

Midterm Exam  
CS 314, Fall '17  
October 27  
SAMPLE SOLUTION

**DO NOT OPEN THE EXAM**  
**UNTIL YOU ARE TOLD TO DO SO**

Name: \_\_\_\_\_

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**Assignment Project Exam Help**

Section: \_\_\_\_\_

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**RIGHT CORNER**  
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**Instructions**

We have tried to provide enough information to allow you to answer each of the questions. If you need additional information, make a *reasonable* assumption, write down the assumption with your answer, and answer the question. There are **5** problems, and the exam has **8** pages. Make sure that you have all pages. The exam is worth **250** points. You have **80 minutes** to answer the questions. Good luck!

This table is for grading purposes only

1	/ 80
2	/ 60
3	/ 30
4	/ 20
5	/ 60
total	/ 250

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## Problem 1 – Regular Expressions, DFA and Context Free Grammars ( 80 pts)

The context-free grammar  $G$  is specified in Backus-Naur-Form as follows, with  $D$  as the start symbol:

1:  $D ::= a D \mid$   
2:      $b E$   
3:  $E ::= b E \mid$   
4:      $F$   
5:  $F ::= c$

1. Give a leftmost derivation ( $\Rightarrow_L$ ) for the string a b c given the grammar above. (15 pts)

$D \Rightarrow_L aD \Rightarrow_L abE \Rightarrow_L abF \Rightarrow_L abc$

2. Give the LL(1) parse table for the grammar  $G$ . Insert the rule number or leave an entry empty. (35 pts)

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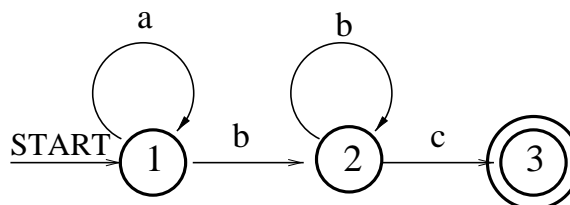
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	a	b	c	eof
F			5	
E		3	4	
D	1	2		

3. Give a regular expression for the language generated by the grammar  $G$ . (15 pts)

$a^*b^+c$

4. Specify a DFA by extending the state transition diagram below. The start state is **state 1**, and the final (accepting) state is **state 3**. You are only allowed to add edges with their appropriate labels, i.e., valid labels are  $a$ ,  $b$ , and  $c$ . Note that an edge may have more than one label. You **must not** add any states. (15 pts)



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## Problem 2 – Context Free Grammars (60 pts)

A *context-free language* is a language that can be specified using a context-free grammar. A *regular language* is a language that can be specified using a regular expression.

For the three languages given below, if the language is context-free, give a compact context-free grammar in Backus-Naur-Form (BNF). If the language is regular, give a compact regular expression using the regular expression syntax introduced in class. If a language is context-free and regular, give both specifications, a BNF and a regular expression. You do not have to justify why you believe a language is not context-free or not regular.

1.  $\{ 0^n 1^{3n} \mid n \geq 0 \}$ , with alphabet  $\Sigma = \{0, 1\}$

$S ::= 0S111 \mid \epsilon$

2.  $\{ a^n b^{3m} c^n \mid n \geq 0, m > 0 \}$ , with alphabet  $\Sigma = \{a, b, c\}$

$S ::= aS \mid B$   
 $B ::= bbbB \mid bbb$

3.  $\{ w \mid w \text{ has at least 2 symbols, but no more than 5} \}$ , with alphabet  $\Sigma = \{0, 1\}$

