

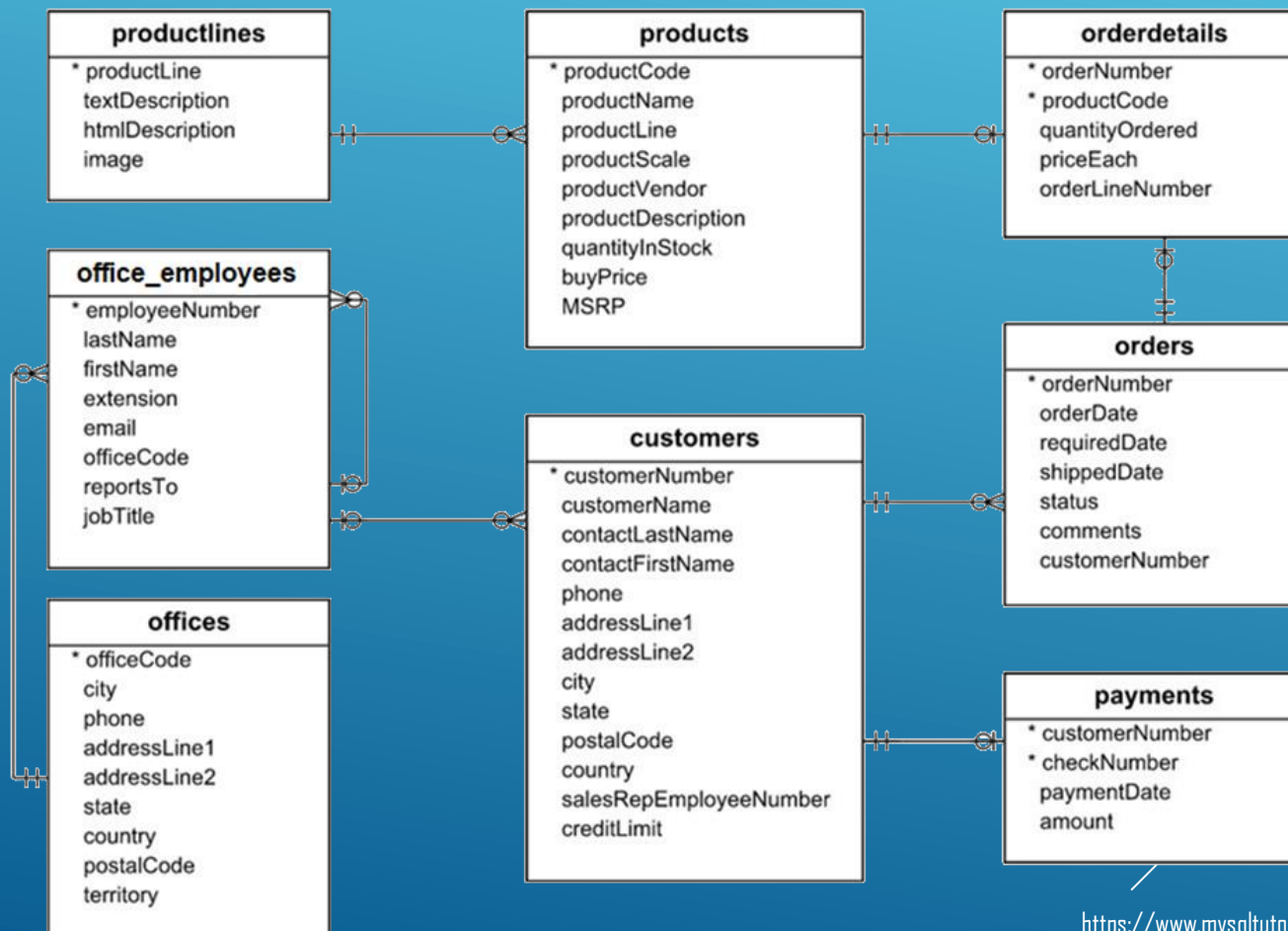
SQL

PART III (JOINS)

SQL

Data Manipulation Language (DML)

Sample Models Schema. Describes an automotive models manufacturer and its sales.



SQL

Data Manipulation Language (DML)

Cross Product (\times). Table aliasing

Return all pairs of student_id's and book_id's

$\pi_{\text{students.student_id, books.book_id}} (\text{students} \times \text{books})$.



SELECT S.student_id, B.book_id **FROM** students S, books B



Except



SELECT S.student_id, B.book_id
FROM students S **CROSS JOIN** books B



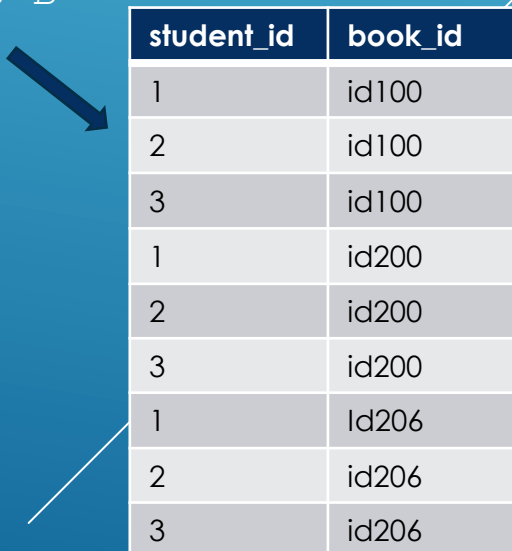
SELECT S.student_id, B.book_id
FROM students S **FULL JOIN** books B

STUDENTS

student_id	name	gender
1	John	M
2	Mike	M
3	Marry	F

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.



student_id	book_id
1	id100
2	id100
3	id100
1	id200
2	id200
3	id200
1	id206
2	id206
3	id206

SQL

Data Manipulation Language (DML)

JOINS ☒

Creating toy schema and data.

