

Ex 2.24

Algorithm 1 Ex 2.24

Input: number of integers n

Output: print all the permutations for integers $1..n$

```
1: procedure PermuteDriver( $n$ ; ;)  
2:   Create empty array  $Arr[1..n]$   
3:   for  $i \in [1, n]$  do  
4:      $Arr[i] \leftarrow i$   
5:   end for  
6:   Permute( $n, Arr$ )  
7: end procedure  
8: procedure Permute( $k$ ; ;  $Arr[1..n]$ )  
9:   if  $k == 1$  then  
10:    print( $Arr[1..n]$ )  
11:   else  
12:    for  $i \in [1, k]$  do  
13:      Swap( $Arr[i], Arr[k]$ )  
14:      Permute( $k - 1, Arr$ )  
15:      Swap( $Arr[i], Arr[k]$ )  
16:    end for  
17:   end if  
18: end procedure
```

I got this question CORRECT

Ex 2.31

At beginning we call $DFSPath(T, 1)$

Algorithm 2 Ex 2.31

Input: Current vertex v to be visit, the index idx that name v should be written in to the array $Names$

Output: Print the path from root to v

```
1: global  $n$  set to a suitably large value, the array  $Names[1..n]$   
2: global the vertices in tree  $T$ , the lists  $child[v]$  for each  $v$  in  $T$   
3: procedure DFSPath( $v, idx$ ; ;)  
4:   if  $v$  is a leaf then  
5:      $Names[idx] \leftarrow v$   
6:     print( $Names[1..idx]$ )  
7:   else  
8:      $Names[idx] \leftarrow v$   
9:     for each child  $c$  in  $child[v]$  do  
10:      DFSPath( $c, idx + 1$ )  
11:    end for  
12:   end if  
13: end procedure
```

I got this question CORRECT. Notes : Line 5 and line 8 can be merged before line 4

Ex 2.32

a)

Algorithm 3 Ex 2.32a

Input: T , root of the tree

Output: colored tree as required

```
1: procedure ColorDriver( $T$ ; ;)  
2:   for each child  $c$  in  $\text{child}T$  do  
3:     Color( $c$ )  
4:   end for  
5: end procedure  
6: procedure Color( $v$ ; ;)  
7:   if  $v$  is a leaf then  
8:      $v.\text{color} \leftarrow \text{Green}$   
9:   else  
10:     $v.\text{color} \leftarrow \text{Green}$   
11:    for each child  $c$  in  $\text{child}[v]$  do  
12:      Color( $c$ )  
13:      if  $c.\text{color} == \text{Green}$  then  
14:         $v.\text{color} \leftarrow \text{Red}$   
15:      return  
16:    end if  
17:  end for  
18: end if  
19: end procedure
```

I got this question CORRECT. Notes: Line 8 and line 10 can be merged before line 7.

b)

Algorithm 4 Ex 2.32b

Input: T , root of the tree

Output: colored(boolean) tree as required

```
1: procedure BooleanColorDriver( $T$ ; ;)  
2:   for each child  $c$  in  $\text{child}T$  do  
3:     BooleanColor( $c$ )  
4:   end for  
5: end procedure  
6: procedure BooleanColor( $v$ ; ;)  
7:   if  $v$  is a leaf then  
8:      $v.\text{bool} \leftarrow \text{True}$   
9:   else  
10:     $v.\text{bool} \leftarrow \text{True}$   
11:    for each child  $c$  in  $\text{child}[v]$  do  
12:      BooleanColor( $c$ )  
13:       $v.\text{bool} \leftarrow v.\text{bool} \text{ And } \text{NOT } c.\text{bool}$   
14:    end for  
15:  end if  
16: end procedure
```

I got this question CORRECT. Notes: Line 8 and line 10 can be merged before line 7.

c)

Algorithm 5 Ex 2.32c, two function version

Input: T , root of the tree

Output: colored(boolean) tree in A's perspective

```
1: procedure BeginGame( $T$ ; ;)  
2:   for each child  $c$  in  $child[T]$  do  
3:      $AtoMove(c)$   
4:   end for  
5: end procedure  
6: procedure AtoMove( $v$ ; ;)  
7:   if  $v$  is a leaf then  
8:      $v.bool \leftarrow True$   
9:   else  
10:     $v.bool \leftarrow True$   
11:    for each child  $c$  in  $child[v]$  do  
12:       $BtoMove(c)$   
13:       $v.bool \leftarrow v.bool \text{ And } c.bool$   
14:    end for  
15:  end if  
16: end procedure  
17: procedure BtoMove( $v$ ; ;)  
18:   if  $v$  is a leaf then  
19:      $v.bool \leftarrow False$   
20:   else  
21:      $v.bool \leftarrow False$   
22:     for each child  $c$  in  $child[v]$  do  
23:        $AtoMove(c)$   
24:        $v.bool \leftarrow v.bool \text{ Or } c.bool$   
25:     end for  
26:   end if  
27: end procedure
```

