

# Lecture 15: Duality, Arbitrage, and

## Asset Pricing

ISyE 6073: Financial Optimization

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## Midterm survey recap

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- 36% response rate
- Thank you for all the feedback!
- Top comments:
  - Recording lectures: difficult to do without more AV tech support
  - Typos: I will do better, please keep me accountable on Piazza
  - Workload/difficulty: exam 1 was too long, some HWs are too long; I will try to shorten homeworks and create some extra credit opportunities

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# Duality Recap

- Lecture 10: Lagrangian duality
  - Lagrangian relaxation
  - Dual problem is finding the best Lagrangian relaxation
- Lecture 11: Sensitivity analysis
  - Origin of shadow prices
  - Generalized dependence on the right-hand side  $b$
- Lecture 12: Application to A/B testing
- Lecture 13: Geometry of duality
  - Normal cones
  - Farkas lemma
- Lecture 14: Large-scale optimization
  - Column generation and cutting planes

Today

Applications of duality in asset pricing and relationship to arbitrage!

## Today

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- A market model
- What are desirable properties of a market?
  - Completeness
  - Law of one price
  - Fundamental theorem(s) of asset pricing

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Two-period market model

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## Market setup

- We consider a portfolio of  $n$  assets (could be bonds, stocks, currencies, etc.)
- Two-period model:
  - At time 0 (today), we know the *price* of each asset, represented by a vector  $\mathbf{p}$

$$\mathbf{p} = \begin{bmatrix} p_1 \\ \vdots \\ p_n \end{bmatrix}$$

- At time 1 (tomorrow), the asset prices are *uncertain*; the world could be in one of  $m$  scenarios
- Scenario matrix:

$$\mathbf{S} = \begin{bmatrix} S_{11} & \cdots & S_{1n} \\ \vdots & \ddots & \vdots \\ S_{m1} & \cdots & S_{mn} \end{bmatrix}, \quad S_{ij} \text{ is the price of asset } j \text{ in scenario } i$$

## Market setup

### Example

Consider the following market

- We have  $n = 3$  assets
- Prices at time 0 (today) are given by  $\mathbf{p} = \begin{bmatrix} 1 \\ 2 \\ 7 \end{bmatrix}$
- At time 1 (tomorrow) the world could be in one of 4 scenarios:

$$\mathbf{S} = \begin{bmatrix} 1 & 3 & 9 \\ 1 & 1 & 5 \\ 1 & 5 & 10 \\ 1 & 2 & 8 \end{bmatrix}$$

**Q:** Are all these assets equally risky?

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# Constructing portfolios

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A portfolio  $\mathbf{x} = \begin{bmatrix} x_1 \\ \vdots \\ x_n \end{bmatrix}$  counts shares purchased at time 0

- $x_i > 0$  means a long position in asset  $i$
- $x_i < 0$  means a short position in asset  $i$
- Price of portfolio at time 0 is  $\mathbf{p}^\top \mathbf{x} = \sum_{j=1}^n p_j x_j$
- Payoff of portfolio at time 1 is

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$$\mathbf{S}\mathbf{x} = \sum_{j=1}^n \mathbf{S}_j x_j = \begin{bmatrix} \sum_{j=1}^n S_{1j} x_j \\ \vdots \\ \sum_{j=1}^n S_{mj} x_j \end{bmatrix} \in \mathbb{R}^m$$

- One payoff value for each scenario



# Constructing portfolios

## Example

- Market setup:  $\mathbf{p} = \begin{bmatrix} 1 \\ 2 \\ 7 \end{bmatrix}$ ,  $\mathbf{S} = \begin{bmatrix} 1 & 3 & 9 \\ 1 & 1 & 5 \\ 1 & 5 & 10 \\ 1 & 2 & 8 \end{bmatrix}$

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- Portfolio  $\mathbf{x} = \begin{bmatrix} 1 \\ -1 \\ 1 \end{bmatrix}$
- Price:  $\mathbf{p}^T \mathbf{x} = 1 - 2 + 7 = 6$

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- Payoff:  $\mathbf{S}\mathbf{x} = \begin{bmatrix} 1 \\ 1 \\ 1 \\ 1 \end{bmatrix} \times 1 + \begin{bmatrix} 3 \\ 1 \\ 5 \\ 2 \end{bmatrix} \times (-1) + \begin{bmatrix} 9 \\ 5 \\ 10 \\ 8 \end{bmatrix} \times 1 = \begin{bmatrix} 7 \\ 5 \\ 6 \\ 7 \end{bmatrix}$

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Complete markets  
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# Completeness

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### Definition

A market is *complete* if any payoff  $\mathbf{z} \in \mathbb{R}^m$  can be achieved by some portfolio  $\mathbf{x}$ .

