OCaml Datatypes

CSI 3120

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OCaml So Far

- We have seen a number of basic types:
  - int
  - float
  - char
  - string
  - bool
- We have seen a few structured types:
  - pairs
  - tuples
  - options
  - lists
- In this lecture, we will see some more general ways to define our own new types and data structures

Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```

Type Abbreviations

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```
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- These abbreviations can be helpful documentation:


```
let distance (p1:point) (p2:point) : float =
  let square x = x *. x in
  let (x1,y1) = p1 in
  let (x2,y2) = p2 in
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```


- But they add nothing of *substance* to the language
  - they are equal in every way to an existing type

Type Abbreviations

- We have already seen some type abbreviations:

```
type point = float * float
```


- As far as OCaml is concerned, you could have written:


```
let distance (p1:float*float)
             (p2:float*float) : float =
  let square x = x *. x in
  let (x1,y1) = p1 in
  let (x2,y2) = p2 in
  sqrt (square (x2 -. x1) +. square (y2 -. y1))
```

- Since the types are equal, you can *substitute* the definition for the name wherever you want
  - we have not added any new data structures

DATA TYPES

Data types

- OCaml provides a general mechanism called a data type for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```


a value with type my_bool is one of two things:

- Tru, or
- Fal

read the "|" as "or"

Data types

- OCaml provides a general mechanism called a data type for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```


Tru and Fal are called "constructors"

a value with type my_bool is one of two things:

- Tru, or
- Fal

read the "|" as "or"

Data types

- OCaml provides a general mechanism called a data type for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```


there's no need to stop at 2 cases; define as many alternatives as you want

Data types

- OCaml provides a general mechanism called a data type for defining new data structures that consist of many alternatives

```
type my_bool = Tru | Fal
```

```
type color = Blue | Yellow | Green | Red
```


- Creating values:

```
let b1 : my_bool = Tru
let b2 : my_bool = Fal
let c1 : color = Yellow
let c2 : color = Red
```

use constructors to create values

Data types

```
type color = Blue | Yellow | Green | Red

let c1 : color = Yellow
let c2 : color = Red
```


- Using data type values:


```
let print_color (c:color) : unit =
  match c with
   | Blue ->
   | Yellow ->
   | Green ->
   | Red ->
```


use pattern matching to determine which color you have; act accordingly

Data types

```
type color = Blue | Yellow | Green | Red

let c1 : color = Yellow
let c2 : color = Red
```


- Using data type values:


```
let print_color (c:color) : unit =
  match c with
  | Blue -> print_string "blue"
  | Yellow -> print_string "yellow"
  | Green -> print_string "green"
  | Red -> print_string "red"
```

Data types

```
type color = Blue | Yellow | Green | Red

let c1 : color = Yellow
let c2 : color = Red
```


- Using data type values:


```
let print_color (c:color) : unit =
  match c with
   | Blue -> print_string "blue"
   | Yellow -> print_string "yellow"
   | Green -> print_string "green"
   | Red -> print_string "red"
```


Why not just use strings to represent colors instead of defining a new type?

Data types

```
type color = Blue | Yellow | Green | Red
```

oops!:

```
let print_color (c:color) : unit =
  match c with
   | Blue -> print_string "blue"
   | Yellow -> print_string "yellow"
   | Red -> print_string "red"
```


Warning 8: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
Green

Data types

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type color = Blue | Yellow | Green | Red
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oops!:

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let print_color (c:color) : unit =
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```


Warning 8: this pattern-matching is not exhaustive.
Here is an example of a value that is not matched:
Green

OCaml's datatype mechanism allow you to create types that contain *precisely* the values you want!

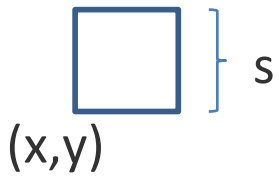
Data Types Can Carry Additional Values

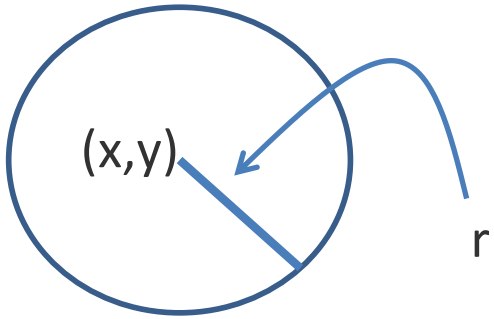
- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float
```