$S ::= A A B B B$

$A ::= 0 \mid 1$

$B ::= 0 \mid 1 \mid \epsilon$

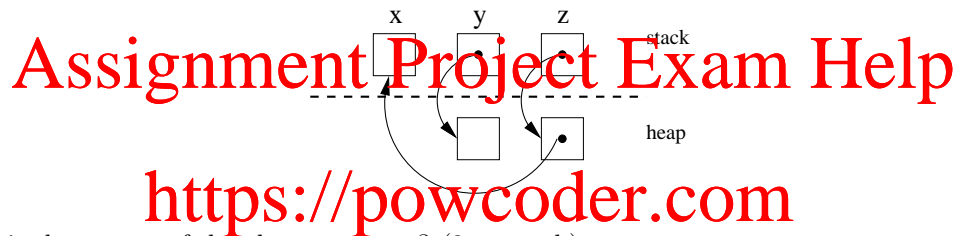
Regular expression:  $(0 \mid 1) (0 \mid 1) (0 \mid 1 \mid \epsilon) (0 \mid 1 \mid \epsilon) (0 \mid 1 \mid \epsilon)$

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### Problem 3 – Pointers and Memory Allocation in C (30 pts)

```
int main() {
    int x;
    int *y;
    int **z;

    z = (int **) malloc (sizeof(int *));
    y = (int *) malloc (sizeof(int));
    x = 3;
    *z = &x;
    *y = x;
    x = x + 3;
    **z = *y + 5;
    printf("x=%d, *y=%d, **z=%d\n", x, *y, **z);
    return 0;
}
```



1. What is the output of the above program? (3 pts each)

x= 8, \*y= 3, \*\*z= 8

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2. Specify, whether the following program objects are allocated on the **stack** (includes global variables), on the **heap**, or **not defined** (2 pts each).

\*x is allocated on the not defined

\*y is allocated on the heap

z is allocated on the stack

x is allocated on the stack

y is allocated on the stack

\*z is allocated on the heap

\*\*y is allocated on the not defined

\*\*z is allocated on the stack

3. Assume the following code segment:

```
int *x;
*x = 5;
printf("%d\n", *x);
```

Is there a problem with this code? Assume that when you ran the code a couple of times, it printed "5". If you believe there is a problem, give a possible "fix" for the problem? (5 pts)

The content of variable `x` is not initialized. However, its content is used as an address of a memory location, and that memory location is assigned the value 5.

To fix the problem, the pointer `x` should be initialized to `NULL` in its declaration, and `x` must point to an object on the heap before it is dereferenced. This object is allocated as follows:

```
x = (int *) malloc(sizeof(int *))
```

This statement should be placed before statement `*x = 5`, i.e., before the dereference operation on `x`.

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## Problem 4 – Compilers vs. Interpreters (20 pts)

To answer this question, please use the following definitions.

**Definition A compiler** maps an input program to a semantically equivalent output program. Note that the input and output language may be the same. □

**Definition An interpreter** maps an input program to the answers it computes; In other words, it executes the program. □

As part of the C project, we used and/or wrote the following programs/commands:

- `gcc` – *usage*: `gcc <program>`
- `compile` – *usage*: `compile <program>`
- `constfolding` – *usage*: `constfolding < <program>`
- `sim` – *usage*: `sim <program>`

Under Unix/Linux, commands can be entered on a single command line if they are separated by a semicolon. For instance, saying `cd foo; ls` will change the current directory to subdirectory `foo`, and will list its files.

In the project, we used several languages, mainly `tinyL`, `RISC` machine code and `ilab` machines object code (executables). **Classify the following commands or the entire sequence of commands** as either compiler or interpreter. Note that you should treat a sequence of commands as a single unit, i.e., as a single composed command. If the single or composed command is a compiler, give its input and output language (e.g.: input language: `C` output language: `tinyL`). For an interpreter, just give its input language.

