

LEACH protocol and LEACH-C protocol simulation using Contiki-NG

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Project Goal

To compare the performance of the original LEACH protocol to one of the LEACH successor protocols

Used network simulator:

Contiki-NG in Cooja

Chosen LEACH successor protocol:

LEACH-C (LEACH Centralized)

LEACH Protocol

- It is an energy-efficient communication protocol for wireless microsensor networks.
- A clustering-based protocol that choose different CH randomly every round. This protocol evenly distributes the energy load among the sensors in the network.

LEACH-C (Centralized)

- Produces excellent clusters by scattering the CH throughout the sensor network
- Since set-up phase is completely executed at the edge gateway (decides which nodes become the CH at the first advertisement phase).
- The steady-state phase of LEACH-C is the same as the original LEACH protocol.

Motivation

LEACH is the first protocol proposing the idea of the clustering-based routing for wireless sensor networks

However...

LEACH has a problem where it does not guarantee the position and number of the CHs in each round

Therefore...

By comparing the succeeded model of LEACH to the original LEACH using a network simulator, visualize how the protocol has improved in data freshness, overhead, and energy efficiency.

Related Work (Papers)

- LEACH (2000)
- LEACH-C (2002)
- "A Survey on Successors of LEACH Protocol,"
- " Clustering objectives in wireless sensor networks: A survey and research direction analysis “
- HEED (2004)
- PEGASIS (2016)

Approach- LEACH Specifications

- All nodes are homogeneous, so it runs the same program
- Only the edge gateway (sink) runs its own program
- All nodes are one hop away from the edge gateway, so it only uses the one-hop unicast and broadcast
- NullNet for communication
- 100 nodes and 1 gateway
- Probability $P = 0.05$
- Transmission range: 150 m

Approach- LEACH Implementation

1. The edge gateway sends out a broadcast advertisement to all nodes within the range (Only once)

```
00:00.458      ID:55      [INFO: LEACH Node] [0] Received packet: address 101
```

2. Starts round: 0
3. Cluster heads are decided

```
1:45:00.501    ID:99      [INFO: LEACH_Node] [0] Cluster head process  
1:45:00.501    ID:99      [INFO: LEACH_Node] random value: 0.050000 0.020432  
1:45:00.501    ID:99      [INFO: LEACH_Node] [0] is_ch: 1  
1:45:00.501    ID:99      [INFO: LEACH Node] [0] Sending out broadcast
```


Approach- LEACH Implementation

LEACH

$$T(n) = \begin{cases} \frac{P}{1 - P * (r \bmod \frac{1}{P})} & \text{if } n \in G \\ 0 & \text{otherwise} \end{cases}$$

LEACH-C

$$K = \sqrt{\frac{N}{2\Pi} \frac{\epsilon_{fs}}{\epsilon_{mp}} \frac{M}{d_{toBS}^2}} \quad P_i(t) = \begin{cases} \frac{k}{N - k * (r \bmod \frac{N}{k})} : & C_i(t) = 1 \\ 0 : & C_i(t) = 0 \end{cases} .$$

Approach- LEACH Implementation

4. CHs send out the advertisement broadcast, and non-CHs keep the address of the CH with the strongest signal(RSSI)

```
1:45:00.537 ID:5 [INFO: LEACH_Node] [0] Receiving advertisement broadcast from 99
```

5. After receiving all the advertisements, non-CHs send out a unicast message to the strongest neighbor CH

```
1:50:00.062 ID:85 [INFO: LEACH_Node] [0] Sending unicast response
```

```
1:50:00.062 ID:85 [INFO: LEACH_Node] [0] sent out unicast to 99
```

Approach- LEACH Implementation

6. Upon receiving unicast from the neighbors, the CH will create a TDMA schedule for the non-CH cluster members

```
2:00:00.501 ID:99 [INFO: LEACH_Node] [0] Sending TDMA schedule
```

7. The non-CH will receive the broadcast only from their strongest neighbor CHs, and start sending data using the received TDMA schedule

```
2:10:00.062 ID:85 [INFO: LEACH_Node] [0] Sending with TDMA process
```

Approach- LEACH Implementation

8. After CH receives all the data from its cluster members, it will fuse the data and send it to the edge gateway

```
2:20:00.690 ID:26 [INFO: LEACH_Node] [0] Sending fused data
```

9. Edge gateway will receive the data

```
2:20:00.717 ID:101 [INFO: App ] Received data from cluster head: 26
```

10. All the collected data will be freed

```
2:36:40.062 ID:85 [INFO: LEACH_Node] [0] Freeing cluster neighbor information
```

Experimental Methodology

- This experiment will be done multiple times using different random seeds and placement.
- It will collect the energy consumption for 100 rounds and compare the mean of the observed values.
- The energy consumption will be observed using the energest API from Contiki-NG
- By comparing the total seconds used for listening to the radio, transmission time, and node sleep time.

Results

After running 100 times:

listen/total time	transmit/total time	off/total time
98.9114%	0.0001%	1.0886%
169:16:40.013 ID:69	Energest:	
169:16:40.013 ID:69	CPU	606300000000s LPM 0s DEEP LPM 0s Total time 606300000000s
169:16:40.013 ID:69	Radio LISTEN	599699602848s TRANSMIT 374152s OFF 6600023000s

It is listening during the time it should be sleeping, so it is not following the LEACH concept of nodes sleeping outside of their transmission time.

As I have not implemented LEACH-C protocol, I do not know the result for comparing between LEACH and LEACH-C performances.

Conclusion/Future Work

- Cooja was good for implementing LEACH, because it had an API for NullNet, which allowed simple communication between the nodes (unlike IPv6 which operates with the MAC layer)
- I would like to implement the correct energy calculation closer to the approach done by the paper
- Fix the LEACH sleep function
- Implement the LEACH-C protocol for comparison

Citations

- Amin Shahraki, Amir Taherkordi, Øystein Haugen, Frank Eliassen, Clustering objectives in wireless sensor networks: A survey and research direction analysis, Computer Networks, Volume 180, 2020, 107376, ISSN 1389-1286, <https://doi.org/10.1016/j.comnet.2020.107376>Links to an external site..
(<https://www.sciencedirect.com/science/article/pii/S1389128620303121>Links to an external site.)
- . W. R. Heinzelman, A. Chandrakasan and H. Balakrishnan, "Energy-efficient communication protocol for wireless microsensor networks," Proceedings of the 33rd Annual Hawaii International Conference on System Sciences, Maui, HI, USA, 2000, pp. 10 pp. vol.2-, doi: 10.1109/HICSS.2000.926982.
- . S. K. Singh, P. Kumar and J. P. Singh, "A Survey on Successors of LEACH Protocol," in IEEE Access, vol. 5, pp. 4298-4328, 2017, doi: 10.1109/ACCESS.2017.2666082.
- W. B. Heinzelman, A. P. Chandrakasan and H. Balakrishnan, "An application-specific protocol architecture for wireless microsensor networks," in IEEE Transactions on Wireless Communications, vol. 1, no. 4, pp. 660-670, Oct. 2002, doi: 10.1109/TWC.2002.804190.
- keywords: {Wireless application protocol;Microsensors;Energy efficiency;Delay;Signal processing algorithms;Remote monitoring;Intelligent networks;Robustness;Wireless communication;Access protocols},

