

Numerical Optimisation
Nonsmooth optimisation

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Lecture 16

Subgradient

For convex differentiable function $f : \mathbb{R}^n \rightarrow \mathbb{R}$ it holds

$$f(y) \geq f(x) + \nabla f(x)^T(y - x).$$

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A vector $g \in \mathbb{R}^n$ is a subgradient of a function f at $x \in \text{dom } f$ if

$$f(y) \geq f(x) + g^T(y - x) \quad \forall y \in \text{dom } f.$$

- $f(x) + g^T(y - x)$ is affine global underestimator
- g is a subgradient of f at x if $(g, -1)$ supports the epigraph of f at $(x, f(x))$

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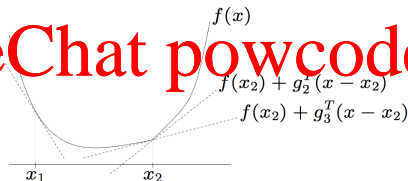


Figure: $\partial f(x_1) = \{\nabla f(x_1)\} = \{g_1\}$, $\partial f(x_2) = [g_3, g_2]$. Fig. from S. Boyd, EE364b, Stanford University.

Subdifferential

A function f is called **subdifferentiable** at x if there exists at least one subgradient at x .

Subdifferential of f at x $\partial f(x)$, is the set of all subgradients of f at x .

$\partial f(x)$ is a closed convex set (can be empty) even if f is not convex.

Proof: It follows from it being intersection of infinite set of halfspaces

$$\partial f(x) = \bigcap_{z \in \text{dom } f} \{g : f(z) \geq f(x) + g^T(z - x)\}.$$

If f is continuous at x , then the $\partial f(x)$ is bounded.

If $f(x)$ is convex and $x \in \text{relint dom } f$

- $\partial f(x)$ is nonempty and bounded
- $\partial f(x) = \{\nabla f(x)\}$ iff f differentiable at x (abuse of notation!)

Minimum of nondifferentiable function (unconstraint)

A point x^* is a minimiser of a function f (not necessarily convex)

iff f is subdifferentiable at x^* and

$$0 \in \partial f(x^*),$$

i.e. $g = 0$ is a subgradient of f at x^* .

Proof: This follows directly from $f(x) \geq f(x^*)$ for all $x \in \text{dom } f$.

f is subdifferentiable at x^* with $0 \in \partial f(x^*)$ is equivalent to

$f(x) \geq f(x^*) + 0^T(x - x^*)$ for all $x \in \text{dom } f$.

The condition $0 \in \partial f(x^*)$ reduces to $\nabla f(x^*) = 0$ when f is convex and differentiable at x^* . Note, that in that case also it is a necessary and sufficient condition.

Minimum of nondifferentiable function (constraint)

Convex constraint optimisation problem

$$\min_{x \in \mathbb{R}^n} f(x) \quad (\text{COP})$$

subject to $f_i(x) \leq 0, \quad i = 1, \dots, m,$

where

- $f, f_i : \mathbb{R}^n \rightarrow \mathbb{R}$ is convex hence subdifferentiable
- strict feasibility holds (Slater's conditions)

Generalised KKT conditions:

x^* is primal optimal and λ^* dual optimal iff

$$\begin{aligned} f(x^*) &\leq 0, \\ \lambda_i^* &\geq 0, \end{aligned} \quad (\text{KKT})$$

$$0 \in \partial f(x^*) + \sum_{i=1}^m \lambda_i^* \partial f_i(x^*),$$

$$\lambda_i^* f_i(x^*) = 0$$

Directional derivatives and subdifferential

For a convex function the *directional derivative* at x in the direction v is

$$f'(x; v) = \lim_{t \rightarrow 0} \frac{f(x + tv) - f(x)}{t}$$

The limit always exists for a convex function, though it can be $\pm\infty$. If f is finite in a neighbourhood of x , then $f'(x; v)$ exists.

f is differentiable at x if for some g (which is $\nabla f(x)$) and all $v \in \mathbb{R}^n$ we have $f'(x; v) = g^T v$ ($f'(x; v)$ is a linear function of v).

