

Chapter 3: Processes

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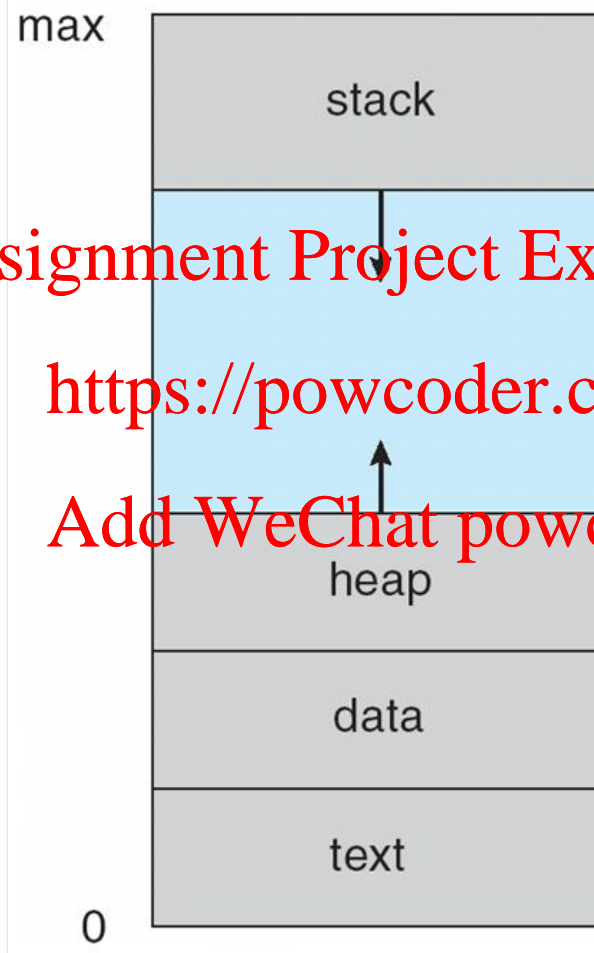
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Process Concept

- **Process** — a program in execution; process execution must progress in sequential fashion
- Multiple parts
 - The program code, also called **text section**
 - Current activity including **program counter**, processor registers
 - **Stack** containing temporary data
 - Function parameters, return addresses, local variables
 - **Data section** containing global variables
 - **Heap** containing memory dynamically allocated during run time
- Program is *passive* entity stored on disk (**executable file**), process is *active*
 - Program becomes process when executable file loaded into memory
- One program can be several processes
 - Consider multiple users executing the same program

Process in Memory



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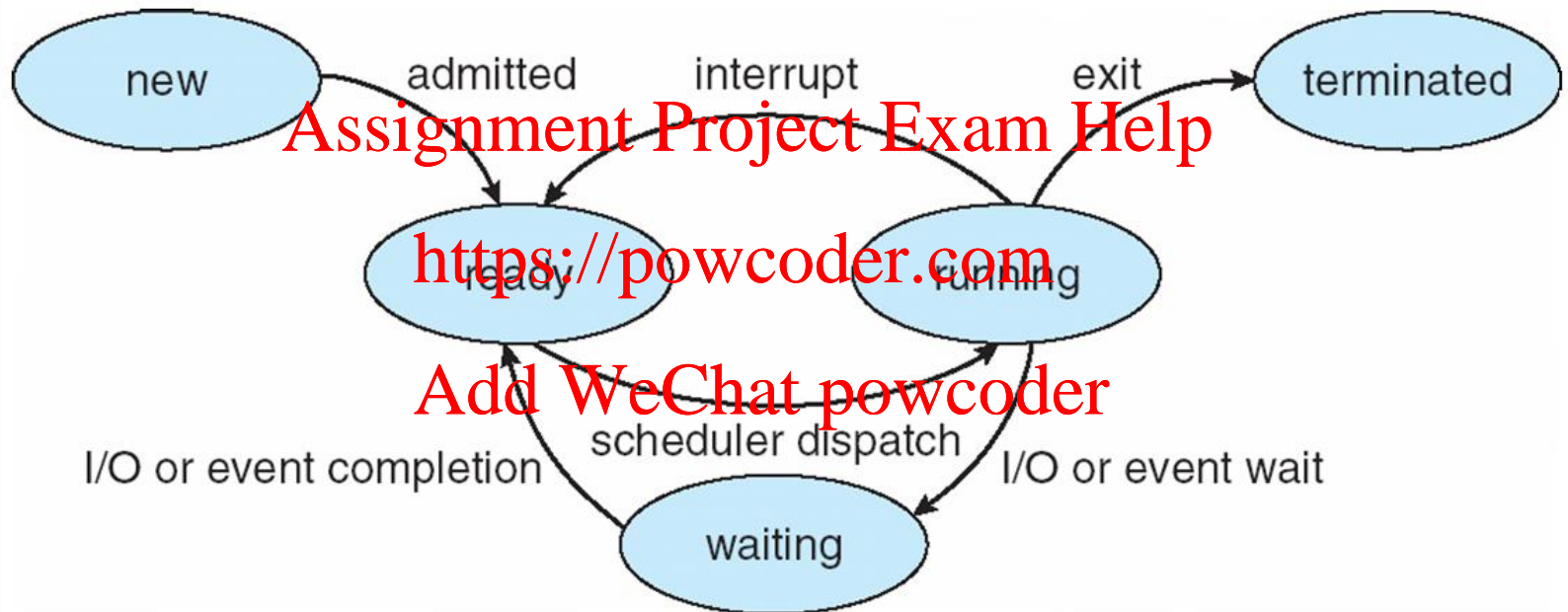
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Process State

- As a process executes, it changes **state**
 - **new**: The process is being created
 - **running**: Instructions are being executed
 - **waiting**: The process is waiting for some event to occur <https://powcoder.com>
 - **ready**: The process is waiting to be assigned to a processor [Add WeChat powcoder](#)
 - **terminated**: The process has finished execution