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	id206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

```
CREATE TABLE Students(student_id INT, name VARCHAR(20), gender CHAR(1));
CREATE TABLE Borrowed(student_id INT, book_id VARCHAR(20));
CREATE TABLE Books(book_id VARCHAR(20), author VARCHAR(20), title VARCHAR(20));

INSERT INTO Students(student_id, name, gender) VALUES (1, 'John', 'M');
INSERT INTO Students(student_id, name, gender) VALUES (2, 'Adam', 'M');
INSERT INTO Students(student_id, name, gender) VALUES (3, 'Sandra', 'F');

INSERT INTO Books(book_id, author, title) VALUES ('id100', 'Ullman', 'DBMS');
INSERT INTO Books(book_id, author, title) VALUES ('id200', 'Linz', 'Automata');
INSERT INTO Books(book_id, author, title) VALUES ('id206', 'Baader', 'Term Rew.');
```

```
INSERT INTO Borrowed(student_id, book_id) VALUES (1, 'id100');
INSERT INTO Borrowed(student_id, book_id) VALUES (1, 'id200');
INSERT INTO Borrowed(student_id, book_id) VALUES (3, 'id200');
INSERT INTO Borrowed(student_id, book_id) VALUES (1, 'id206');
INSERT INTO Borrowed(student_id, book_id) VALUES (3, 'id200');
```

SQL

Data Manipulation Language (DML)

JOINS ⋈

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	ld206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

Return the list of student names who borrowed books written by "Linz".

SQL

Data Manipulation Language (DML)

JOINS ☒

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	id206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

Return the list of student names who borrowed books written by "Linz".

```
SELECT S.Name
FROM Books B, Borrowed R, Students S
WHERE S.student_id=R.student_id AND
      R.book_id=B.book_id AND
      B.Author ='Linz'
```



name
John
Sandra
Sandra

SQL

Data Manipulation Language (DML)

JOINS ⋈

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	id206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B, Borrowed R, Students S
WHERE S.student_id=R.student_id AND
      R.book_id=B.book_id AND
      B.Author ='Linz'
```



name
John
Sandra

SQL

Data Manipulation Language (DML)

JOINS ⋈ . INNER JOIN

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	id206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B INNER JOIN Borrowed R ON R.book_id=B.book_id
           INNER JOIN Students S ON S.student_id=R.student_id
WHERE B.Author ='Linz'
```



name

John

Sandra

SQL

Data Manipulation Language (DML)

JOINS ⋈. INNER JOIN

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

BORROWED

student_id	book_id
1	id100
1	id200
3	id200
1	id206
3	id200

BOOKS

book_id	author	title
id100	Ullman	DBMS
id200	Linz	Automata
id206	Baader	Term Rew.

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B INNER JOIN Borrowed R USING(book_id)
INNER JOIN Students S USING(student_id)
WHERE B.Author = 'Linz'
```

In case the joining column names is the same (Natural Join)

ORACLE®

MySQL®

name
John
Sandra

SQL

Data Manipulation Language (DML)

JOINS ☒

STUDENTS

student_id	name	gender
1	John	M
2	Adam	M
3	Sandra	F

student_id	book_id	book_id	author	title
1	id100	id100	Ullman	DBMS
1	id200	id200	Linz	Automata
3	id200	id200	Linz	Automata
1	id206	id206	Baader	Term Rew.
3	id200	id200	Linz	Automata

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B INNER JOIN Borrowed R ON R.book_id=B.book_id
      INNER JOIN Students S ON S.student_id=R.student_id
WHERE B.Author ='Linz'
```



SQL

Data Manipulation Language (DML)

JOINS ☒

student_id	name	gender	student_id	book_id	book_id	author	title
1	John	M	1	id100	id100	Ullman	DBMS
1	John	M	1	id200	id200	Linz	Automata
3	Sandra	F	3	id200	id200	Linz	Automata
1	John	M	1	id206	id206	Baader	Term Rew.
3	Sandra	F	3	id200	id200	Linz	Automata

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B INNER JOIN Borrowed R ON R.book_id=B.book_id
      INNER JOIN Students S ON S.student_id=R.student_id
WHERE B.Author ='Linz'
```



SQL

Data Manipulation Language (DML)

JOINS ☒

student_id	name	gender	student_id	book_id	book_id	author	title
1	John	M	1	id200	id200	Linz	Automata
3	Sandra	F	3	id200	id200	Linz	Automata
3	Sandra	F	3	id200	id200	Linz	Automata

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name
FROM Books B INNER JOIN Borrowed R ON R.book_id=B.book_id
           INNER JOIN Students S ON S.student_id=R.student_id
WHERE B.Author ='Linz'
```




SQL

Data Manipulation Language (DML)

JOINS ☒

Return the list of student names who borrowed books written by "Linz".

```
SELECT DISTINCT S.Name  
  FROM Books B INNER JOIN Borrowed R ON R.book_id=B.book_id  
        INNER JOIN Students S ON S.student_id=R.student_id  
 WHERE B.Author ='Linz'
```



name
John
Sandra

SQL

Data Manipulation Language (DML)

