

## SMART CONTRACT AUDIT REPORT

for

PolkaCipher

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PeckShield May 21, 2021

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## 1 Introduction

Given the opportunity to review the design document and related source code of the PolkaCipher protocol and the associated token contract, we outline in the report our systematic method to evaluate potential security issues in the smart contract implementation, expose possible semantic inconsistency between smart contract code and the documentation, and provide additional suggestions or recommendations for improvement. Our results show that the given version of smart contracts can be further improved due to the presence of several issues related to either security or performance. This document outlines our audit results.

#### 1.1 About PolkaCipher

PolkaCipher is a cross-chain privacy preserving project in the Polkadot ecosystem. The primary aim is to expand the applications of NFTs & DeFi and related use cases in business settings as well as to increase the reach of Web3 applications in blockchain economies via a seamless integration with PolkaCipher. To achieve that, there is a need to bridge the gap between an off-chain Internet economy and on-chain blockchain decentralized applications using PolkaCipher that relies on verified vendors to provide accurate and trusted data. By design, PolkaCipher, once ready, will be community owned and fully decentralized.

The basic information of PolkaCipher is as follows:

Table 1.1: Basic Information of PolkaCipher

Item Description

ltem	Description
Issuer	PolkaCipher
Туре	Ethereum ERC20 Token Contract
Platform	Solidity
Audit Method	Whitebox
Audit Completion Date	May 21, 2021

In the following, we show the Git repository and the commit hash value used in this audit:

https://github.com/polkacipher/Contracts.git (c79c731)

And this is the commit ID after all fixes for the issues found in the audit have been checked in:

https://github.com/polkacipher/Contracts.git (b936867)

#### 1.2 About PeckShield

PeckShield Inc. [7] is a leading blockchain security company with the goal of elevating the security, privacy, and usability of current blockchain ecosystem by offering top-notch, industry-leading services and products (including the service of smart contract auditing). We are reachable at Telegram (https://t.me/peckshield), Twitter (http://twitter.com/peckshield), or Email (contact@peckshield.com).

### 1.3 Methodology

To standardize the evaluation, we define the following terminology based on OWASP Risk Rating Methodology [6]:

- <u>Likelihood</u> represents how likely a particular vulnerability is to be uncovered and exploited in the wild;
- Impact measures the technical loss and business damage of a successful attack;
- Severity demonstrates the overall criticality of the risk;

Likelihood and impact are categorized into three ratings: *H*, *M* and *L*, i.e., *high*, *medium* and *low* respectively. Severity is determined by likelihood and impact and can be classified into four categories accordingly, i.e., *Critical*, *High*, *Medium*, *Low* shown in Table 1.2.

We perform the audit according to the following procedures:

- <u>Basic Coding Bugs</u>: We first statically analyze given smart contracts with our proprietary static code analyzer for known coding bugs, and then manually verify (reject or confirm) all the issues found by our tool.
- <u>ERC20 Compliance Checks</u>: We then manually check whether the implementation logic of the audited smart contract(s) follows the standard ERC20 specification and other best practices.
- Additional Recommendations: We also provide additional suggestions regarding the coding and development of smart contracts from the perspective of proven programming practices.

High Critical High Medium

High Medium

Low

Medium Low

High Medium

Low

High Medium

Low

Likelihood

Table 1.2: Vulnerability Severity Classification

To evaluate the risk, we go through a list of check items and each would be labeled with a severity category. For one check item, if our tool does not identify any issue, the contract is considered safe regarding the check item. For any discovered issue, we might further deploy contracts on our private testnet and run tests to confirm the findings. If necessary, we would additionally build a PoC to demonstrate the possibility of exploitation. The concrete list of check items is shown in Table 1.3.

#### 1.4 Disclaimer

Note that this security audit is not designed to replace functional tests required before any software release, and does not give any warranties on finding all possible security issues of the given smart contract(s) or blockchain software, i.e., the evaluation result does not guarantee the nonexistence of any further findings of security issues. As one audit-based assessment cannot be considered comprehensive, we always recommend proceeding with several independent audits and a public bug bounty program to ensure the security of smart contract(s). Last but not least, this security audit should not be used as investment advice.

