

451 Let a and b be binary interactive variables. Define

$loop = \text{if } b \text{ then } loop \text{ else } ok \text{ fi}$

Add a time variable according to any reasonable measure, and then express

$b := \perp \parallel loop$

as an equivalent program but without using \parallel .

§ The left process owns b . Variable a could belong to either process; let's give it to the right process. Let assignment take time 1. Then the left process is

$\neg b(t+1) \wedge t' = t+1$

Add recursive time to $loop$, and the right process is

$loop$

$= \text{if } b \text{ then } t := t+1. loop \text{ else } ok \text{ fi}$ unroll

$= \text{if } b \text{ then } t := t+1. \text{if } b \text{ then } t := t+1. loop \text{ else } ok \text{ fi else } ok \text{ fi}$ Substitution Law

$= \text{if } b \text{ then if } b(t+1) \text{ then } t := t+2. loop \text{ else } t := t+1 \text{ fi else } ok \text{ fi}$

The left process gives us $\neg b(t+1)$

$= \text{if } b \text{ then if } \perp \text{ then } t := t+2. loop \text{ else } t := t+1 \text{ fi else } ok \text{ fi}$

$= \text{if } b \text{ then } t := t+1 \text{ else } ok \text{ fi}$

$= \text{if } b \text{ then } a t' = a t \wedge t' = t+1 \text{ else } t' = t \text{ fi}$

$t' = t$ implies $a t' = a t$

$= a t' = a t \wedge \text{if } b \text{ then } t' = t+1 \text{ else } t' = t \text{ fi}$

The independent composition is

$\exists tL, tR. \neg b(t+1) \wedge tL = t+1 \wedge a tR = a t \wedge \text{if } b \text{ then } tR = t+1 \text{ else } tR = t \text{ fi}$

$\wedge t' = \max tL, tR$ The left process takes time 1 and the right process takes time 0 or 1, so the maximum is 1

$= \neg b(t+1) \wedge a(t+1) = a t \wedge t' = t+1$

We no longer have an independent

composition, so a and b are both variables

$= b := \perp$ <https://powcoder.com>

Add WeChat powcoder