JOINS ⋈. Join Types

A

key	value
1	A1
2	A2
3	A3

B

key	value
1	B1
4	B4
5	B5



```
INSERT INTO A(`key`, value) VALUES (1, 'A1');  
INSERT INTO A(`key`, value) VALUES (2, 'A2');  
INSERT INTO A(`key`, value) VALUES (3, 'A3');  
INSERT INTO B(`key`, value) VALUES (1, 'B1');  
INSERT INTO B(`key`, value) VALUES (4, 'B4');  
INSERT INTO B(`key`, value) VALUES (5, 'B5');
```



```
CREATE TABLE A(`key` INT, value VARCHAR(10));  
CREATE TABLE B(`key` INT, value VARCHAR(10));
```



```
CREATE TABLE A([key] INT, value VARCHAR(10));  
CREATE TABLE B([key] INT, value VARCHAR(10));
```



ORACLE



```
CREATE TABLE A("key" INT, value VARCHAR(10));  
CREATE TABLE B("key" INT, value VARCHAR(10));
```



if ANSI mode is of " otherwise

```
SET GLOBAL sql_mode = 'ANSI';  
SET SESSION sql_mode = 'ANSI';
```

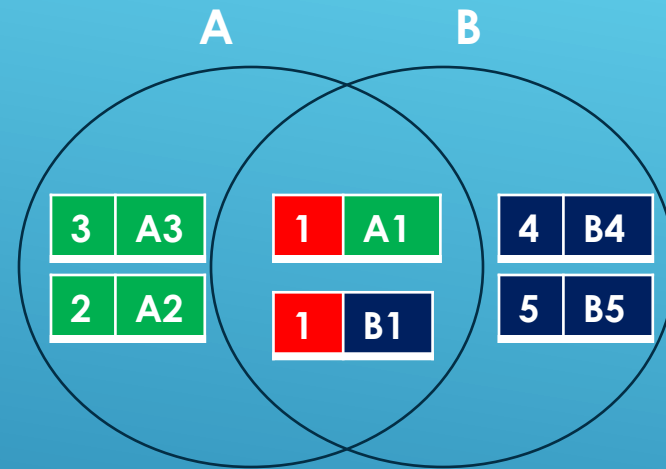
```
SELECT @@GLOBAL.sql_mode;  
SELECT @@SESSION.sql_mode;
```

SQL

Data Manipulation Language (DML)

JOINS ⋈. Join Types

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5





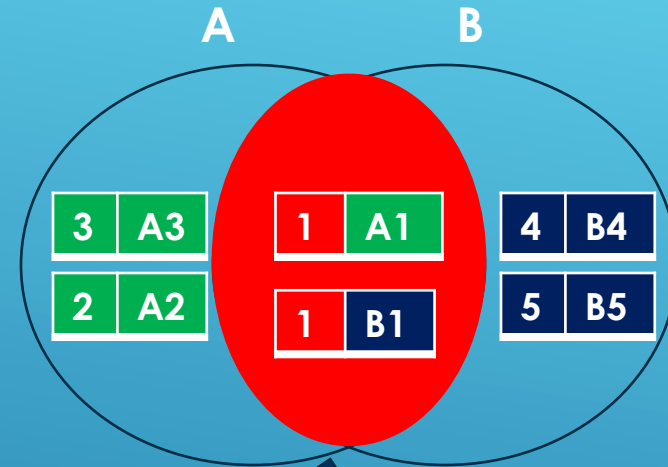
For keywords escaping

SQL

Data Manipulation Language (DML)

JOINS ☒. Inner Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1:

```
SELECT A.value, B.value  
FROM A INNER JOIN B ON A.`key`=B.`key`
```

Version 2:

```
SELECT A.value, B.value  
FROM A, B  
WHERE A.`key`=B.`key`
```

A.value	B.Value
A1	B1



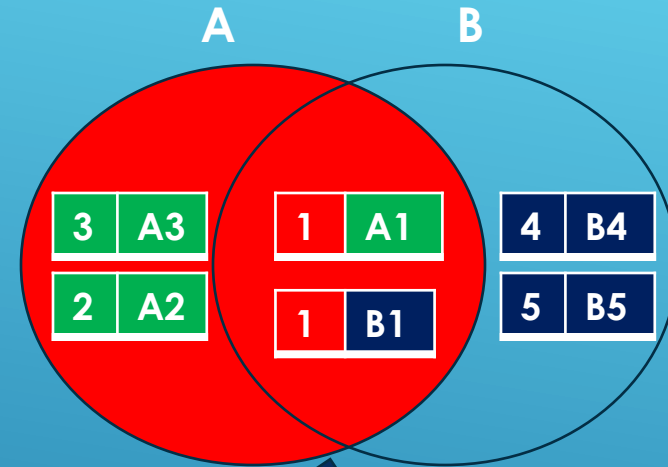
For keywords escaping

SQL

Data Manipulation Language (DML)

JOINS ☒ . Left Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1:



```
SELECT A.value, B.value  
FROM A LEFT JOIN B ON A.`key`=B.`key`
```

Version 2:



```
SELECT A.value, B.value  
FROM A LEFT OUTER JOIN B ON A.`key`=B.`key`
```

