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**SUBJECT: DESIGN AND ANALYSIS OF
ALGORITHM**

CODE: CS 351

STRING MATCHING ALGORITHM

PRACTICAL-8.1

AIM:

Suppose you are given a source string $S[0..n-1]$ of length n , consisting of symbols a and b . Suppose that you are given a pattern string $P[0..m-1]$ of length $m < n$, consisting of symbols a , b , and $*$, representing a pattern to be found in string S . The symbol $*$ is a “wild card” symbol, which matches a single symbol, either a or b . The other symbols must match exactly. The problem is to output a sorted list M of valid “match positions”, which are positions j in S such that pattern P matches the substring $S[j..j+|P|-1]$.

For example, if $S = ababbab$ and $P = ab*$, then the output M should be $[0, 2]$. Implement a straightforward, naive algorithm to solve the problem.

PROGRAM CODE:

```
#include<iostream>
#include <string.h>

using namespace std;

int main()
{
    char t[100], p[100];
    int tn, pn, shift[20] = {0}, s = 0, i, j = 0, count = 0, m = 0;
    cout<<"ENTER THE TEXT : ";
    cin>>t;

    cout<<"ENTER THE PATTERN : ";
    cin>>p;

    tn = strlen(t);
    pn = strlen(p);

    while (s != (tn - pn + 1))
    {
        j = 0;
        for (i = s; i < pn + s; i++)
        {
            if (p[j] == t[i])
            {
                count++;
                if (count == pn)
                {
                    count = 0;
                }
            }
        }
    }
}
```

```
        shift[m] = s;
        m++;
    }
}
else
{
    count = 0;
    break;
}
j++;
}
s++;
}
if (m > 0)
{
    printf("\n\nVALID SHIFTS : ");
    for (i = 0; i < m; i++)
        printf("%d \n", shift[i]);
}
else
{
    printf("\n\n NO VALID SHIFTS AVAILABLE");
}

cout<<"PARTH PATEL\n19DCS098"<<endl;
return 0;
}
```

OUTPUT:

```
ENTER THE TEXT : parthpatel
ENTER THE PATTERN : rt

VALID SHIFTS : 2
PARTH PATEL
19DCS098
```

Conclusion:

- Naive pattern searching is the simplest method among other pattern searching algorithms.
- It checks for all character of the main string to the pattern.
- This algorithm is helpful for smaller texts.
- It does not need any pre-processing phases.
- The time complexity is $O(m*n)$. The m is the size of pattern and n is the size of the main string

PRACTICAL-8.2**AIM:**

Implement Rabin karp algorithm and test it on the following test cases:

Test Case	String	Pattern
1	2359023141526739921	31415 q=13
2	ABAAABCDBBABCDDDEBCABC	ABC q=101

PROGRAM CODE:

```
#include <iostream>
#include <string.h>
#define d 256
using namespace std;
void search(char pattern[], char text[], int q)
{
    int M = strlen(pattern);
    int N = strlen(text);
    int i, j;
    int p = 0;
    int t = 0;
    int h = 1;
    for (i = 0; i < M - 1; i++)
        h = (h * d) % q;
    for (i = 0; i < M; i++)
    {
        p = (d * p + pattern[i]) % q;
        t = (d * t + text[i]) % q;
    }
    for (i = 0; i <= N - M; i++)
    {
        if (p == t)
        {
            for (j = 0; j < M; j++)
            {
                if (text[i + j] != pattern[j])
                    break;
            }
        }
    }
}
```

```
        if (j == M)

            cout<<"PATTERN FOUND AT INDEX : "<<i<<endl;

        }
        if (i < N - M)
        {
            t = (d * (t - text[i] * h) + text[i + M]) % q;
            if (t < 0)
                t = (t + q);
        }
    }
}

int main()
{

    char text[100];
    char patterntern[100];
    int q;

    cout<<"ENTER THE TEXT : ";
    cin>>text;
    cout<<"ENTER THE PATTERN : ";
    cin>>patterntern;

    cout<<"ENTER THE VALUE OF q : ";
    cin>>q;
    search(patterntern, text, q);

    cout<<"PARTH PATEL\n19DCS098"<<endl;
    return 0;
}
```


OUTPUT:**Test Case-1:**

```
ENTER THE TEXT : 2359023141526739921
ENTER THE PATTERN : 31415
ENTER THE VALUE OF q : 13
PATTERN FOUND AT INDEX : 6
PARTH PATEL
19DCS098
```

Test Case-2:

```
ENTER THE TEXT : ABAAABCDBBABCDDDEBCABC
ENTER THE PATTERN : ABC
ENTER THE VALUE OF q : 101
PATTERN FOUND AT INDEX : 4
PATTERN FOUND AT INDEX : 10
PATTERN FOUND AT INDEX : 18
PARTH PATEL
19DCS098
```

CONCLUSION:

- For text of length n and p patterns of combined length m , its average and best case time complexity is $O(n+m)$, but its worst-case time complexity is $O(nm)$