

# APPLYING THE BORSUK-ULAM THEOREM

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ABSTRACT. This writeup acts as a sketch of a presentation on topological combinatorics given for Math 2240 at Cornell. We hope present one potential answer to the question of “why topology?” by connecting two seemingly disparate areas of mathematics: combinatorics (where we work with the discrete) and topology (where we work with the continuous). Broadly, the talk begins with a general introduction to topology, setting up the necessary machinery to state and prove the famous *Borsuk-Ulam* theorem (Theorem 2.1), and continues to discuss its various combinatorial applications.

We assume no more than familiarity with the standard Euclidean topology in  $\mathbb{R}^n$ .

## 1. SETUP

## 2. BORSUK-ULAM

**Theorem 2.1** (Borsuk-Ulam). *Given a continuous map  $f : \mathbb{S}^n \rightarrow \mathbb{R}^n$ ,  $f$  identifies two antipodal points: i.e.  $\exists \mathbf{x} \in \mathbb{S}^n$  such that  $f(\mathbf{x}) = f(-\mathbf{x})$ .*

### 3. HEX GAME

The game of Hex is played between two players, say Red and Green. Each of them takes turns coloring a finite hexagonal grid. After all hexagons are colored, Red wins if there's a red path connecting the top and bottom, and Green wins if there's a green path connecting the left and right.

**Remark.** Play a game after introducing Hex give a feel for the problem (if time permits).

**Theorem 3.1.** *Hex can never end in a tie.*

We won't assume that

## 4. FAIR DIVISION

Suppose you and a friend steal an open necklace, engraved with  $m$  precious stones. There are  $d$  kinds of stones (labelled  $1, \dots, d$ ), an even number of each kind. Neither you nor your co-conspirator know the values of the different stone types. Hence, you decide to split the stones between each of you such that you both have the same number of jewels of each kind.

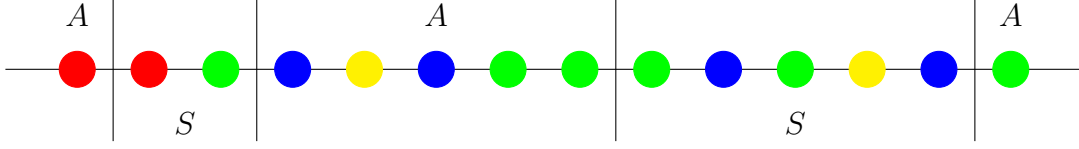


FIGURE 1. Example necklace and cut: segments labelled  $A$  go to Akash and those labelled  $S$  go to Sid. Our necklace has 14 stones of 4 types and was split using 4 cuts.

Instead of removing the individual stones one by one, you decide to split the jewels by making as few cuts as possible to the necklace, divvying out the remaining strands between yourselves.

**Theorem 4.1.** *The minimal number of cuts needed is at most  $d$ .*

*Proof.* At its current state, there isn't any good way to apply Borsuk-Ulam on our puzzle. Let's change that! We'll consider the interval  $[0, m]$  as our necklace. Break  $[0, m]$  into the union of subintervals

$$\bigcup_{k=0}^{m-2} [k, k+1) \cup [m-1, m].$$

The  $k$ th stone corresponds to the  $k$ th subinterval. Then, define characteristic functions for the stones as  $f_i : [0, m] \rightarrow \{0, 1\}$  for  $i \in [d]$  such that

$$f_i(x) = \begin{cases} 1 & x \in [k-1, k) \text{ and the } k\text{th stone is of type } i \\ 0 & \text{otherwise} \end{cases}.$$

Fix  $0 = z_0 \leq z_1 \leq z_2 \leq \dots \leq z_d \leq m = z_{d+1}$  to act as the cuts on our string. If we wish to assign the part  $[z_i, z_{i-1}]$  ( $i = 1, \dots, d+1$ ) to the first thief, set  $x_i = (\sqrt{z_i - z_{i-1}})/m$  and  $x_i = -(\sqrt{z_i - z_{i-1}})/m$  otherwise. The tuple  $(x_1, x_2, \dots, x_{d+1}) \in \mathbb{S}^d$  encodes a cut of the necklace. Noting that  $z_j = \sum_{i=1}^j m^2 x_i^2$ , consider the continuous map

$$(x_1, x_2, \dots, x_{d+1}) \rightarrow \left( \sum_{j=1}^{d+1} \text{sign}(x_j) \int_{[z_j, z_{j-1}]} f_1(x) \, dx, \dots, \sum_{j=1}^{d+1} \text{sign}(x_j) \int_{[z_j, z_{j-1}]} f_d(x) \, dx \right).$$

Since  $f$  is antipodal, theorems 2.1 and ?? tell us  $\exists \mathbf{x} \in \mathbb{S}^n$  such that  $f(\mathbf{x}) = 0$ . The cut associated  $\mathbf{x}$  is a fair division.

However, we might run into a problem translating this back into the discrete case! What if some  $z_i$  associated with  $\mathbf{x}$  lies in  $(k, k+1)$  for some  $k \in \mathbb{Z}$ ? (I.e. we cut partially into one of the jewels.) Given a non-integral cut subdividing a stone of type  $i$ , where a portion of stone  $i$  is assigned to the first thief, we know our cut is either unnecessary or there exists at least one other partial cut into stones of type  $i$  assigned to thief one since the sums of lengths of stone  $i$  intervals is an integer. In the latter case, we can move the first “non-integral” cut to the right and the remaining to the left, without changing the loot for each thief.  $\square$

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