

Lecture - 1

Digital Circuit

- Number systems

- Logic gates

- Logic family

(combinational & sequential circuit)

→ Number system

1) Binary No. System

base = 2 or B

digits = 0, 1 ex. $(101)_B$ or $(101)_2$

2) Octal No. System

base = 8 or O

digits = 0, 1, 2 ... 7

ex. $(345)_8$

3) Decimal

base = 10 or D

digits = 0, 1, ..., 9

ex. $(189)_{10}$

4) Hexadecimal No. system

base - 16 or H

Digits - 0 ~ 9

Symbol	A	B	C	D	E	F
10, 11, 12, 13, 14, 15	1	1	1	1	1	1

5) Base - 4

Digits 0, 1, 2, 3

ex. $(12)_4$

0, 1, 2, 3

10, 11, 12, 13

20, 21, 22, 23

In

Hexadecimal number system

0, 1, 2, 3, 4, 5, 6, 7, 8, 9, A, B, C, D, E, F, 10, 11, 12,

13, 14, 15, 16, 17, 18, 19, 1A, 1B, 1C, 1D, 1E, 1F, 20, 21, 22, 23, 24, 25, 26, 27, 28, 29, 2A, 2B, 2C, 2D, 2E, 2F

Base - 4

0, 1, 2, 3

Base - 10

0, 1, 2, 3, 4, 5, 6, 7, 8, 9

Base - 8

0, 1, 2, 3, 4, 5, 6, 7

Note:-

Largest number in any number system is 1 less than its base.

Conversion:- 0.625_{10} (Binary part)

Hence 21 - 3rd

① Division by 2
Quotient 130321
Dividend 730321
 $(25)_{10} = (11001)_2$

Quotient	2	15	→ Remainder
divisor	2	12	1
	2	6	0
	2	3	0
	2	1	1
	0	1	

$$(25)_{10} = (11001)_2$$

$$\rightarrow (0.625)_{10} = (0.1101)_2$$

② $0.625 \times 2 = 1.25$ 1 AS
 $0.25 \times 2 = 0.5$ 0
 $0.5 \times 2 = 1.0$ 1

$$(0.625)_{10} = (0.1101)_2$$

$$(15.625)_{10} = (11001.101)_2$$

* Octal to binary :-

$$(73)_8 = (111001)_2$$

* Decimal - octal. :- 9 15 8

$$(73)_{10} = (111)_8$$

$$\begin{array}{r} 8 | 73 \\ 8 | 9 \quad 1 \\ 8 | 1 \quad 1 \\ \hline 0 \quad 1 \end{array}$$

$$(73)_{10} = (111)_8$$

$$* (73)_{10} = (111)_8 = 1 \cdot 8^2 + 1 \cdot 8^1 + 1 \cdot 8^0$$

$$\begin{array}{r} 16 | 73 \\ 16 | 24 \quad 9 \\ \hline 0 \quad 4 \end{array}$$

$$(73)_{10} = (49)_{16} = (111)_8$$

Note:- If base is greater, than number will be smaller, If base is lesser than number will be bigger.

Binary to decimal.

$$\text{MSB} \curvearrowleft (1001)_2 = (\quad)_{10} \quad \text{LSB} \quad \text{- least significant bit}$$

$$\begin{pmatrix} 1 & 0 & 0 & 1 \end{pmatrix}_2 \quad q=7$$

$$= 1x^2^0 + 0x^2^1 + 0x^2^2 + 1x^2^3$$

$$= 1 + 0 + 0 + 8$$

$$(1001)_2 = (9)_{10}$$

→ binary to decimal:

$$(101.101)_2 = (1.75)_{10.8}$$

$$\begin{pmatrix} 1 & 0 & 1 \\ 0 & 1 & 0 \end{pmatrix}_2 = \begin{pmatrix} 1 & -2 & 3 \\ 0 & 1 & 0 \end{pmatrix} \in \mathcal{O}_1(\mathbb{R})$$

$$= 0 \times 2^0 + 1 \times 2^1 + 1 \times 2^{-1} + 0 \times 2^{-2} + 1 \times 2^{-3}$$

$$(10.101)_2 = (2.625)_{10}$$

adhesive went + stopper in end of

$$j_0 = 1 \quad (1, 1, 0, 1, 0, 1, 0)_2 = 1011010_2$$

Specid 3d Missed today with

will reach the sea

$$\underbrace{001}_{1} \quad \underbrace{101}_5 \quad \underbrace{010}_2 = (152)_8$$