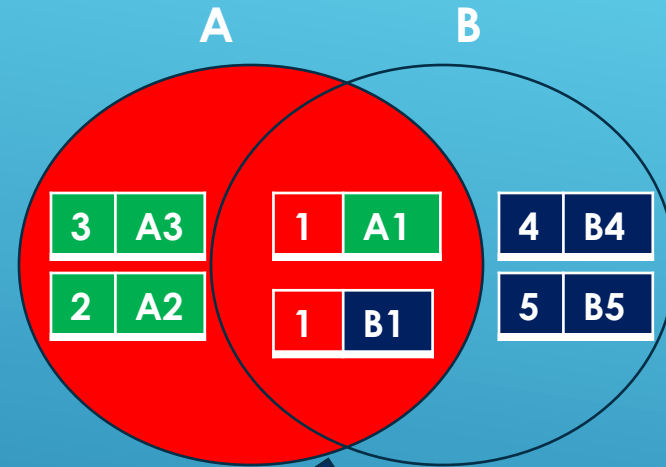
A.value	B.Value
A1	B1
A2	NULL
A3	NULL

SQL

Data Manipulation Language (DML)

JOINS ☒ . Left Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1*:



```
SELECT A.value, B.value  
FROM A, B WHERE A.[key] *= B.[key]
```

Version 2*:

ORACLE

```
SELECT A.value, B.value  
FROM A, B WHERE A."key"=B."key" (+)
```

A.value	B.Value
A1	B1
A2	NULL
A3	NULL



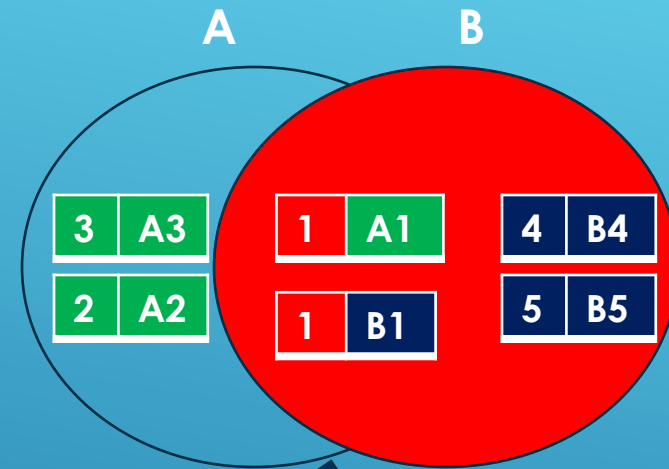
For keywords escaping

SQL

Data Manipulation Language (DML)

JOINS ☒. Right Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1:



```
SELECT A.value, B.value  
FROM A RIGHT JOIN B ON A.`key`=B.`key`
```

Version 2:



```
SELECT A.value, B.value  
FROM A RIGHT OUTER JOIN B ON A.`key`=B.`key`
```

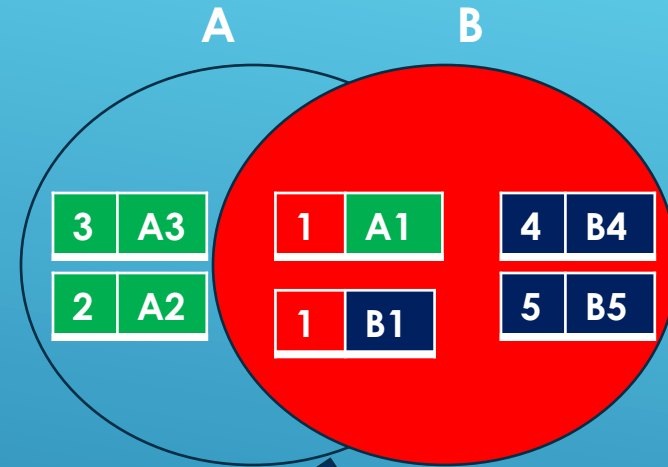
A.value	B.Value
A1	B1
NULL	B4
NULL	B5

SQL

Data Manipulation Language (DML)

JOINS ☒. Right Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1*:



```
SELECT A.value, B.value  
FROM A, B WHERE A.[key]=*B.[key]
```

Version 2*:

ORACLE

```
SELECT A.value, B.value  
FROM A, B WHERE A."key" (+)=B."key"
```

A.value	B.Value
A1	B1
NULL	B4
NULL	B5

SQL

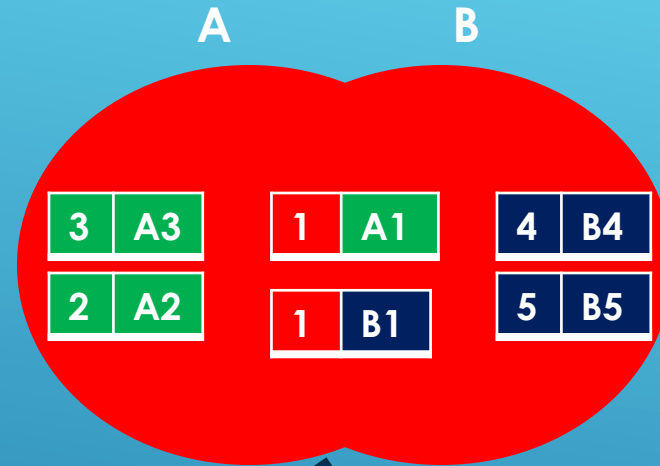


For keywords escaping

Data Manipulation Language (DML)

JOINS ☒. Full Join

A		B	
key	value	key	value
1	A1	1	B1
2	A2	4	B4
3	A3	5	B5



Version 1:



```
SELECT A.value, B.value  
FROM A FULL JOIN B ON A.`key`=B.`key`
```

Version 2:



Except



```
SELECT A.value, B.value  
FROM A FULL OUTER JOIN B ON A.`key`=B.`key`
```

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A3	NULL

SQL

Data Manipulation Language (DML)

JOINS ☒ . Full Join

How to simulate full joins?



