# **Complex Dynamical Networks:**

## **Lecture 6a: Community Structures**

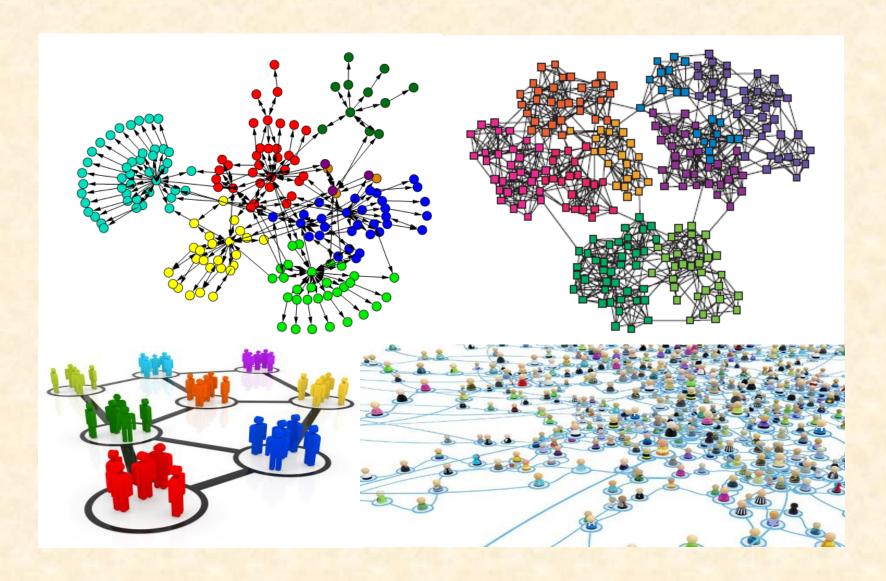
**EE 6605** 

Instructor: G Ron Chen

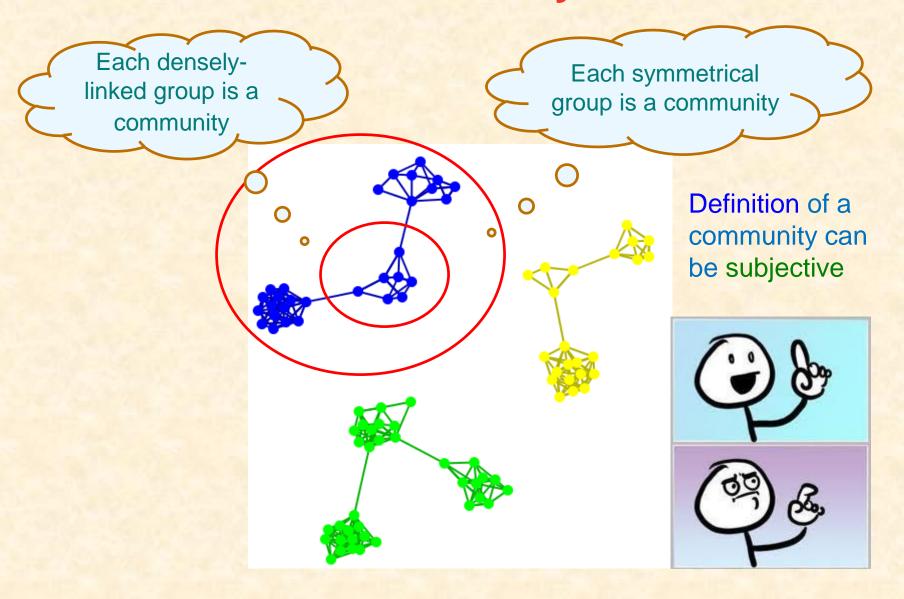


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## **Community Structure in Complex Networks**



## Community



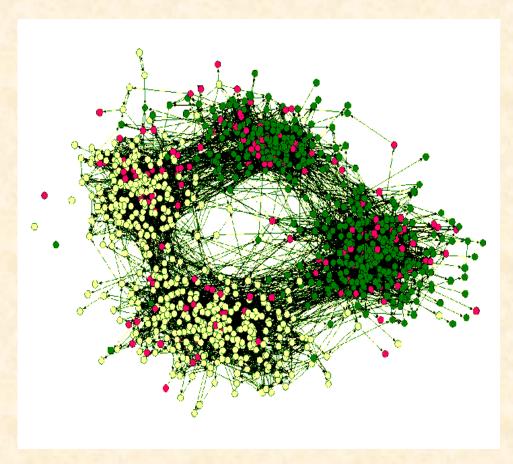
## **Community Detection**

- Common Definition: A community is a set of nodes among which the connections are relatively dense and strong, or interactions are relatively frequent
- A network has a community structure if the network can naturally be divided into clusters of nodes with dense internal connections and sparse external connections
- Community Detection (Clustering, Grouping)
   To find cohesive subgroups from a given graph
- Applications

Understanding interactions between people (or systems)
Visualizing and navigating on large-scale networks
Forming bases for other tasks such as data mining

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### **Example:** USA school integration and friendship segregation





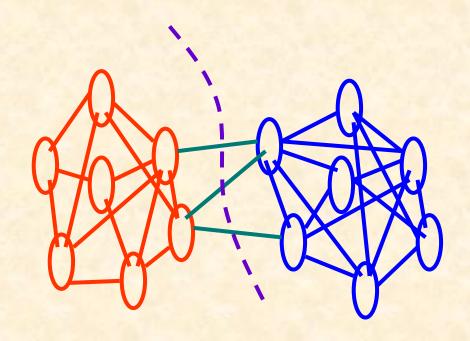
Race: left (white) to right (black)

Grade: up (junior) to down (senior)

J. Moody, Amer. J. of Sociology, 2001

## **Community Detection**

Many real networks have a natural community structure, and we want to discover this structure



For a large-scale network, how can we discover community structure in an automated way?

### Criteria: Which partition is better?

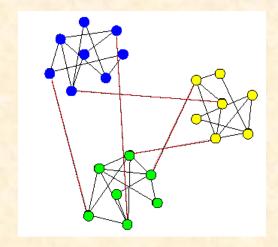
### Quality Measures:

How to evaluate an algorithm's performance while the community structure is unknown?

