

Ans-(i) The given transition table:-

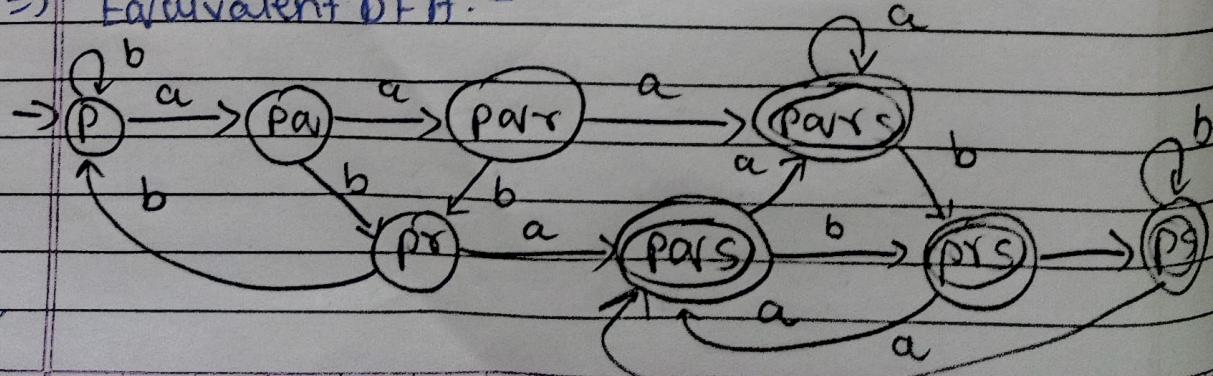
State \ Σ	a	b	Final state: S initial state: P
P	{p, a τ }	p	
a τ	τ	τ	
r	s	ϕ	
s	s	s	

=> This is clearly NFA. So, we need to make transition table for DFA with the help of above transition.

=> Initial state: P.

State \ Σ	a	b
\rightarrow P	[pa]	p
[pa τ]	[pa τ s]	[p τ s]
[pa τ s]	[pa τ ss]	[p τ s]
[p τ s]	[pa τ s]	p
*[pa τ s]	[pa τ ss]	[p τ s]
*[pa τ ss]	[pa τ ss]	[p τ ss]
*[p τ ss]	[p τ ss]	[ps]
*[ps]	[pa τ s]	[ps]

=> Equivalent DFA:-



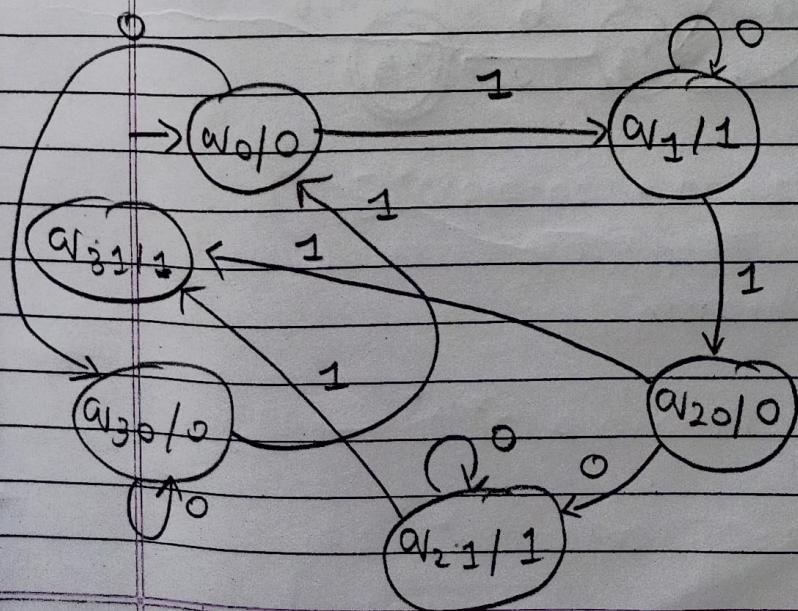
Ans-(2)

* Convert mealy to moore.

⇒ From the given transition table, we will construct transition table for Moore Machine.

