4.1.1 Class Discussion

Given these two relations:

CUSTOMER (<u>cust_id</u>, cust_name, cust_address)

ORDER (<u>order_id</u>, order_date, cust_id)

| CUST_ID | CUST_NAME | CUST_ADDRESS | CUST_GENDER |
|---------|-----------|-------------------------|-------------|
| 1 | Jack | 6 Jalan Jaya 2 | М |
| 2 | Jill | 185 Jalan Ampang | F |
| 3 | John | 16-02, Jalan Permas 9/2 | М |

| ORDER_ID | ORDER_DATE | CUST_ID |
|----------|------------|---------|
| 1 | 5/3/2021 | 1 |
| 2 | 5/4/2021 | 1 |
| 3 | 6/4/2021 | 1 |
| 4 | 5/4/2021 | 2 |
| 5 | 6/4/2021 | 3 |
| 6 | 8/4/2021 | 3 |

1. Identify the following terms based on the above CUSTOMER and ORDER relations:

- a) Relation customer and order
- b) Attribute cust_id, 4 for customer, 3 for order
- c) Domain CUST_NAME 20 character string
- d) Tuple {1, Jack, 6 Jalan Jaya 2, M}, 6, 3 tuples
- e) Degree of a relation and 4,3
- f) Cardinality 3 for CUSTOMER, 6 for ORDER
- g) Primary key CUST_ID for CUSTOMER and ORDER_ID for ORDER
- h) Foreign key no foreign key for CUSTOMER and CUST_ID for ORDER

4.1.2 Choosing the Primary key

1. In any relation, tuples must be unique. However, in many cases, the set of all the attributes in a relation is not considered a candidate key. Why not?

Table: Employee

| Emp_SSN | Emp_Number | Emp_FName E | Emp_Lname |
|-----------|------------|-------------|-----------|
| | | | |
| 123456789 | 226 | Steve | John |
| 999999321 | 227 | Steve | Smith |
| 888997212 | 228 | Sherry | Smith |
| 777778888 | 229 | Robert | John |

Super Key:

Emp_SSN

Emp_Number

Emp_SSN, Emp_Number, Emp_FName, Emp_Lname

Emp_SSN, Emp_Number

Emp_SSN, Emp_Number, Emp_FName

Emp FName, Emp Lname

Candidate key:

Emp_SSN,

Emp_Number

Emp_FName, Emp_Lname

Primary Key:

- Emp_SSN
- Emp_no

2. On the other hand, suppose we do have a relation where the set of all attributes is a candidate key. In this case, show that this set must, therefore, be the only candidate key and hence the primary key.

| PET ID | VISIT DATE | PROCEDURE |
|-----------|-------------|----------------------------|
| 246ROVER | JAN 13/2019 | 01 RABIES VACCINATION |
| 246ROVER | MAR 27/2019 | 10 EXAMINE AND TREAT WOUND |
| 246ROVER | APR 02/2019 | 05 HEART WORM TEST |
| 298SPOT | JAN 21/2019 | 08 TETANUS VACCINATION |
| 298SPOT | MAR 10/2019 | 05 HEART WORM TEST |
| 341MORRIS | JAN 23/2019 | 01 RABIES VACCINATION |
| 341MORRIS | JAN 13/2019 | 01 RABIES VACCINATION |
| 519TWEEDY | APR 30/2019 | 20 ANNUAL CHECK UP |
| 519TWEEDY | APR 30/2019 | 12 EYE WASH |

Superkey:

- PET ID, VISIT DATE, PROCEDURE

Candidate Key:

PET ID, VISIT DATE, PROCEDURE

Primary Key:

PET ID, VISIT DATE, PROCEDURE

3. Identify the primary key and foreign key for these three relations:

ORDER (ORDER_ID, ORDER_DATE, CUST_ID)

ORDERLINE (ORDER_ID, PROD_NO, OL_QTYORDERED, OL_LINEPRICE)

PRODUCT (PROD_NO, PROD_DESC, PROD_UNITPRICE)

ORDER

PK: ORDER_ID

FK: CUST_ID

ORDERLINE

PK: ORDER_ID and PROD_NO together

FK: ORDER_ID, PROD_NO

PRODUCT

PK: PROD_NO

FK: None

Relationships:

Order - Product (Many-to-many)

Order - Orderline (one to many)

Product - Orderline (one to many)

4. Identify the superkey(s), candidate key(s) and the primary key for the relation if the following business rules are applicable:

- A dentist can only see a single patient at a particular date and time
- A dentist treats a patient in a particular surgery room, and
- A patient can see the same dentist multiple times

APPOINTMENT

| dentist_i | dentist_nam e | patient_i d | patient_nam e | appointment_dateti me | surgery_roomn o |
|-----------|------------------|----------------|------------------|--------------------------|--------------------|
| D1 | Jack | P1 | Jill | 5/4/2020 10am | 1 |
| D1 | Jack | P2 | Smith | 5/4/2020 12pm | 1 |
| D2 | John | P3 | Mary | 5/4/2020 10pm | 2 |
| D2 | John | P3 | Mary | 6/4/2020 2pm | 2 |
| D3 | Will | P2 | Smith | 6/4/2020 4pm | 1 |
| D3 | Will | P4 | Doe | 7/4/2020 6pm | 3 |
| | | | | | |