#### Benchmarks:

Which algorithm is the best to characterize a given network with a known community structure?

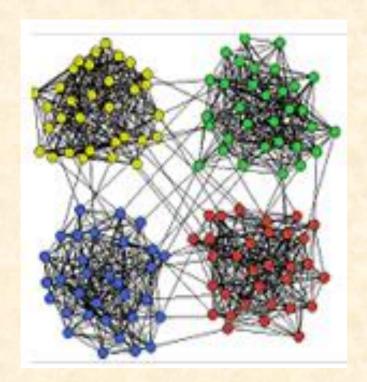




## **Benchmark Examples**

A real network: Zachary's Karate Network

An artificial network: Planted *L*-partition model



W. W. Zachary, J. Anthropological Research 33: 452-473, 1977

A. Condon and R. M. Karp, Random Structures & Algorithms 18(2): 116-140, 2001

## **Zachary's Karate Network**

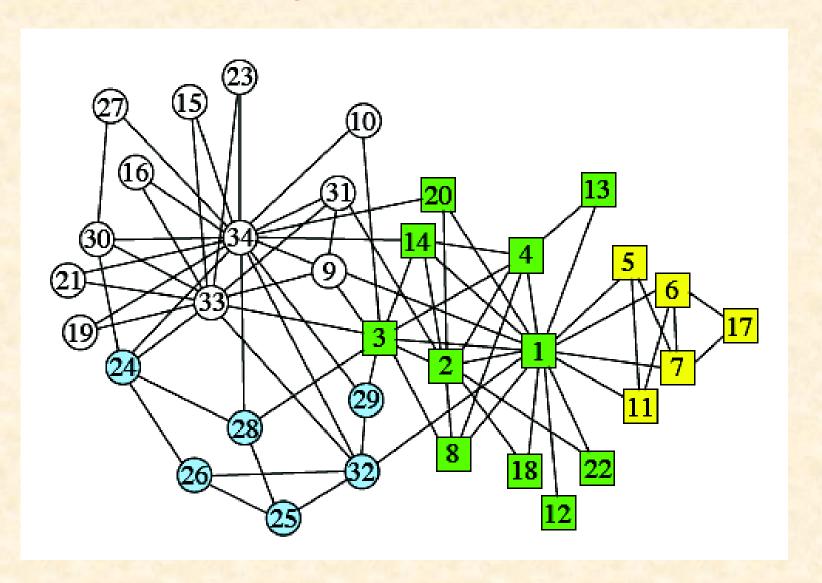






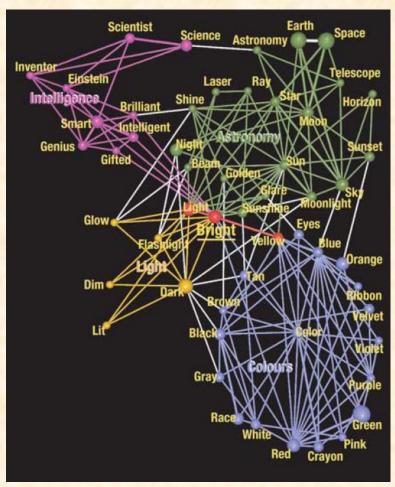
Story Background / Dataset

### **Zachary's Karate Network**



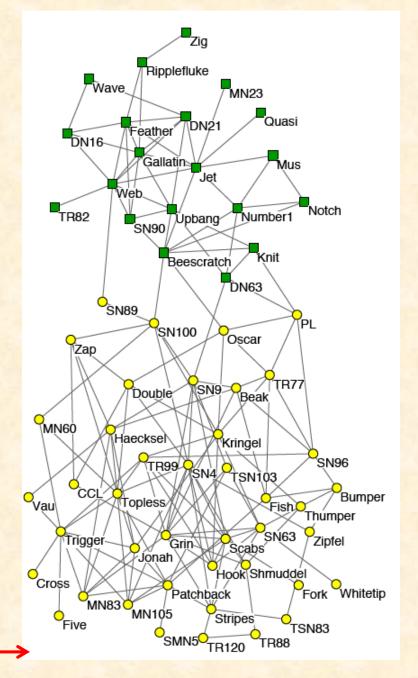
#### Other

### **Benchmark Examples**

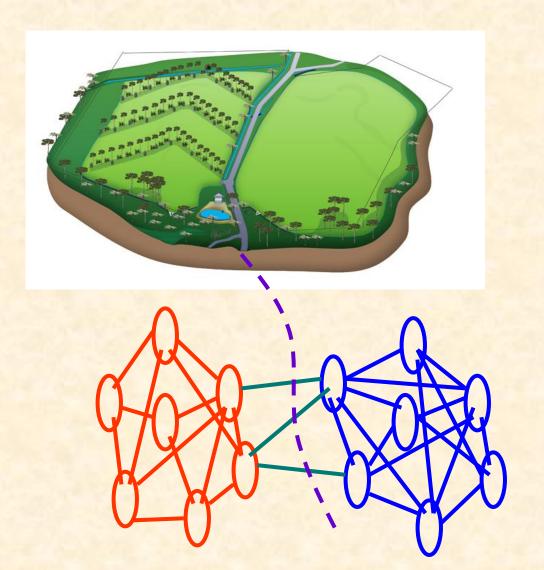


word network

Lusseau's network of bottlenose dolphins →



## **Example:** Planted 2-Partition Model



L=2

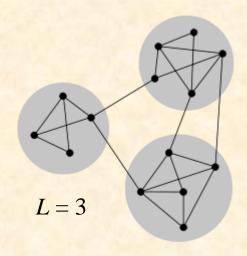
### Planted L-Partition Model

Partition the graph into L parts (for example, L = 3)

Each node has a probability (or, degree)  $P_{in}$  of being connected to nodes inside its group and a probability  $P_{out}$  of being connected to nodes outside its group