- Read as: a simple_shape is either:
  - a Circle, which contains a pair of a point and float, or
  - a Square, which contains a pair of a point and float





Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float

let origin : point = (0.0, 0.0)

let circ1  : simple_shape = Circle (origin, 1.0)
let circ2  : simple_shape = Circle ((1.0, 1.0), 5.0)
let square : simple_shape = Square (origin, 2.3)
```

Data Types Can Carry Additional Values

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```

Compare

- Data types are more than just enumerations of constants:

```
type point = float * float

type simple_shape =
  Circle of point * float
| Square of point * float

let simple_area (s:simple_shape) : float =
  match s with
  | Circle (_, radius) -> 3.14 *. radius *. radius
  | Square (_, side) -> side *. side
```



```
type my_shape = point * float

let simple_area (s:my_shape) : float =
  (3.14 *. radius *. radius)  ?? or ??  (side *. side)
```

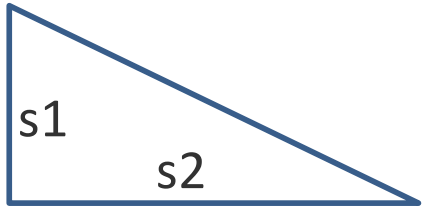
More General Shapes

```
type point = float * float

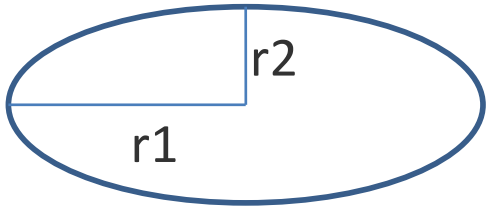
type shape =
    Square of float
  | Ellipse of float * float
  | RtTriangle of float * float
  | Polygon of point list
```


Square s =

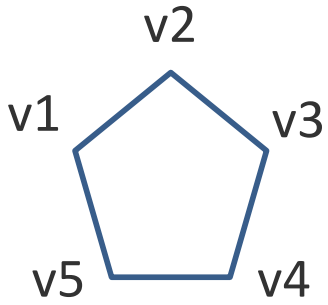
RtTriangle (s1, s2) =



Ellipse (r1, r2) =



Polygon [v1; ...;v5] =



More General Shapes

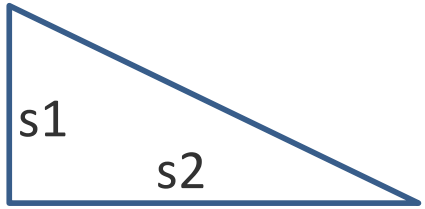
```
type point = float * float
type radius = float
type side = float

type shape =
    Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

Type abbreviations can aid readability

Square s =

RtTriangle (s1, s2) =



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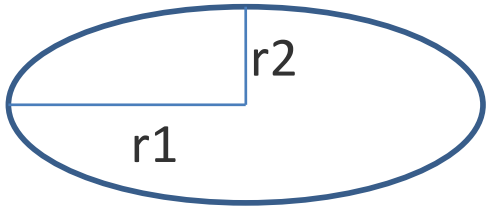
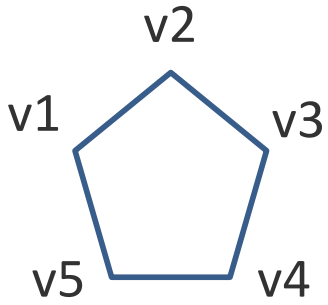


