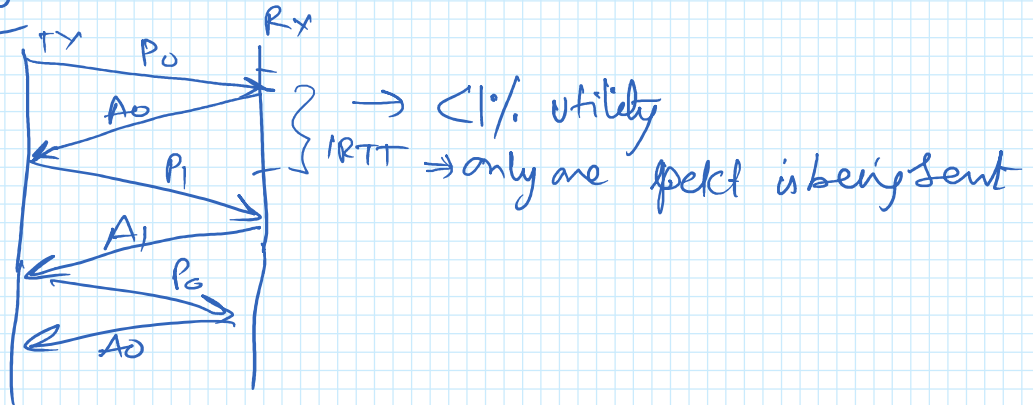


# Lecture 9

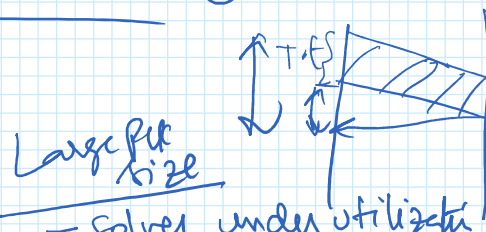
07 September 2021 17:04

- Recap - stop & go
- Pipelining based protocol  $\begin{cases} \rightarrow \text{Go Back N} \\ \rightarrow \text{Selective Repeat} \end{cases}$
- + dual implementation/simulations of  $\rightarrow$  also in book's website
- Congestion control.

## stop & go



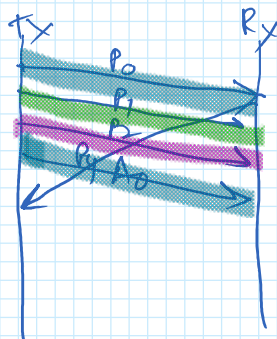
Improve stop & go Perf? ①  $\rightarrow$  Increase packet size  $\approx RTT$



