Efficient Synchronization of Linux Memory Regions over a Network: A Comparative Study and Implementation (Notes)

A user-friendly approach to application-agnostic state synchronization

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Structure

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Introduction

- Memory management in Linux
- Memory as the universal storage API
- What would be possible if memory would be the universal way to access resources?
- Why efficient memory synchronization is the missing key component
- High-level use cases for memory synchronization in the industry today
- Pull-Based Memory Synchronization with userfaultfd
 - · Plain language description of userfaultfd (what are page faults)
 - Exploring an alternative method by handling page faults using signals
 - Handlers and registration
 - History of userfaultfd
 - · Allocating the shared region
 - · Maximum shared region size is limited by available physical memory
 - Transferring handler sockets between processes
 - · Implementing userfaultfd bindings in Go

Content

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- Pull-Based Memory Synchronization with userfaultfd
 - Page faults occur when a process tries to access a memory region that has not yet been mapped into a process' address space
 - By listening to these page faults, we know when a process wants to access a specific piece of memory
 - We can use this to then pull the chunk of memory from a remote, map it to the address on which the page fault occured, thus only fetching data when it is required
 - · Usually, handling page faults is something that the kernel does
 - In our case, we want to handle page faults in userspace
 - In the past, this used to be possible by handling the SIGSEGV signal in the process
 - In our case however, we can use a recent system called userfaultfd to do this in a more elegant way (available since kernel 4.11)
 - userfaultfd allows handling these page faults in userspace
 - Implementing this in Go was quite tricky, and it involves using unsafe