RtTriangle [v1; ...;v5] =



More General Shapes

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type radius = float
type side = float

type shape =
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  | Polygon of point list
```

Square builds a shape from a single side

Ellipse builds a shape from a pair of radiuses

RtTriangle builds a shape from a pair of sides


```
let sq   : shape = Square 1.0
let ell  : shape = Ellipse (1.0, 2.0)
let rt   : shape = RtTriangle (1.0, 1.0)
let poly : shape = Polygon [(0., 0.); (1., 0.); (0.; 1.)]
```

they are all shapes;
they are constructed in different ways

Polygon builds a shape from a list of points (where each point is itself a pair)

More General Shapes

```
type point = float * float
type radius = float
type side = float

type shape =
    Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

a data type also defines
a pattern for matching


```
let area (s : shape) : float =
  match s with
  | Square s ->
  | Ellipse (r1, r2)->
  | RtTriangle (s1, s2) ->
  | Polygon ps ->
```

More General Shapes

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type point = float * float
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let area (s : shape) : float =
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  | Polygon ps ->
```

Square carries a value with type float so s is a pattern for float values

RtTriangle carries a value with type float * float so (s1, s2) is a pattern for that type

More General Shapes

```
type point = float * float
type radius = float
type side = float

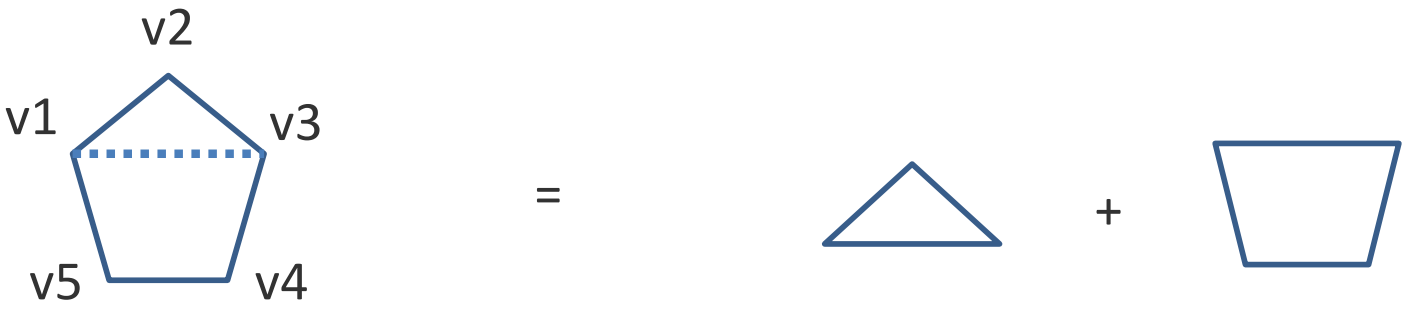
type shape =
    Square of side
  | Ellipse of radius * radius
  | RtTriangle of side * side
  | Polygon of point list
```

a data type also defines
a pattern for matching