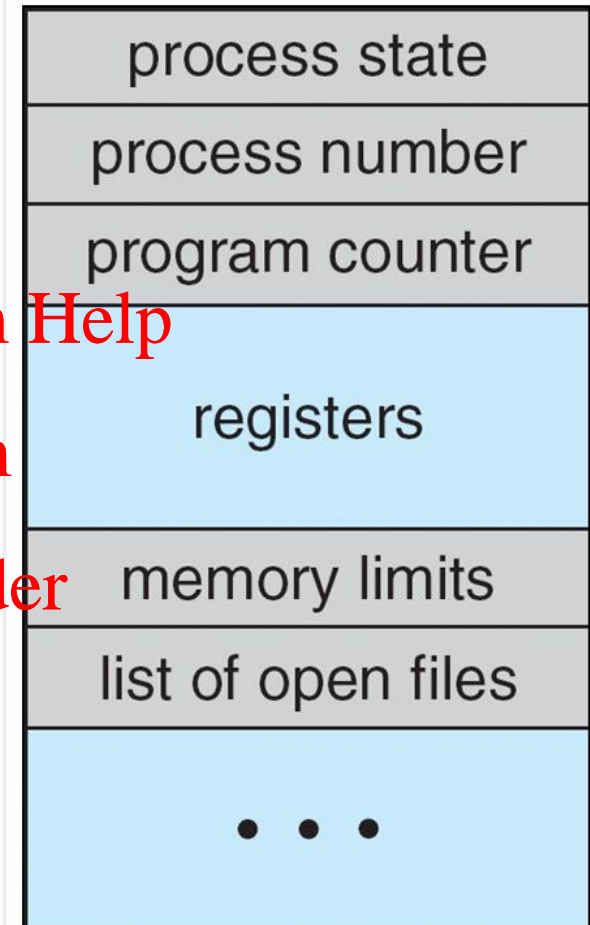
Diagram of Process State



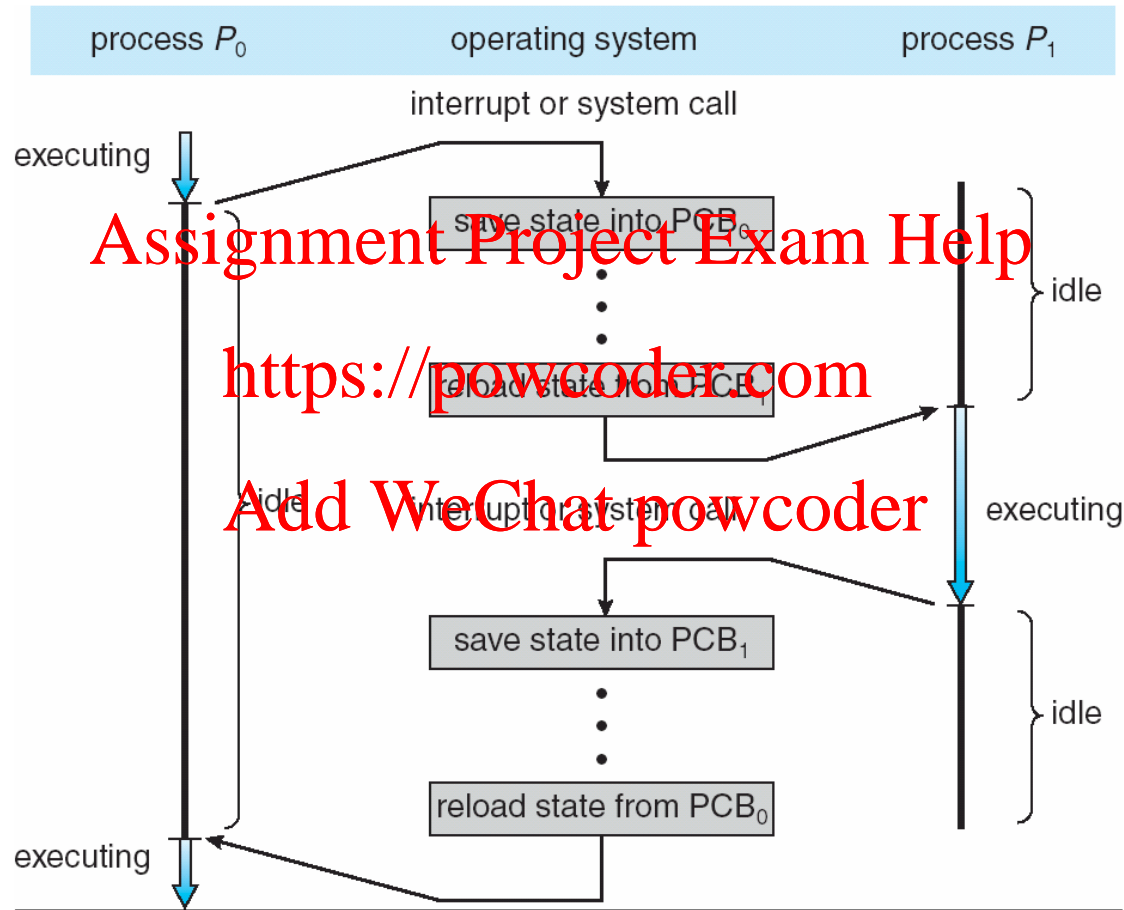
Process Control Block (PCB)

Information associated with each process
(also called **task control block**)

- Process state – running, waiting, etc
- Program counter – location of instruction to next execute
- CPU registers – contents of all process-centric registers
- CPU scheduling information- priorities, scheduling queue pointers
- Memory-management information – memory allocated to the process
- I/O status information – I/O devices allocated to process, list of open files