1. `compile test1`

*Answer (mark one):* compiler: ☒ or interpreter: ☐

input language: `tinyL`, output language: `ILOC RISC machine code`

2. `compile test1; sim tinyL.out`

*Answer (mark one):* compiler: ☐ or interpreter: ☒

input language: `tinyL`, output language: \_\_\_\_\_

3. `gcc Compiler.c`

*Answer (mark one):* compiler: ☒ or interpreter: ☐

input language: `C`, output language: `ilab machine code (executable)`

4. `constfolding < tinyL.out`

*Answer (mark one):* compiler: ☒ or interpreter: ☐

input language: `ILOC RISC machine code`, output language: `ILOC RISC machine code`

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## Problem 5 – Syntax-Directed Translation (60 pts)

Assume the following partial expression grammar:

$$\begin{aligned} \langle \text{expr} \rangle &::= + \langle \text{expr} \rangle \langle \text{expr} \rangle \mid \\ &\quad * \langle \text{expr} \rangle \langle \text{expr} \rangle \mid \\ &\quad \langle \text{const} \rangle \\ \langle \text{const} \rangle &::= 1 \mid 2 \end{aligned}$$

instr. format	description	semantics
<b>memory instructions</b>		
loadI $\langle \text{const} \rangle \Rightarrow_L r_x$	load constant value $\langle \text{const} \rangle$ into register $r_x$	$r_x \leftarrow \langle \text{const} \rangle$
<b>arithmetic instructions</b>		
add $r_x, r_y \Rightarrow_L r_z$	add contents of registers $r_x$ and $r_y$ , and store result into register $r_z$	$r_z \leftarrow r_x + r_y$
mult $r_x, r_y \Rightarrow_L r_z$	multiply contents of registers $r_x$ and $r_y$ , and store result into register $r_z$	$r_z \leftarrow r_x * r_y$

Here is a recursive descent parser that implements a compiler for the above grammar. Here is the important part of the code:

```
int expr() {
    int reg, left_reg, right_reg;
    switch (token) {
        case '+': next_token();
                    left_reg = expr(); right_reg = expr(); reg = next_register();
                    CodeGen(ADD, left_reg, right_reg, reg);
                    return reg;
        case '*': next_token();
                    left_reg = expr(); right_reg = expr(); reg = next_register();
                    CodeGen(MULT, left_reg, right_reg, reg);
                    return reg;
        case '1':
        case '2': return const();
    }
}

int const() {
    int reg;
    switch (token) {
        case '1': next_token(); reg = next_register();
                    CodeGen(LOADI, 1, reg);
                    return reg;
        case '2': next_token(); reg = next_register();
                    CodeGen(LOADI, 2, reg);
                    return reg;
    }
}
```

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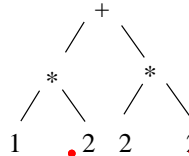
Make the following assumptions:

- The value of variable “token” has been initialized correctly.
- The function `CodeGen` is the one from our first project.
- **The first call to function `next_register()` the shown parser returns integer value “1”.** In other words, the first register that the generated code will be using is register  $r_1$ .
- Your parser “starts” by calling function `expr()` on the entire input.

1. Show the code that the recursive descent parser generates for input

<code>+ * 1 2 * 2 2</code>
----------------------------

will produce:



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loadI 1  $\Rightarrow$  r1 <https://powcoder.com>

loadI 2  $\Rightarrow$  r2

mult r1, r2  $\Rightarrow$  r3

loadI 2  $\Rightarrow$  r4

loadI 2  $\Rightarrow$  r5

mult r4, r5  $\Rightarrow$  r6

add r3, r6  $\Rightarrow$  r7



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2. Change the basic recursive-descent parser to implement an interpreter for our example language. You may insert pseudo code in the spaces marked by \_\_\_\_\_. No error handling is necessary.

```
int  expr() {

    int a, b;
    switch (token) {
    case '+': next_token();
              a = expr(); b = expr();
              return (a + b);

    case '*': next_token();
              a = expr(); b = expr();
              return (a * b);

    case '1':
    case '2': return const();
    }

}

int  const() {
    switch (token) {
    case '1': next_token();
              return 1 ;

    case '2': next_token();
              return 2 ;
    }

}
```

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