Midterm Exam CS 314, Fall '17 October 27 SAMPLE SOLUTION

DO NOT OPEN THE EXAM UNTIL YOU ARE TOLD TO DO SO

Name:	
Rutgers ID number:	
Assignment Project Exam Help Section:	
https://powcoder.com	

WRITE YOUR NAME ON EACH PAGE IN THE UPPER Add WIGHT CORPUS WCODE!

Instructions

We have tried to provide enough information to allow you to answer each of the questions. If you need additional information, make a *reasonable* assumption, write down the assumption with your answer, and answer the question. There are 5 problems, and the exam has 8 pages. Make sure that you have all pages. The exam is worth 250 points. You have 80 minutes to answer the questions. Good luck!

This table is for grading purposes only

1	/ 80
2	/ 60
3	/ 30
4	/ 20
5	/ 60
total	/ 250

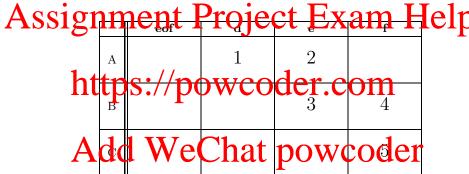
Problem 1 – Regular Expressions, DFA and Context Free Grammars (80 pts)

The context-free grammar ${\sf G}$ is specified in Backus-Naur-Form as follows, with ${\sf A}$ as the start symbol:

1. Give a rightmost derivation (\Rightarrow_R) for the string d e f given the grammar above. (15 pts)

$$A \Rightarrow_R dA \Rightarrow_R deB \Rightarrow_R deC \Rightarrow_R def$$

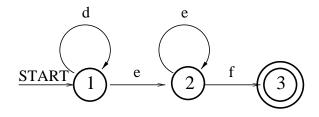
2. Give the LL(1) parse table for the grammar G. Insert the rule number or leave an entry empty. (35 pts)



3. Give a regular expression for the language generated by the grammar ${\sf G.}$ (15 pts)

$$d^*e^+f$$

4. Specify a DFA by extending the state transition diagram below. The start state is **state 1**, and the final (accepting) state is **state 3**. You are only allowed to add edges with their appropriate labels, i.e., valid labels are a, b, and c. Note that an edge may have more than one label. You **must not** add any states. (15 pts)



NAME:	

Problem 2 – Context Free Grammars (60 pts)

A context-free language is a language that can be specified using a context-free grammar. A regular language is a language that can be specified using a regular expression.

For the three languages given below, if the language is context-free, give a compact context-free grammar in Backus-Naur-Form (BNF). If the language is regular, give a compact regular expression using the regular expression syntax introduced in class. If a language is context-free and regular, give both specifications, a BNF and a regular expression. You do not have to justify why you believe a language is not context-free or not regular.

1. $\{ w \mid w \text{ has at least 3 symbols, but no more than 6} \}$, with alphabet $\Sigma = \{ a, b \}$

$$S ::= A A A B B B$$

$$A ::= a \mid b$$

$$\mathbf{B} ::= \mathbf{a} \mid \mathbf{b} \mid \ \epsilon$$

 $^{\text{S} ::= aaSb \mid \epsilon}$ https://powcoder.com

3. { $a^nb^mc^n \mid {\bf n} \geq 0, \, {\bf m} > 0$ }, with alphabet $\Sigma = \{{\bf a}, \, {\bf b}, \, {\bf c}\}$

 $s := aSc \mid B$ Add WeChat powcoder

Problem 3 – Pointers and Memory Allocation in C (30 pts)

```
int main() {
  int     x;
  int *y;
  int **z;

z = (int **) malloc (sizeof(int *));
  y = (int *) malloc (sizeof(int));
  x = 0;
  *z = &x;
  *y = x;
  x = x + 4;
  **z = *y + 5;
  printf("x=%d, *y=%d, **z=%d\n", x, *y, **z);
  return 0;
}
```

Assignment Project Exam Help heap https://powcoder.com

1. What is the output of the above program? (3 pts each)

x= 5, *y= 0, Azdd WeChat powcoder

- 2. Specify, whether the following program objects are allocated on the **stack** (includes global variables), on the **heap**, or **not defined** (2 pts each).
 - x is allocated on the stack
- **y is allocated on the <u>not defined</u>
- z is allocated on the stack
- *x is allocated on the <u>not defined</u>
- y is allocated on the stack
- **z is allocated on the stack
- *y is allocated on the heap
- *z is allocated on the heap
- 3. Assume the following code segment:

```
int *x;
*x = 5;
printf("%d\n", *x);
```

Is there a problem with this code? Assume that when you ran the code a couple of times, it printed "5". If you believe there is a problem, give a possible "fix" for the problem? (5 pts)

The content of variable x is not initialized. However, its content is used as an address of a memory location, and that memory location is assigned the value 5.

To fix the problem, the pointer x should be initialized to NULL in its declaration, and x must point to an object on the heap before it is dereferenced. This object is allocated as follows:

```
x = (int *) malloc(sizeof(int *))
```

This statement should be placed before statement *x = 5, i.e., before the dereference operation on x.