Make group of 3 bit (for octal) (bcos using 3 bit you can represent any number in octal)

* Binary to hexadecimal

$$(10101011)_2 = (?)_{16}$$

make group of 4 bit for hex-decimal (bcos using 4 bit you can represent any number in hexadecimal number system)

$$\underbrace{11010}_A \quad \underbrace{1011}_B = (AB)_{16}$$

* Convert $(4F)_{16} = (?)_8$

$$= 4 \quad 1 \quad 0 \quad F$$

$$= (0100 \quad 1111)_2$$

$$= \underbrace{001}_1 \quad \underbrace{001}_1 \quad \underbrace{111}_7 = (117)_8$$

$$(6A)_{16} = ()_{10}$$

$(6A)_{16}$ not yet to work 2016
 word (1000 up tide 1000)
 type conversion not up tide 1000
 $= A \times 16^0 + 6 \times 16^1$ the parallel
 $= 10 + 96 = (106)_{10}$

* gate question

Q The octal equivalent of binary

$$(AB \cdot CD)_{16}$$

$$\therefore (A \cdot B \cdot C \cdot D)_{16}$$

$$(1010 \ 1011 \cdot 1100 \ 1101)_2$$

↔

$$010 \ 101 \ 011 \cdot 110 \ 011 \ 010$$

$$2 \ 5 \ 3 \cdot 6 \ 3 \ 2$$

$$(253.632)_8$$

will reach the sea.

Q The binary representation of decimal number (1.375) is

$$(1)_{10} = ()_2$$

$$\begin{array}{r} | \\ 2 \mid 1 \\ \hline 0 \quad 1 \end{array}$$

$$= (1)_2$$

$$01 \text{ (ES)}$$

$$01 \text{ (22)}$$

$\times 2$, second

$$01 \times 2 = 10$$

second

$$(0.375)_{10} = ()_2$$

$$\begin{aligned} 0.375 \times 2 &= 0.75 && \xrightarrow{\text{top}} 01 < 10 \\ 0.75 \times 2 &= 1.5 && \xrightarrow{\text{middle}} 1 \\ 0.5 \times 2 &= 1.0 && \xrightarrow{\text{bottom}} 01 \end{aligned}$$

$$(0.375)_{10} = (0.011)_2$$

$$\checkmark (1.375)_{10} = (1.011)_2$$

double verifying

$$(1.011)_2 = ()_{10}$$

$$(1.011)_2 = 1 \times 2^0 + 0 \times 2^1 + 1 \times 2^{-2}$$

$$0 \xrightarrow{-1-2-3} + 1 \times 2^{-3}$$

$$= 1 + 0 + 0.25 + 0.125$$

$$= 1 + 0.375 = (1.375)_{10}$$

Arithmetic operations

Addition carried out

$$\begin{array}{r} (23)_{10} \\ + (55)_{10} \\ \hline 78 \end{array}$$

$$\begin{array}{r} 21 (282.1) \\ + 55 \\ \hline 78 \end{array}$$

if $\frac{8}{\text{base}} = \text{Valid decimal number}$,
No carry.

$$\begin{array}{r} 68 \\ + 46 \\ \hline 114 \end{array}$$

$14 > 10$ It's not
base

single digit valid decimal

number so it written as

$$14 = 10 \times 1 + 4$$

Base → Carry → Result

Base × Carry + Result

$$11 = 10 \times 1 + 1$$

1 → result

$$0.1(282.1) = 282.0 \times 1 + 1$$

addition with Binary

$$\begin{array}{r} \text{0 1}_2 \\ + \text{0 0}_2 \\ \hline \text{1 0 0 1}_2 \end{array}$$

base 2

2 < base (2)