Table 1.3: The Full List of Check Items

Category	Check Item
	Constructor Mismatch
	Ownership Takeover
	Redundant Fallback Function
	Overflows & Underflows
	Reentrancy
	Money-Giving Bug
	Blackhole
	Unauthorized Self-Destruct
Basic Coding Bugs	Revert DoS
Dasic Coung Dugs	Unchecked External Call
	Gasless Send
	Send Instead of Transfer
	Costly Loop
	(Unsafe) Use of Untrusted Libraries
	(Unsafe) Use of Predictable Variables
	Transaction Ordering Dependence
	Deprecated Uses
	Approve / TransferFrom Race Condition
ERC20 Compliance Checks	Compliance Checks (Section 3)
	Avoiding Use of Variadic Byte Array
	Using Fixed Compiler Version
Additional Recommendations	Making Visibility Level Explicit
	Making Type Inference Explicit
	Adhering To Function Declaration Strictly
	Following Other Best Practices

# 2 | Findings

#### 2.1 Summary

Here is a summary of our findings after analyzing the PolkaCipher. During the first phase of our audit, we study the smart contract source code and run our in-house static code analyzer through the codebase. The purpose here is to statically identify known coding bugs, and then manually verify (reject or confirm) issues reported by our tool. We further manually review business logics, examine system operations, and place ERC20-related aspects under scrutiny to uncover possible pitfalls and/or bugs.

Severity	# of Findings
Critical	1
High	1
Medium	3
Low	2
Informational	0
Total	7

Moreover, we explicitly evaluate whether the given contracts follow the standard ERC20 specification and other known best practices, and validate its compatibility with other similar ERC20 tokens and current DeFi protocols. The detailed ERC20 compliance checks are reported in Section 3. After that, we examine a few identified issues of varying severities that need to be brought up and paid more attention to. (The findings are categorized in the above table.) Additional information can be found in the next subsection, and the detailed discussions are in Section 4.

#### 2.2 Key Findings

Overall, no ERC20 compliance issue was found and our detailed checklist can be found in Section 3. Note that the smart contract implementation can be improved by resolving the identified issues (shown in Table 2.1), including 1 critical-severity vulnerability, 1 high-severity vulnerability, 3 medium-severity vulnerabilities, and 2 low-severity vulnerabilities.

ID	Severity	Title	Category	Status
PVE-001	Low	Susceptible Replays Against Multisig-	Coding Practices	Fixed
		Wallet::execute()		
PVE-002	High	Denial-Of-Service Against MultisigWal-	Coding Practices	Fixed
		let::execute()		
PVE-003	Medium	Reentrancy Risk In MultisigWal-	Coding Practices	Fixed
		let::execute()		
PVE-004	Critical	Funds Can Be Fully Drained From Vest	Security Features	Fixed
PVE-005	Low	safeTransfer()/safeTransferFrom() Re-	Coding Practices	Fixed
		placement		
PVE-006	Medium	Logic Error of getClaimable()	Business Logic	Fixed
PVE-007	Medium	Assumed Trust on Admin Keys	Security Features	Fixed

Table 2.1: Key PolkaCipher Audit Findings

As a common suggestion, due to the fact that compiler upgrades might bring unexpected compatibility or inter-version consistencies, it is always preferred to use fixed compiler versions whenever possible. As an example, we highly encourage to explicitly indicate the Solidity compiler version, e.g., pragma solidity 0.8.0; instead of pragma solidity ^0.8.0;.

Besides recommending specific countermeasures to mitigate these issues, we emphasize that it is always important to develop necessary risk-control mechanisms and make contingency plans, which may need to be exercised before the mainnet deployment. The risk-control mechanisms need to kick in at the very moment when the contracts are being deployed in mainnet. Please refer to Section 3 for our detailed compliance checks and Section 4 for elaboration of reported issues.