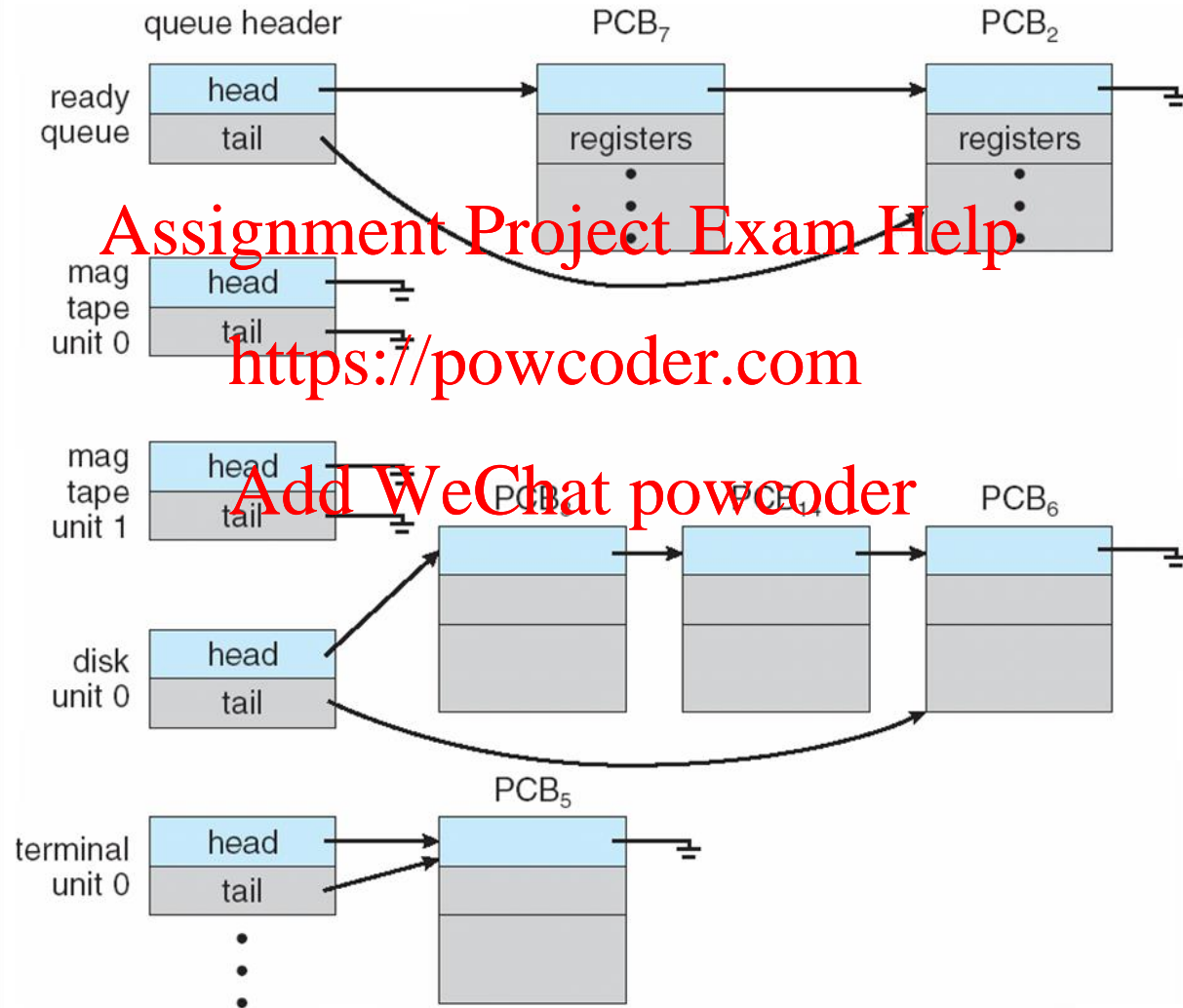
CPU Switch From Process to Process



Process Scheduling

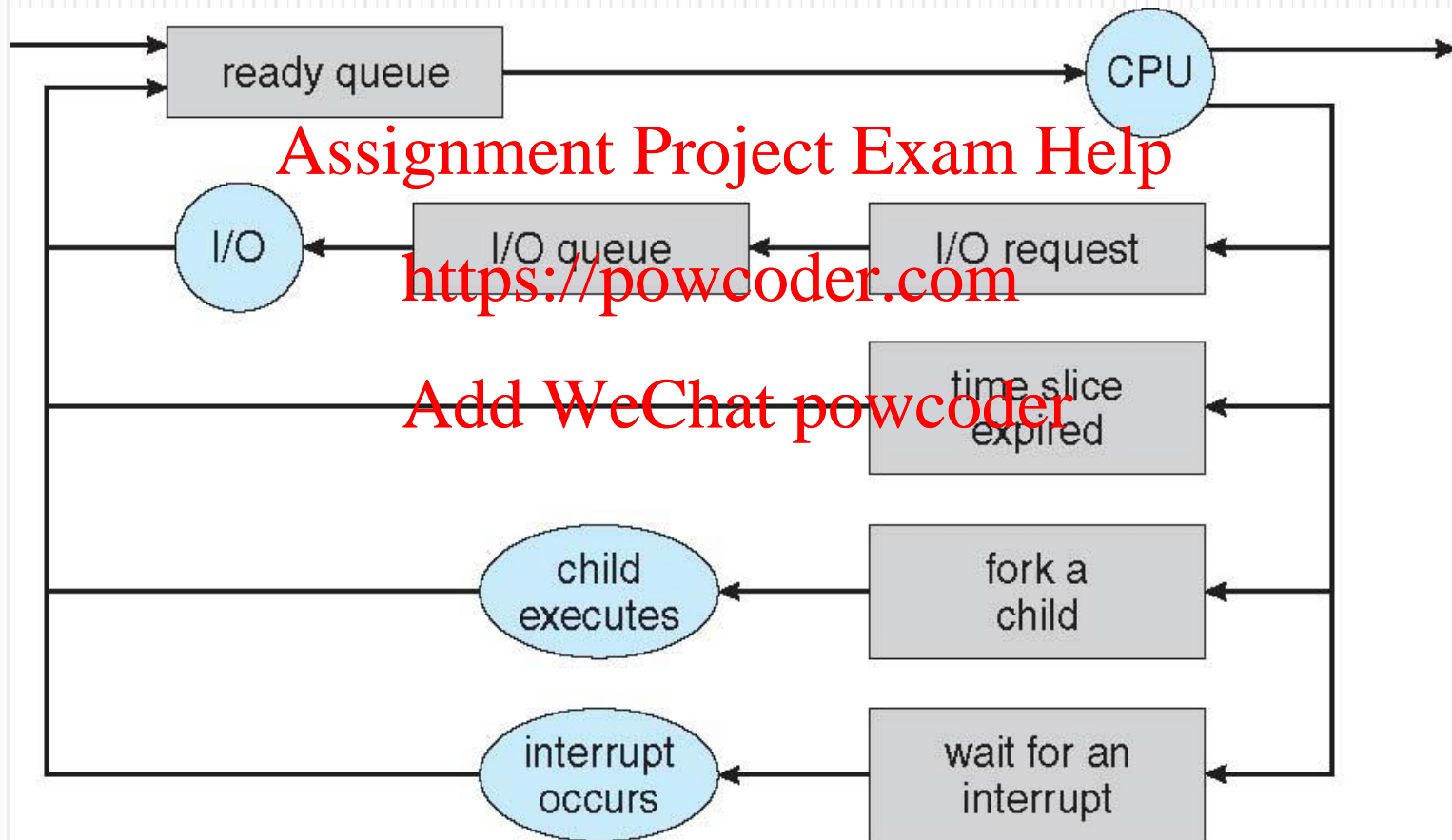
- Maximize CPU use, quickly switch processes onto CPU for time sharing
- **Process scheduler** selects among available processes for next execution on CPU
- Maintains **scheduling queues** of processes
 - **Job queue** – set of all processes in the system
 - **Ready queue** – set of all processes residing in main memory, ready and waiting to execute
 - **Device queues** – set of processes waiting for an I/O device
 - Processes migrate among the various queues