2 is not valid binary num= digit

$$2 = 2 \times 1 + 0$$

base carry result

Octal 1
6 7

$$\begin{array}{r} 2 \\ + 1 \\ \hline 7 \end{array}$$

$$7 + 2 = 9 \quad \times \text{Valid}$$

$$7 + 2 = 9$$

$$9 = 8 \times 1 + 1$$

Hexadecimal

$$\begin{array}{r} 2 F 1 \\ + A 3 A \\ \hline D 2 B \end{array}$$

2 + 1 × F = 7
3 + 1 × F = 8

$$F + 3 = 15 + 3 = 18 \quad \times \text{Valid}$$

$$18 = 16 \times 1 + 2$$

$$F(021) = \boxed{\text{Valid}}$$

$$\text{Add } (63)_7 + (56)_8$$

will give different → direct add them

So we have to convert both into some number system other than $(0, 1)$

Can (100 others) \rightarrow

$$(56)_8 = ((\text{base})_{10} = (1)_{10})_2$$

$$(56)_8 \quad 7^0 + 1 \times 8^1 = 8$$

$$6 \times 8^0 + 5 \times 8^1$$

$$= 6 + 40$$

$$= (46)_{10}$$

$$\begin{array}{r} 7 \\ \overline{)46} \\ 7 \\ \overline{)6} \\ 7 \\ \overline{)0} \\ 6 \end{array}$$

$$= (64)_7$$

$\boxed{1}$	$\boxed{(64)_7}$
$\boxed{+}$	$\boxed{(63)_7}$
\hline	
160	

$$\boxed{\text{Ans}} = (160)_7$$

$$7 = 7 \times 1 + 0$$

$$13 = 7 \times 1 + 6$$

Subtraction

~~base~~
10+2

Base

$$\begin{array}{r} 56 \\ - 29 \\ \hline 27 \end{array}$$

math 586 39

- 123
29

$$\frac{-1.45}{11.5}$$

Binary

11 ♂

$$\begin{array}{r} \textcircled{101} \\ \underline{- 11} \\ \hline 00 \end{array}$$

Oct 91 :-

$$\begin{array}{r}
 8+2 \\
 62 \\
 -\underline{45} \\
 \hline
 33
 \end{array}$$

~~100~~

Hexadecimal :-

16+A

$$\frac{1}{0} \frac{F + 1 \times 8 - 8}{B}$$

Lecture 2] $(15)_7 - (26)_7 + (53)_7$

Wrong

explanations

* Multiplication No system

$$\begin{array}{r}
 \overset{201}{\cancel{211}} \\
 \times 6 \\
 \hline
 66
 \end{array}
 \quad
 \begin{array}{r}
 58 \\
 \times 36 \\
 \hline
 30
 \end{array}
 \quad
 \begin{array}{r}
 + 150 \\
 \hline
 180
 \end{array}$$

* Binary:-

$$\begin{array}{r}
 101 \\
 \times 11 \\
 \hline
 101 \\
 1010 \\
 \hline
 1111
 \end{array}$$

* Octal :-

$$\begin{array}{r}
 74 \\
 \times 21 \\
 \hline
 74 \\
 148 \\
 \hline
 170
 \end{array}$$

$$8 = 8 \times 1 + 0$$

$$15 = 8 \times 1 + 7$$

Hexadecimal :- (1-Second)

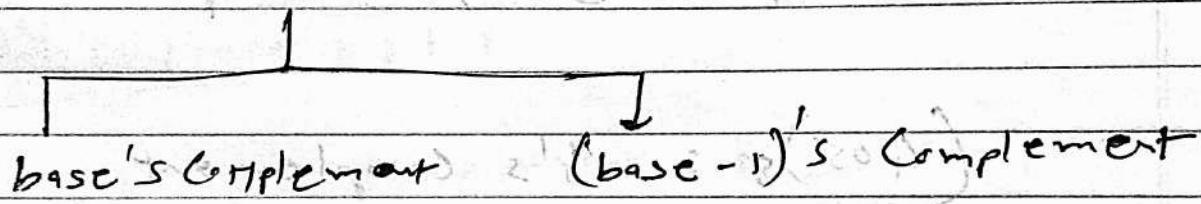
HexaDecimal LFS \leftrightarrow 8(2nd)

$$\begin{array}{r} AB \\ \times 16 \\ \hline 22 FFF \\ \hline 156203 \\ 156 \times 16 \\ \hline 16 B6 \end{array}$$