A	
key	value
1	A1
2	A2
3	A2

B	
key	value
1	B1
4	B4
5	B5

Expected result

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL

SQL

Data Manipulation Language (DML)

JOINS ☒ . Full Join

How to simulate full joins?



A	
key	value
1	A1
2	A2
3	A2

B	
key	value
1	B1
4	B4
5	B5

Solution 1:

```
SELECT A.value, B.value
  FROM A LEFT JOIN B ON A.`key`=B.`key`
UNION
SELECT A.value, B.value
  FROM A RIGHT JOIN B ON A.`key`=B.`key`;
```

Expected result

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL



A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL

SQL

Data Manipulation Language (DML)

JOINS ☒ . Full Join

How to simulate full joins?



A	
key	value
1	A1
2	A2
3	A2

B	
key	value
1	B1
4	B4
5	B5

Solution 2:

```
SELECT A.value, B.value
  FROM A LEFT JOIN B ON A.`key`=B.`key`
UNION ALL
SELECT A.value, B.value
  FROM A RIGHT JOIN B ON A.`key`=B.`key`;
```

Expected result

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL



A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL
A1	B1

SQL

Data Manipulation Language (DML)

JOINS ☒ . Full Join

How to simulate full joins?



A	
key	value
1	A1
2	A2
3	A2

B	
key	value
1	B1
4	B4
5	B5

Solution 3:

```
SELECT A.value, B.value
  FROM A LEFT JOIN B ON A.`key`=B.`key`
UNION ALL
SELECT A.value, B.value
  FROM A RIGHT JOIN B ON A.`key`=B.`key`
 WHERE A.`key` IS NULL;
```

Expected result

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL

||

A.value	B.Value
A1	B1
NULL	B4
NULL	B5
A2	NULL
A2	NULL

SQL

Data Manipulation Language (DML)

JOINS ⋈. Counting joins

T1		T2		T3	
A	B	B	C	C	D

|T1| - number of records in T1;
|T2| - number of records in T2;
|T3| - number of records in T3;

Fill the question marks with the correct numbers or records.

```
SELECT * FROM T1 CROSS JOIN T2;
```

Cross Product	
Min	Max
?	?

```
SELECT * FROM T1 INNER JOIN T2 ON T1.B=T2.B;
```

Inner Join	
Min	Max
?	?

```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B;
```

Left Join	
Min	Max
?	?

```
SELECT * FROM T1 FULL JOIN T2 ON T1.B=T2.B;
```

Full Join	
Min	Max
?	?

```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B  
RIGHT JOIN T3 ON T2.C=T3.C;
```

Multiple Outer Joins	
Min	Max
?	?

SQL

Data Manipulation Language (DML)

JOINS \bowtie . Counting joins

T1		T2		T3	
A	B	B	C	C	D

$|T1|$ - number of records in T1;
 $|T2|$ - number of records in T2;
 $|T3|$ - number of records in T3;

Fill the question marks with the correct numbers or records.

```
SELECT * FROM T1 CROSS JOIN T2;
```

Cross Product	
Min	Max
$ T1 * T2 $	$ T1 * T2 $

Number of records returned by the cross product is always the multiplication between the sizes of the two tables.

SQL

Data Manipulation Language (DML)

JOINS ⋈. Counting joins

T1		T2		T3	
A	B	B	C	C	D

|T1| - number of records in T1;
|T2| - number of records in T2;
|T3| - number of records in T3;

Fill the question marks with the correct numbers or records.

```
SELECT * FROM T1 INNER JOIN T2 ON T1.B=T2.B;
```

Inner Join	
Min	Max
0	T1 * T2

Minimum records are obtained in case no combination of records make the join condition **True**.
Maximum number of records can be obtained in case the join condition is **True** for any pairs of records. For example join on Gender where all students are Females.

SQL

Data Manipulation Language (DML)

JOINS \bowtie . Counting joins

T1		T2		T3	
A	B	B	C	C	D

|T1| - number of records in T1;
|T2| - number of records in T2;
|T3| - number of records in T3;

Fill the question marks with the correct numbers or records.

```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B;
```

Left Join	
Min	Max
T1	T1 * T2

Minimum records are obtained in case no combination of records make the join condition **True**. In this case only the records from T1 are returned.

Maximum number of records can be obtained in case the join condition is **True** for any pairs of records (we assume that T2 has at least one record, if T2 has no records then maximum is |T1|).

SQL

Data Manipulation Language (DML)

JOINS \bowtie . Counting joins

T1		T2		T3	
A	B	B	C	C	D

$|T1|$ - number of records in T1;
 $|T2|$ - number of records in T2;
 $|T3|$ - number of records in T3;

Fill the question marks with the correct numbers or records.

Minimum records are obtained in case the following holds:

1. **SELECT DISTINCT B FROM T1** = **SELECT B FROM T1**
2. **SELECT DISTINCT B FROM T2** = **SELECT B FROM T2**
3. **SELECT B FROM T1** \subseteq **SELECT B FROM T2** OR **SELECT B FROM T1** \subseteq **SELECT B FROM T2**

For the maximum number of records returned we have to keep in mind that there are cases when $|T1| + |T2| > |T1| * |T2|$, for example 1 and 3. $|T1| + |T2|$ is obtained where the join condition is **False** for all records and $|T1| * |T2|$ when the join is **True** for all records.

