## COMP2610/COMP6261 - Information Theory

Tutorial 9: Stream and Noisy Channel Coding

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Week 11 (16th – 20th Oct), Semester 2, 2017

- 1. Complete arithmetic coding (Question 4, Tutorial 8) from previous tutorial if you have not completed.
- 2. Consider a channel with inputs  $\mathcal{X} = \{a, b, c\}$ , outputs  $\mathcal{Y} = \{a, b, c, d\}$ , and transition matrix

$$Q = \begin{bmatrix} 0.5 & 0 & 0 \\ 0.5 & 0.5 & 0 \\ 0 & 0.5 & 0.5 \\ 0 & 0 & 0.5 \end{bmatrix}.$$

- (a) Asyming p. 1027 (25714), Paris by munatinformation I(Y) between the input and output of the channel.
- (b) Assuming  $p_{\chi} = (0.25, 0.25, 0.5)$ , what is the average probability of error of the channel?
- (c) Calvin claims that he has constructed a block code for Q with rate 0.01 bits per transmission and maximal block error probability 8%, as his aim was tible? Ustry you can so transmission.
- (d) Hobbes claims that he has constructed a block code for Q with rate 100 bits per transmission and maximal block error probability 1%. Is his claim possible? Justify your answer.
- 3. Noisy Coding (Exclared in Well-Chat powcoder
  - (a) A binary erasure channel with input  $x \in \{0,1\}$  and output  $y \in \{0,?,1\}$  has transition matrix

$$Q_E = \begin{bmatrix} 1 - q & 0 \\ q & q \\ 0 & 1 - q \end{bmatrix}.$$

Find the mutual information I(X;Y) between the input and output for a general distribution  $\mathbf{p}_X=(p_0,p_1)$  over inputs. Show that the capacity of this channel is  $C_E=1-q$  bits.

(b) A Z-channel has transition probability matrix

$$Q_Z = \begin{bmatrix} 1 & q \\ 0 & 1 - q \end{bmatrix}.$$

Show that, using a (2,1) code, that two uses of a Z-channel can be made to emulate one use of an erasure channel, and state the erasure probability of that erasure channel. Hence show the capacity of the Z-channel  $C_Z \geq \frac{1}{2}(1-q) = \frac{1}{2}C_E$  bits.

Explain why this result is an inequality rather than an equality.