- solves underutilization
- But increases packet error probability

## pipeline

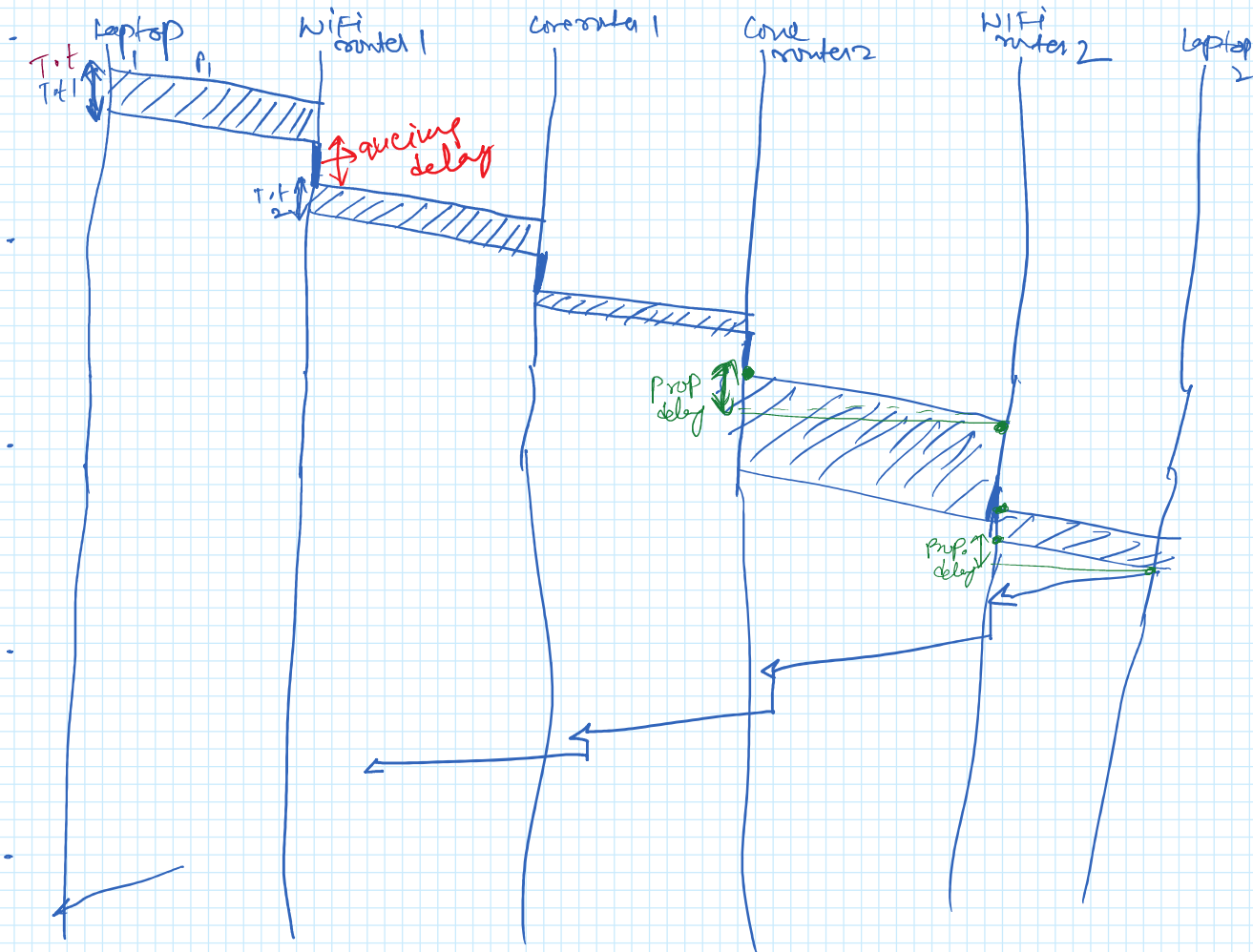
②



$T \cdot t \cdot$  = transmit time of a packet  
 $RTT$

$$\left. \begin{array}{l} \# \text{ Pkts} \\ \text{without} \\ \text{ack} \\ \text{to} \\ \text{send} \end{array} \right\} = \frac{RTT}{T \cdot t \cdot}$$

In Reality timing diagram:



- window size

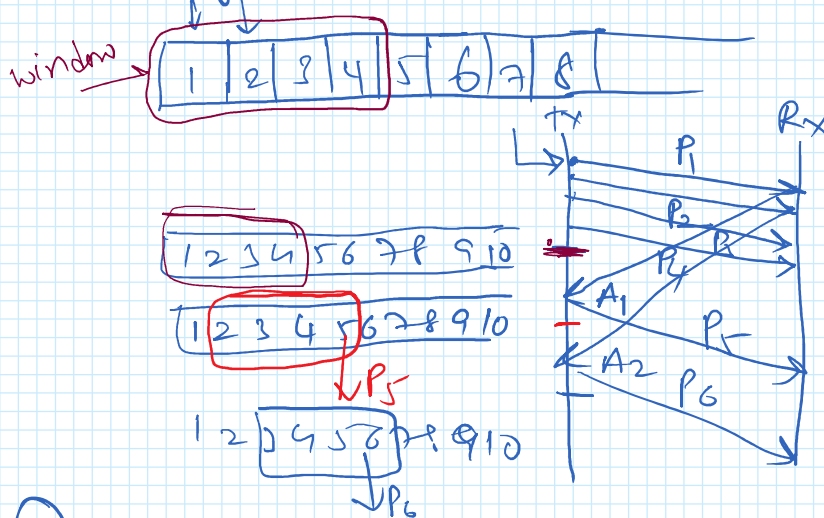
Max

- No. of unacknowledged packets that the sender is willing to tolerate

- In real protocols like TCP, the window size changes (adaptive) based on the current traffic situation

