CS 314 Principles of Programming Languages

Lecture 15: Functional Programming Assignment Project Exam Help

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Class Information

- HW5 posted in Sakai.
- HW6 will be posted by the end of today.
- Next week's recitation is a review for midterm. Attendance is encouraged!
- Reminder: **Midterm exam** 11/7, in class, closed book, closed notes. Midterm covers lecture 1-12, hw 1-5, and corresponding book chapters.
- There will be extended office hours next week.

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 Pay attention to Sakai announcements!

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Scheme - Functions as First Class Values (Higher-Order)

Functions as arguments:

(define f (lambda (g x) (g x))

- $(f number? 0) \Rightarrow (number? 0) \Rightarrow \#t$
- (f length '(1 2)) \Rightarrow (length '(1 2)) \Rightarrow 2
- $(f (lambda (x) (* 2 3)) 3) \Rightarrow ((lambda (x) (* 2 3)) 3) \Rightarrow (* 2 3) \Rightarrow 6$

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Scheme - Functions as First Class Values (Higher-Order)

Computation, i.e., **function application** is performed by reducing the initial S-expression (program) to an S-expression that represents a value.

Reduction is performed by **substitution**, i.e., replacing formal by actual parameters in the function body.

Examples for S-expressions that directly represent values, i.e., cannot be further reduced:

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- function values (e.ghttdampodacoderecom
- constants (e.g.: 3, #tadd WeChat powcoder

Computation completes when S-expression cannot be further reduced

Review: Higher-Order Functions (Cont.)

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• (plusn 5) evaluates to antiportion that adds 5 to its argument:

Question: How would you write down the value of (plusn 5)?

$$(lambda (x) (+ 5 x))$$

• ((plusn 5) 6) = ?

Review: Higher-Order Functions (Cont.)

```
In general, any n-ary function

(lambda (x_1 x_2 ... x_n) e)

can be rewritten as a nest of n unary functions:

(lambda (x_1)

(lambda (x_2)

(... (lambda (x_n)) reject Exam Help

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```

Review: Higher-Order Functions (Cont.)

In general, any n-ary function

(lambda
$$(x_1 x_2 ... x_n) e$$
)

can be rewritten as a nest of *n* unary functions:

This translation process istpalled wurding out means that having functions with multiple parameters do not add anything to the expressiveness of the language:

Review: Higher-order Functions: map

```
(define map (lambda (f) 1)

(if (null? 1) '() list

(cons (f (car 1)) (map f (cdr 1)))

) Apply f to the first resemble for Project Exam proof to the rest of l

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```

- map takes two arguments: a function f and a list l
- map builds a new list by applying the function to every element of the (old) list

Review: Higher-Order Functions: map

• Example:

$$(\text{map abs '}(-1\ 2\ -3\ 4)) \Rightarrow (1\ 2\ 3\ 4)$$

 $(\text{map (lambda } (x)\ (+\ 1\ x))\ '(-1\ 2\ 3) \Rightarrow (0\ 3\ 4)$

• Actually, the built-in **map** can have more than two arguments:

$$(map + '(1\ 2\ 3)\ '(4\ 5\ 6)) \Rightarrow (5\ 7\ 9)$$
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Review: More on Higher-Order Functions

reduce

Higher order function that takes a binary, associative operation and uses it to "roll up" a list

```
(define reduce
 (lambda (op l id)
             ( if (null? 1)
                id
                (op (car Assignment Broject Exam) Help
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              (reduce + '(10 20 30) 0) ⇒
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(+ 10 (reduce + '(20 30) 0)) ⇒
Example:
                (+ 10 (+ 20 (reduce + '(30) 0))) \Rightarrow
                (+ 10 (+ 20 (+ 30 (reduce + '() 0)))) \Rightarrow
                (+ 10 (+ 20 (+ 30 0))) \Rightarrow
                60
```

Review: Higher-Order Functions

Compose higher order functions to form compact powerful functions

```
(define sum

(lambda (f l)

(reduce + ( map f l ) 0) ) )

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(sum (lambda (x) (* 2 xh))tp(sl/2p3)) eder.com

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(reduce (lambda (x y) (+ 1 y)) '(a b c) 0) \Rightarrow
```