I got this question CORRECT. Notes: Line 8 and line 10 can be merged before line 7. Line 19 and line 21 can be merged before line 18.

Ex 2.57

Algorithm 6 Ex 2.57

Input: number of vertices n , child lists $Child[1..n]$

Output: Return the name of the Root

```
1: function FindRoot( $n, Child[1..n]$ ; ;)  
2:   Create empty Marker Array visited[1.. $n$ ] initialize all elements to False  
3:   for  $i \in [1, n]$  do  
4:     for each  $c$  in  $Child[i]$  do  
5:        $visited[c] \leftarrow True$   
6:     end for  
7:   end for  
8:   for  $i \in [1, n]$  do  
9:     if  $visited[i] == False$  then  
10:      return  $i$   
11:    end if  
12:  end for  
13: end function
```

I got this question CORRECT.

Ex 3.2

a)

$$n > 2^{32}$$

I got this question CORRECT.

b)

$$n > 32$$

I got this question CORRECT.

c)

$$n > 2^8$$

I got this question CORRECT.

Ex 3.3

In increasing order:

$$\left(\frac{1}{2}\right)^n < n^{\frac{1}{\log_2 n}} < 2 < 1001n^{10010001} < (\log_2 n)^{\log_2 n} = n^{\log_2 \log_2 n} < n^{\log_2 n} < 10010n + 1.0000000000001^n < 1.0000001^n < (\log_2 n)^n$$

I got this question CORRECT. Notes: We should use << instead of <

Ex 3.4

a

$$\begin{cases} S(1) = 1 \\ S(n) = S(n-1) + \Theta(1) \end{cases}, n > 1 \quad (1)$$

I got this question CORRECT.

b

$$\begin{cases} S(1) = 1 \\ S(n) = S(n-1) + \Theta(1) \end{cases}, n > 1 \quad (2)$$

I got this question CORRECT.

c

$$\begin{cases} S(1) = 1 \\ S(n) = 2S(n-1) + \Theta(1) \end{cases}, n > 1 \quad (3)$$

I got this question CORRECT.

d

$$\begin{cases} S(1) = 1 \\ S(2) = 1 \\ S(n) = S(n-1) + S(n-2) + \Theta(1) \end{cases}, n > 2 \quad (4)$$

I got this question CORRECT.

Ex 3.5

True: a, b1, b, c1, c, d, f1, f, g

False: e1, e

I got this question CORRECT.

Ex 3.6

$$f(n) = (\lg n)^{\lg n}$$

As I check, I think I got this question CORRECT. However this answer didn't listed on solution. Here's how we prove it:

$$\log(n^k) = k \times \log n \quad (1)$$

$$\log((\log n)^{\log n}) = \log \log n \times \log n \quad (2)$$

$$\log(c^n) = n \times \log c \quad (3)$$

Apparently, (1) << (2) << (3)

Ex 3.13

a)

$$\begin{cases} T(1) = 1 \\ T(n) = \frac{T(n)}{3} + \frac{T(n-1)}{3} + \frac{2T(n-1)}{3} + \Theta(1) = \frac{T(n)}{3} + T(n-1) + \Theta(1) \end{cases}, n > 1 \quad (5)$$

I got this question CORRECT.

b)

$$\begin{cases} Y(1) = 1 \\ Y(n) = \frac{31Y(n)}{3} + \frac{41Y(n) + 1}{3} + \frac{59Y(n-1) + 26Y(n-1) + 1}{3} = \frac{31Y(n)}{3} + 42Y(n-1) + \frac{2}{3} \end{cases}, n > 1 \quad (6)$$

I got this question CORRECT.