If  $p_{in} > p_{out}$  then the graph has a community structure

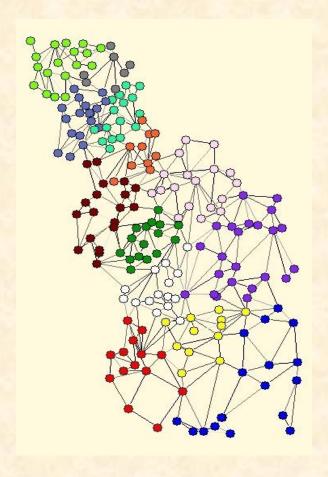
Otherwise, it is a homogeneous (e.g., random) graph



## **Methods and Algorithms**

- 1) Minimum-cut method
- 2) Hierarchical clustering
- 3) GN algorithm
- 4) Clique-based methods
- 5) Modularity maximization
- 6) Information-based algorithms

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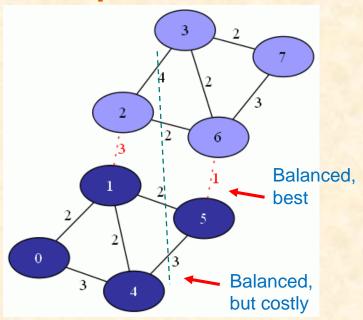
### 1) Minimum-Cut Method

One of the oldest algorithms for dividing networks into parts

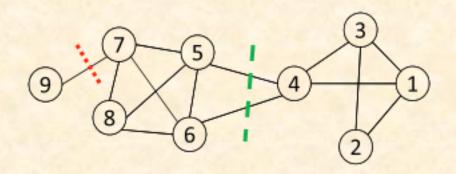
The network is divided into a pre-determined number of groups, in approximately the same size ("balanced"), chosen such that the number or cost (total weights) of edges between groups is minimized

### Karger's algorithm

### **Example:**



### **Ratio Cut and Normalized Cut**



- If cost (total weights) of edges between groups are not chosen appropriately, the corresponding minimum-cost cut may yield an unbalanced partition (e.g., with one set being a singleton)
- Other objective functions:

Ratio 
$$\operatorname{Cut}(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{\operatorname{cut}(C_i, \bar{C}_i)}{|C_i|},$$

Normalized 
$$\operatorname{Cut}(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{\operatorname{cut}(C_i, \bar{C}_i)}{\operatorname{vol}(C_i)}$$

 $C_i$ : community i

 $\overline{C}_i$ : complementary community

 $|C_i|$ : number of nodes in  $C_i$ 

 $vol(C_i)$ : sum of degrees in  $C_i$ 

 $cut(C_i, \overline{C_i}) =$ degrees of the cut

### **Example**

### For partition by red: $\pi_1$

Ratio 
$$Cut(\pi_1) = \frac{1}{2} \left( \frac{1}{1} + \frac{1}{8} \right) = 9/16 = 0.56$$

Normalized Cut(
$$\pi_1$$
) =  $\frac{1}{2} \left( \frac{1}{1} + \frac{1}{27} \right) = 14/27 = 0.52$ 

### For partition by green: $\pi_2$

Ratio 
$$Cut(\pi_2) = \frac{1}{2} \left( \frac{2}{4} + \frac{2}{5} \right) = 9/20 = 0.45 < Ratio  $Cut(\pi_1)$$$

Normalized 
$$\operatorname{Cut}(\pi_2) = \frac{1}{2} \left( \frac{2}{12} + \frac{2}{16} \right) = 7/48 = 0.15 < \operatorname{Normalized } \operatorname{Cut}(\pi_1)$$

### $\rightarrow$ We should cut $\pi_2$ as comparing to $\pi_1$

Ratio Cut
$$(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{cut(C_i, \bar{C}_i)}{|C_i|},$$

Normalized 
$$\operatorname{Cut}(\pi) = \frac{1}{k} \sum_{i=1}^{k} \frac{\operatorname{cut}(C_i, \bar{C}_i)}{\operatorname{vol}(C_i)}$$

 $|C_i|$ : number of nodes in  $C_i$   $vol(C_i)$ : sum of degrees in  $C_i$  $cut(C_i, \overline{C_i})$  = degrees of the cut

### 2) Node-Similarity-Based Clustering

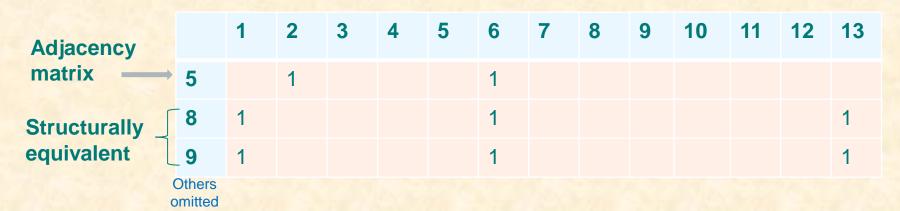
Put all nodes with the same (or close) similarity into the same community

- For large-scale networks:
  - Consider the connections as features
  - Use <u>Cosine similarity</u> or <u>Jaccard similarity</u> to compute node similarity
  - Apply the classical K-means Clustering Algorithm
- K-means Clustering Algorithm
  - Each cluster is associated with a centroid (center point)
  - Each node is assigned to the cluster with the closest centroid

#### **Algorithm 1** Basic K-means Algorithm.

- 1: Select K points as the initial centroids.
- 2: repeat
- 3: Form K clusters by assigning all points to the closest centroid.
- 4: Recompute the centroid of each cluster.
- 5: until The centroids don't change

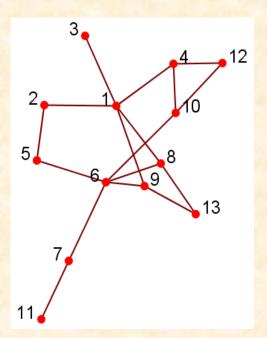
### **Node Similarity**



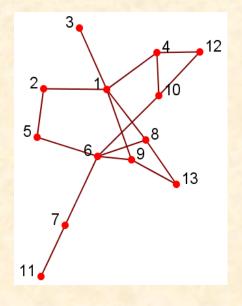
Cosine Similarity:  $similarity = cos(\theta) = \frac{A \cdot B}{\|A\| \|B\|}$ .