HexaDecimal 2nd (8)

Complements :- 1's & 2's

HexaDecimal L(1-based) To HexaDec



ex. L + HexaDecimal 2nd =

binary

\equiv 2's 1's

Decimal 10's 9's

How to find complement :- 1's & 2's

① (base-1)'s complement

Subtracting each digit from the largest digit in the number system.

(base-1)'s complement.

ex $(405)_8 \Rightarrow 7^1$'s complement

$$\begin{array}{r}
 2+1 \times 2^1 \\
 2+1 \times 2^1 \quad 777\ldots \\
 \hline
 405 \quad 2^1 \\
 \hline
 372 \quad 2^1
 \end{array}$$

2) base's complement:-

by Simply adding 1 to the result of (base-1)'s complement

$$\begin{aligned}
 (405)_8 &= 8^1 \text{ complement} \\
 &= 7^1 \text{ complement} + 1
 \end{aligned}$$

$$\begin{array}{r}
 2 = 372 \\
 + 1 \\
 \hline
 373
 \end{array}$$

* Calculate 3^1's complement for $(301)_4$

(Base-1)'s complement

$$\begin{array}{r}
 3^1 \text{ complement} \quad 333 \\
 301 \\
 \hline
 032
 \end{array}$$

3^1 complement

* Calculate 2's complement for

$$(1001)_2 \quad | \quad 0 \quad 0 \quad 1$$

$$\begin{array}{r} \text{1's complement} \\ \hline 1 \quad 1 \quad 1 \quad 1 \\ 1 \quad 1 \quad 1 \quad 1 \\ \hline 1 \quad 0 \quad 0 \quad 1 \\ \hline 0 \quad 1 \quad 1 \quad 0 \end{array}$$

$$\begin{array}{r} 0 \quad 1 \quad 1 \quad 0 \\ + \quad \quad \quad \quad 1 \\ \hline 0 \quad 1 \quad 1 \quad 1 \end{array}$$

direct 1's complement & 2's complement

$1001 \Rightarrow$ Reverse all bit
1's complement $\rightarrow 0110$

$$1001 \Rightarrow$$

 \downarrow
LSB

Start from LSB
Look for first "1"
keep that "1" as it is, ~~go~~ after first "1" reverse all bit

2's Complement \leftarrow ~~first~~ 1

$$\begin{array}{r} 1 \quad 0 \quad 1 \quad 0 \\ \downarrow \downarrow \downarrow \downarrow \\ 0 \quad 1 \quad 1 \quad 0 \end{array}$$