Ready Queue And Various I/O Device Queues



Representation of Process Scheduling

□ Queueing diagram represents queues, resources, flows



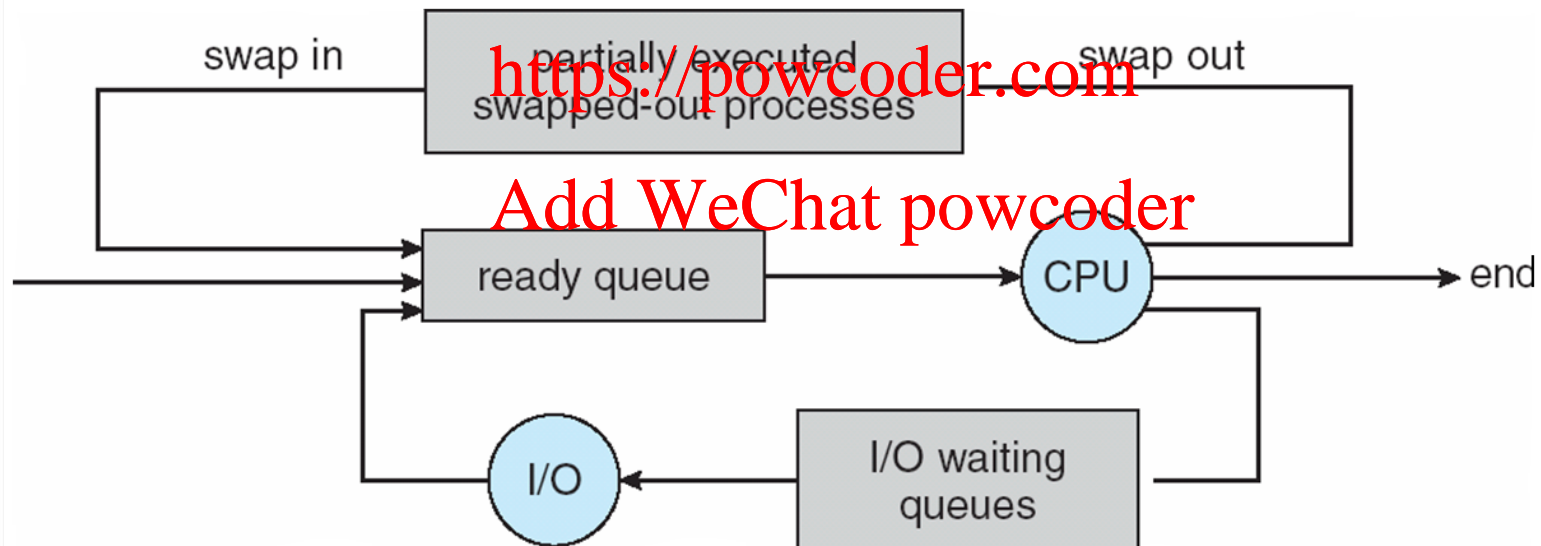
Schedulers

- **Long-term scheduler** (or **job scheduler**) – selects which processes should be brought into the ready queue - The long-term scheduler controls the **degree of multiprogramming** - Long-term scheduler is invoked very infrequently (seconds, minutes) \Rightarrow (may be slow)
- **Short-term scheduler** (or **CPU scheduler**) – selects which process should be executed next and allocates CPU - Sometimes the only scheduler in a system - Short-term scheduler is invoked very frequently (milliseconds) \Rightarrow (must be fast)
- Processes can be described as either:
 - **I/O-bound process** – spends more time doing I/O than computations, many short CPU bursts
 - **CPU-bound process** – spends more time doing computations; few very long CPU bursts
- Long-term scheduler strives for good *process mix*

Addition of Medium Term Scheduling

- **Medium-term scheduler** can be added if degree of multiple programming needs to decrease
 - Remove process from memory, store on disk, bring back in from disk to continue execution: **swapping**

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Context Switch

- When CPU switches to another process, the system must **save the state** of the old process and load the **saved state** for the new process via a **context switch**
- **Context** of a process represented in the PCB
- Context-switch time is overhead; the system does no useful work while switching
 - The more complex the OS and the PCB -> longer the context switch
- Time dependent on hardware support
 - Some hardware provides multiple sets of registers per CPU -> multiple contexts loaded at once