```
let area (s : shape) : float =
  match s with
  | Square s -> s *. s
  | Ellipse (r1, r2) -> pi *. r1 *. r2
  | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
  | Polygon ps -> ???
```

Computing Area

- How do we compute polygon area?
- For convex polygons:
  - Case: the polygon has fewer than 3 points:
    - it has 0 area! (it is a line or a point or nothing at all)
  - Case: the polygon has 3 or more points:
    - Compute the area of the triangle formed by the first 3 vertices
    - Delete the second vertex to form a new polygon
    - Sum the area of the triangle and the new polygon



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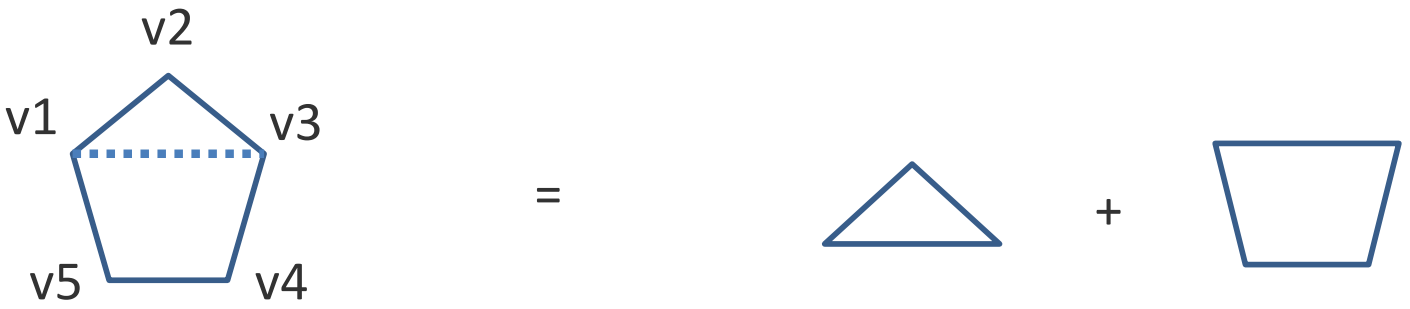
    - Compute the area of the triangle formed by the first 3 vertices
    - Delete the second vertex to form a new polygon
    - Sum the area of the triangle and the new polygon
- Note: This is a beautiful inductive algorithm:
  - the area of a polygon with n points is computed in terms of a smaller polygon with only n-1 points!

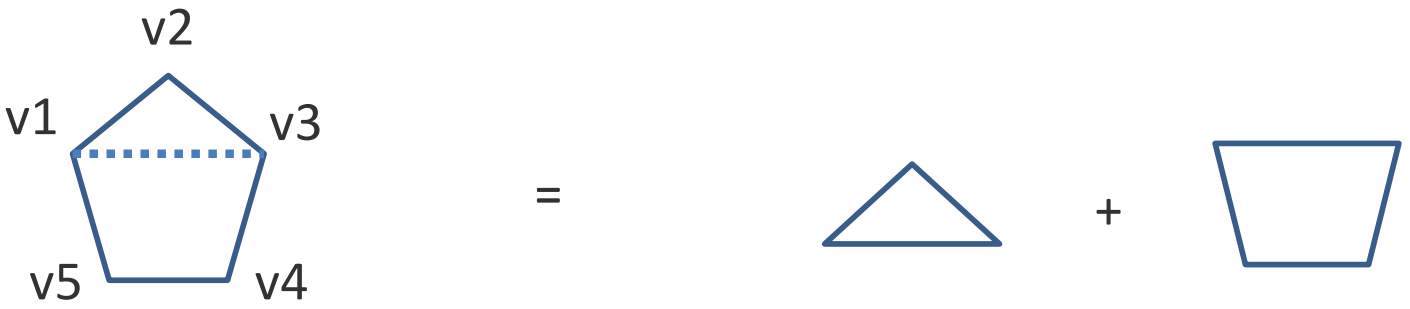


Computing Area

```
let area (s : shape) : float =
  match s with
   | Square s -> s *. s
   | Ellipse (r1, r2) -> r1 *. r2
   | RtTriangle (s1, s2) -> s1 *. s2 /. 2.
   | Polygon ps -> poly_area ps
```

This pattern says the list has at least 3 items


```
let rec poly_area (ps : point list) : float =
  match ps with
   | p1 :: p2 :: p3 :: tail ->
      tri_area p1 p2 p3 +. poly_area (p1::p3::tail)
   | _ -> 0.
```



Computing Area

```
let tri_area (p1:point) (p2:point) (p3:point) : float =
  let a = distance p1 p2 in
  let b = distance p2 p3 in
  let c = distance p3 p1 in
  let s = 0.5 *. (a +. b +. c) in
  sqrt (s *. (s -. a) *. (s -. b) *. (s -. c))
```



```
let rec poly_area (ps:point list) : float =
  match ps with
  | p1 :: p2 :: p3 :: tail ->
      tri_area p1 p2 p3 +. poly_area (p1::p3::tail)
  | _ -> 0.
```