Present State	Next State		Output
	$a=0$	$a=1$	
v_0	v_3	v_1	0
v_1	v_1	v_0	1
v_2	v_{21}	v_{31}	0
v_{21}	v_{21}	v_{31}	1
v_3	v_{30}	v_0	0
v_{31}	v_{30}	v_0	1

* moore machine transition diagram:-



Ans-(3) From the given DFA table, we will minimize it by writing the equivalence form of it.

$$\pi_0 = \{v_0, v_1, v_2, v_3, v_5\} \setminus \{v_4\}$$

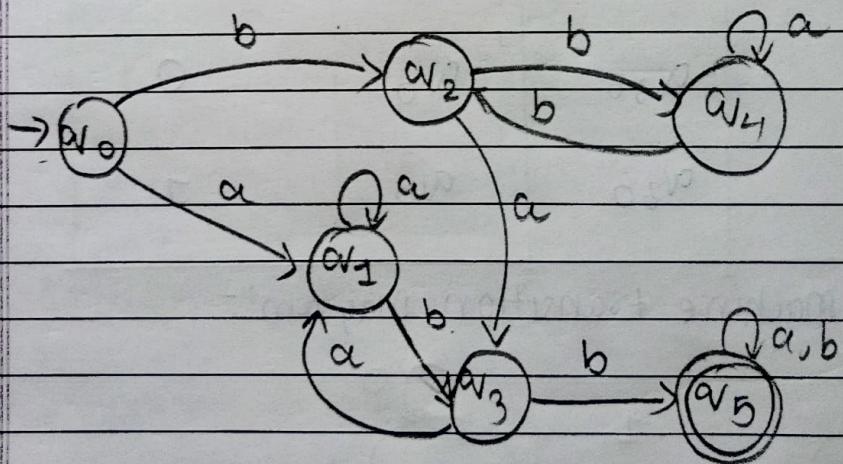
$$\Rightarrow \pi_1 = \{v_0, v_1, v_2, v_4\} \setminus \{v_3\} \setminus \{v_5\}$$

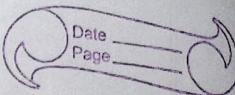
$$\Rightarrow \pi_2 = \{v_0, v_1, v_4\}, \{v_2\}, \{v_3\}, \{v_5\}$$

$$\Rightarrow \pi_3 = \{v_0, v_4\}, \{v_1\}, \{v_2\}, \{v_3\}, \{v_5\}$$

$$\Rightarrow \pi_4 = \{v_0\}, \{v_1\}, \{v_2\}, \{v_3\}, \{v_4\}, \{v_5\}$$

\Rightarrow Transition diagram of minimized DFA:-





Ans - (4)

$$\alpha v_1 = \varepsilon + \alpha v_1 0 + \alpha v_4 0 + \alpha v_3 0 \quad (1)$$

$$\alpha v_2 = \alpha v_2 1 + \alpha v_4 1 + \alpha v_1 1 \quad (2)$$

$$\alpha v_3 = \alpha v_2 0 \quad (3)$$

$$\alpha v_4 = \alpha v_2 1 \quad (4)$$

$$\Rightarrow \text{By (3) & (4); } \alpha v_4 = \alpha v_2 0 1 \quad (5)$$

 \Rightarrow In eqn - (1);

$$\alpha v_1 = \varepsilon + \alpha v_2 0 0 + \alpha v_1 0 + \alpha v_2 0 1 0$$

$$\alpha v_1 = \underbrace{\varepsilon + \alpha v_2 (00 + 010)}_{Q}, \underbrace{\alpha v_1 0}_{R P}$$

 \Rightarrow By applying Arden's theorem;

$$\alpha v_1 = (\varepsilon + \alpha v_2 (00 + 010)) 0^* \quad (6)$$

 \Rightarrow In eqn (2);

$$\alpha v_2 = \alpha v_2 1 + \alpha v_2 0 1 + [\varepsilon + \alpha v_2 (00 + 010)] 0^*$$

$$\alpha v_2 = \alpha v_2 1 + \alpha v_2 0 1 + \varepsilon 0^* + \alpha v_2 (00 + 010) 0^*$$