Context Switch

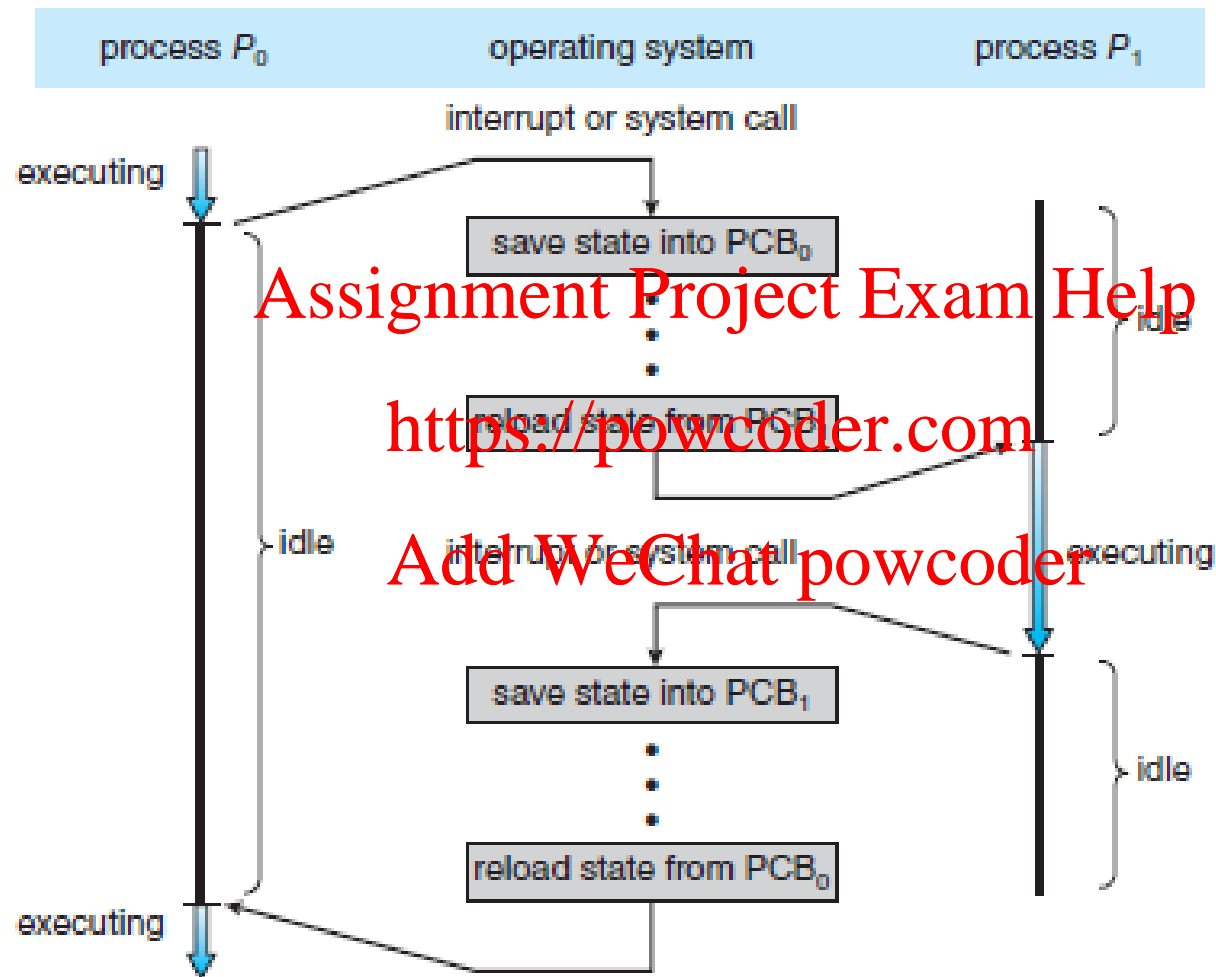


Figure 3.6 Diagram showing context switch from process to process.

Process Creation

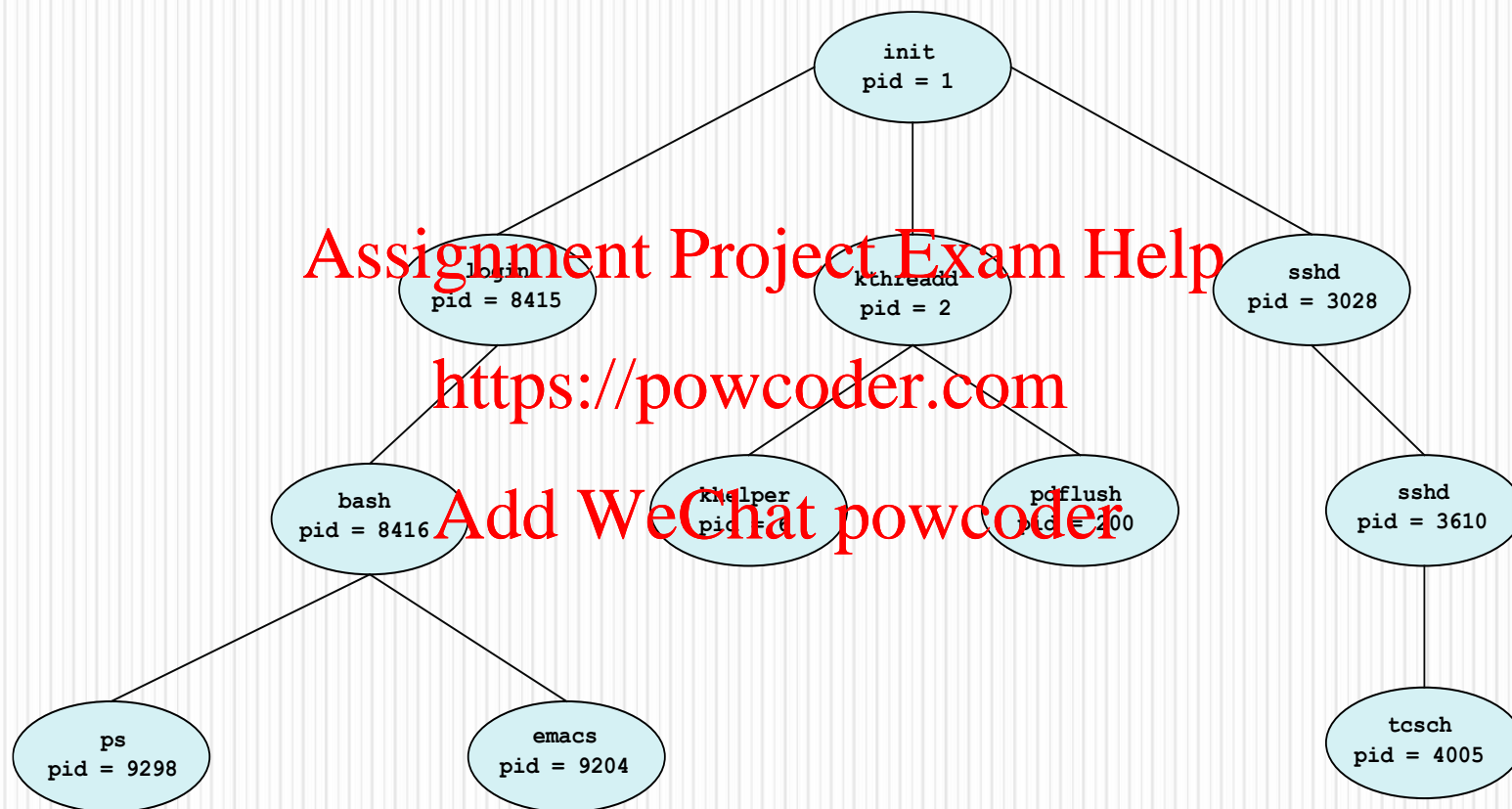
- **Parent** process create **children** processes, which, in turn create other processes, forming a **tree** of processes

- Generally, process identified and managed via a **process identifier (pid)**

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- Resource sharing options
 - Parent and children share all resources
 - Children share subset of parent's resources
 - Parent and child share no resources
- Execution options
 - Parent and children execute concurrently
 - Parent waits until children terminate

A Tree of Processes in Linux

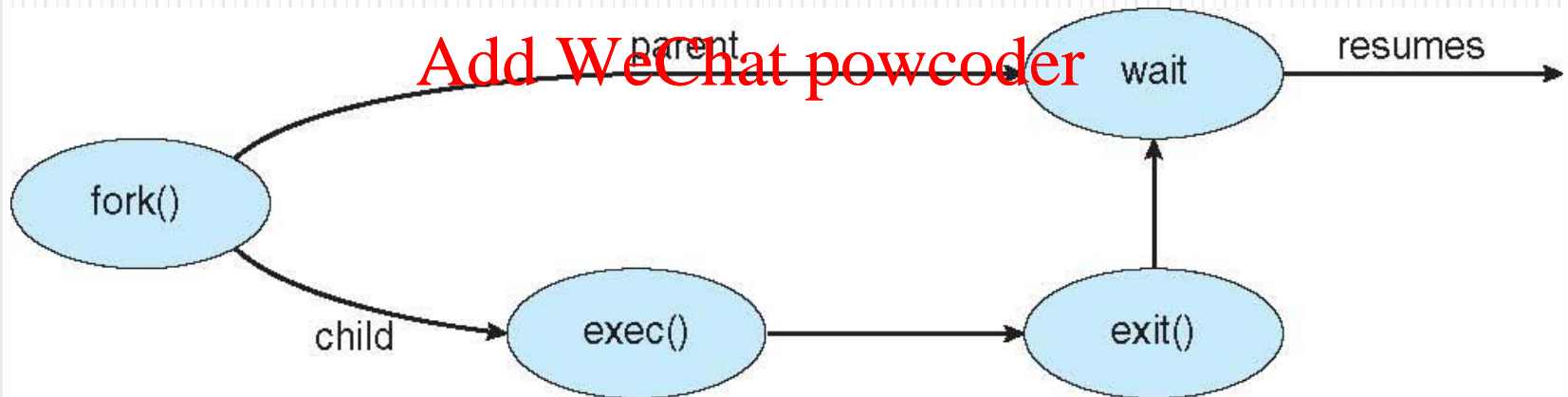


Process Creation (Cont.)