The directional derivative $f'(x; v)$ of a convex function f satisfies

$$f'(x; v) = \sup_{g \in \partial f(x)} g^T v.$$

Proof idea: Note that $f'(x; v) \geq \sup_{g \in \partial f(x)} g^T v$ by the definition of the subgradient $f(x + tv) - f(x) \geq tg^T v$ for any $t \in \mathbb{R}$ and $g \in \partial f(x)$. Other direction: show that all affine functions below $v \rightarrow f'(x; v)$ may be taken to be linear.

Weak subgradient calculus: formulas for finding *one* $g \in \partial f(x)$.

If you can compute f , you can usually compute one subgradient.

Many algorithms require only one subgradient.

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Strong subgradient calculus: formula for finding *the whole* subdifferential $\partial f(x)$

Optimality conditions and some algorithms require the whole differential.

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Basic rules:

- non-negative scaling: for $\alpha > 0$, $\partial(\alpha f) = \alpha \partial f$
- addition: $\partial(f_1 + f_2) = \partial f_1 + \partial f_2$
- affine transformation: $g(x) = f(Ax + b)$,
 $\partial g(x) = A^T \partial f(Ax + b)$
- finite point wise maximum: $f = \max_{i=1, \dots, m} f_i$,
 $\partial f(x) = \text{Co} \bigcup \{ \partial f_i(x) : f_i(x) = f(x) \}$ (convex hull of a union of subdifferentials of active functions at x)

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Subgradient and descent direction

p is a descent direction for f at x if $f'(x; p) < 0$.

If f is differentiable, $-\nabla f$ is always a descent direction (except when it is 0).

For a nondifferentiable convex function f , $p = -g$, $g \in \partial f(x)$ need not to be a descent direction.

Example: $f(x) = |x_1| - 2|x_2|$

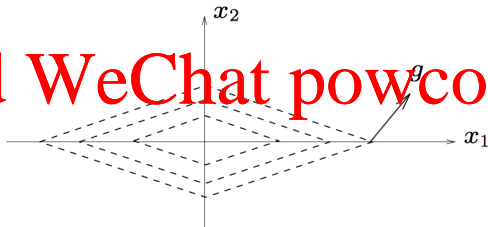


Figure: Fig. from S. Boyd, EE364b, Stanford University.

Subgradient and distance to sublevel set

For a convex f , if $f(z) < f(x)$, $g \in \partial f(x)$, then for small $t > 0$

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Thus $-g$ is descent direction for $\|x - z\|_2$, for any z with $f(z) < f(x)$.

Proof:

$$\begin{aligned}\|x - tg - z\|_2^2 &= \|x - z\|_2^2 - 2tg^T(x - z) + t^2\|g\|_2^2 \\ &\leq \|x - z\|_2^2 - \underbrace{2t(f(x) - f(z))}_{>0} + \underbrace{t^2\|g\|_2^2}_{t: \frac{t}{2}\|g\|_2^2 < f(x) - f(z)}\end{aligned}$$

In particular, choosing $z = x^*$, we obtain that the negative subgradient is a descent direction for distance to optimal point x^* .

Proximal operator of $f : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$

$$\text{prox}_{\lambda f}(v) := \arg \min_x (f(x) + 1/(2\lambda) \|x - v\|_2^2), \lambda > 0 \quad (\text{PROX})$$

Evaluating $\text{prox}_{\lambda f}$ involves solving a convex optimisation problem

Can evaluate numerically via e.g. BFGS, but often the convex problem (PROX) has an analytical solution or at least a specialised linear time algorithm.

Indicator function of a closed convex set, $C \neq \emptyset$

$$I_C(x) = \begin{cases} 0 & x \in C \\ \infty & \text{otherwise} \end{cases}$$

Proximal operator of I_C is the Euclidean projection

$$\text{prox}_{\lambda I_C}(v) = \arg \min_{x \in C} \|x - v\|_2 = \Pi_C(v)$$

Many properties of projection carry over to proximal operator.