Superkey:

dentist_id, dentist_name, appointment_datetime, surgery_room

Dentist_id, patient_id, appointment_datetime

dentist_id, appointment_datetime, surgery_roomno

dentist_id, appointment_datetime

Dentist_id, dentist_name, patient_id, patient_name, appointment_datetime

Dentist_id, dentist_name, patient_id, patient_name, appointment_datetime, surgery_room

patient_id, appointment_datetime

Candidate key:

dentist_id, patient_id, appointment_datetime?

dentist_id, appointment_datetime

appointment datetime, surgery roomno

patient_id, appointment_datetime

4.2.1 Relational Algebra Exercise

HOTEL (HOTEL_NO , HOTEL_NAME, HOTEL_CITY)

ROOM (ROOM_NO , HOTEL_NO , ROOM_TYPE, ROOM_PRICE)

BOOKING (HOTEL NO , GUEST NO , BDATE FROM , BDATE_TO, ROOM NO)

GUEST (GUEST_NO , GUEST_NAME, GUEST_ADDRESS)

πσΜ

1. List the names and cities of all hotels

 $\pi(HOTEL_NAME, HOTEL_CITY)(HOTEL)$

2. List all single rooms with a price below \$50

σ ROOM_TYPE = 'SINGLE' and ROOM_PRICE < \$50 (ROOM)

3. List the names of all hotels in Melbourne

 $\pi(HOTEL_NAME)(\sigma HOTEL_CITY="Melbourne"(HOTEL))$

4. List all names of hotels which have presidential suite room

 $\pi(HOTEL_NAME)(\sigma ROOM_TYPE = 'presidential suite'(HOTEL))$

S5. List the price and type of all rooms at the Grosvenor Hotel

 π (ROOM_TYPE, ROOM_PRICE) σ (HOTEL_NAME = 'Grosvenor Hotel'(HOTEL MROOM))

6. List all names and addresses of guests currently staying in deluxe room of any hotel (assume that if the guest has a tuple in the BOOKING relation, then they are currently staying in the hotel)

TTGUEST_NAME, GUEST_ADDRESS(OROOM_TYPE = deluxe room(ROOM ⋈ GUEST ⋈ BOOKING))

7) List all names and addresses of guests currently staying at the Grosvenor Hotel (assume that if the guest has a tuple in the BOOKING relation, then they are currently staying in the hotel)

π(names, address)(Booking⋈(Guest number = guest number)Guest), using right outer join

4.2.2. Advanced Relational Algebra Exercise

Considers these four relations:

CUSTOMER (cust_id, cust_name, cust_address)

PRODUCT (prod_no, prod_desc, prod_unitprice, prod_stock)

STAFF(staff_name, staff_position)

SALE (cust_id, sale_date, prod_no, sale_qty, sold_by)

*Note that sold_by value is the staff who made the sale

πσΜ

1. List names of customers and descriptions of products bought by the customer. How many tuples will be returned by the relational algebra query that you have constructed as your answer?

P1 =
$$\pi$$
(prod_no, prod_desc)(PRODUCT) #retrieve all prod_desc

C1 =
$$\pi$$
(cust id, cust name)(CUSTOMER) #retrieve all cust_name

$$S1 = \pi(cust_id, prod_no)$$
 #filter out unnecessary attributes from SALE

$$\mathsf{SP} = \mathsf{P1} \bowtie \mathsf{S1} \bowtie \mathsf{C1}$$
 #add prod_desc and cust_name to SALE

 $A1 = \pi(cust_name, prod_desc)(SPC)$ #project the cust_name and prod_desc for each tuple in sale

5 tuples will be returned

2. List all ∩names which are shared by customers and staff

$$N1 = \pi(cust_name)$$
 (CUSTOMER)

N2 =
$$\pi$$
(staff_name) (STAFF)
N3 = N1 \cap N2

- 3. List descriptions of products that haven't been sold
 - $S1 = \pi(prod_no)(SALE)$
 - $P1 = \pi(prod_no)(PRODUCT)$
 - PD1 = π (prod_no,prod_desc)(PRODUCT)
 - NS1 = P1-S1
 - $\pi(\text{prod_desc})(\text{NS1} \bowtie \text{PD1})$
- 4. List names of clerks who don't have any sales yet

$$\pi(\text{staff_name})(\sigma_{\text{staff_position=clerk AND sold_by=null}}(\text{SALE} \bowtie \text{STAFF}))$$

a =
$$\pi$$
 (staff_name (STAFF) \approx sold_by (SALE))
out = π a (σ staff_position = clerk and sale_qty= 0)

5. List categories (positions) of staff who have made sales