# 3 | ERC20 Compliance Checks

The ERC20 specification defines a list of API functions (and relevant events) that each token contract is expected to implement (and emit). The failure to meet these requirements means the token contract cannot be considered to be ERC20-compliant. Naturally, as the first step of our audit, we examine the list of API functions defined by the ERC20 specification and validate whether there exist any inconsistency or incompatibility in the implementation or the inherent business logic of the audited contract(s).

Table 3.1: Basic View-Only Functions Defined in The ERC20 Specification

Item	Description	Status
name()	Is declared as a public view function	✓
name()	Returns a string, for example "Tether USD"	✓
symbol()	Is declared as a public view function	✓
Symbol()	Returns the symbol by which the token contract should be known, for	✓
	example "USDT". It is usually 3 or 4 characters in length	
decimals()	Is declared as a public view function	✓
uecimais()	Returns decimals, which refers to how divisible a token can be, from 0	✓
	(not at all divisible) to 18 (pretty much continuous) and even higher if	
	required	
totalSupply()	Is declared as a public view function	✓
total Supply()	Returns the number of total supplied tokens, including the total minted	✓
	tokens (minus the total burned tokens) ever since the deployment	
balanceOf()	Is declared as a public view function	✓
balanceO1()	Anyone can query any address' balance, as all data on the blockchain is	✓
	public	
allowance()	Is declared as a public view function	1
anowance()	Returns the amount which the spender is still allowed to withdraw from	✓
	the owner	

Our analysis shows that there is no ERC20 inconsistency or incompatibility issue found in the audited PolkaCipher. In the surrounding two tables, we outline the respective list of basic viewonly functions (Table 3.1) and key state-changing functions (Table 3.2) according to the widely-

Table 3.2: Key State-Changing Functions Defined in The ERC20 Specification

ltem	Description	Status
	Is declared as a public function	✓
	Returns a boolean value which accurately reflects the token transfer status	✓
tuomafau()	Reverts if the caller does not have enough tokens to spend	<b>√</b>
transfer()	Allows zero amount transfers	✓
	Emits Transfer() event when tokens are transferred successfully (include 0	✓
	amount transfers)	
	Reverts while transferring to zero address	<b>✓</b>
	Is declared as a public function	✓
	Returns a boolean value which accurately reflects the token transfer status	✓
	Reverts if the spender does not have enough token allowances to spend	✓
	Updates the spender's token allowances when tokens are transferred suc-	✓
transferFrom()	cessfully	
	Reverts if the from address does not have enough tokens to spend	✓
	Allows zero amount transfers	✓
	Emits Transfer() event when tokens are transferred successfully (include 0	✓
	amount transfers)	
	Reverts while transferring from zero address	✓
	Reverts while transferring to zero address	✓
	Is declared as a public function	/
	Returns a boolean value which accurately reflects the token approval status	<b>√</b>
approve()	Emits Approval() event when tokens are approved successfully	<b>√</b>
	Reverts while approving to zero address	<b>√</b>
Tue n efe n()	Is emitted when tokens are transferred, including zero value transfers	<b>√</b>
Transfer() event	Is emitted with the from address set to $address(0x0)$ when new tokens	<b>√</b>
	are generated	
Approval() event	Is emitted on any successful call to approve()	<b>√</b>

adopted ERC20 specification. In addition, we perform a further examination on certain features that are permitted by the ERC20 specification or even further extended in follow-up refinements and enhancements (e.g., ERC777/ERC2222), but not required for implementation. These features are generally helpful, but may also impact or bring certain incompatibility with current DeFi protocols. Therefore, we consider it is important to highlight them as well. This list is shown in Table 3.3.

