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```
print(np.mean(X, axis=0))
```

What is the expected runtime of this algorithm? Well, it will be the number of 1s in the array, plus one, unless we have $X_n = 1$, in which case we will know that we don't have to generate an additional geometric random variable.

So this has expectation: $E[\sum X_i + 1 | X_n = 0]P(X_n = 0) + E[\sum X_i | X_n = 1]P(X_n = 1) = p(n-1) + 1 = pn + (1-p)$. This suggests a small optimization: if $p < 0.5$, we proceed as before; if $p > 0.5$, use $q = (1-p)$ as the parameter for the geometric random variable, and swap the 1s and 0s in the above algorithm.

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Sampling a Binomial Random

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Variable

In the last section, we showed a method for generating n independent $\text{Bernoulli}(p)$ random variables. Clearly, we can use the method to generate a $\text{Binomial}(n, p)$ by taking their sum.

However, we could also directly exploit the inverse cdf transformation to sample Binomials. There is, of course, no closed-form formula for the binomial cdf, so we will have to do something like this:

- Step 1: $U \sim \text{Unif}(0, 1)$
- Step 2: $i = 0$, let $pr = \binom{n}{i} p^i * (1-p)^{n-i}$, $F = pr$
- Step 3: If $U \leq F$ return $X = i$
- Step 4: $i = i + 1$, $pr = pr + \binom{n}{i} p^i * (1-p)^{n-i}$, $F = F + pr$
- Step 5: go to step 3

Using this method, we make $X + 1$ comparisons, giving us an expected number of iterations of $np + 1$.