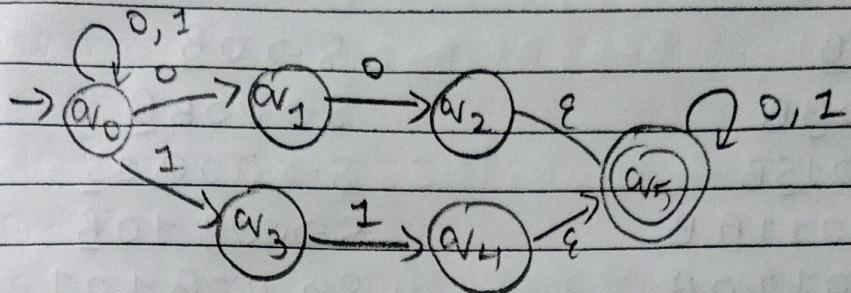
$$\alpha v_2 = 0^* + \alpha v_2 [1 + 01 + (00 + 010) 0^*]$$

 \rightarrow from Arden's theorem;

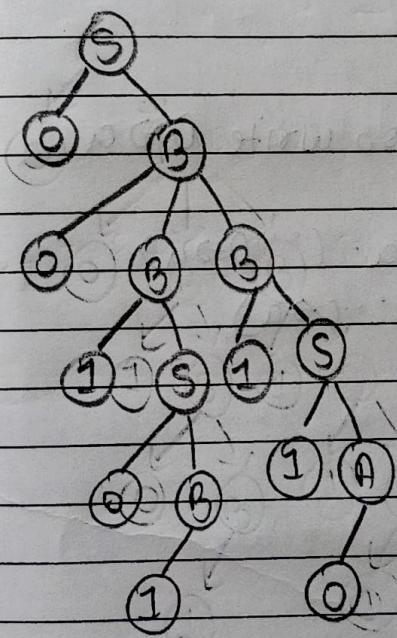
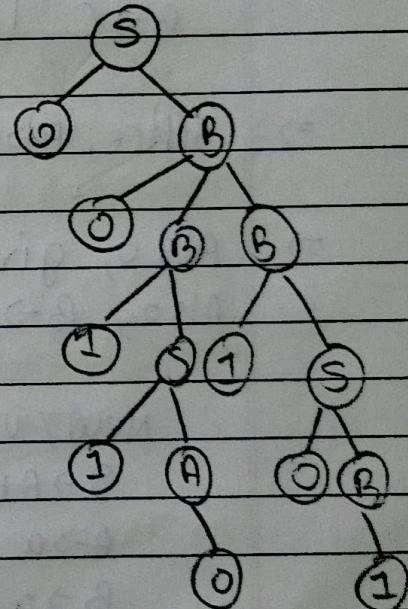
$$\alpha v_2 = 0^* (1 + 01 + (00 + 010) 0^*)^* \quad (7)$$

$$\therefore RF = 0^* (1 + 01 + (00 + 010) 0^*)^* 0 1$$

Ans-(5)

Given RE: $(0+1)^*(00+11)(0+1)^*$ 

Ans-(6)

 $S \rightarrow 0B|1A$ $A \rightarrow 0|0S|1AA$ $B \rightarrow 1|1S|0BB$ Left DerivationRight Derivation

Leftmost

$S \rightarrow OB$
 $S \rightarrow OOB_B$
 $S \rightarrow OOB_S B$
 $S \rightarrow OOB_1 A B$
 $S \rightarrow OOB_1 B$
 $S \rightarrow OOB_1 S$
 $S \rightarrow OOB_1 10B$
 $S \rightarrow OOB_1 101$
 $S \rightarrow OOB_1 101$

Right most

$S \rightarrow OB$
 $S \rightarrow OOB_B$
 $S \rightarrow OOB_1 S$
 $S \rightarrow OOB_1 10B$
 $S \rightarrow OOB_1 101$
 $S \rightarrow OOB_1 101$
 $S \rightarrow OOB_1 101$

Ans - 18) \equiv	$S \rightarrow AB \quad C \rightarrow D$ $A \rightarrow a \quad D \rightarrow E$ $B \rightarrow C \quad E \rightarrow a$
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\Rightarrow since, $E \rightarrow a \therefore$ we can also write $D \rightarrow a$

\Rightarrow Also; given $C \rightarrow D \therefore C \rightarrow a (\because D \rightarrow a)$