```
let area (s : shape) : float =
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  | Polygon ps -> poly_area ps
```

INDUCTIVE DATA TYPES

Inductive data types

- We can use data types to define inductive data
- A binary tree is:
  - a Leaf containing no data
  - a Node containing a key, a value, a left subtree and a right subtree

Inductive data types

- We can use data types to define inductive data
- A binary tree is:
  - a Leaf containing no data
  - a Node containing a key, a value, a left subtree and a right subtree


```
type key = int
type value = string
```



```
type tree =
  Leaf
| Node of key * value * tree * tree
```

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```



```
let rec insert (t:tree) (k:key) (v:value) : tree =
```

Inductive data types

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type key = int
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```



```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf ->
  | Node (k', v', left, right) ->
```

Again, the type definition specifies the cases you must consider

Inductive data types

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type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```



```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
```

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
| Node of key * value * tree * tree
```



```
let rec insert (t:tree) (k:key) (v:value) : tree =
  match t with
  | Leaf -> Node (k, v, Leaf, Leaf)
  | Node (k', v', left, right) ->
      if k < k' then
        Node (k', v', insert left k v, right)
      else if k > k' then
        Node (k', v', left, insert right k v)
      else
        Node (k, v, left, right)
```

Inductive data types

```
type key = int
type value = string

type tree =
  Leaf
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```
let rec insert (t:tree) (k:key) (v:value) : tree =
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      else
        Node (k, v, left, right)
```

Inductive data types: Another Example

- We can use the type "int" to represent natural numbers
  - but that is kind of broken: it also contains negative numbers
  - we have to use a dynamic test and a default value:

```
let double (n : int) : int =
  if n < 0 then
    0
  else
    double_nat n
```

Inductive data types: Another Example

- We can use the type "int" to represent natural numbers
  - but that is kind of broken: it also contains negative numbers
  - we have to use a dynamic test and a default value:
  - or raise an exception

```
let double (n : int) : int =
  if n < 0 then
    raise (Failure "negative input!")
  else
    double_nat n
```


  - it would be nice if there was a way to define the natural numbers exactly, and use OCaml's type system to guarantee no client ever attempts to double a negative number

Inductive data types

- Recall, a natural number n is either:
  - zero, or
  - m + 1
- We use a data type to represent this definition exactly:

Inductive data types

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type nat = Zero | Succ of nat
```

Inductive data types

- Recall, a natural number n is either:
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type nat = Zero | Succ of nat
```



```
let rec nat_to_int (n : nat) : int =
 match n with
   Zero -> 0
 | Succ n -> 1 + nat_to_int n
```

Inductive data types

- Recall, a natural number n is either:
  - zero, or
  - m + 1
- We use a data type to represent this definition exactly:


```
type nat = Zero | Succ of nat

let rec nat_to_int (n : nat) : int =
 match n with
   Zero -> 0
 | Succ n -> 1 + nat_to_int n

let rec double_nat (n : nat) : nat =
  match n with
   | Zero -> Zero
   | Succ m -> Succ (Succ (double_nat m))
```

Summary

- OCaml data types: a powerful mechanism for defining complex data structures:
  - They are precise
    - contain exactly the elements you want, not more elements
  - They are general
    - recursive, non-recursive (mutually recursive and polymorphic)
  - The type checker helps you detect errors
    - missing cases in your functions

PARAMETERIZED TYPE DEFINITIONS

```
type ('key, 'val) tree =
    Leaf
  | Node of 'key * 'val * ('key, 'val) tree * ('key, 'val) tree


type 'a inttree = (int, 'a) tree


type istree = string inttree
```


Note: `istree` is an abbreviation for `(int,string) tree`

General form:

definition:

type 'x f = body

use:

arg f

A more conventional notation would have been (but is not ML):

definition:

type f x = body

use:

f arg

Take-home Message

- Think of parameterized types like functions:
  - a function that take a type as an argument
  - produces a type as a result

- Theoretical basis:
  - System F-omega
  - a typed lambda calculus with general type-level functions as well as value-level functions