$$sim(5,8) = \frac{0 \times 1 + 1 \times 0 + 1 \times 1 + 0 \times 1}{\sqrt{2} \times \sqrt{3}} = \frac{1}{\sqrt{6}}$$

Jaccard Similarity:  $J(A,B) = \frac{|A \cap B|}{|A \cup B|}$ .  $J(5,8) = \frac{|\{6\}|}{|\{1,2,6,13\}|} = \frac{1}{4}$ 



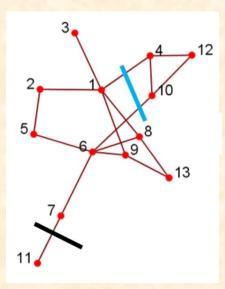
## **Node-Similarity-Based Clustering**

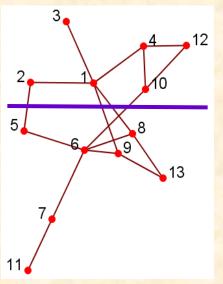


Cosine, Jaccard

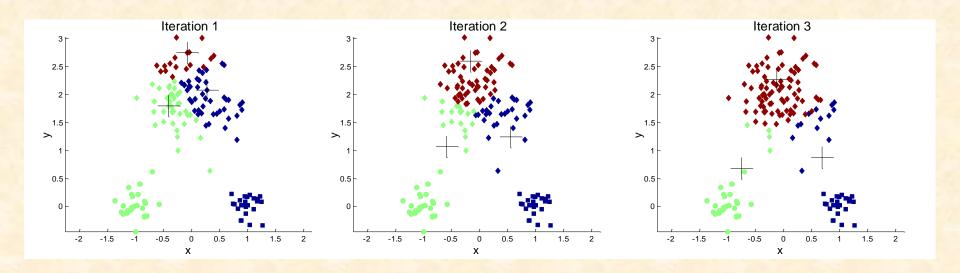


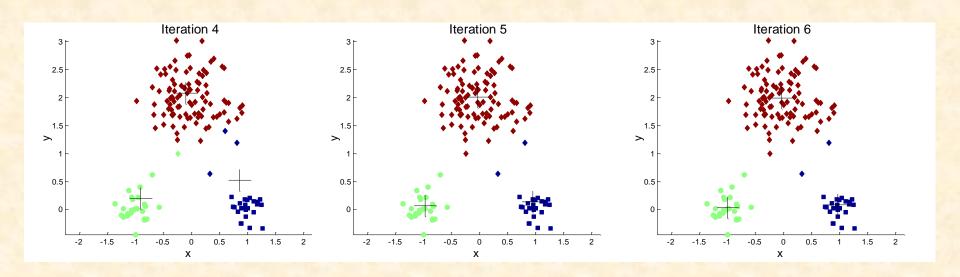
K-means





## **Illustration of K-means Clustering**





### 2) Hierarchical Clustering Method

A similarity measure is used to quantify node pairs

Commonly used measures include: cosine similarity, Jaccard similarity, and <u>Hamming distance</u> between rows of the adjacency matrix, etc.

Then, according to any of such measures, similar nodes are grouped into communities

**Example:** Hamming distance

between: toned and roses

ls: 3

**Example:** Hamming distance

between: 100100 and 011011

ls: 6

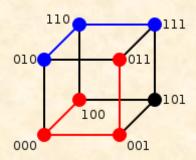
### 2) Hierarchical Clustering Method

In information theory, the **Hamming distance** between two strings of equal length is the number of positions at which the corresponding symbols are different

### **Hamming distance:**

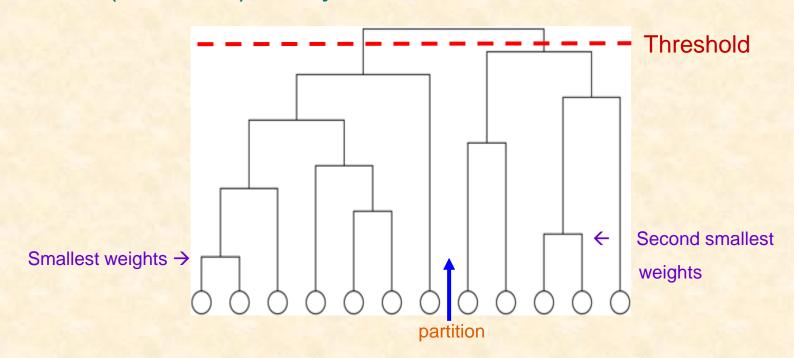
100 → 011 has distance 3 (red path)

 $010 \rightarrow 111$  has distance 2 (blue path)



### **Hierarchical Clustering Method**

- Calculate a "weight" (e.g., similarity value) for every pair of nodes, which represents how closely connected this pair of nodes is
- Starting with all disconnected nodes, add edges between pairs, one by one, in increasing order of their weights
- Result: Some nested components, where one can take a "slice" (threshold) at any level of the tree



## 3) GN Algorithm

Girvan–Newman (GN) algorithm identifies "heavy" edges in a network that lie between communities and then removes them, leaving only the communities

This is a betweenness-based clustering method: Identification is performed by employing the edgebetweenness, yielding results of reasonable quality

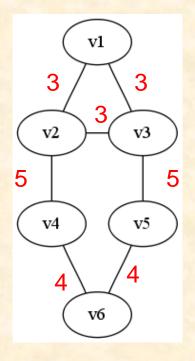
Drawbacks: High computational complexity