SELECT * FROM T1 FULL JOIN T2 ON T1.B=T2.B;

Full Join	
Min	Max
$\text{Max}(T1 , T2)$	$\text{Max}(T1 + T2 , T1 * T2)$

SQL

Data Manipulation Language (DML)

JOINS \bowtie . Counting joins

T1		T2		T3	
A	B	B	C	C	D

|T1| - number of records in T1;
|T2| - number of records in T2;
|T3| - number of records in T3;

Fill the question marks with the correct numbers or records.

Minimum records are obtained when all join conditions are false. Maximum is obtained when all join conditions are true for all records (**same assumption that each table has at least 1 row**).

```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B  
           RIGHT JOIN T3 ON T2.C=T3.C;
```

Multiple Outer Joins	
Min	Max
T3	T1 * T2 * T3

SQL

Data Manipulation Language (DML)

JOINS Counting joins

T1		T2		T3	
A	B	B	C	C	D

|T1| - number of records in T1;
|T2| - number of records in T2;
|T3| - number of records in T3;

Fill the question marks with the correct numbers or records.

```
SELECT * FROM T1 CROSS JOIN T2;
```

Cross Product	
Min	Max
T1 * T2	T1 * T2

```
SELECT * FROM T1 INNER JOIN T2 ON T1.B=T2.B;
```

Inner Join	
Min	Max
0	T1 * T2

```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B;
```

Left Join	
Min	Max
T1	T1 * T2

```
SELECT * FROM T1 FULL JOIN T2 ON T1.B=T2.B;
```

Full Join	
Min	Max
Max(T1 , T2)	Max(T1 + T2 , T1 * T2)

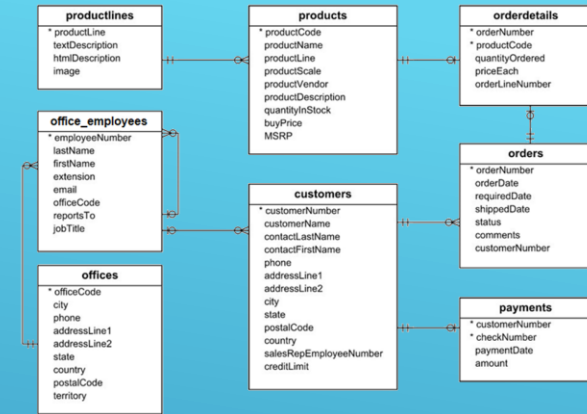
```
SELECT * FROM T1 LEFT JOIN T2 ON T1.B=T2.B  
RIGHT JOIN T3 ON T2.C=T3.C;
```

Multiple Outer Joins	
Min	Max
T3	T1 * T2 * T3

SQL

Data Manipulation Language (DML)

JOINS ∞ . Mimic Difference



Return the codes for those "Motorcycles" (productLine) products that were never ordered in a quantity greater than 50. Schema details can be found [here](#).



```
SELECT productCode FROM products WHERE productLine='Motorcycles'
EXCEPT
SELECT productCode FROM orderDetails WHERE quantityOrdered>50;
```

How about

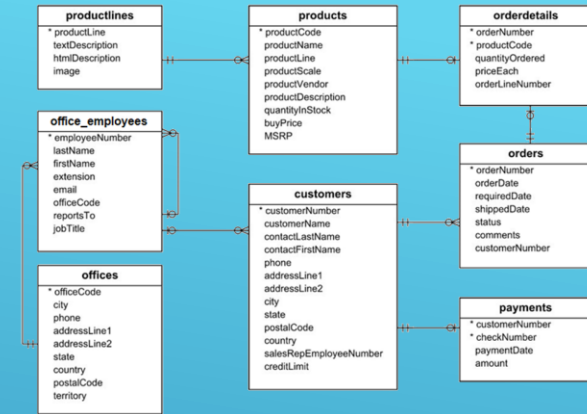


SQL

Data Manipulation Language (DML)

JOINS ∞. Mimic Difference

Problem started [here](#).



Return the codes for those "Motorcycles" (productLine) products that were never ordered in a quantity greater than 50. Schema details can be found [here](#).

```
SELECT productCode FROM products WHERE productLine='Motorcycles'
EXCEPT
SELECT productCode FROM orderDetails WHERE quantityOrdered>50;
```

How about     Access ?

```
SELECT P.productCode FROM products P LEFT JOIN orderDetails OD
      ON P.productCode=OD.productCode AND quantityOrdered>50
WHERE P.productLine='Motorcycles' AND OD.orderNumber IS NULL;
```

As we will see this can be done also using **NOT IN** or **NOT EXISTS** expressions.

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

```
SELECT *  
FROM accounts A,  
      orders O,  
      deals D,  
      dealDetails DD,  
      ordersDemand OD,  
      dealClients DC,  
      clients C  
WHERE O.accountId=A.accountId AND O.dealId=D.dealId AND  
O.dealId=DD.dealId AND OD.orderId=O.orderId AND  
DC.dealId=D.dealId AND C.clientId=DC.clientId AND OD.qty>50 AND  
O.status='In Progress' AND C.name='John';
```

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

```
SELECT *  
FROM accounts A INNER JOIN orders O ON O.accountId=A.accountId  
      INNER JOIN deals D ON O.dealId=D.dealId  
      INNER JOIN dealDetails DD ON O.dealId=DD.dealId  
      INNER JOIN ordersDemand OD ON OD.orderId=O.orderId  
      INNER JOIN dealClients DC ON DC.dealId=D.dealId  
      INNER JOIN clients C ON C.clientId=DC.clientId  
WHERE OD.qty>50 AND O.status='In Progress' AND C.name='John';
```

