2]$ is stored in $r2$, $A[i+1]$ is stored in $r3$ and $A[i]$ is stored in $r4$. Now $A[i+5]$ is stored in $r5$. Using these registers, S_2 is calculated and stored in $r10$. S_2 is compared with S minimum. If S_2 is smaller than S , it is stored as minimum and corresponding $r8$ is stored in $r9$ and $r8$ decrements. If not smaller, then just $r8$ decrements.

Thus after each decrement in $r8$, the values $A[i]$ to $A[i+5]$ are stored cyclically in registers $r1-r6$. This process continues until $r8$ becomes less than zero. Once $r8$ is less than zero it breaks out of the loop and branches to X . $r9$ contains $n-6-i_0$, i_0 is the index corresponding to S which is minimum. Using $r11$ which is stored with $n-6$, $r9$ is subtracted from $r11$ and stored in $r9$. Now $r9$

is stored with i_0 which is the index corresponding to minimum S . This value is moved to r_0 which returns this value to the main program.

The total number of clock cycles taken by the function = $116 * L + 966$

Here L is clock cycles taken by LDR

Number of clock cycles required for an array of size 100 = 1546

Number of registers required to be saved = 8 ($r_4, r_5, r_6, r_7, r_8, r_9, r_{10}, r_{11}$)

Cycles taken by LDR/STR	1	5	10	20
Cycles taken by the function	1082	1546	2126	3286

Question 2)

In the second question, we are required to minimize the total clock cycles given that LDR instruction takes equal time as other instructions. In this question, we require 9 registers. Initially r_0 is stored with address of $A[0]$, r_1 is stored with n .

Registers :

r_0 : To store address of $A[i-1]$

r_1 : To store n

r_2 : To store $n-6-i_0$. i_0 is the index corresponding to minimum S .

r_3 : To store address of $A[i+5]$

r_4 : To store S of previous iteration

r_5 : To store S of present iteration

r_6 : Is used as scratch

r_7 : To store minimum S

r_8 : To store index which decrements every iteration to keep track of them

We use the fact that

$$S_i = A[i+5] - A[i-1] - S_{i-1}$$

Initially r_4 is stored with $A[5]$. Using r_6 to load values of elements of array temporarily, we use r_6 and add to or subtract it from r_4 to get S_0 in r_4 . This r_4 is moved into r_7 , i.e., it is stored as S minimum as it is the first S we calculated. r_8 is stored with $n-7$ and r_2 is stored with $n-6$.

Now, in the loop, $A[i+5]$ is stored in r5, and $A[i-1]$ is stored in r6. Now subtracting r6 and r4 from r5 we get S in r5. This S is compared with the minimum S till now. If the present S is less than the minimum S, it is stored in minimum S and corresponding r8 is also stored in r2 and r8 is decremented. If it is not less, only r8 is decremented.

The loop continues till r8 is positive or zero. Once r8 is negative, it breaks from the loop.

r2 is now stored with $n-6-i0$. Using the value in r1 and subtracting r2 from $n-6$ and storing it in r2, we get i0 in r2. This r2 is stored in r0 which returns the value to main function.

This scheme can't be used in first problem as the no of clock cycles taken by LDR is 5 times more than normal instructions which drastically increases the total clock cycles if we use two or more LDRs in the loop.

The total number of clock cycles taken by function in question 2 = $204 * L + 765$
Here L is clock cycles taken by LDR

No. of clock cycles taken by the function in question 2 = 969

No. of registers to be saved = 5

Cycles for LDR	1	5	10	20
Total Cycles for Q1	1082	1546	2126	3286
Total Cycles for Q2	969	1173	1377	1581