Table 3.3: Additional Opt-in Features Examined in Our Audit

Feature	Description	Opt-in
Deflationary	Part of the tokens are burned or transferred as fee while on trans-	_
	fer()/transferFrom() calls	
Rebasing	The balanceOf() function returns a re-based balance instead of the actual	_
	stored amount of tokens owned by the specific address	
Pausable	The token contract allows the owner or privileged users to pause the token	_
	transfers and other operations	
Blacklistable	The token contract allows the owner or privileged users to blacklist a	
	specific address such that token transfers and other operations related to	
	that address are prohibited	
Mintable	The token contract allows the owner or privileged users to mint tokens to	✓
	a specific address	
Burnable	The token contract allows the owner or privileged users to burn tokens of	_
	a specific address	

## 4 Detailed Results

### 4.1 Susceptible Replays Against MultisigWallet::execute()

• ID: PVE-001

• Severity: Low

Likelihood: Low

Impact: Medium

• Target: MultisigWallet

• Category: Coding Practices [5]

CWE subcategory: CWE-1109 [1]

#### Description

The PolkaCipher protocol has a core MultisigWallet contract that allows approvers to collect their signatures for a given action. The MultisigWallet contract supports three approvers: Once the collected number of (verified) signatures is larger than the required threshold (saved in \_required), the intended action will be executed.

```
function execute(address to, uint256 value, bytes memory data, bytes memory sig1,
37
            bytes memory sig2, bytes memory sig3) public payable returns (bool success) {
38
            bytes32 hash = keccak256(abi.encodePacked(nonce, to, value ,data));
39
            uint8 approvals;
40
            approvals += voters[hash.recover(sig1)];
            approvals += voters[hash.recover(sig2)];
41
42
            approvals += _voters[hash.recover(sig3)];
43
            if (approvals >= required){
44
                (success,) = to.call{value: value}(data);
45
           }
46
            nonce++;
47
```

Listing 4.1: MultisigWallet :: execute()

In particular, we show above the <code>execute()</code> function. The function generates the <code>hash</code> using required fields, i.e., <code>nonce</code>, <code>to</code>, <code>value</code>, and <code>data</code>. However, this calculation assumes all signatures produced by different signers are in the same chain. The absence of a <code>EIP-712 domainSeparator</code> with the <code>EIP-155 chainID</code> in current calculation makes signature validation susceptible to replay across different chains.

Moreover, it is also possible to replay sig1 in sig2 and/or sig3. By doing so, we can readily replay one valid signature to pass the required threshold and execute the transaction.

**Recommendation** Add the EIP-712 domainSeparator with the chainID into calculation. Also, prevent the same signature from being replayed.

Status The issue has been addressed by the following commits: 09d1f0b and b936867.

### 4.2 Denial-Of-Service Against MultisigWallet::execute()

• ID: PVE-002

Severity: High

• Likelihood: Medium

Impact: High

• Target: MultisigWallet

• Category: Coding Practices [5]

• CWE subcategory: CWE-1109 [1]

#### Description

In the MultisigWallet contract, there is a global nonce state that is designed to record the total number of successfully executed transactions. With that, it is used to prevent unwanted replays by being included in the computation of the generated hash for signature validation. However, we notice that nonce can be advanced by anyone even when the signature validation fails. With that, a denial-of-service situation can be introduced to block the transaction execution.

To elaborate, we list below the related <code>execute()</code> function. As mentioned in Section 4.1, this function supports three voters: Once the collected (verified) signatures is larger than the required threshold, the intended transaction will be executed. It comes to our attention that <code>nonce</code> is included in the <code>hash</code> calculation, but it can be advanced regardless of the signature validation result. Therefore, a malicious actor can simply advance to foil the signature validation, hence blocking the transaction execution.

```
function execute(address to, uint256 value, bytes memory data, bytes memory sig1,
37
            bytes memory sig2, bytes memory sig3) public payable returns (bool success) {
38
            bytes32 hash = keccak256(abi.encodePacked(nonce, to, value ,data));
39
            uint8 approvals;
40
            approvals += voters[hash.recover(sig1)];
41
            approvals += _voters[hash.recover(sig2)];
42
            approvals += _voters[hash.recover(sig3)];
            if (approvals >= required){
43
44
                (success,) = to.call{value: value}(data);
45
            }
46
            nonce++;
```

