Assignment 4 v1

Assignment 4 requires **one project** and **one brief essay**.

Project

An implementation of the **synchronous EIG-based Byzantine agreement** algorithm, working on the complete graph.

The project shall consist of a single source file, containing definitions for **Node** and for **Arcs**, an adapted version of our previous Arcs (not part of the algorithm).

In the submitted version, your own **standard output stream** (Console.Out) must fully conform to the sample output described below. This is very important for our automarker.

The **appendix** contains snapshots with our visual Byzantine demo running a sample scenario.

Essay

Select and critically **review** any interesting paper(s) which discuss(es) one of the following topics (your choice):

- **Byzantine** algorithms and **blockchains**.
- **Authenticated Byzantine** algorithms.

- **Asynchronous Byzantine** algorithms.
- **Practical Byzantine** algorithms (less costly).

Your review should be structured and formatted as a short conference paper, well-argued and supported by additional evidence, e.g. references/citations. It must adhere to a strict page limit, **up to 3 pages** (without references).

Note: Such topics appear under various names, such as: Byzantine / distributed / fault-tolerant generals / algorithm / agreement / consensus / protocol … not all for the same problem, but many are. E.g. just from the Dijkstra prize list: 2001, 2005, 2007, 2010, 2015, 2017 (https://en.wikipedia.org/wiki/Dijkstra_Prize)

If needed, please don't hesitate to check with us about the suitability of your selected topic / papers.

Config

Arcs reads the config file from **stdin**, e.g.:

```
4 2 0                                   // N L V0
1 0 1 0 0 1 1 011 011 011 011           // P V F script
2 0 0                                   // P V F
4 1 0
3 1 0
```

There are *N+1* lines in total, where *N* = # of process nodes, *N* in 1..9. End -of-line comments are optional and here indicate the meaning of each value: *N* = # of process nodes; *L* = # of EIG levels; V_0 (well known); *P* = process number (1..*N*); *V* = initial choice; *F* = failure flag (0: loyal, 1: faulty).

Each faulty process line *P* continues with a script indicating the messages that that process *P* is expected to send. Conceptually, there are *L*·*N* blocks of values, in order:

- *S*=1: *N* blocks of length 1
- *S*=2: *N* blocks of length *N*-1
- *S*=3: *N* blocks of length *N*-2
- …
- *S*=*L*: *N* blocks of length *N*-*L*+1

For each given level *S*, the *N* blocks are to be sent to processes 1, 2, …, *N*, in this order. Each block represents the level *S*-1 values that need to be sent, in left-to-right order (considering an ordered EIG tree) - more precisely, what values that faulty process claims to have recorded at level *S*-1.

Faulty-script conventions (similar to our **demo**):

- legal sequence digits are **0**, **1**, **2, 3**
- **0** and **1** indicate the two values exchanged by loyal participants
- **2** indicates **null**, missing/corrupt messages – to be replaced by $\mathbf{v_0}$
- **3** indicates that the faulty participant sends the **correct** value at this position
- **illegal** characters are assumed to be **2**
- **missing** characters are assumed to a **repeat** of the **last** preceding digit (0, 1, 2, 3)
- internal faulty script **spaces** are **optional** and only relevant for readability
- **extra** characters are **discarded**

For example, the following faulty scripts are equivalent:

```
0 0 0 1 011 021 112 222
0001 0110211122
000101102111xxxxxxxxxxxxxxxxx
```

Workflow

There are $2 \cdot L+1$ steps, numbered $S = 0, 1, 2, \ldots, 2 \cdot L$. Step $S = 0$ is the initialisation; steps $S = 1, 2, \ldots, L$ are the top-down messaging rounds; steps $S = L+1, L+2, \ldots, 2 \cdot L$ are the bottom-up evaluation.

S=0: Arcs sends init messages to each of the *N* processes, in order *P* = 1, 2, …, *N*. Each init message contains *N*, *P*, *V*, *V0*, *L*. For a faulty node, the message also includes the required faulty script. Process *P* sets its EIG root (level 0) value to *V*, and then prints this value in the format: [0 *P V*]. Finally, its responds by sending its level *S*=1 messages to Arcs.

S=1, 2, … *L*-1: Arcs sends all due level *S* messages to each of the *N* processes, in order *P* = 1, 2, …, *N*. Process *P* fills its level *S* top-down values, and then prints these, in left-to-right order, in the format [*S P values*], separating sibling groups by single spaces. Finally, it responds by sending its level *S*+1 messages to Arcs. Note that each sibling group has *N*+1-*S* values.

S=*L*. Same as above, but we are at the leaves level, where there are no more messages to send. Conceptually, top-down values become bottom-up values. Processes respond with "empty" messages.

S=*L*+1, *L*+2, …, 2·*L*. Arcs send "empty" messages to each of the *N* processes, in order *P* = 1, 2, …, *N*. Process *P* evaluates its level 2·*L*-bottom-up values, using the same format as above.

Obviously, for *S*=2·*L*, processes print their final decisions, in the format [*S P V*].

Sample printout corresponding to the above config file (*N*=4, etc):

```
0 1 0
0 2 0
0 3 1
0 4 1
1 1 0011
1 2 0011
1 3 1011
1 4 1011
2 1 011 000 111 111
2 2 011 000 111 111
2 3 011 000 111 111
2 4 011 000 111 111
3 1 1011
3 2 1011
3 3 1011
3 4 1011
4 1 1
4 2 1
4 3 1
4 4 1
```


Notes

You are free to select actual message format that suits you best. Important is to design a credible simulation, where all contacts between nodes is realised via messages gathered and forwarded by Arcs.

Please don't hesitate to ask if you need clarifications. Thanks!

Submit electronically:
(1) Your programs, as single source file, to the COMPSCI automarker, and

(2) Your essay as PDF, to Canvas.

Deadline: Monday 12 October, 2020, 23:00.

Do not leave it for the last minute, please. Remember that you can resubmit and, by default, we only consider your last submission.

Late submissions will incur penalties, 0.25% off for each hour late, for up to eight days.

APPENDIX