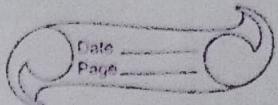
Also, $B \rightarrow C \therefore B \rightarrow a (\because C \rightarrow a)$

\therefore New / Updated Rules:

$S \rightarrow AB$ $A \rightarrow a$ $B \rightarrow a$	$C \rightarrow a$ $D \rightarrow a$ $E \rightarrow a$
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\Rightarrow AS; $S \rightarrow AB; \therefore C, D, E$ are useless symbol

$S \rightarrow AB$
 $A \rightarrow a$
 $B \rightarrow a$



Ans-(a)

$$\begin{array}{l|l} S \rightarrow aAD & B \rightarrow b \\ A \rightarrow aB/bAB & D \rightarrow d \end{array}$$

→ Taking $X_1 \rightarrow aA$; $X_2 \rightarrow bA$; $X_3 \rightarrow a$

$$\therefore S \rightarrow X_1 D$$

$$A \rightarrow X_3 B \mid X_2 B \quad (\because \text{CNF Form})$$

$$B \rightarrow b$$

$$D \rightarrow d \quad | \quad X_2 \rightarrow bA$$

$$X_1 \rightarrow aA \quad | \quad X_3 \rightarrow a$$

Ans-(b) $S \rightarrow AA|a$

$$A \rightarrow SS|b$$

→ Replace S with A_1 | $A_1 \rightarrow A_2 A_2 | a$
 Replace A with A_2 | $A_2 \rightarrow A_1 A_1 | b$

A_2 can be written as:

$$A_2 \rightarrow A_2 A_2 A_1 | b | a A_1 \quad (\because A_1 \rightarrow A_2 A_2)$$

- Here, the first production having left recursion
 → We will remove left Recursion

$$A_2 \rightarrow A_2 A_2 A_1$$

$$\therefore \text{let } Z_2 \rightarrow A_2 A_1 | A_2 A_1 Z_2$$

$A_2 \rightarrow b | a A_1$ are in CNF

Now, $A_2 \rightarrow A_2 A_2 A_1$

$A_2 \rightarrow bA_2 A_1 | aA_1 A_2 A_1$

$A_2 \rightarrow bZ_2 | aA_1 A_2 Z_2$

$\therefore A_2 \rightarrow bA_2 b | aA_1 bZ_2 | aA_1 Z_2$

Now, for $A_1 \rightarrow A_2 A_2 | a$

$\Rightarrow A_1 \rightarrow bA_2 | aA_1 A_2 | bZ_2 A_2 | aA_1 A_2 Z_2$

Final production list: -

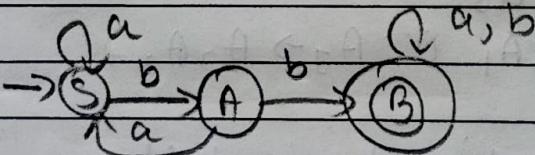
$A \cdot A_1 \rightarrow bA_2 | aA_1 A_2 | bZ_2 A_2 | aA_1 Z_2 A_2 | a$

$A_2 \rightarrow bA_1 | bZ_2 | aA_1 Z_2$

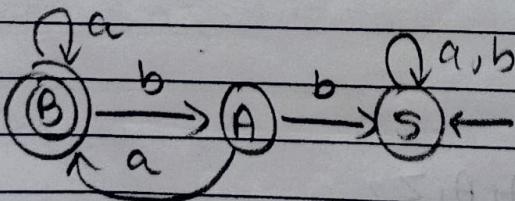
$Z_2 \rightarrow bA_1 | aA_1 A_1 | bZ_2 A_1 | aA_1 Z_2 A_1 | bA_2 A_1 Z_2 |$
 $aA_1 A_1 Z_2 | bZ_2 A_1 Z_2 | aA_1 A_1 Z_2 Z_2$

Ans-(iii) $S \rightarrow aS | bA \quad | \quad B \rightarrow aB | bB | \epsilon$
 $A \rightarrow aS | bB$

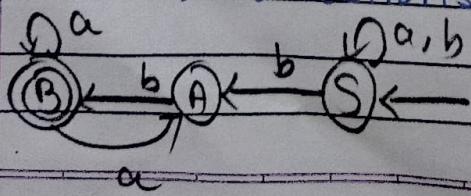
S1:

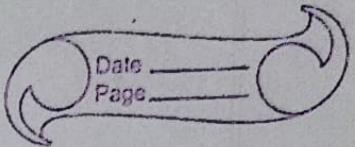


S2: Interchange initial & final state



S3: Reverse directions





2) Final left linear grammar: -

$$S \rightarrow aS \mid bS \mid bA$$

$$A \rightarrow bB$$

$$B \rightarrow ab \mid aA \mid \epsilon$$

Ans(12)
=

$$S \rightarrow (L) \mid \epsilon$$

$$L \rightarrow LSIS$$

One production has left recursion.

$$L \rightarrow LS \mid S \overline{A} \overline{d}$$

$$L \rightarrow SA^I$$

$$A^I \rightarrow SA^I \mid \epsilon$$

(: for form $A \rightarrow A\alpha \mid B$)

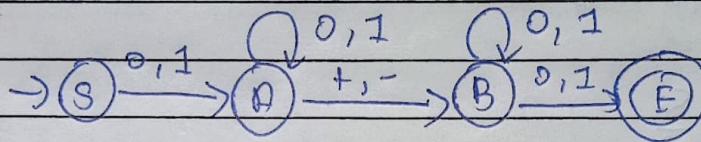
$$A \rightarrow BA^I$$

$$A \rightarrow \alpha A^I \mid \epsilon$$

Ans $\Rightarrow S \rightarrow iE + S1 | iE + S e/a, E \rightarrow b$

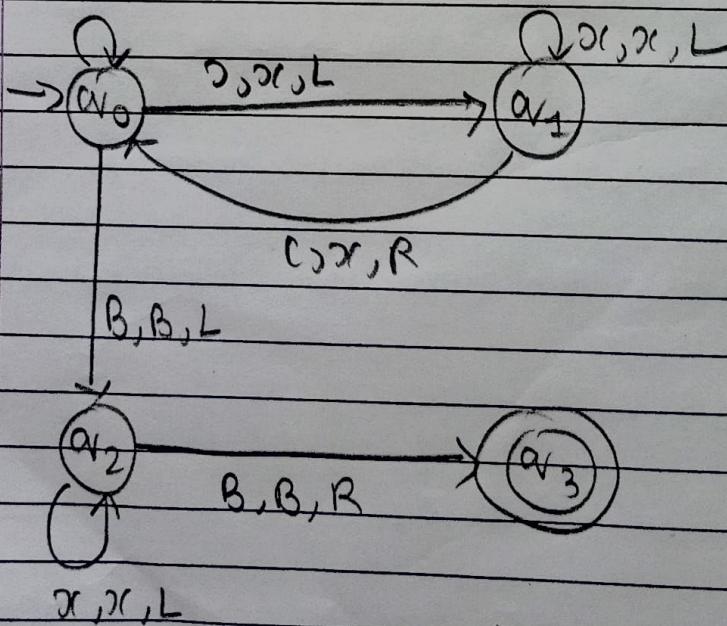
$S \rightarrow iE + S S' / a$
 $S1 \rightarrow e / eS$
 $E \rightarrow b$

Ans $\Rightarrow S \rightarrow OA | 1A$
 $A \rightarrow OA | 1A | +B | -B$
 $B \rightarrow OB | 1B | O | 1$



Ans $\Rightarrow \dots B C C D C D D C D B \dots$

\Rightarrow Turing machine :-



- $\Rightarrow S(aV_0, C) = (aV_0, C, R)$
 $S(aV_0, X) = (aV_1, X, L)$
 $S(aV_1, C) = (aV_0, X, R)$
 $S(aV_0, X) = (aV_1, X, R)$
 $S(aV_1, X) = (aV_1, X, L)$
 $S(aV_0, Y) = (aV_2, Y, L)$
 $S(aV_2, X) = (aV_2, X, L)$
 $S(aV_2, Y) = (aV_2, Y, H)$
 $S(aV_2, C) = (aV_2, N, H)$
 $S(aV_1, Y) = (aV_1, W, H)$

Ans) Constructing turing machine odd palindrome over $\{a, b\}$.

