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Here is an informal explanation of the one-point laws. Let's start with
        \exists v: D \cdot v = x \wedge b
and we are given x: D. Let's suppose D is nat. The existential quantification is an infinite
disjunction.
        (0=x \land b) \lor (1=x \land b) \lor (2=x \land b) \lor ...
We can use context in each of these disjuncts.
          (0=x \land (b \text{ but replace } x \text{ with } 0))
        \vee (1=x \land (b \text{ but replace } x \text{ with } 1))
        \vee (2=x \wedge (b \text{ but replace } x \text{ with } 2))
Since x is in nat, exactly one of 0=x, 1=x, 2=x, ... is \top and the others are \bot. That's why it's
called "one-point". Let's suppose x is 1.
          (\perp \land (b \text{ but replace } x \text{ with } 0))
        \vee (\top \land (b \text{ but replace } x \text{ with } 0))
        \vee (\perp \land (b \text{ but replace } x \text{ with } 0))
        \perp v ( b but replace x with 1 ) v \perp v ...
        (b but replace x with 1)
Now the other signment Project Exam Help
and we are given x: D. Let's suppose D is nat. The universal quantification is an infinite
conjunction.
        (0=x ⇒ b) ∧ (https:///xpowcoder.com
We can use context in each of these conjuncts.
          (0=x \Rightarrow (b \text{ but replace } x \text{ with } 0))
        Since x is in nat, exactly one of 0=x, 1=x, 2=x, ... is \top and the others are \bot. That's why it's
called "one-point". Let's suppose x is 1.
          (\bot \Rightarrow (b \text{ but replace } x \text{ with } 0))
        \land (\top \Rightarrow (b \text{ but replace } x \text{ with } 0))
        \land (\bot \Rightarrow (b \text{ but replace } x \text{ with } 0))
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 $\top \land (b \text{ but replace } x \text{ with } 1) \land \top \land ...$

(b but replace x with 1)