- Address space
 - Child duplicate of parent
 - Child has a program loaded into it

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- UNIX examples <https://powcoder.com>



C Program Forking Separate Process

```
#include <sys/types.h>
#include <stdio.h>
#include <unistd.h>

int main()
{
    pid_t pid;

    /* fork a child process */
    pid = fork();

    if (pid < 0) { /* error occurred */
        fprintf(stderr, "Fork Failed");
        return 1;
    }
    else if (pid == 0) { /* child process */
        execlp("/bin/ls", "ls", NULL);
    }
    else { /* parent process */
        /* parent will wait for the child to complete */
        wait(NULL);
        printf("Child Complete");
    }

    return 0;
}
```

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Process Termination

- Process executes last statement and asks the operating system to delete it (**exit()**)
 - Output data from child to parent (via **wait()**)
 - Process' resources are deallocated by operating system
- Parent may terminate execution of children processes (**abort()**)
 - Child has exceeded allocated resources
 - Task assigned to child is no longer required
 - If parent is exiting
 - Some operating systems do not allow child to continue if its parent terminates
 - All children terminated - **cascading termination**

Process Termination

- Wait for termination, returning the pid:

```
pid_t t_pid; int status;
```

```
pid = wait(&status);
```

- If no parent waiting, then terminated process is a **zombie**
- If parent terminated, processes are **orphans**

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Interprocess Communication

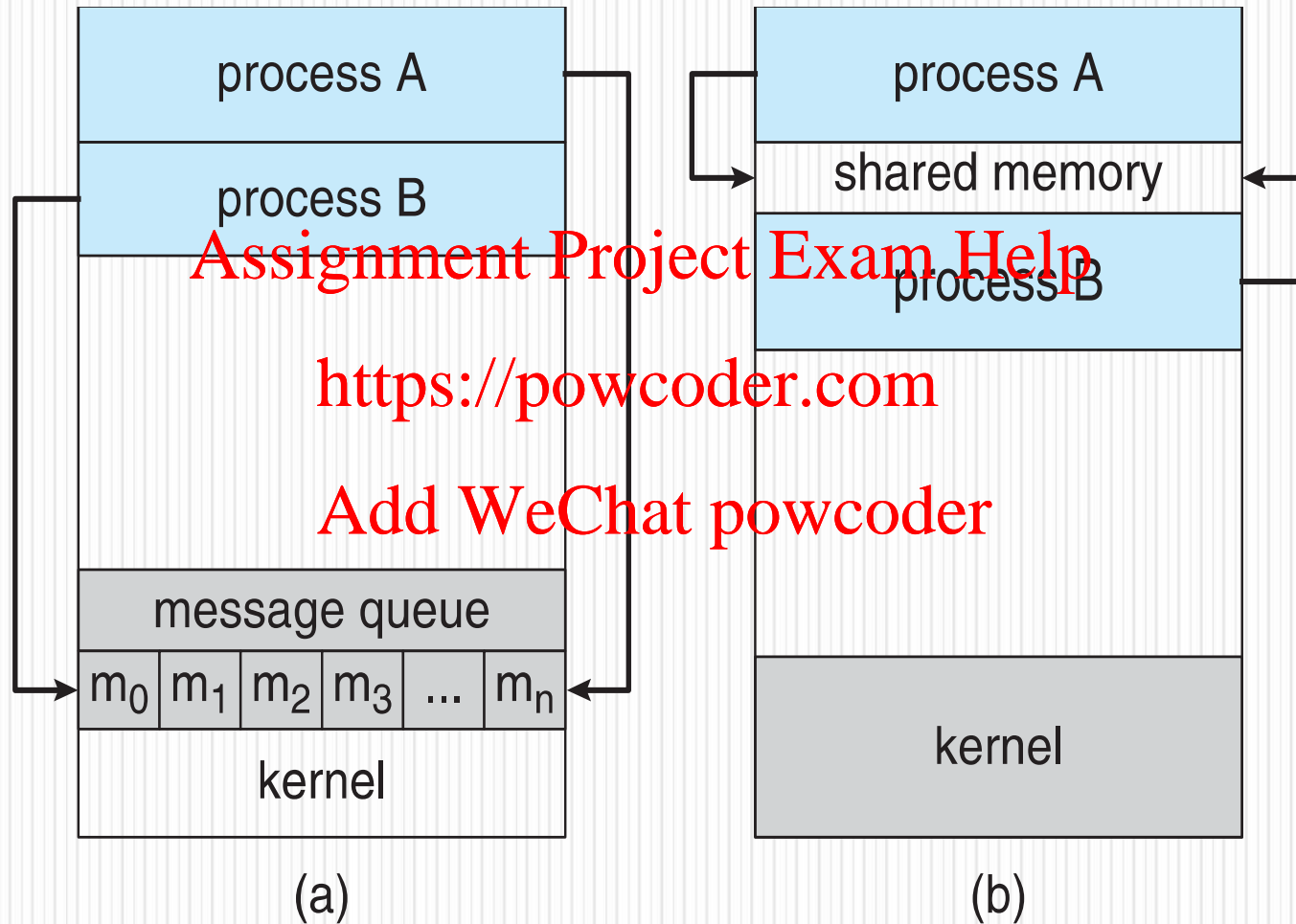
- Processes within a system may be *independent* or *cooperating*
- Cooperating process can affect or be affected by other processes, including sharing data
- Reasons for cooperating processes:
 - Information sharing
 - Computation speedup
 - Modularity
 - Convenience
- Cooperating processes need **interprocess communication (IPC)**
- Two models of IPC
 - **Shared memory**
 - **Message passing**

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Communications Models



Message Passing Vs. Shared Memory

Message passing:

- Useful for exchanging smaller amounts of data
- Easier to implement
- Implemented using system calls
- Every message requires kernel intervention - slower

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Shared memory: Add WeChat powcoder