Now add 1 to get 2's complement

~~ex~~

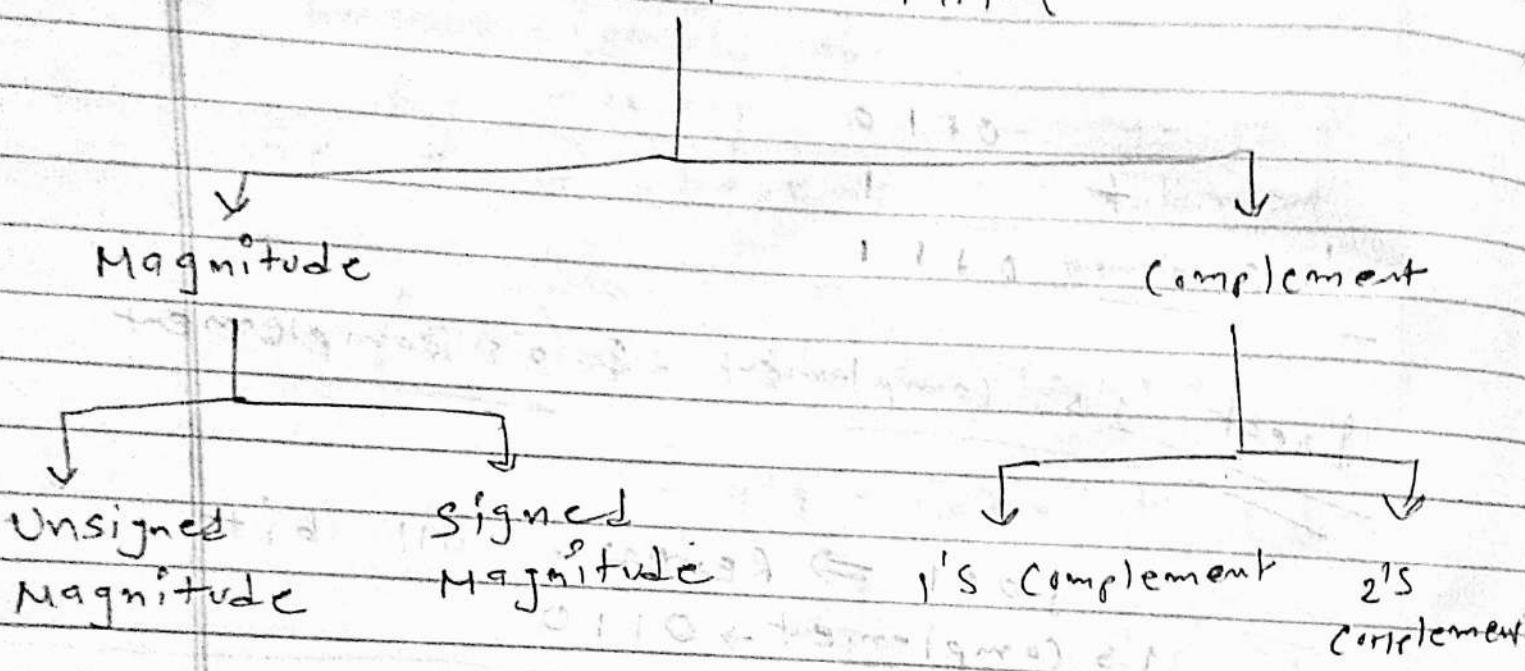
Hint
2

$$\begin{array}{cccccc} & 1 & 0 & 0 & 1 & (1001) \\ \downarrow & \downarrow & \downarrow & \downarrow & \downarrow \\ 0 & 1 & 1111 & 1 & 1001 \end{array}$$

2's complement



Number Representation



→ we can → Both +ve → Both → Both
 represent & -ve & +ve & -ve → +ve &
 only +ve & -ve → +ve & -ve

$$\begin{array}{l} \rightarrow +6 \\ = 110 \end{array} \rightarrow +6 \rightarrow \begin{array}{c} +6 \\ 0110 \end{array} \rightarrow +6$$

~~0110~~ ↓ ↓ ↓ ↓

+ve 0 1 1 0

representation of positive number are same in all case

In signed magnitude	Signed magnitude	I's Complement	2's complement
-6			
X			
+6 in 6bit	+6 in 6bit	+6 in 6bit	+6 in 6bit
0'00110	0'00110	0'00110 → +6	0'00110 → +6
-6 in 6bit	-6 in 6bit	-6 in 6bit	-6 in 6bit
1'11010	1'11010	1'11010 → -6	1'11010 → -6
F -	Copy MSB & -6 representation of I's complement	Copy MSB & -6 representation of I's complement	Copy MSB & -6 representation of 2's complement
0 -	11'001	11'001	11'010