Listing 4.2: MultisigWallet :: execute()

**Recommendation** Move the nonce update inside the if-branch.

**Status** The issue has been addressed by the following commit: 09d1f0b.

### 4.3 Reentrancy Risk In MultisigWallet::execute()

• ID: PVE-003

• Severity: High

• Likelihood: Medium

• Impact: High

• Target: Multiple Contracts

• Category: Coding Practices [5]

• CWE subcategory: CWE-561 [3]

#### Description

A common coding best practice in Solidity is the adherence of checks-effects-interactions principle. This principle is effective in mitigating a serious attack vector known as re-entrancy. Via this particular attack vector, a malicious contract can be reentering a vulnerable contract in a nested manner. Specifically, it first calls a function in the vulnerable contract, but before the first instance of the function call is finished, second call can be arranged to re-enter the vulnerable contract by invoking functions that should only be executed once. This attack was part of several most prominent hacks in Ethereum history, including the DAO [9] exploit, and the recent Uniswap/Lendf.Me hack [8].

We notice there are several occasions the <code>checks-effects-interactions</code> principle is violated. Using the <code>MultisigWallet</code> as an example, the <code>execute()</code> function (see the code snippet below) is provided to externally call a token contract to transfer assets. However, the invocation of an external contract requires extra care in avoiding the above <code>re-entrancy</code>.

Apparently, the interaction with the external contract (line 44) starts before effecting the update on internal states (line 46), hence violating the principle. In this particular case, if the external contract has certain hidden logic that may be capable of launching re-entrancy via the very same execute() function.

```
37
       function execute(address to, uint256 value, bytes memory data, bytes memory sig1,
            bytes memory sig2, bytes memory sig3) public payable returns (bool success) {
38
            bytes32 hash = keccak256(abi.encodePacked(nonce, to, value ,data));
39
            uint8 approvals;
40
            approvals += voters[hash.recover(sig1)];
41
            approvals += voters[hash.recover(sig2)];
            approvals += voters[hash.recover(sig3)];
42
43
            if (approvals >= required){
44
                (success,) = to.call{value: value}(data);
45
           }
46
            nonce++;
47
```

Listing 4.3: MultisigWallet :: execute()

Moreover, there is a need to validate the return boolean success to be true. Also, the initiateVest () routine also shares the same reentrancy risk.

**Recommendation** Validate the action execution result returned in success and add the nonReentrant modifier to prevent reentrancy.

**Status** The issue has been addressed by the following commit: ffa42f7.

### 4.4 Funds Can Be Fully Drained From Vest

ID: PVE-004

Severity: Critical

Likelihood: High

• Impact: High

• Target: Vest

• Category: Security Features [5]

• CWE subcategory: CWE-287 [2]

#### Description

The PolkaCipher protocol has a vesting contract that allows investors to set their own custom vesting arrangements. For each arrangement, the investor is able to specify the entire vesting period, the total amount, the beneficiary, as well as the initial amount of tokens that can be immediately available for withdrawal.

In the following, we show below the responsible initiateVest() function. Our analysis exposes a critical problem in this function: it firstly transfers the entire amount from the investor into this contract for custody. Afterwards, it transfers the initial amount to the intended beneficiary. However, there is a lack of validation on the given initial. A malicious actor can simply set initial = \_token .balanceOf(vest) to drain all the funds in this contract.

```
62
        // v.initiateVest(advisor, amount, 0, 30 days, 180 days);
       function initiateVest (address owner, uint256 amount, uint256 initial, uint256 cliff,
63
            uint256 linear) public virtual returns (bytes32){
64
            token.transferFrom( msgSender(), address(this), amount);
65
           bytes32 index = getVestId(owner, nonce[owner]);
66
            vests[index] = Vest(owner, amount, block.timestamp, initial, cliff, linear, 0);
            vests[index].claimed = vests[index].claimed.add( vests[index].initial);
67
            _token.transfer(_vests[index].owner, _vests[index].initial);
68
69
            nonce [owner]++;
70
           return index;
71
```