- Useful for exchanging larger amounts of data
- Requires synchronization
- System calls are only required to set up shared-memory region
- Once set up, all accesses are routine memory accesses, no further kernel assistance required - faster

Cooperating Processes

- *Independent* process cannot affect or be affected by the execution of another process

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- *Cooperating* process can affect or be affected by the execution of another process

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- Advantages of process cooperation
 - Information sharing
 - Computation speed-up
 - Modularity
 - Convenience

Producer-Consumer Problem

- Paradigm for cooperating processes, *producer* process produces information that is consumed by a *consumer* process
 - **unbounded-buffer** places no practical limit on the size of the buffer
 - **bounded-buffer** assumes that there is a fixed buffer size

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Bounded-Buffer – Shared-Memory Solution

- Shared data

```
#define BUFFER_SIZE 10
```

```
typedef struct {
```

```
    . . .
```

```
} item;
```

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```
item buffer[BUFFER_SIZE];
```

```
int in = 0;
```

```
int out = 0;
```

- Solution is correct, but can only use BUFFER_SIZE-1 elements

Bounded-Buffer – Producer

```
item next_produced;
while (true) {
    /* produce an item in next produced */
    while (((in + 1) % BUFFER_SIZE) == out)
        ; /* do nothing */
    buffer[in] = next_produced;
    in = (in + 1) % BUFFER_SIZE;
}
```

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Bounded Buffer – Consumer

```
item next_consumed;  
while (true) {  
    while (in == out)  
        ; /* do nothing */  
    next_consumed = buffer[out];  
    out = (out + 1) % BUFFER_SIZE;  
  
    /* consume the item in next  
consumed */  
}
```

Interprocess Communication – Message Passing

- Mechanism for processes to communicate and to synchronize their actions
- Message system – processes communicate with each other without resorting to shared variables
- IPC facility provides two operations:
 - **send**(*message*) – message size fixed or variable
 - **receive**(*message*)
- If *P* and *Q* wish to communicate, they need to:
 - establish a **communication link** between them
 - exchange messages via send/receive
- Implementation of communication link
 - physical (e.g., shared memory, hardware bus)
 - logical (e.g., direct or indirect, synchronous or asynchronous, automatic or explicit buffering)

Implementation Questions

- How are links established?
- Can a link be associated with more than two processes?
- How many links can there be between every pair of communicating processes?
- What is the capacity of a link?
- Is the size of a message that the link can accommodate fixed or variable?
- Is a link unidirectional or bi-directional?

Direct Communication

- Processes must name each other explicitly:
 - send** ($P, message$) – send a message to process P
 - receive** ($Q, message$) – receive a message from process Q

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- Properties of communication link
 - Links are established automatically
 - A link is associated with exactly one pair of communicating processes
 - Between each pair there exists exactly one link
 - The link may be unidirectional, but is usually bi-directional

Indirect Communication

- Messages are directed and received from mailboxes (also referred to as ports)
 - Each mailbox has a unique id
 - Processes can communicate only if they share a mailbox

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- Properties of communication link
 - Link established only if processes share a common mailbox
 - A link may be associated with many processes
 - Each pair of processes may share several communication links
 - Link may be unidirectional or bi-directional

Indirect Communication

- Operations
 - create a new mailbox
 - send and receive messages through mailbox
 - destroy a mailbox

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- Primitives are defined as:

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send($A, message$) – send a message to mailbox A

receive($A, message$) – receive a message from mailbox A

Indirect Communication

- Mailbox sharing

- P_1 , P_2 , and P_3 share mailbox A
- P_1 sends; P_2 and P_3 receive
- Who gets the message?

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- Solutions

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- Allow a link to be associated with at most two processes
- Allow only one process at a time to execute a receive operation
- Allow the system to select arbitrarily the receiver. Sender is notified who the receiver was.

Synchronization

- Message passing may be either blocking or non-blocking
- **Blocking** is considered **synchronous**
 - **Blocking send** has the sender block until the message is received
 - **Blocking receive** has the receiver block until a message is available
- **Non-blocking** is considered **asynchronous**
 - **Non-blocking send** has the sender send the message and continue
 - **Non-blocking receive** has the receiver receive a valid message or null

Solution to Producer-Consumer Problem using Blocking Send() and Receive()

```
message next_produced;
while (true){
    /* produce an item in next produced */
    send(next_produced);
}