The subtle error of this problem is: The factor for $Y(n)$ on the right side of the equation is greater than 1, plus both $Y(n)$ and $Y(n-1)$ cannot be negative, thus $Y(n)$ have to be $+\infty$ to satisfy this equation.

c)

$$\begin{cases} S(1) = 0 \\ S(n) = \frac{S(n) + 1}{3} + \frac{S(n-1) + 1}{3} + \frac{S(n-1) + 1 + S(n-1) + 1}{3} = \frac{S(n)}{3} + S(n-1) + \frac{4}{3} \end{cases}, n > 1 \quad (7)$$

I got this question CORRECT.

Ex 3.17

$$\begin{cases} Name(1) = 1 \\ Name(2n) = 2Name(n) - 1, n \geq 1 \\ Name(2n+1) = 2Name(n) + 1, n \geq 1 \end{cases} \quad (8)$$

I got this question CORRECT.

Ex 3.37

a)

$$U(99) = 99 + U(89.1) = 188.1 + U(80.19) = 268.29 + U(72.171) = 340.461 + U(64.9539) = 1314.7695 \quad (9)$$

$$Z(99) = 99 + Z(90) = 189 + Z(81) = 270 + Z(73) = 343 + Z(66) = 409 + Z(60) = 1309 \quad (10)$$

$$U(99) > Z(99) \quad (11)$$

I got this question CORRECT.

b)

The reason that $U(99) > Z(99)$ is, for some x , we have:

$$U(x) = x + \left(\frac{9}{10}\right)x + \left(\frac{9}{10}\right)^2x + \dots + \left(\frac{9}{10}\right)^{n-1}x + 15 \times \left(\frac{9}{10}\right)^n x \quad (12)$$

$$Z(x) = x + \left\lceil \left(\frac{9}{10}\right)x \right\rceil + \left\lceil \frac{9}{10} \left\lceil \left(\frac{9}{10}\right)x \right\rceil \right\rceil + \dots + \left\lceil \left(\frac{9}{10}\right) \dots \left\lceil \left(\frac{9}{10}\right)x \right\rceil \dots \right\rceil + 15 \times \left\lceil \left(\frac{9}{10}\right) \dots \left\lceil \left(\frac{9}{10}\right)x \right\rceil \dots \right\rceil \quad (13)$$

$n \text{ of } \frac{9}{10}s$ $n+1 \text{ of } \frac{9}{10}s$

Where in equation (12), $\left(\frac{9}{10}\right)^{n-1}x > 65$ and $\left(\frac{9}{10}\right)^n x \leq 65$, and in equation (13), $\left\lceil \left(\frac{9}{10}\right) \dots \left\lceil \left(\frac{9}{10}\right)x \right\rceil \dots \right\rceil > 65$
 $n \text{ of } \frac{9}{10}s$

For the first n items in $U(x)$ and $Z(x)$, they are approximately the same, although $Z(x)$ is a little greater(or equal) than $U(x)$ for each of these items.

For the last 1 item in $U(x)$ and last 2 items in $Z(x)$, there's a constant factor 15, thus the last item in $U(x)$ is larger than the sum of last two items in $Z(x)$. This is what happened in $U(99)$ and $Z(99)$.

Note: Apparently, x cannot be too large since the first n elements also slightly affect the sums of the equations. Actually there are only 7 ranges that satisfy $U > Z$

[80.00000, 80.24690]

[88.88890, 89.16322]

[98.44878, 99.07026]

[109.51730, 110.07806]

[121.72730, 122.30896]

[135.27650, 135.89884]

[150.64110, 150.99872]

I got this question CORRECT. Notes : I didn't figure out how to proof that there are no more numbers(other than the ones in these 7 ranges) that satisfy $U > Z$. Just the program cannot find anymore. The intuition only give us that x cannot be very large, however dealing with real numbers is much more tougher with *ceil*

c)

The answer is no.

Under this condition, the first n items in $U(x)$ and $Z(x)$ are not changed. For the sum of these items, $Z(x)$ is slightly larger(or equal, for some x) than $U(x)$.

The rest item in $U(x)$ is C . The rest items in $Z(x)$ is $\left\lceil \left(\frac{9}{10}\right) \dots \left\lceil \left(\frac{9}{10}\right)x \right\rceil \dots \right\rceil + C$. Apparently, the later one is larger.
 $n \text{ of } \frac{9}{10}s$

Thus, under this condition, for all x , $U(x) \leq Z(x)$

I got this question CORRECT. Notes: Each item in the equations is corresponding to each level in the recursion tree.