

143 In natural variables s and n prove

$P \Leftarrow \text{if } n=0 \text{ then } ok \text{ else } n:=n-1. s:=s+2^n-n. t:=t+1. P \text{ fi}$
where $P = s' = s + 2^n - n \times (n-1)/2 - 1 \wedge n'=0 \wedge t' = t+n$.

§ Proof by parts (3 of them) and by cases (2 of them), so 6 things to prove. First part, first case, starting with the right side:

$n=0 \wedge ok$ expand ok
 $= n=0 \wedge n'=n \wedge s'=s \wedge t'=t$ arithmetic and specialization
 $\Rightarrow s' = s + 2^n - n \times (n-1)/2 - 1$

First part, last case:

$n>0 \wedge (n:=n-1. s:=s+2^n-n. t:=t+1. s' = s + 2^n - n \times (n-1)/2 - 1)$ Substitution Law 3 times
 $= n>0 \wedge s' = s + 2^{n-1} - (n-1) + 2^{n-1} - (n-1) \times (n-1-1)/2 - 1$ arithmetic
 $= n>0 \wedge s' = s + 2^n - n \times (n-1)/2 - 1$ specialization
 $\Rightarrow s' = s + 2^n - n \times (n-1)/2 - 1$

Middle part, first case:

$n=0 \wedge ok$ expand ok
 $= n=0 \wedge n'=n \wedge s'=s \wedge t'=t$ transitivity and specialization
 $\Rightarrow n'=0$

Middle part, last case:

$n>0 \wedge (n:=n-1. s:=s+2^n-n. t:=t+1. n'=0)$ Substitution Law 3 times
 $= n>0 \wedge n'=0$ specialization
 $\Rightarrow n'=0$

Last part, first case:

$n=0 \wedge ok$ expand ok
 $= n=0 \wedge n'=n \wedge s'=s \wedge t'=t$ arithmetic and specialization
 $\Rightarrow t' = t+n$

Last part, last case:

$n>0 \wedge (n:=n-1. s:=s+2^n-n. t:=t+1. t' = t+n)$ Substitution Law 3 times
 $= n>0 \wedge t' = t+n-1$ arithmetic and specialization
 $\Rightarrow t' = t+n$

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