Examples of proximal operators

Important special choices of f , for which $\text{prox}_{\lambda f}$ has a closed form:

- $f(x) = \frac{1}{2} \|Px - q\|_2^2$,

$$\text{prox}_{\lambda f}(v) = (P^T P + \lambda^{-1} I)^{-1} (P^T q + \lambda^{-1} v).$$
$$(P^T P + Q)^{-1} = Q^{-1} - Q^{-1} P^T (I + P Q^{-1} P^T)^{-1} P Q^{-1}.$$

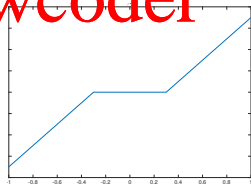
- f is separable i.e. $f(x) = \sum_i^n f_i(x_i)$, proximal operator acts componentwise

$$(\text{prox}_{\lambda f}(v))_i = \text{prox}_{\lambda f_i}(v_i), \quad i = 1, \dots, n$$

- $f(x) = \|x\|_1$

$\text{prox}_{\lambda}(v) = S_{\lambda}(v)$
with elementwise soft thresholding

$$S_{\delta}(x) = \begin{cases} x - \delta & x > \delta \\ 0 & x \in [-\delta, \delta] \\ x + \delta & x < -\delta \end{cases}$$



Examples of proximal operators

Another important example which does not admit close form is
Total Variation, $f(x) = TV(x)$, defined as follows

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$$TV(x) := \sum_{i=1}^{m-1} \sum_{j=1}^{n-1} \sqrt{(x_{i,j} - x_{i+1,j})^2 + (x_{i,j} - x_{i,j+1})^2}$$

$$+ \sum_{i=1}^{m-1} |x_{i,n} - x_{i+1,n}| + \sum_{j=1}^{n-1} |x_{m,j} - x_{m,j+1}|$$

assuming standard reflexive boundary conditions

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$$x_{m+1,j} = x_{m,j}, \quad x_{i,n+1} = x_{i,n}.$$

The proximal operator has to be computed iteratively using
e.g. Chambolle-Pock algorithm (primal dual proximal gradient).

Resolvent of subdifferential operator

Proximal operator

$$\text{prox}_{\lambda f}(v) = \arg \min_x (f(x) + 1/(2\lambda) \|x - v\|_2^2).$$

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The first order condition for the minimiser reads

$$\begin{aligned} 0 &\in \partial f(x) + 1/\lambda(x - v) \\ v - x &\in \lambda \partial f(x) \\ v &\in \lambda \partial f(x) + x \end{aligned}$$

$$\text{prox}_{\lambda f}(v) = (I + \lambda \partial f)^{-1} v$$

Mapping $(I + \lambda \partial f)^{-1}$ is called **resolvent** of operator ∂f .

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x^* minimises f iff x^* is a fixed point

$$x^* = \text{prox}_f(x^*)$$

Moreau envelope or Moreau-Yosida regularisation of f

$$M_{\lambda f}(v) = \inf_x (f(x) + 1/(2\lambda) \|x - v\|_2^2).$$

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$M_{\lambda f}$ is a smoothed (regularised) version of f .

- always has full domain
- always continuously differentiable
- has the same minimisers as f

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Can show that $M_f = (f^* + 1/2 \|\cdot\|_2^2)^*$.

Example: Moreau envelope of $|\cdot|$ is the Huber function

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$$M_{|\cdot|}(x) = \begin{cases} \frac{\lambda}{2} x^2 & |x| \leq 1 \\ |x| - \frac{1}{2} & |x| > 1 \end{cases}$$

Moreau decomposition: $v = \text{prox}_f(v) + \text{prox}_{f^*}(v)$ is
generalisation of orthogonal decomposition $v = \Pi_W(v) + \Pi_{W^\perp}(v)$.
It follows from Moreau decomposition that $(I_W)^* = I_{W^\perp}$.

$$\min_x f(x) + g(x) \quad \text{subject to } x \in \mathbb{E} \quad (1)$$

- \mathbb{E} : finite dimensional Euclidean space with inner product $\langle \cdot, \cdot \rangle$ and a self dual norm $\|\cdot\| = \langle \cdot, \cdot \rangle^{1/2} = \|\cdot\|_*$, e.g. space of $n \times m$ images, $\mathbb{R}^{n \times m}$
- $f: \mathbb{E} \rightarrow \mathbb{R}$ continuously differentiable with Lipschitz continuous gradient, $\|\nabla f(x) - \nabla f(y)\| \leq L(f)\|x - y\|, \forall x, y \in \mathbb{E}$.
- $g: \mathbb{E} \rightarrow (-\infty, \infty]$ proper closed convex.