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

Using below schema return all triples (clientName, accountId, amount) for all clients which has in their account less than 1000\$.

clients	
clientName	accountId
Rob	1
Adam	2
John	3



accounts	
accountId	amount
1	10000
2	500

Schema/data definition.

```
CREATE TABLE clients (clientName VARCHAR(20), accountId INT);
```

```
CREATE TABLE accounts (accountId INT, amount INT);
```

```
INSERT INTO clients (clientName, accountId) VALUES ('Rob', 1);
```

```
INSERT INTO clients (clientName, accountId) VALUES ('Adam', 2);
```

```
INSERT INTO clients (clientName, accountId) VALUES ('John', 3);
```

```
INSERT INTO accounts (accountId, amount) VALUES (1, 10000);
```

```
INSERT INTO accounts (accountId, amount) VALUES (2, 500);
```

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

Using below schema return all triples (clientName, accountId, amount) for all clients which has in their account less than 1000\$.

clients	
clientName	accountId
Rob	1
Adam	2
John	3



accounts	
accountId	amount
1	10000
2	500

```
SELECT C.*, ISNULL(A.amount,0)
FROM clients C, accounts A
WHERE C.accountId *= A.accountId AND ISNULL(A.amount,0)<1000;
```



clientName	accountId	amount
Rob	1	0
Adam	2	500
John	3	0

Is it what expected?

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

Using below schema return all triples (clientName, accountId, amount) for all clients which has in their account less than 1000\$.

clients	
clientName	accountId
Rob	1
Adam	2
John	3



accounts	
accountId	amount
1	10000
2	500

```
SELECT C.*, ISNULL(A.amount,0)
FROM clients C, accounts A
WHERE C.accountId *= A.accountId AND ISNULL(A.amount,0)<1000;
```



clientName	accountId	amount
Rob	1	0
Adam	2	500
John	3	0

Why is Rob returned?

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

Using below schema return all triples (clientName, accountId, amount) for all clients which has in their account less than 1000\$.

clients	
clientName	accountId
Rob	1
Adam	2
John	3



accounts	
accountId	amount
1	10000
2	500

```
SELECT C.*, ISNULL(A.amount,0)
FROM clients C, accounts A
WHERE C.accountId *= A.accountId AND ISNULL(A.amount,0)<1000;
```



```
SELECT C.*, ISNULL(A.amount,0)
FROM clients C INNER JOIN accounts A
ON C.accountId = A.accountId AND ISNULL(A.amount,0)<1000;
```



clientName	accountId	amount
Rob	1	0
Adam	2	500
John	3	0

SQL

Data Manipulation Language (DML)

JOINS ☒. Writing joins

Using below schema return all triples (clientName, accountId, amount) for all clients which has in their account less than 1000\$.

clients	
clientName	accountId
Rob	1
Adam	2
John	3



accounts	
accountId	amount
1	10000
2	500

Correct version.

```
SELECT C.*, ISNULL(A.amount,0)
FROM clients C LEFT JOIN accounts A
      ON C.accountId = A.accountId
WHERE ISNULL(A.amount,0)<1000;
```

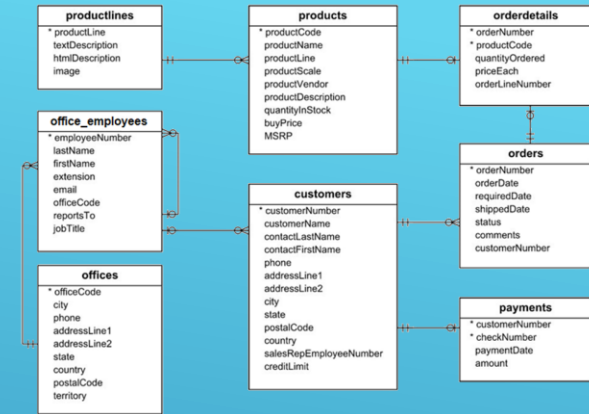


clientName	accountId	amount
Adam	2	500
John	3	0

SQL

Data Manipulation Language (DML)

Updating Joins



Using consultants table below update the email address in the office_employees table with the vendorEmail from the consultants table based on the employee number. Schema details can be found [here](#).

employeeNumber	vendorEmail
1102	gbondur@vendors.com
1337	lbondur@vendors.com
1611	afixter@vendors.com
1625	ykato@vendors.com



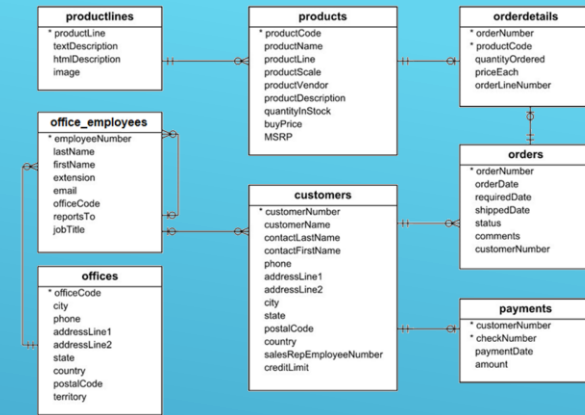
```
CREATE TABLE consultants (employeeNumber INT NULL, vendorEmail VARCHAR(100) NULL);
```

```
INSERT INTO consultants (employeeNumber, vendorEmail) VALUES (1102, 'gbondur@vendors.com');
INSERT INTO consultants (employeeNumber, vendorEmail) VALUES (1337, 'lbondur@vendors.com');
INSERT INTO consultants (employeeNumber, vendorEmail) VALUES (1611, 'afixter@vendors.com');
INSERT INTO consultants (employeeNumber, vendorEmail) VALUES (1625, 'ykato@vendors.com');
```