message next_consumed;
while (true) {
    receive(next_consumed);
    /* consume the item in next consumed */
}
```

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Buffering

- Queue of messages attached to the link; implemented in one of three ways
 1. Zero capacity – 0 messages
Sender must wait for receiver (rendezvous)
 2. Bounded capacity – finite length of n messages
Sender must wait if link full
 3. Unbounded capacity – infinite length
Sender never waits

Communications in Client-Server Systems

- Sockets
- Remote Procedure Calls

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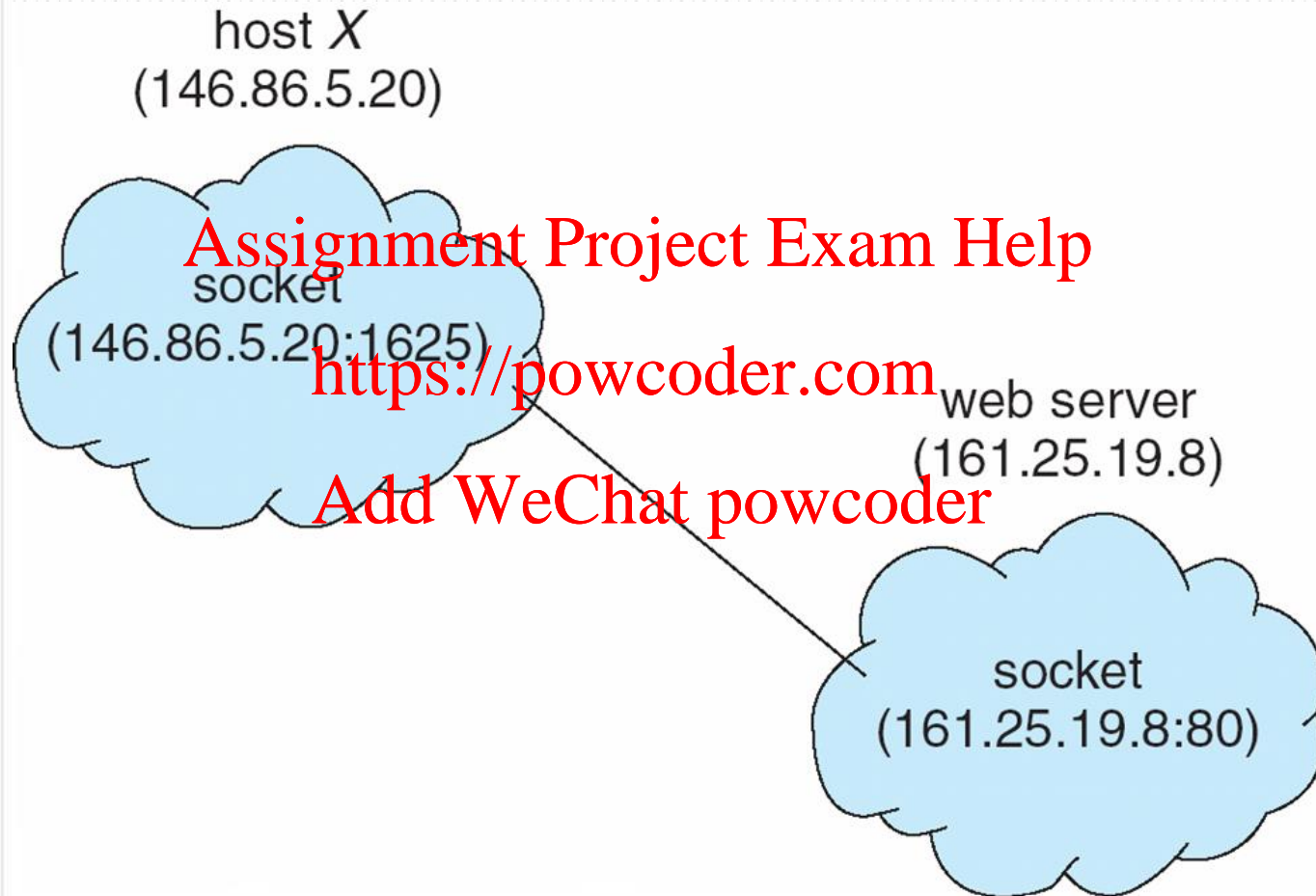
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Sockets

- A **socket** is defined as an endpoint for communication
- Concatenation of IP address and **port** – a number included at start of message packet to differentiate network services on a host
- The socket **161.25.19.8:1625** refers to port **1625** on host **161.25.19.8**
- Communication consists between a pair of sockets
- All ports below 1024 are *well known*, used for standard services

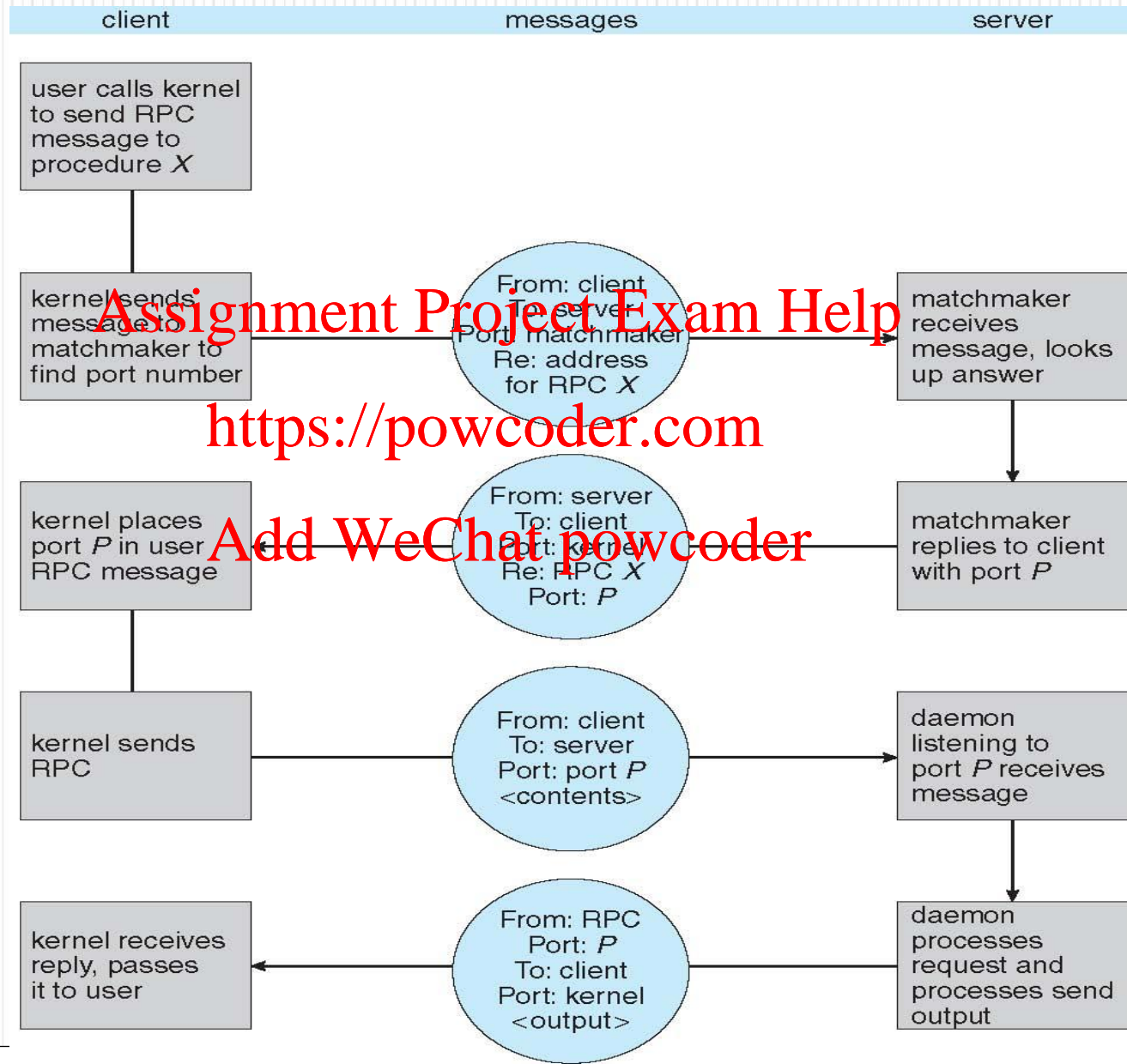
Socket Communication



Remote Procedure Calls

- Remote procedure call (RPC) abstracts procedure calls between processes on networked systems
 - Again uses ports for service differentiation
- **Stubs** – client-side proxy for the actual procedure on the server
- The client-side stub locates the server and **marshalls** the parameters
- The server-side stub receives this message, unpacks the marshalled parameters, and performs the procedure on the server
- Data representation handled via **External Data Representation (XDL)** format to account for different architectures
- Remote communication has more failure scenarios than local
 - Messages can be delivered *exactly once* rather than *at most once*
- OS typically provides a rendezvous (or **matchmaker**) service to connect client and server

Execution of RPC



End of Chapter 3

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