Signed Magnitude

0 0 0 0
0 0 0 1 ←
0 1 1 0 0 0 0
0 1 1 1 1 1 1

1 0 0 0 1 0 0

↑ 0 0 0 1 0 0

1 1 1 0 0 0

↑ 1 1 1 0 0 0

disadvantages →

→ 2 different representation for 0

→ Range

$$\begin{array}{c} \text{1 bit} \\ -7 \text{ to } +7 \text{ (4 bits)} \\ -(2^{4-1}-1) \text{ to } +(2^{4-1}-1) \\ -7 \text{ to } +7 \end{array}$$

$$n \text{ bit } -(2^{n-1}-1) \text{ to } +(2^{n-1}-1)$$

* 1's Complement

Decimal

0 0 0 0 0 0 0

+0

0 0 0 0 0 1 1

+1

0 1 1 1 0 0 0

+6

0 1 1 1 1 0 0

+7

1 0 0 0 0 0 0 1

-7

1 0 0 0 0 0 1 0

-6

1 1 1 1 1 0 0 0

-1

0 1 0 1 1 1 1 1 1

-0

Decimal

+0

+1 ←

+6 ←

+7

-0

-1

-6

-7

short form

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

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1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7

$-(2^{n-1}-1)$ to $+(2^{n-1}-1)$

1's Complement

Decimal

+0

+1

+6

+7

-7

-6

-1

-0

size

disadvantages

→ 2 different representation for 0

→ Range

1 bit

-7 to +7

($2^{4-1}-1$) to $+(2^{4-1}-1)$

-7 to +7</p

disadvantages:

→ 2 different representations for 0

→ long delay

$1101 +$

$0 = 7 \quad 7 + 7 \quad 4 \text{ bit}$

$-(2^4 - 1) \text{ to } +(2^{4-1} - 1) \in \beta -$

$-7 \rightarrow 7$

$n \text{ bit } -(2^{n-1} - 1) \text{ to } +(2^{n-1} - 1)$

2's complement

$000001 +$

001001

0110

0111

1000

1001 biased

1110

decimal = 2

$+00010 = 2 +$

$+1010$

$+0011$

$+7$

-8

-7

-2

$01111011 + 01110 = 2$

with more zeros, it's more precise + 1/2

to add $-8 + 7 \rightarrow +7$ (4 bit)

$-(2^4 - 1) \rightarrow +(2^{4-1} - 1)$

$n \text{ bit } -(2^{n-1}) \rightarrow +(2^{n-1} - 1)$

000100

100101

111001

111000

Arithmetic operation using Complement

Q) $5 - 4 = ?$

1's complement

$$+5 \Rightarrow 0101$$

$$+4 \Rightarrow 0100$$

$$\begin{array}{r} \cancel{1} \\ +1 \\ \hline 1000 \end{array}$$

$$-4 \Rightarrow (+0100) + \cancel{0} + (-1100)$$

$$\begin{array}{r} \cancel{1} \\ +0 \\ \hline 0001 \end{array} \checkmark +1$$

8's complement: $+0 + (-1) \Rightarrow This \text{ is in } 1\text{'s complement form}$

Q) $5 - 4$

$$+5 = 0101$$

$$+4 = 0100$$

$$\begin{array}{r} \cancel{1} \\ +1 \\ \hline 0001 \end{array}$$

$$-4 = 1100$$

$$\begin{array}{r} 1 \\ +0101 \\ \hline 0101 \end{array} 5$$

$$+1100 \quad 4$$

$$\begin{array}{r} \cancel{1} \\ +1 \\ \hline 0001 \end{array} \checkmark +1$$

\Rightarrow This result is
in 2's complement
form

discard it

gate $x = 01110$ & $y = 11001$ are 2
5 bit binary number represented in
2's complement form. The sum of
 $x + y$ represented in 2's complement
format using 6 bit is

