The blind-spot of dynamically regularized factorizations

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 Had just finished preliminary tests for distributed (C++11 over MPI) sparse primal-dual IPM for

$$\min_{x} \frac{1}{2} x^{T} Q x + c^{T} x$$

s.t. $Ax = b$, $Gx + s = h$, $s \ge 0$

when I signed up for this talk.

- Had tested synthetic BP, BPDN, CP, DS, EN, LAV, NNLS, SVM, and TV with great success
- ▶ Decided that solving $\min_{x} ||Ax b||_{\infty}$ meant that I should be able to solve

$$Ax = b$$

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- ► Sparse Cholesky can be vastly less expensive than Bunch-Kaufman (same for LU w/ and w/o pivoting)
- In a distributed-memory context, dynamic pivoting prevents a priori load balancing
- ► Forming A^HA or AA^H should be avoided for both stability and sparsity reasons
- Many sparse-direct solvers (e.g., Pardiso, SuperLU_Dist, and WSMP) support dynamically regularized pivots
- ► Each diagonal modification implies a small **and** rank-one perturbation of the original problem, and so, with exact arithmetic, we could expect a great preconditioner
- ► Interior Point solvers (e.g., [Altman/Gondzio-1998]) often intelligently choose the sign of the perturbation of quasi-semidefinite matrix [Vanderbei-1993]

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Suppose someone handed you

and asked you to solve Ax = b.

Despite

$$diag(L) = ones(n, 1), ||L||_{max} = 2$$

for even n = 100, the residual will tend to be O(1).

Because A is extremely ill-conditioned, there is no contradiction with the good backwards stability implied by |A| and $|L||L^H|$.

While the $\lambda(L) = \{1\}$, the ϵ -pseudospectrum is a large disk containing the origin

Suppose someone handed you

$$A = \begin{pmatrix} 1 & 2 & & & & \\ 2 & 5 & 2 & & & \\ & 2 & 5 & \ddots & & \\ & & \ddots & \ddots & & \\ & & & & 5 & 2 \\ & & & & & 2 & 5 \end{pmatrix} = \begin{pmatrix} 1 & & & & & \\ 2 & 1 & & & & & \\ & 2 & 1 & & & & \\ & & \ddots & \ddots & & & \\ & & & & 1 & 2 & \\ & & & & & 2 & 1 \end{pmatrix} \begin{pmatrix} 1 & 2 & & & & \\ 1 & 2 & & & & \\ & 1 & 2 & & & \\ & & & \ddots & & & \\ & & & & 1 & 2 \\ & & & & & 1 \end{pmatrix}$$

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A well-conditioned extension

Consider

$$K = \begin{pmatrix} 4J_{1/2}(n)^TJ_{1/2}(n) & I\\ I & -I \end{pmatrix},$$

which, for n = 100, has cond(K) = 14.719

Then LDL^T factorization without pivoting will fail catastrophically before invoking regularization

K is an example of a well-conditioned quasi-definite matrix where dynamic regularization dramatically fails.

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The stability of quasi-definite *LDL*^H factorization

Suppose

$$K = \begin{pmatrix} G & A \\ A^H & -H \end{pmatrix}, \quad G, H \succ 0.$$

[Gill et al.-1996] showed that

Econd(K) =
$$\left(1 + \frac{\max\{\|A^HG^{-1}A\|_2, \|AH^{-1}A^H\|_2\}}{\|K\|_2}\right)$$
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which allows for the failure of the previous example since $\|G^{-1}\|_2$ was very large.

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Preconditioning with regularization

[Saunders-1995,1996] suggests instead using a factorization of the matrix

$$\mathcal{K}_{\delta,\delta} \equiv egin{pmatrix} G + \delta I & A \ A^H & -H - \delta I \end{pmatrix}$$

as a preconditioner, where $\delta = 10^{-4}$ is typical.

He showed that, for $\sigma_{\min}(A) < \delta < ||A||_2$, G = H = 0,

$$\mathsf{Econd}(K_{\delta,\delta}) pprox \left(rac{\|A\|_2}{\delta}
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which might often be 10⁸. When *A* has a dense row, the former is often much preferred to $A^{H}A + \delta^{2}I$.

Elemental's IPMs now default to $\delta = 10^{-4}$ and use the iteratively refined solution against $K_{\delta,\delta}$ as a preconditioner for FGMRES(10).

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There is a long history [Siegel-1965,Björck-1967,Hatchel-1974,Björck-1992] of

$$\min_{x} \|Ax - b\|_2 \Rightarrow A^H(Ax - b) = 0$$

$$K_{\alpha} \begin{pmatrix} s \\ x \end{pmatrix} \equiv \begin{pmatrix} \alpha I & A \\ A^H & 0 \end{pmatrix} \begin{pmatrix} r/\alpha \\ x \end{pmatrix} = \begin{pmatrix} b \\ 0 \end{pmatrix}.$$

- ▶ cond(K_{α}) roughly varies between cond(A) and cond(A)²
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- The (carefully-scaled!) augmented system is likely to have a sparser factorization than the normal equations (and they share the same *effective* conditioning)
- Given a Lanczos procedure for estimating σ_{min}(A), we have a powerful tool for distributed sparse-direct linear and least squares problems
- Augmented systems should be revived, though a careful usage of extended precision may be needed; perhaps computation of |A| and |L||D||L|^T could predict the need for refactoring in extended precision

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