SQL

Data Manipulation Language (DML)

Updating Joins



Using consultants table below update the email address in the office_employees table with the vendorEmail from the consultants table based on the employee number. Schema details can be found [here](#).



```
UPDATE office_employees E INNER JOIN consultants C ON E.employeeNumber=C.employeeNumber
SET E.email = C.vendorEmail;
```



```
UPDATE office_employees E SET E.email = C.vendorEmail
FROM consultants C WHERE E.employeeNumber=C.employeeNumber;
```



```
UPDATE office_employees SET email = C.vendorEmail
FROM office_employees E INNER JOIN consultants C ON E.employeeNumber=C.employeeNumber;
```

How about

ORACLE

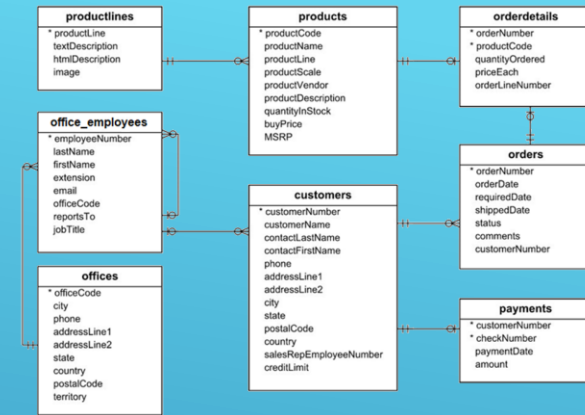


?

SQL

Data Manipulation Language (DML)

Merge Statement



Using consultants table below update the email address in the office_employees table with the vendorEmail from the consultants table based on the employee number. Schema details can be found [here](#).

How about   ? Could be done with bit more complex nested queries.

MERGE

```
INTO target_table
USING table_source
ON (merge_search_condition)
[ WHEN MATCHED
  THEN (UPDATE SET set_clause | DELETE) ]
[ WHEN NOT MATCHED
  THEN INSERT (column_list) VALUES (values_list) ]
```

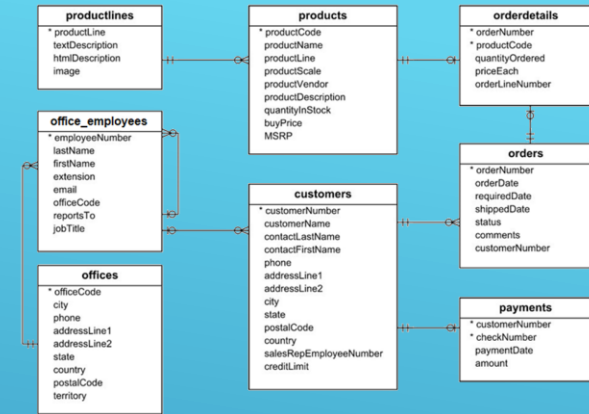


Merge does INSERT/UPDATE/DELETE of rows in a single statement. When `merge_search_condition` is true, than that row can be either updated or deleted (**WHEN MATCHED**) if false, then we can insert a new row (**WHEN NOT MATCHED**)

SQL

Data Manipulation Language (DML)

Merge Statement



Using consultants table below update the email address in the office_employees table with the vendorEmail from the consultants table based on the employee number. Schema details can be found [here](#).

Solution.

MERGE

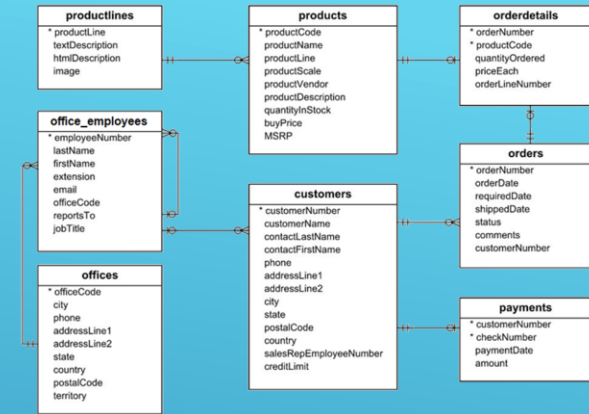
```
INTO office_employees E
USING consultants C
ON (E.employeeNumber=C.employeeNumber)
WHEN MATCHED
    THEN UPDATE SET E.email = C.vendorEmail
```



SQL

Data Manipulation Language (DML)

Merge Statement



Using consultants table below update the email address in the office_employees table with the vendorEmail from the consultants table based on the employee number. Schema details can be found [here](#).

Not so orthodox solutions but helps knowing the syntax in case you want to insert data that may violate the key constraints. Note that the constant values in the below select statements do not play any role because we know that all data is duplicated. This is not always the case.



```
INSERT INTO office_employees (employeeNumber, email, firstName, lastName, extension, officeCode, jobTitle)
SELECT employeeNumber, vendorEmail, 'x', 'x', 'x', 'x', 'x' FROM consultants
ON DUPLICATE KEY UPDATE email=consultants.vendorEmail;
```



```
INSERT INTO office_employees (employeeNumber, email, firstName, lastName, extension, officeCode, jobTitle)
SELECT employeeNumber, vendorEmail, 'x', 'x', 'x', 'x', 'x' FROM consultants
ON CONFLICT (employeeNumber) DO SET email=consultants.vendorEmail;
```