From first order optimality condition we have

$$\begin{aligned} 0 &\in \nabla f(x^*) + \partial g(x^*) \\ 0 &\in \tau \nabla f(x^*) + \tau \partial g(x^*) - x^* + x^* \\ (I + \tau \partial g)(x^*) &\in (I - \tau \nabla f)(x^*) \\ x^* &= (I + \tau \partial g)^{-1}(I - \tau \nabla f)(x^*) \end{aligned} \quad (2)$$

$$x_k = \text{prox}_{\tau_k g}(x_{k-1} - \tau_k \nabla f(x_{k-1}))$$

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- **Gradient Projection:** $g(x) = l_C(x)$: smooth constrained minimisation, $\tau_k \in (0, 2/L(f))$

$$x_k = \Pi_C(x_{k-1} - \tau_k \nabla f(x_{k-1})).$$

- **Proximal Minimization:** $f(x) = 0$: non-smooth convex minimisation

$$x_k = \arg \min_x \left\{ g(x) + \frac{1}{2\tau_k} \|x - x_{k-1}\|^2 \right\}.$$

- **Iterative Shrinkage Thresholding Algorithm (ISTA):**

$$g(x) = \|x\|_1, f(x) = \|Ax - b\|^2, \tau_k \in (0, 2/L(f))$$

$$x_k = S_{\tau_k}(x_{k-1} - \tau_k \nabla f(x_{k-1})).$$

$$x_k = \text{prox}_{\tau_k g}(x_{k-1} - \tau_k \nabla f(x_{k-1}))$$

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- **Gradient Projection:** $g(x) = I_C(x)$: smooth constrained minimisation, $\tau_k \in (0, 2/L(f))$

$$x_k = \Pi_C(x_{k-1} - \tau_k \nabla f(x_{k-1})).$$

- **Iterative Shrinkage Thresholding Algorithm (ISTA):**

$$g(x) = \|\lambda\|_1, f(x) = \|Ax - b\|^2, \tau_k \in (0, 2/L(f))$$

$$x_k = S_{\tau_k}(x_{k-1} - \tau_k \nabla f(x_{k-1})).$$

Slow convergence, if $\tau_k = \tau = 1/L$, $L \geq L(f)$, $F(x) := f(x) + g(x)$

$$F(x_k) - F^* \leq \frac{L\|x_0 - x^*\|^2}{2k}.$$

Fast Iterative Shrinkage Thresholding Algorithm (FISTA):

Initialize: $y_1 := x_0 \in \mathbb{E}$, $\tau_1 = 1$.

$$\text{Step } k : \quad x_k = \text{prox}_{1/L}(g) \left(y_k - \frac{1}{L} \nabla f(y_k) \right)$$

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$$\tau_{k+1} = \frac{1 + 4F_k^2}{2}$$

$$y_{k+1} = x_k + \frac{\tau_k - 1}{\tau_{k+1}} (x_k - x_{k-1}).$$

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Convergence, if $\tau_k = \tau = 1/L$, $L \geq L(f)$, $F(x) := f(x) + g(x)$

$$F(x_k) - F^* \leq \frac{2L \|x_0 - x^*\|^2}{(k+1)^2}.$$

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Fast Projection Gradient [Nesterov'83]: $g(x) = l_c(x)$

$$x_k = \Pi_C \left(y_k - \frac{1}{L} \nabla f(y_k) \right).$$

More details on Nesterov algorithm see e.g. [http:](http://www.seas.ucla.edu/~vandenbe/236C/lectures/fgrad.pdf)

[//www.seas.ucla.edu/~vandenbe/236C/lectures/fgrad.pdf](http://www.seas.ucla.edu/~vandenbe/236C/lectures/fgrad.pdf)

Review: Optimisation with equality constraints

Let $f : \mathbb{R}^n \rightarrow \mathbb{R} \cup \{+\infty\}$, closed, proper and convex.