**GN** Algorithm

M. Girvan and M. E. J. Newman, PNAS, 99(12): 7821-7826, 2002

### **GN** Algorithm

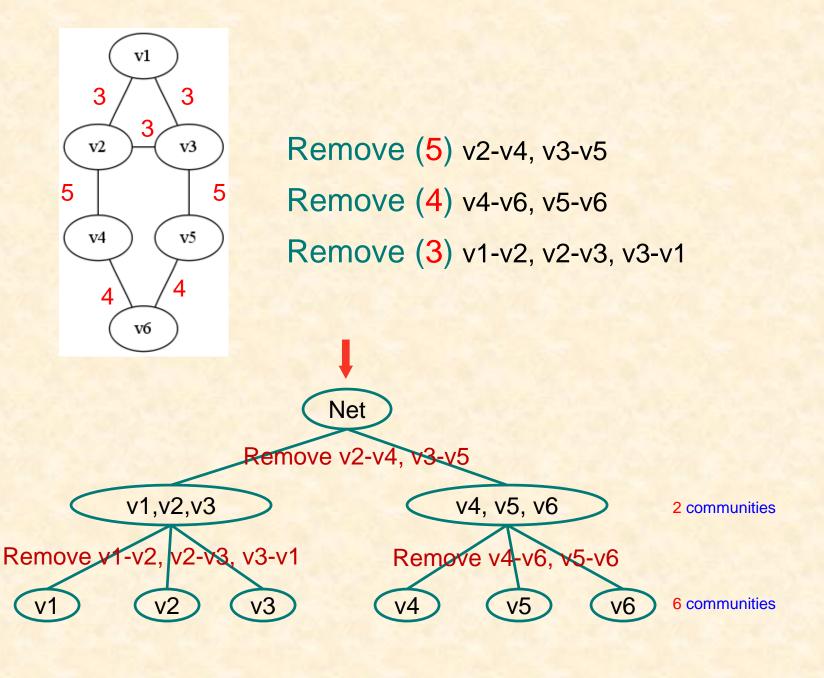
- Calculate all the edge-betweenness in the network
- Remove the edge with the highest betweenness
- Re-calculate all the edge-betweennesses for the resulting (smaller) network
- Repeat the above, until no edge is left