Assignment Project Exam Help https://powcoder.com Add WeChat powcoder

NAME:	

Problem 4 – Compilers vs. Interpreters (20 pts)

To answer this question, please use the following definitions.

Definition A <u>compiler</u> maps an input program to a semantically equivalent output program. Note that the input and output language may be the same. \Box

Definition An <u>interpreter</u> maps an input program to the answers it computes; In other words, it executes the program. \Box

As part of the C project, we used and/or wrote the following programs/commands:

- gcc usage: gcc program>
- compile usage: compile cprogram>
- ullet constfolding usage: constfolding < <pre>cprogram>
- \bullet sim -usage: sim <program>

Under Unix/Linux, commands can be entered on a single command line if they are separated by a semicolon. For instance, saying cd foo; 1s will change the current directory to subdirectory foo, and will list its files.

In the project System of the plantages, man(h) Lew Rec Madifecode and plantages object code (executables). Classify the following commands or the entire sequence of commands as either compiler or interpreter. Note that you should treat a sequence of commands as a single unit, i.e., as a single composed command. If the single or composed command is a compiler, give its input and output language (e.g.: input language: Cautius against 1971 Wor an interpreter us the its input language.

1.	Answer (mark one): Answer (mark one): Workstein powcoder
	Answer (mark one): compiler: X for interplated POWCOUCI
	input language: $\underline{\mathbf{C}}$, output language: $\underline{\mathbf{ilab}}$ machine code (executable)
2.	compile test1
	Answer (mark one): compiler: X or interpreter
	input language: $\underline{\operatorname{tinyL}}$, output language: $\underline{\operatorname{ILOC}}$ RISC machine code
3.	<pre>constfolding < tinyL.out</pre>
	Answer (mark one): compiler: X or interpreter
	input language: $\underline{\text{ILOC RISC machine code}}, \text{output language: } \underline{\text{ILOC RISC machine code}}$
4.	compile test1; sim tinyL.out
	Answer (mark one): compiler: or interpreter X
	input language: tinyL, output language:

NAME:	

Problem 5 – Syntax-Directed Translation (60 pts)

Assume the following partial expression grammar:

instr. format	$\operatorname{description}$	semantics
memory instructions		
$\texttt{loadI} \; <\! \texttt{const}\! > \Rightarrow_R \; r_x$	load constant value $<$ const $>$ into register r_x	$r_x \leftarrow < \text{const} >$
arithmetic instructions		
add r_x , r_y \Rightarrow_R r_z	add contents of registers r_x and r_y , and	$r_z \leftarrow r_x + r_y$
	store result into register r_z	
mult r_x , r_y $\Rightarrow_R r_z$	multiply contents of registers r_x and r_y , and	$r_z \leftarrow r_x * r_y$
	store result into register r_z	

Here is a recurs vs. is in the input of the code:

Here is a recurs vs. is in the input of the code:

```
int expr() {
       int reg, left_left[hts:eg/powcoder.com] switch (token) {
       case '+': next_token();
                   left_reg = expr(); right_reg = expr(); reg = next_register();
                   code len appl laft leg cight regt reg Dowcoder
       case '*': next_token();
                   left_reg = expr(); right_reg = expr(); reg = next_register();
                   CodeGen(MULT, left_reg, right_reg, reg);
                   return reg;
       case '1':
       case '2': return const();
}
int const() {
       int reg;
       switch (token) {
       case '1': next_token(); reg = next_register();
                     CodeGen(LOADI, 1, reg);
                     return reg;
       case '2': next_token(); reg = next_register();
                     CodeGen(LOADI, 2, reg);
                     return reg;
       }
}
```

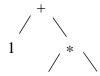
NAME:	
-------	--

Make the following assumptions:

- The value of variable "token" has been initialized correctly.
- The function CodeGen is the one from our first project.
- The first call to function next_register() the shown parser returns integer value "1". In other words, the first register that the generated code will be using is register r_1 .
- Your parser "starts" by calling function expr() on the entire input.
- 1. Show the code that the recursive descent parser generates for input

+ 1 * 2 1

will produce:



Assignment Project Exam Help

 $loadI 1 \Rightarrow r1$ https://powcoder.com

 $loadI 2 \Rightarrow r2$

 $loadI \stackrel{?}{1} \Rightarrow r^3Add WeChat powcoder$

mult r2, r3 \Rightarrow r4

add r1, r4 \Rightarrow r5

NAME:	

2. Change the basic recursive-descent parser to implement an interpreter for our example language. You may insert pseudo code in the spaces marked by ______. No error handling is necessary.

```
int expr() {
      int a, b;
      switch (token) {
      case '+': next_token();
                 a = expr(); b = expr();
             return (a + b);
      case '*': next_token();
                 a = expr(); b = expr();
             return (a * b);
      case '1':
      case '2': return const();
}
 int const() {
               nment Project Exam Help
              next_token();
      case '1':
              return 1;
             https://powcoder.com
      }
}
            Add WeChat powcoder
```