Primal problem

$\min_x f(x) \quad \text{subject to } Ax = b$

Lagrangian

$$\mathcal{L}(x, y) = f(x) + y^T (Ax - b)$$

Dual function

$$g(y) = \inf_x \mathcal{L}(x, y) = -f^*(-A^T y) - b^T y$$

y : dual variable (Lagrange multiplier),

f^* : convex conjugate of f (f^* is convex and closed even if f is not).

Dual problem (always concave, $y^* \leq x^*$, $y^* = x^*$ if strong duality holds)

$$\max_y g(y). \quad (3)$$

Gradient descent for primal problem (assuming f continuously differentiable)

$$x_{k+1} = x_k - \tau_k \nabla f(x_k)$$

Gradient ascent for dual problem (assuming g continuously differentiable)

$$x_{k+1} = \arg \min_x L(x, y_k)$$

$$y_{k+1} = y_k + \tau_k \underbrace{(Ax_{k+1} - b)}_{=\nabla g(y_k)}$$

Remark: the primal update is part of evaluation of $\nabla g(y_k)$

- + for separable f it leads to a parallel algorithm.
- various conditions necessary for convergence e.g. strict convexity of f , $f(x) < \infty, \forall x$.

Augmented Lagrangian

Augmented Lagrangian

$$\mathcal{L}_\rho(x, y) = f(x) + y^T(Ax - b) + \rho/2 \|Ax - b\|_2^2, \quad \rho > 0 \quad (\text{AL})$$

Equivalent to Lagrangian of an equivalent problem (for all feasible x the quadratic term equals 0)

$$\min_x f(x) + \rho/2 \|Ax - b\|_2^2, \quad \text{subject to } Ax = b.$$

Method of multipliers (MM): dual ascent applied to (AL)

$$x_{k+1} = \arg \min \mathcal{L}_\rho(x, y_k)$$

$$y_{k+1} = y_k + \rho \underbrace{(Ax_{k+1} - b)}_{= \nabla_y \mathcal{L}_\rho(x_{k+1}, y_k)}$$

Using ρ as the step size guarantees dual feasibility of (x_{k+1}, y_{k+1}) :

$$0 = \nabla_x \mathcal{L}_\rho(x_{k+1}, y_k) = \nabla f(x) + A^T y_k + \rho A^T (Ax - b) \Big|_{x=x_{k+1}} = \\ \nabla f(x_{k+1}) + A^T y_{k+1} =: s_{k+1} = 0.$$

+ converges under more general conditions

– augmented Lagrangian is non-separable.

Alternating Directions Methods of Multipliers (ADMM)

Blend separability of dual ascent with superior convergence of MM:

$$\min_{x \in \mathbb{R}^n, z \in \mathbb{R}^m} f(x) + g(z) \quad \text{subject to } Ax + Bz = c \quad (4)$$

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with $A \in \mathbb{R}^{p \times n}$, $B \in \mathbb{R}^{p \times m}$, $c \in \mathbb{R}^p$ and $g: \mathbb{R}^m \rightarrow \mathbb{R} \cup \{\infty\}$ closed, proper and convex.

The equality constraint comes from the split of the variable into x and z with the objective function separable across the splitting.

Augmented Lagrangian

$$L_\rho(x, z, y) = f(x) + g(z) + y^T(Ax + Bz - c) + \rho/2 \|Ax + Bz - c\|_2^2,$$

ADMM

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$$x_{k+1} = \arg \min_x L_\rho(x, z_k, y_k)$$

$$z_{k+1} = \arg \min_z L_\rho(x_{k+1}, z, y_k)$$

$$y_{k+1} = y_k + \rho(Ax_{k+1} + Bz_{k+1} - c).$$

Alternating Directions Methods of Multipliers (ADMM)

Blend separability of dual ascent with superior convergence of MM:

$$\min_{x \in \mathbb{R}^n, z \in \mathbb{R}^m} f(x) + g(z) \quad \text{subject to } Ax + Bz = c \quad (4)$$

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The equality constraint comes from the split of the variable into x and z with the objective function separable across the splitting.