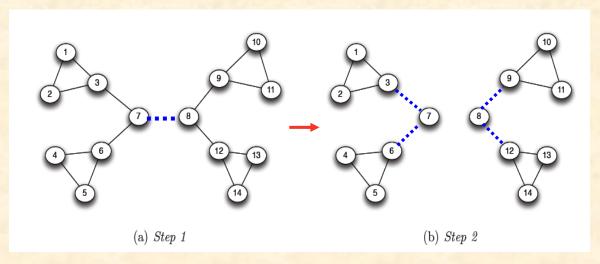
Remove (5) v2-v4, v3-v5

Remove (4) v4-v6, v5-v6

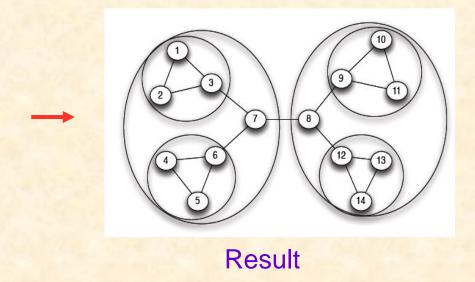
Remove (3) v1-v2, v2-v3, v3-v1



### **Example:**

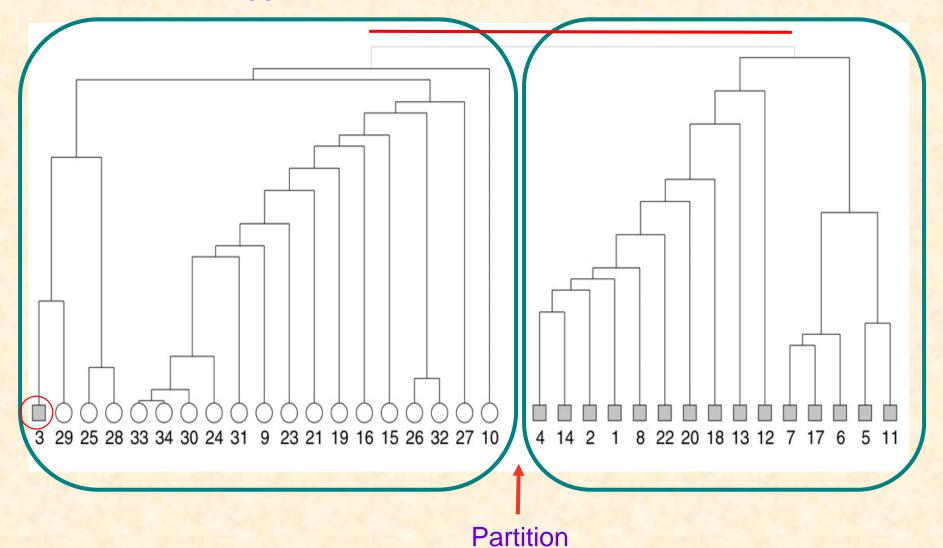


Step 1: remove 7-8 Step 2: remove 3-7, 6-7; 8-9, 8-12

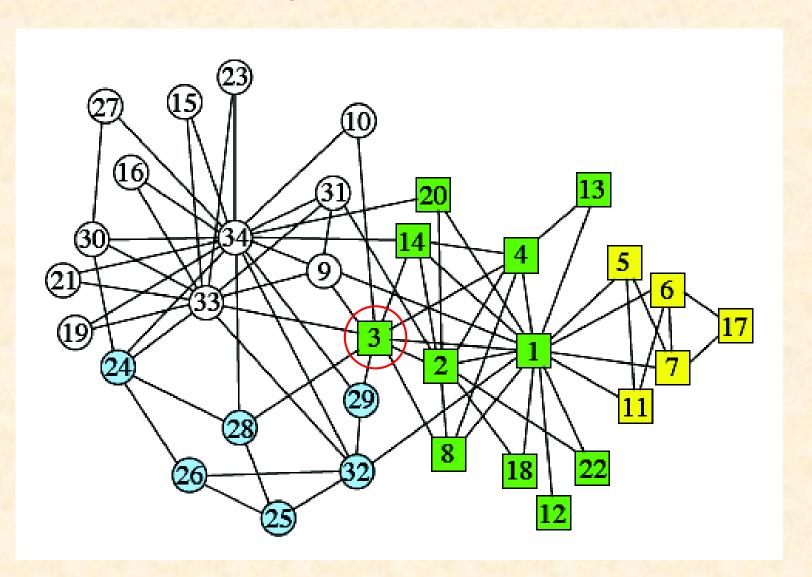


### **GN** Algorithm

### **Applied to the Karate Club Dataset**

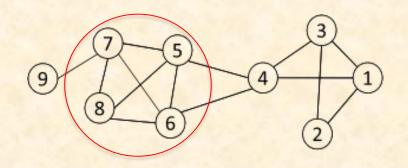


### **Zachary's Karate Network**



### 4) Clique-Based Methods

- > A clique is a fully-connected subgraph
- ➤ There are several clique-based community detection algorithms (computationally NP-hard)
- Since a node can be a member of more than one clique, these methods usually yield overlapping community structures



For example:

Nodes (5, 6, 7, 8) form a clique

Nodes (1, 2, 3), (1, 3, 4), (4, 5, 6), also form a clique, respectively

### **Clique-Based Methods**

- One approach is to find all the maximal cliques, which are cliques that are not a subgraph of any other clique
- > To find maximum cliques: Bron-Kerbosch Algorithm

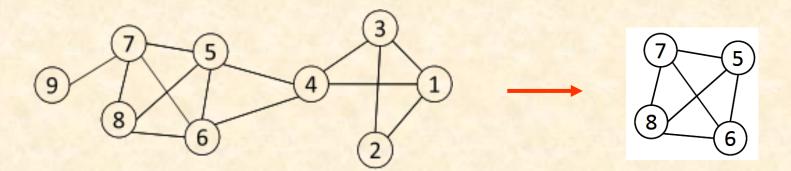
Recursively apply the following pruning procedure:

Sample a (large) subgraph from the given network, and find a clique of size k in it (say, by a greedy algorithm)

To find a larger clique, all nodes with degree  $\leq k-1$  are removed from the whole network

Repeat the above, until the network is small enough

### **Example**



- First, suppose we sampled a sub-network with nodes {1, 2, 3, 4, 5} and found a clique {1, 2, 3} of size 3
   [This is used as a reference to search for a possible larger clique]
- Next, to find a clique of size > 3, remove all nodes with degrees
   1 and 2
  - > Remove nodes 2, 9
  - > Remove nodes 1, 3, 4
- Result is a larger clique: {5, 6, 7, 8}
- If the network is huge, continue: size > 4, size > 5, ...

### **Clique Percolation Method**

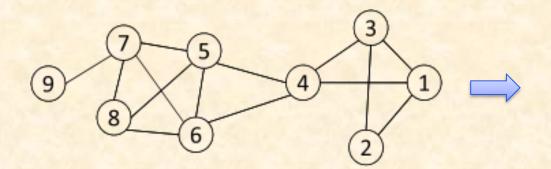
- A node can be a member of more than one clique. These methods usually yield overlapping community structures
- CPM is a method to find overlapping communities

**Input:** A network, and a parameter *k* 

#### **Procedure:**

- > Find all cliques of size k in the given network
- Construct a clique graph:
  - Two cliques are adjacent if they share k-1 nodes
- > Each connected cluster in the clique graph is a community

### **Example**



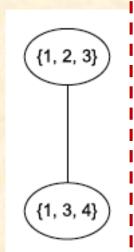
### **Cliques of size 3:**

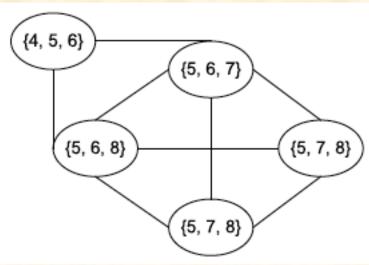
{1, 2, 3}, {1, 3, 4}, {4, 5, 6}, {5, 6, 7}, {5, 6, 8}, {5, 7, 8}, {6, 7, 8}

k = 3k - 1 = 2 shared nodes

#### **Communities:**

overlapping: node 4

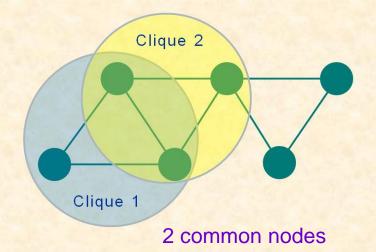




## **k-Clique Communities**

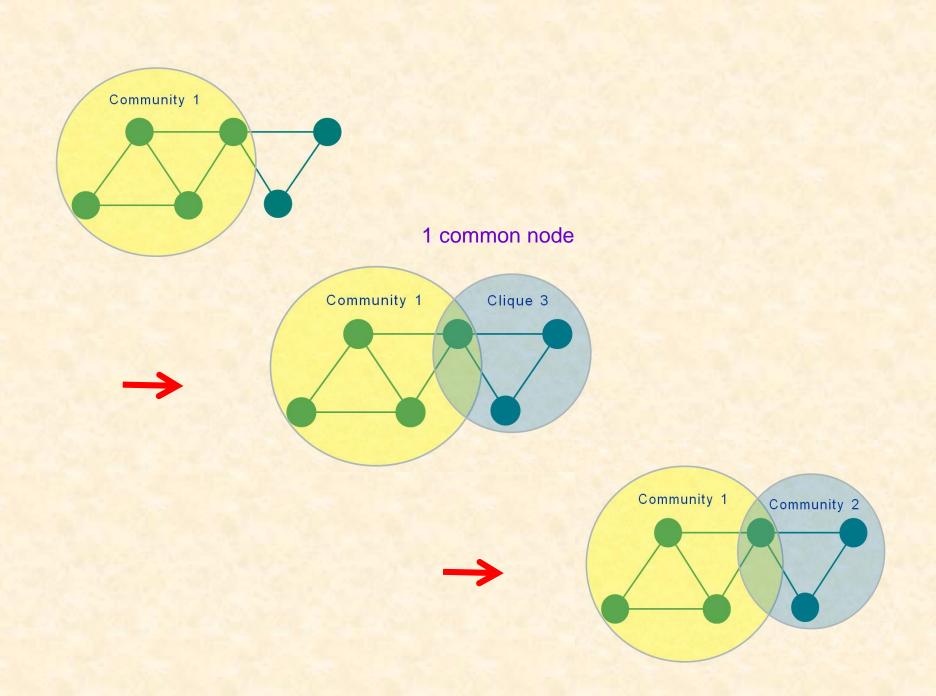
Union of all k-cliques that can be reached from each other through a series of adjacent k-cliques