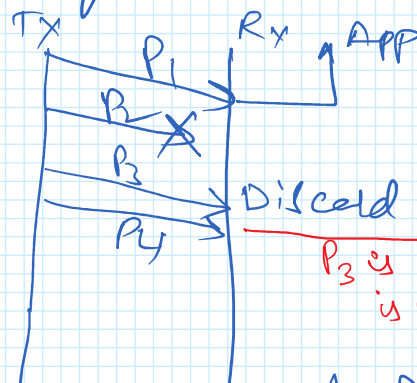
# GoBack N

① - Sliding window based protocol.



Window size = 4.

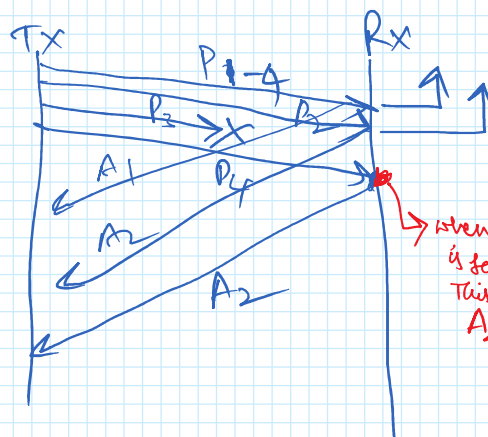
② - Rx. accepts only in-order packets.



$P_2$  is discarded as the Rx is expecting  $P_2$ .  $P_3$  is received "out-of-order".

③ Cumulative Acknowledgment

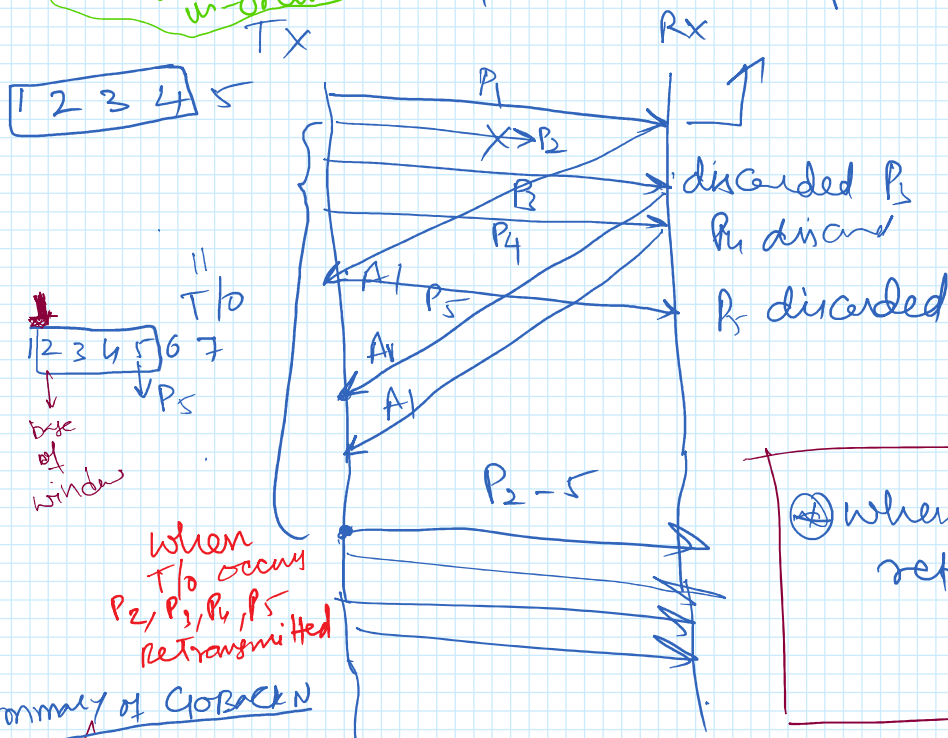
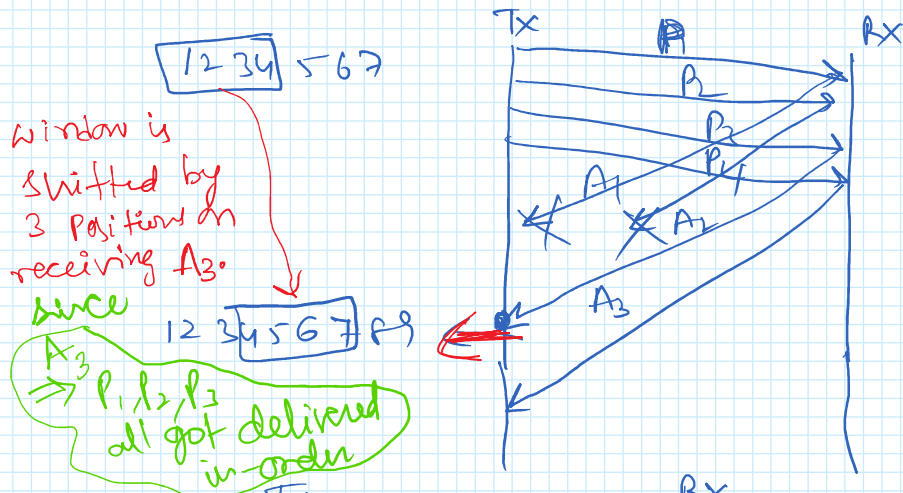
$A_4 \Rightarrow$  All packets till  $P_4$  are received in-order