Augmented Lagrangian

$$L_\rho(x, z, y) = f(x) + g(z) + y^T(Ax + Bz - c) + \rho/2 \|Ax + Bz - c\|_2^2,$$

Dual ascent on \mathcal{L}_ρ (joint minimisation)

$$(x_{k+1}, z_{k+1}) = \arg \min_{x, z} L_\rho(x, z, y_k)$$

$$y_{k+1} = y_k + \rho(Ax_{k+1} + Bz_{k+1} - c).$$

ADMM: scaled form

Augmented Lagrangian

$$\begin{aligned} L_{\rho}(x, z, y) &= f(x) + g(z) + y^T (Ax + Bz - c) + \rho/2 \|Ax + Bz - c\|_2^2, \\ &= f(x) + g(z) + y^T r + \rho/2 \|r\|_2^2, \\ &= f(x) + g(z) + \rho/2 \|r + u\|_2^2 - \rho/2 \| \underbrace{u}_{u = (1/\rho)y} \|_2^2, \end{aligned}$$

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with $u = (1/\rho)y$ the scaled dual variable.

ADMM: scaled form

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$$\begin{aligned} x_{k+1} &= \arg \min_x f(x) + \rho/2 \|Ax + Bz_k - c + u_k\|_2^2 \\ z_{k+1} &= \arg \min_z g(z) + \rho/2 \|Ax_{k+1} + Bz - c + u_k\|_2^2 \\ u_{k+1} &= u_k + Ax_{k+1} + Bz_{k+1} - c. \end{aligned}$$

Assume in addition that the unaugmented Lagrangian \mathcal{L} has a saddle point.

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For f, g proper, closed, convex, it follows that strong duality holds (no explicit assumptions on A, B, c).

Under these assumptions the ADMM iterates satisfy

- Residual convergence: $r^k \rightarrow 0$ as $k \rightarrow \infty$ i.e. the iterates approach feasibility.
- Objective convergence: $f(x^k) + g(z^k) \rightarrow p^*$ as $k \rightarrow \infty$ i.e. the objective function of the iterates approach the optimal value
- Dual variable convergence: $y^k \rightarrow y^*$ as $k \rightarrow \infty$, where y^* is a dual optimal point.

Note, that x^k, z^k need not converge to optimal points, although such a result can be shown under additional assumptions.

Optimality conditions

Necessary and sufficient optimality conditions for ADMM

$$Ax^* + Bz^* - c = 0 \quad \text{primal feasibility}$$

$$0 \in \partial f(x^*) + A^T y^* \quad \text{dual feasibility}$$

$$0 \in \partial g(z^*) + B^T y^* \quad \text{dual feasibility}$$

As for MM, it follows from $z_{k+1} = \arg \min_z \mathcal{L}_\rho(x_{k+1}, z, y_k)$ that z_{k+1} and y_{k+1} always satisfy the last equation.

From $x_{k+1} = \arg \min_x \mathcal{L}_\rho(x, z_k, y_k)$ we have

$$0 \in \partial f(x_{k+1}) + A^T y_k + \rho A^T (Ax_{k+1} + Bz_k - c)$$

$$= \partial f(x_{k+1}) + A^T (y_k + \rho r_{k+1} + \rho B(z_k - z_{k+1}))$$

$$= \partial f(x_{k+1}) + A^T y_{k+1} + \rho A^T B(z_k - z_{k+1})$$

or equivalently

$$s_{k+1} := \rho A^T B(z_{k+1} - z_k) \in \partial f(x_{k+1}) + A^T y_{k+1},$$

which can be interpreted as dual feasibility condition and s_{k+1} is the *dual residual* at iteration $k + 1$.

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