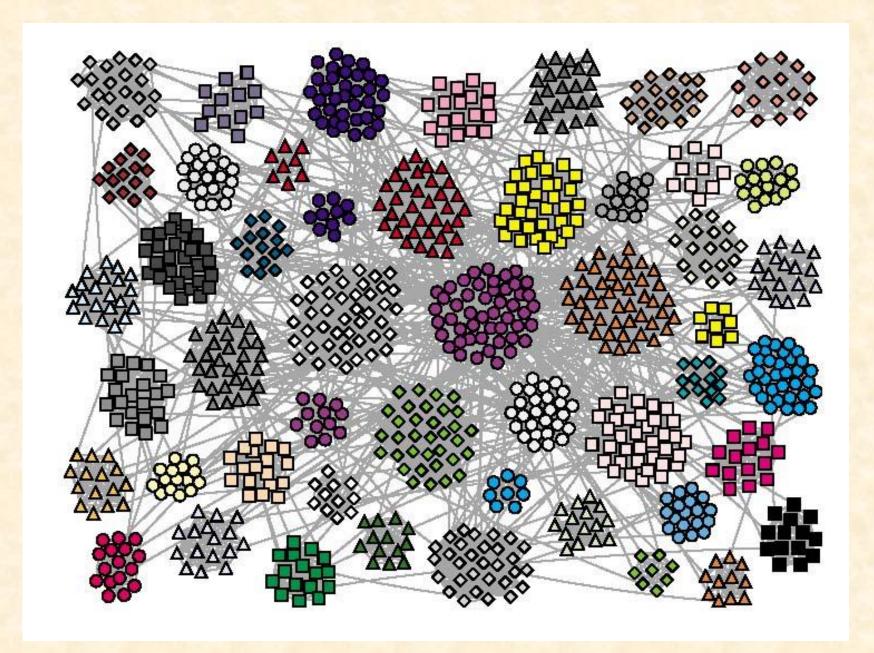
## Example:



k = 3

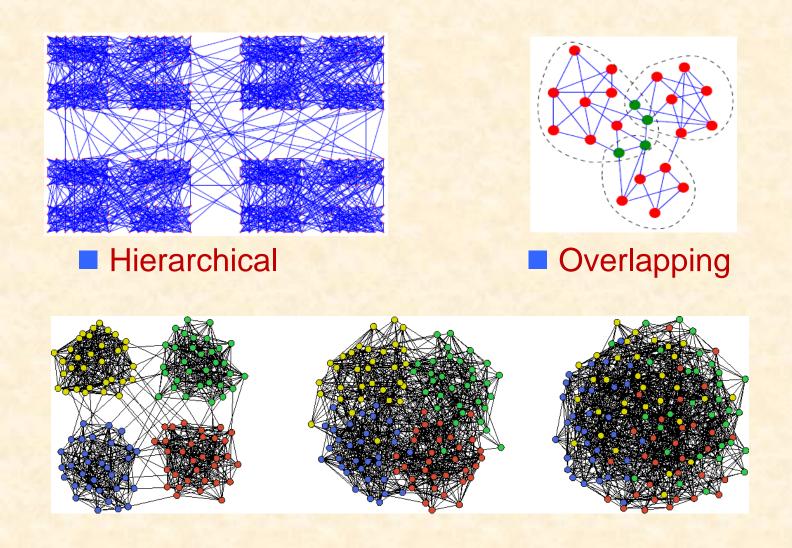
→ They belong to the same community



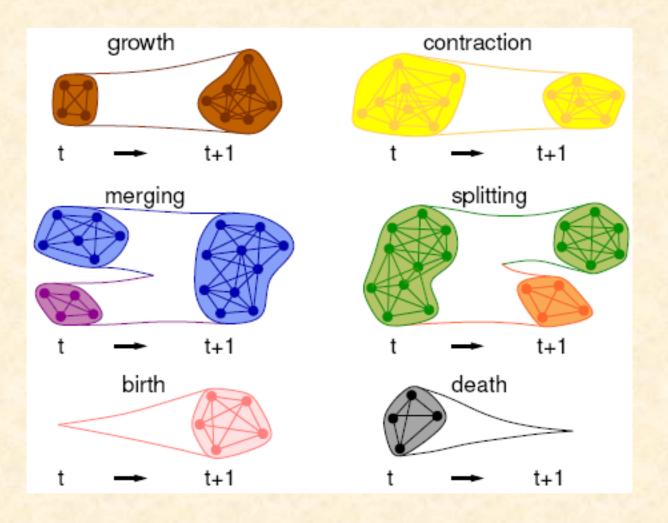


A. Lancichinetti, S. Fortunato, and F. Radicchi, Phys. Rev. E 78, 046110 (2008)

## **Detecting Community Structure: Challenges**

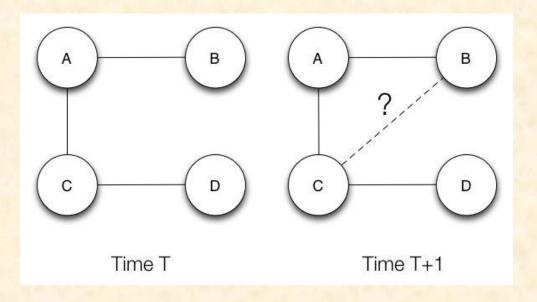


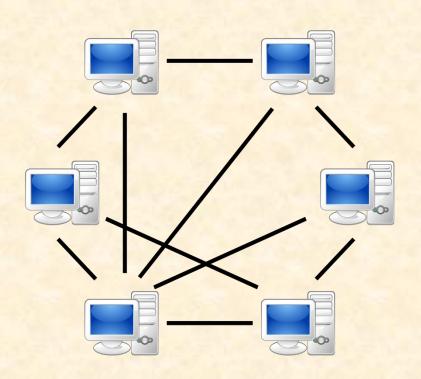
### **Detecting Community Structure: Challenges**



Evolution, Emergence

An emerging community?