only Acks  $\Rightarrow$  cumulative  
- NO NACK are used.

when  $P_4$  received,  $A_2$  is sent instead of  $A_4$ . This is because,  $P_4$  is discarded.  $A_2 \Rightarrow$  till  $P_2$  all are delivered in-order.

## Example 3



⊕ when timeout occurs retransmit all packets in window

## Summary of GoBackN

- GoBack N
- No buffer at Rx. Only in-order packets will be accepted.
  - Tx. maintains a sliding window
  - T/O ⇒ Resend all packets in window
  - cum. Ack ⇒ all packets till that are received in order

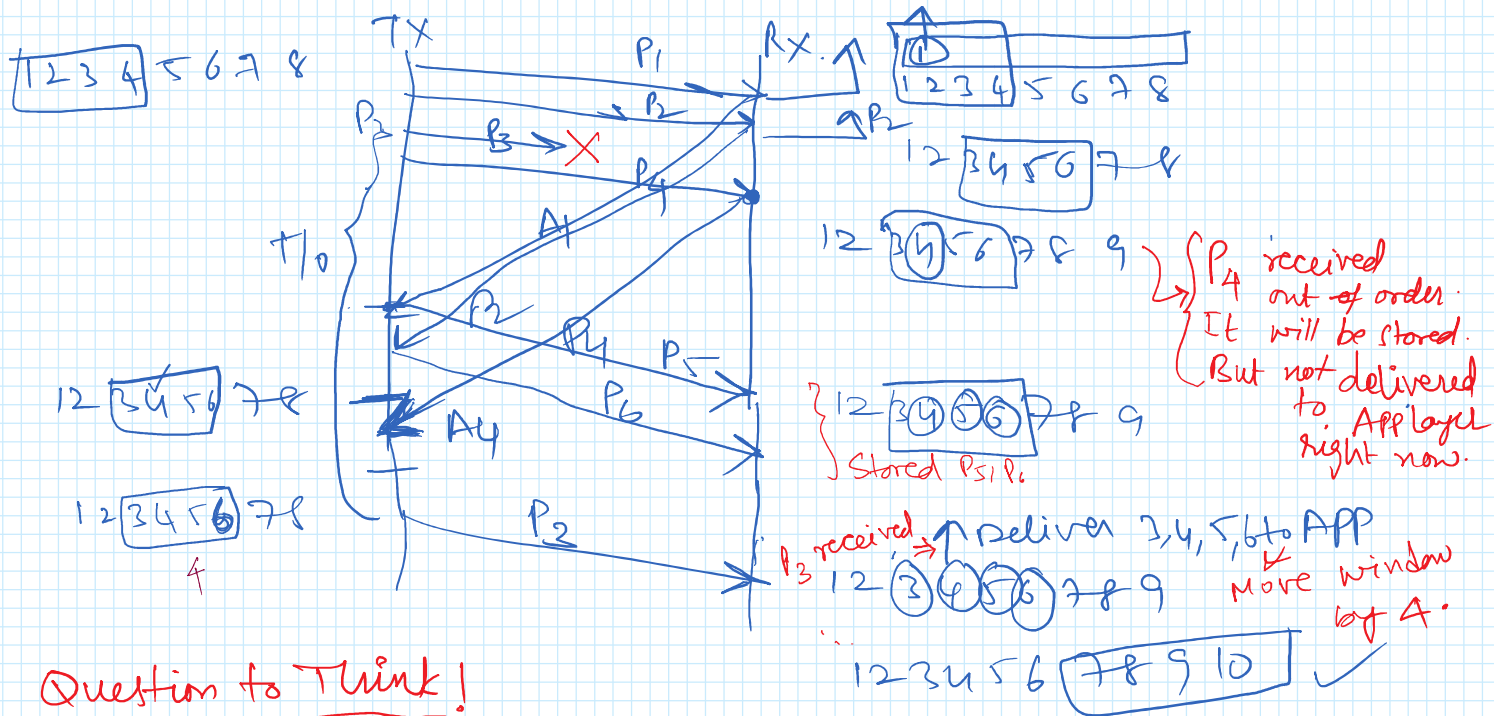
# Selective Repeat Algorithm

- Rx accepts out-of-order packets.

- Rx. <sup>has</sup> buffer and a window

↓  
determines how many out-of-order packets it can store.

- Acts are specific to packet received.  
(i.e., NOT cumulative Ack)



## Question to Think!

- Should Rx & TX window size be same??

Case (i) If Rx window size > Tx window size