Lexical Scoping and let, let*, and letrec

All are variable binding operations:

```
LET = let, let*, letrec

(LET ( (v1 e1) (v2 e2) ... (vn en) Assignment Project Exam Help e https://powcoder.com

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```

Lexical Scoping and let, let*, and letrec

• let:

- binds variables to values (no specific order), and evaluate body **e** using bindings
- new bindings are not effective during evaluation of any e_i

• let*:

- binds variables to values in textual order of write-up
- new binding e_{i-1} is effective for next e_i https://powcoder.com

• letrec:

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 bindings of variables to values in no specific order
- independent evaluations of all e_i to values have to be possible
- new bindings effective for all e_i
- mainly used for recursive function definitions

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```
( let ( (a 5) (b 6) )
 (+ab) ) ;; \Rightarrow 11
```

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(let (
$$(a 5) (b 6)$$
)
 $(+ab)$) ;; $\Rightarrow 11$

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(let (
$$(a 5) (b 6)$$
)
 $(+ab)$) ;; $\Rightarrow 11$

(let ((a 5) (b (+ a 6)))
(+ a b)
$$;; \Rightarrow ERROR; a: undefined$$
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```
( let ( (a 5) (b 6) )
    (+ab) ;; \Rightarrow 11
(let ( (a 5) (b (+ a 6)) )
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(let* ( (a 5) (b (+ a 6)) Add WeChat powcoder
   (+ab)
```

```
( let ( (a 5) (b 6) )
    (+ab) ;; \Rightarrow 11
(let ( (a 5) (b (+ a 6)) )
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(let* ( (a 5) (b (+ a 6)) Add WeChat powcoder
    (+ab) ;; \Rightarrow 16
```

Typically used for local definitions for recursive functions

```
(letrec ( (a 5)

(b ( lambda() (+ a 6) ) ) )

( + a (b) ) ) ;; ⇒ 16
```

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Typically used for local definitions for recursive functions

Typically used for local definitions for recursive functions

```
(letrec ( (even? (lambda (x) (or (= x 0) (odd? (- x 1))))) )

(odd? (lambda (x) (and (not (= x 0)) (even? (- x 1))))))

( list (even? Signment Project Exam) Help ;; ⇒ (#f #t #t) https://powcoder.com

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```

Typically used for local definitions for recursive functions

Lexical (Static) Scoping

- The occurrence of a variable matches the lexically closed **binding** occurrence.
- An occurrence of a variable without a matching binding occurrence is called *free*.
- A variable can occur free and bound in an expression.

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We only substitute frentoscuprences of the formal arguments in the function body when making a function application!

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Free and Bound

Consider the following function definition:

```
(define plusn

(lambda (n)

(lambda (x) (+ n x))

n is free

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```

Free and Bound

Consider the following function definition:

```
(define plusn

(lambda (n)

(lambda (x) (+ n x))

n is bound

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```

Environment and Closure

Environment:

• Record the bindings for the variables. If we need the value that a variable denotes, we just look it up in the environment.

An environment is a finite map from variables to values

ρ ∈ Envsīg Variables ject Values Help

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Closures

Pair the environment with a function (lambda abstraction). The environment must contain values for all free variables of the function. The function can only be evaluated in its attached environment.

Such a pairing is called a Closure.

A closure is a pair consisting of an environment and a function definition.

cl
$$\in$$
 Closure $= \{\langle \lambda, \rho \rangle \mid \text{Free Var}(\lambda) \subseteq \text{DOM}(\rho) \}$
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Closures can be used to implement lexical scoping.

They represent lexically scoped function values.

How to Apply a Closure?

How to apply a closure value to actual argument values?

- Let c_v be the closure value < (lambda (x) e), ρ >
- Apply c_v to a value a_v as follows:

Evaluate the body e of the function in the environment ρ of the closure extended by the suppoint of the formal parameter \mathbf{x} to the actual parameter \mathbf{a}_v ($\rho[\mathbf{x} \longrightarrow a_v]$).

How to Apply a Closure: Examples

```
((lambda(\mathbf{x}))
    ((lambda(z) ((lambda(x)(z x)) 3)) (lambda(y)(+ x y)))) 1)
                               closure interpreter
                                                        \{\ \}
((lambda(z) ((lambda(x)(z x)) 3))
                                                        \{x \rightarrow 1\}
   (lambda(y)(+ x y)) Assignment Project Exam Help
((lambda(\mathbf{x})(z \mathbf{x}))
                                  https://powcoder.xom1,
                                  Add WeChat powcoder (lambda(y)(+xy)), \{x\rightarrow 1\})
 3 )
                                                        \{ x \rightarrow 3,
(z x)
                                                         z \rightarrow \langle (lambda(y)(+xy)), \{x \rightarrow 1\} \rangle \}
(+ x y)
                                                         \{x \rightarrow 1, y \rightarrow 3\}
```

Next Lecture

Reading:

- Scott, Chapter 11.1 11.3
- Scott, Chapter 11.7

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