Listing 4.4: Vest:: initiateVest ()

**Recommendation** Validate the given input by ensuring initial is always smaller than the amount.

**Status** The issue has been addressed by the following commit: 09d1f0b.

### 4.5 safeTransfer()/safeTransferFrom() Replacement

• ID: PVE-005

Severity: LowLikelihood: Low

• Impact: Medium

• Target: Vest.sol

Category: Coding Practices [5]CWE subcategory: CWE-561 [3]

#### Description

Though there is a standardized ERC-20 specification, many token contracts may not strictly follow the specification or have additional functionalities beyond the specification. In this section, we examine the approve() routine and possible idiosyncrasies from current widely-used token contracts.

In particular, we use the popular stablecoin, i.e., USDT, as our example. We show the related code snippet below. On its entry of approve(), there is a requirement, i.e., require(!((\_value != 0) && (allowed[msg.sender][\_spender] != 0))). This specific requirement essentially indicates the need of reducing the allowance to 0 first (by calling approve(\_spender, 0)) if it is not, and then calling a second one to set the proper allowance. This requirement is in place to mitigate the known approve()/transferFrom() race condition (https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729).

```
194
195
        * @dev Approve the passed address to spend the specified amount of tokens on behalf
            of msg.sender.
196
        * Cparam _spender The address which will spend the funds.
197
        * @param _value The amount of tokens to be spent.
198
        function approve(address spender, uint value) public onlyPayloadSize(2 * 32) {
199
200
201
            // To change the approve amount you first have to reduce the addresses'
202
            // allowance to zero by calling 'approve(_spender, 0)' if it is not
203
             // already 0 to mitigate the race condition described here:
204
             // https://github.com/ethereum/EIPs/issues/20#issuecomment-263524729
205
             require(!(( value != 0) && (allowed[msg.sender][ spender] != 0)));
206
207
             allowed [msg.sender] [ _spender] = _value;
208
             Approval (msg. sender, _spender, _value);
209
```

Listing 4.5: USDT Token Contract

Because of that, a normal call to approve() with a currently non-zero allowance may fail. Moreover, it is important to note that for certain non-compliant ERC20 tokens (e.g., USDT), the transfer () function does not have a return value. However, the IERC20 interface has defined the transfer()

interface with a bool return value. As a result, the call to transfer() may expect a return value. With the lack of return value of USDT's transfer(), the call will be unfortunately reverted.

```
62
         // v.initiateVest(advisor, amount, 0, 30 days, 180 days);
63
        function initiateVest(address owner, uint256 amount, uint256 initial, uint256 cliff,
             uint256 linear) public virtual returns (bytes32){
64
            token.transferFrom( msgSender(), address(this), amount);
65
            bytes32 index = getVestId(owner,_nonce[owner]);
            _vests[index]= Vest(owner, amount, block.timestamp, initial, cliff, linear, 0);
66
            \_vests[index].claimed = \_vests[index].claimed.add(\_vests[index].initial);
67
68
            _token.transfer(_vests[index].owner, _vests[index].initial);
69
            nonce [owner]++;
70
            return index;
71
```

Listing 4.6: Vest:: initiateVest ()

Because of that, a normal call to transfer() is suggested to use the safe version, i.e., safeTransfer (), In essence, it is a wrapper around ERC20 operations that may either throw on failure or return false without reverts. Moreover, the safe version also supports tokens that return no value (and instead revert or throw on failure). Note that non-reverting calls are assumed to be successful. To use this library you can add a using SafeERC20 for IERC20. Similarly, there is a safe version of approve()/transferFrom() as well, i.e., safeApprove()/safeTransferFrom().

**Recommendation** Accommodate the above-mentioned idiosyncrasy about ERC20-related approve()/transfer()/transferFrom().

Status The issue has been addressed by the following commit: ffa42f7.