⇒ No use since extra space unused because TX is not going to send any new packets

Case (ii) Tx window size > Rx window size.

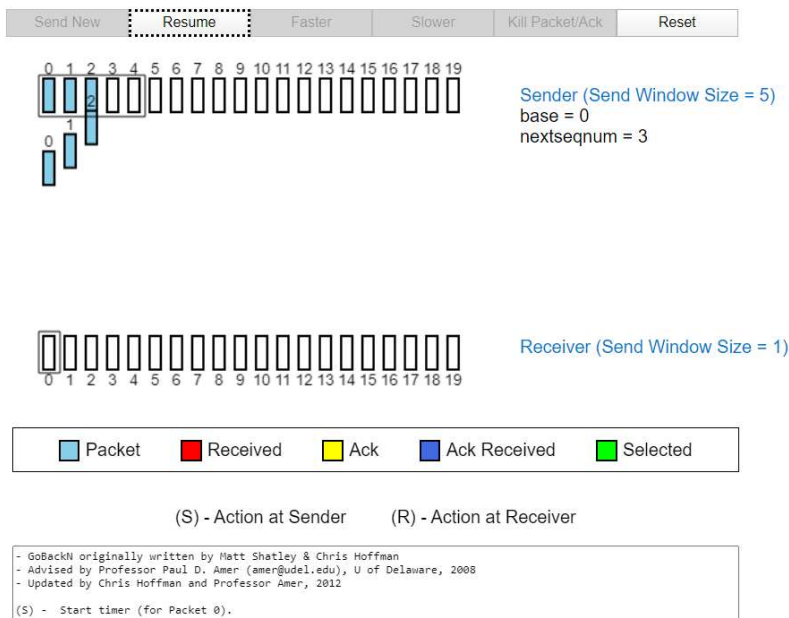
Chapter 3

# Interactive Animations on Book website

URL link  $\Rightarrow$  [https://wps.pearsoned.com/ecs\\_kurose\\_compnetsw/6/216/55463/14198702.cw/index.html](https://wps.pearsoned.com/ecs_kurose_compnetsw/6/216/55463/14198702.cw/index.html)

## Go-Back-N Protocol.

This interactive animation brings to life the Go-Back-N protocol. In this demo, the sending window limit "New". This action will begin moving data packets between sender and receiver. To simulate loss, select easier. Speed up or slow down the simulation by clicking "Faster" or "Slower". BE PATIENT for retransmissions.



Problem in chapter 3 of Kurose & Ross Book 6 edition (Indian version)

R12

Send five packets  
kill first packet.

GoBack N