### **Predict:**

Which computer is likely to connect to which computer



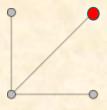
#### **Predict:**

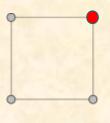
Which person is likely to connect to which friend

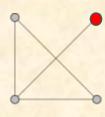
For a large-scale network, this could be challenging



### Simple Cases:







Degree sequences:

(3,1,1,1)

(2,2,2,2)

(3,2,2,1)

#### Criteria:

Similarity: Similar degrees, properties, importance, ...

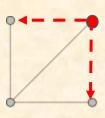
Commonality: Common friends, features, ...

Closeness: Closeness, distances, ...

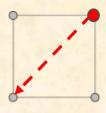
Based on:

**Degree similarity** 

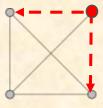
**Closeness/Distance** 



(3,1,1,1)



(2,2,2,2)



(3,2,2,1)

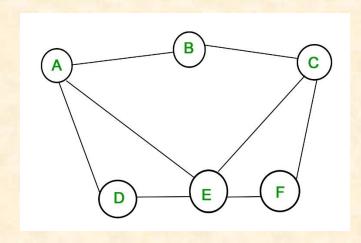
#### Criteria:

Similarity: Similar degrees, properties, importance, ...

Commonality: Common friends, features, ...

Closeness: Closeness, distances, ...

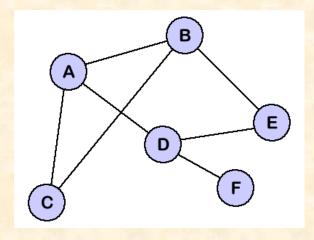
## **Examples**



Predict link(s) from node "B"

Answer: B - D and B - F

(They have degree 2)

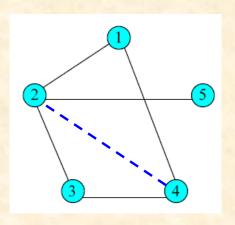


Predict link(s) from node "F"

Answer: F - E and A - F

(There is no other degree-1 node)

(They have a shortest distance)



Criterion:

Based on node-degree average

 $\bullet$  Give every node-pair (i, j) a value (weight):

$$(i,j) = \sqrt{k_i k_j}$$
 (node *i* has degree  $k_i$ )

- Predict a link between two un-connected nodes with a highest node-pair value:
- \*  $(1,2) = \sqrt{6}$ ,  $(1,3) = \sqrt{4}$ ,  $(1,4) = \sqrt{4}$ ,  $(1,5) = \sqrt{2}$  $(2,3) = \sqrt{6}$ ,  $(2,4) = \sqrt{6}$ ,  $(2,5) = \sqrt{3}$ ,  $(3,4) = \sqrt{4}$ ,  $(3,5) = \sqrt{2}$ ,  $(4,5) = \sqrt{2}$
- → Predict a new link: 2 4

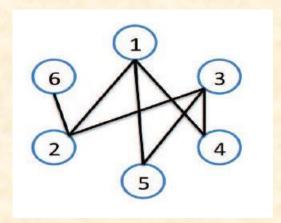
#### **Based on Commonality**

Neighborhood of node x:  $N(x) = \{i: i \text{ connects to } x\}$ 

Neighborhood of node y:  $N(y) = \{i: i \text{ connects to } y\}$ 

Intersect:  $N(x) \cap N(y) = \{i: i \text{ belongs to both neighborhoods}\}$ 

Cardinality:  $|N(x) \cap N(y)| =$  number of nodes in the intersect



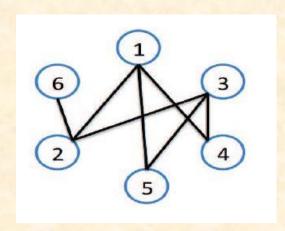
#### Example:

$$N(1) = \{2, 4, 5\}$$

$$N(6) = \{2\}$$

$$\rightarrow N(1) \cap N(6) = \{2\}$$

$$\rightarrow |N(1) \cap N(6)| = 1$$



$$|N(1) \cap N(3)| = 3$$

$$|N(2) \cap N(4)| = 2$$

$$|N(2) \cap N(5)| = 2$$

$$|N(4) \cap N(5)| = 2$$

$$|N(1) \cap N(6)| = 1$$

$$|N(3) \cap N(6)| = 1$$

$$|N(4) \cap N(6)| = 0$$

$$|N(5) \cap N(6)| = 0$$

### **Based on Commonality**

 $\rightarrow$ 

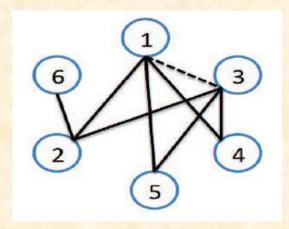
Predict 1st new link: 1 – 3

Predict 2nd new links:

$$2 - 4$$
 or  $2 - 5$  or  $4 - 5$ 

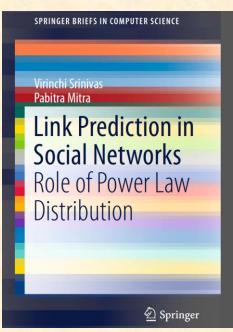
and so on





### Other Link Prediction Methods

- Node-similarity-based methods
- Topology-similarity-based methods
  CN, AA, RA, LP, Katz, LRW, SRW, RWR, ...
- Maximum-likelihood analytic methods Layer-structure modeling, random partitioning, ...
- Machine Learning
- **\*** .......



# BREAK 10 minutes