### 4.6 Logic Error of getClaimable()

• ID: PVE-006

Severity: Medium

Likelihood: High

Impact: Low

• Target: Vest

• Category: Business Logic [5]

CWE subcategory: CWE-561 [3]

#### Description

The vesting contract in PolkaCipher allows users to withdraw vested assets from the contract. It also provides a convenience routine getClaimable() to query current claimable amount. In the following, we examine this getClaimable() implementation.

To elaborate, we show below the getClaimable() routine. It implements a rather straightforward logic by firstly subtracting the initial from the amount, then multiplying the result by timepassed,

and next dividing by linear, and finally subtracting the claimed from it.

```
36
         function getClaimable (address owner, uint256 nonce) public view virtual returns (
             uint256 tokenClaimable) {
37
            bytes32 index = getVestId(owner, nonce);
38
            if (block.timestamp >= vests[index].start.add( vests[index].cliff) && block.
                timestamp<= vests[index].start.add( vests[index].cliff).add( vests[index].</pre>
                linear)) {
39
                uint256 timePassed = block.timestamp.sub(( vests[index].start).add( vests[
                    index].cliff));
40
                tokenClaimable = (( _vests[index].amount-_vests[index].initial).mul(
                    timePassed).div( vests[index].linear)).sub( vests[index].claimed);
41
            } else if (block.timestamp > _vests[index].start.add(_vests[index].cliff).add(
                vests[index].linear)) {
42
                tokenClaimable = vests[index].amount.sub( vests[index].claimed);
43
44
```

Listing 4.7: Vesting :: getClaimable()

It comes to our attention that the claimed has already included the initial, hence effectively subtracting the tokenClaimable by initial twice. It is considered a typical logic error.

Recommendation Add the initial amount back before subtracting claimed.

Status The issue has been addressed by the following commit: a10242f.

### 4.7 Assumed Trust on Admin Keys

ID: PVE-007

Severity: Medium

Likelihood: Low

• Impact: High

• Target: Token.sol

• Category: Security Features [4]

• CWE subcategory: CWE-287 [2]

#### Description

In the PolkaCipher token contract, there is an owner account that plays a critical role in governing and regulating the entire operation and maintenance such as minting a certain amount (1e9 \* 1e18) to the specified manager.

```
function mint(address manager) public virtual onlyOwner{
    _mint(manager,1e9*1e18);
}
```

Listing 4.8: Token::mint()

We note that the above mint() function allows for the owner to mint more tokens into circulation without being capped. We understand the need of the privileged functions for contract operation, but at the same time the extra power to the owner may also be a counter-party risk to the contract users. Therefore, we list this concern as an issue here from the audit perspective and highly recommend making these onlyOwner privileges explicit or raising necessary awareness among contract users.

**Recommendation** Make the list of extra privileges granted to owner explicit to PolkaCipher users.

Status The issue has been addressed by the following commit: 09d1f0b.



# 5 Conclusion

In this security audit, In this security audit, we have examined the design and implementation of the PolkaCipher token contract. During our audit, we first checked all respects related to the compatibility of the ERC20 specification and other known ERC20 pitfalls/vulnerabilities. We then proceeded to examine other areas such as coding practices and business logics. Overall, we identified seven issues of varying severities that were promptly confirmed and fixed by the team. In the meantime, as disclaimed in Section 1.4, we appreciate any constructive feedbacks or suggestions about our findings, procedures, audit scope, etc.



# References

- [1] MITRE. CWE-1109: Use of Same Variable for Multiple Purposes. https://cwe.mitre.org/data/definitions/1109.html.
- [2] MITRE. CWE-287: Improper Authentication. https://cwe.mitre.org/data/definitions/287.html.
- [3] MITRE. CWE-561: Dead Code. https://cwe.mitre.org/data/definitions/561.html.
- [4] MITRE. CWE CATEGORY: 7PK Security Features. https://cwe.mitre.org/data/definitions/ 254.html.
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