(4)

$$100111$$

(5)

$$001000$$

(6)

$$000111$$

(7)

$$101001$$

$$N = 01110 \leftarrow +14$$

$$4 = 111001 \leftarrow -7$$

J. 2's

$$00111 \leftarrow -7$$

16	0	0
-7	1	1
+7	1	0

in 6 bit, 2's complement

norm

100011

Lecture:-3

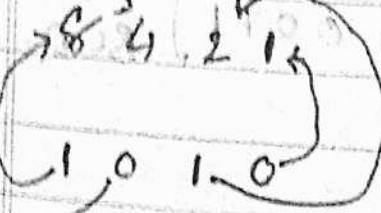
Codes

Weighted

Non-weighted

Code

weight



0100 ← excess

→ BCD 8 + 0001 = 8 → Excess -3
→ gray code

BCD is always 4 bit

BCD

decimal	BCD	Excess-3
0	0 0 0 0	0 0 0 0
1	0 0 0 1	0 0 0 1
2	0 0 1 0	0 0 1 0
3	0 0 1 1	0 0 1 1
4	0 1 0 0	0 1 0 0
5	0 1 0 1	0 1 0 1
6	0 1 1 0	0 1 1 0
7	0 1 1 1	0 1 1 1
8	1 0 0 0	1 0 0 0
9	1 0 0 1	1 0 0 1
10		
11		
12		
13		
14		
15		

Invalid BCD code

to represent greater number

$$(23)_{10} \rightarrow (0010\ 0011)_{BCD}$$

Excess - 3 code

$$\text{Excess - 3} = \text{BCD} + 3$$

Decimal	B ₃	B ₂	B ₁	B ₀	E ₃	E ₂	E ₁	E ₀
0	0	0	0	0	0	0	1	1
1	0	0	1	0	0	1	0	0
2	0	1	0	0	1	0	0	1
3	0	1	0	1	1	0	1	1
4	0	1	1	0	0	1	1	1
5	0	1	1	0	1	1	0	0
6	0	1	1	0	1	1	0	1
7	0	1	1	1	1	0	0	0
8	1	0	0	0	1	0	1	1
9	1	0	0	1	1	1	0	0

It's also called self Complementing code.

→ Gray Code :- (Cyclic Code all unit distance code)

	B ₃	B ₂	B ₁	B ₀	G ₃	G ₂	G ₁	G ₀
0	0	0	0	0	0	0	0	0
1	0	0	0	1	0	0	0	1
2	0	0	1	0	0	0	1	0
3	0	0	1	1	0	0	1	0
4	0	1	0	0	0	1	1	0
5	0	1	0	1	0	1	1	0
6	0	1	1	0	0	1	0	1
7	0	1	1	1	0	1	0	0
8	1	0	0	0	1	0	0	0
9	1	0	0	1	1	0	0	1
10	1	0	1	0	1	1	0	1
11	1	0	1	1	1	1	1	0
12	1	1	0	0	1	0	1	0
13	1	1	0	1	1	0	1	1
14	1	1	1	0	1	0	0	1
15	1	1	1	1	1	0	0	0

BCD Addition

	1	0	1	0	1	0
(a)	1	0	1	0	1	0
83	0	0	0	0	1	0
+ 16	0	0	0	1	1	0
<u>-</u> 39	0	0	1	1	1	0
	0	0	1	1	1	0
	1	0	0	1	3	0
	0	1	0	1	3	0
(b)	1	1	0	1	1	0
66	0	1	1	0	0	1
+ 55	0	1	0	1	0	1
<u>-</u> 121	0	1	1	1	0	1

~~Both
Invalid
BCD code~~

to make it's valid add

$$\begin{array}{r} 10110 = +6 \\ \hline 0110 \end{array}$$

each Invalid 0001 0 0 1 0 0 0

each Invalid

BCD Number: ↓ ↓ ↓ ↓

BCD Number: ↓ ↓ ↓ ↓

1 2 3

0	1	0	1		0	0	0	1
-	-	0	1		0	0	0	-
-	0	0	1		0	0	0	-

0 0 0 1 1 1 1 1

Convert binary to gray code

(e)

$$(1011)_2 = (1110)_G$$

(e)

$$(1111)_2 = (1000)_G$$

First ~~do~~ MSB as it is
then ~~do~~ do \oplus -or operation of 1st
and 2nd bit (from MSB side)
result will be your 2nd bit
of gray code .. same process

(a)

$$(0100)_2 = (0110)_G$$

$$(0100)_2 = (0110)_G$$