ACM Lab Week 4

Complete Search

October 8, 2015

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Backtracking is an important technique for performing a complete search.

Idea: recursively enumerate a set of partial candidates which eventually includes all solution candidates.

Pseudo-code:

```
backtrack(candidate) {
     if reject(candidate) {
2
       return;
     if is_solution(candidate) {
5
       output(candidate);
6
     }
7
     for (c ∈ extensions(candidate)) {
8
       backtrack(c);
9
10
11
```

There are three important design choices:

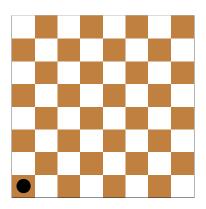
- 1 how do I decide whether I should reject a candidate?
- 2 how do I know whether I have found a solution?
- 3 what are the extensions of a candidate?

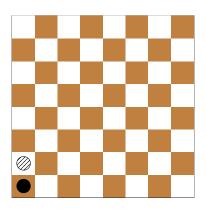
For the performance it is crucial to reject candidates as early as possible.

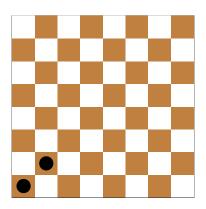
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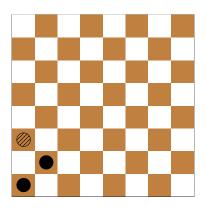
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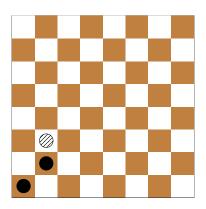
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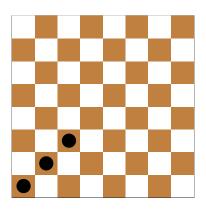


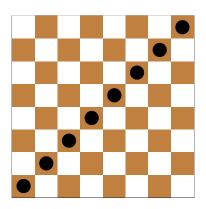


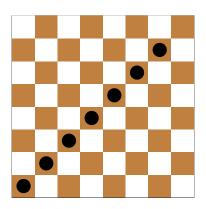


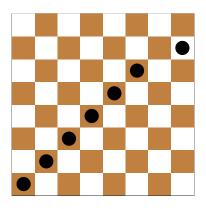


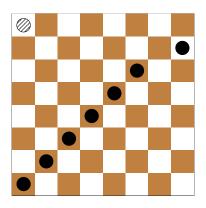


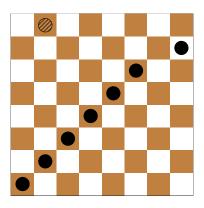


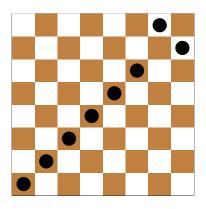


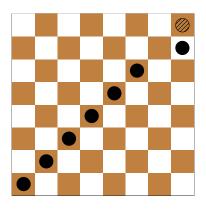


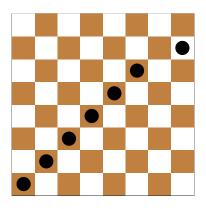


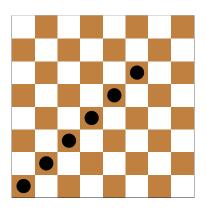


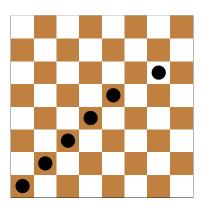


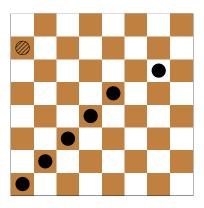


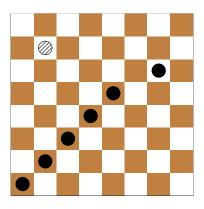


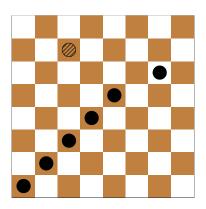


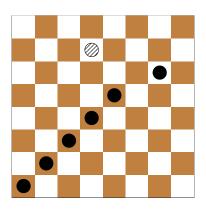


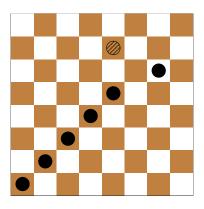


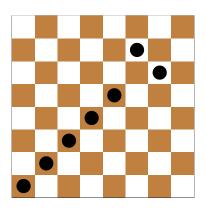


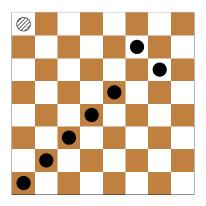


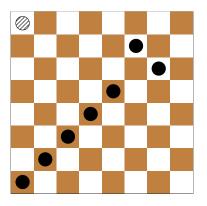












And so on...

When designing a backtracking algorithm, the goal should always be to restrict the set of extensions as much as possible as early as possible.

Don't be afraid to use heuristics to achieve this goal.

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Problem

Given two vertices u and v of a graph G, enumerate all paths between u and v in lexicographic order.

- \blacksquare The partial candidates are paths starting at u.
- The extensions of a path are obtained by adding one edge at the end. Process the extensions in lexicographic order.
- A path is accepted if it ends at v.
- A path can be rejected if there is no way to reach v from the endpoint (check with DFS).

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