- The **asset span**  $\mathcal{M}$  is the set of payoffs achievable by some portfolio  $\mathbf{x}$ :

$$\mathcal{M} = \{\mathbf{z} \in \mathbb{R}^m : \mathbf{z} = \mathbf{S}\mathbf{x} \text{ for some } \mathbf{x} \in \mathbb{R}^n\}$$

- $\mathcal{M}$  is the column space of  $\mathbf{S}$ ; completeness means  $\mathcal{M} = \mathbb{R}^m$

# Completeness

Example

Consider a market with 1 asset and 2 scenarios:

$$p = 1, \quad \mathbf{S} = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$$

- Can we achieve the payoff  $\mathbf{z} = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$ ?
- We cannot! The payoff in scenario 1 must always be the same as the payoff in scenario 2
- This market is not complete!
- The matrix does not have *full row rank*

# Complete markets: linear algebra review!

## Theorem

The following statements are equivalent:

- A market is complete:  $\mathcal{M} = \mathbb{R}^m$
- The column space of  $S$  is  $\mathbb{R}^m$
- $S$  has  $m$  linearly independent columns
- $S$  has  $m$  linearly independent rows (full row rank)
- $S$  has rank  $m$

## Number of assets

A complete market must have more assets than scenarios, i.e.  $m \leq n$ .  
Otherwise,  $S$  has rank at most  $n < m$

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The Law of One Price  
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# Law of One Price

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Theorem

*If two portfolios  $x_1$  and  $x_2$  have the same payoff in all scenarios at time 1, they must have the same price at time 0.*

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$$Sx_1 = Sx_2 \Rightarrow p^T x_1 = p^T x_2$$

Equivalently

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$$Sx = 0 \Rightarrow p^T x = 0$$

(an asset with 0 payoff in all scenarios should have price 0)

## Implications of the Law of One Price

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- $Sx = 0 \Leftrightarrow x \in \text{null}(S)$ 
  - $x$  is orthogonal to each row  $s_i$  of  $S$
  - $x \perp \text{span}(s_1, \dots, s_n)$
- $p^T x = 0 \Leftrightarrow p \perp x$ 
  - Therefore  $p \in \text{null}(S)^\perp$  (orthogonal complement)
  - $p \in \text{span}(s_1, \dots, s_n) = \text{span}(S^T)$
  - There exists some  $q \in \mathbb{R}^n$  such that  $S^T q = p$

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## Example in an incomplete market

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### Example

- $p = 1, S = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$  <https://powcoder.com>
- Law of one price holds trivially:  $S^T x = 0 \Rightarrow x = 0 \Rightarrow p^T x = 0$
- $q^T = [1 \ 0]$  or  $q^T = [0 \ 1]$  both satisfy  $q^T S = p^T$

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# Why do we care about the law of one price?

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- Consider a market where the law doesn't hold
  - There exists a portfolio  $x$  such that
    - $Sx = 0$  (0 payoff in all scenarios)
    - $p^T x = \epsilon > 0$  (nonzero price)
  - Consider the portfolio  $\alpha x$  (scale  $x$  by  $\alpha$ )
    - $S\alpha x = \alpha Sx = 0$  (still 0 payoff in all scenarios)
    - $p^T \alpha x = \alpha \epsilon$  (arbitrary price)
  - Now take an arbitrary portfolio  $y \in \mathbb{R}^n$  with payoff  $Sy = z \in \mathbb{R}^m$ , and augment it by  $\alpha x$ 
    - $S(y + \alpha x) = Sy = z$  (same payoff as  $y$ )
    - $p^T(y + \alpha x) = p^T y + \alpha \epsilon$  (arbitrary price)
  - Any portfolio can be purchased at any price!

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# Payoff pricing

## Payoff pricing

How much do we need to pay at time 0 to secure a particular payoff  $\mathbf{z}$  at time 1?

- Only defined if  $\mathbf{z} \in \mathcal{M} = \text{col}(\mathbf{S})$
- Same payoff can be achieved by multiple portfolios
- Set of prices determined by function  $q : \mathbb{R}^m \rightarrow P(\mathbb{R})$ :

$$q(\mathbf{z}) = \{\mathbf{p}^\top \mathbf{x} : \mathbf{S}\mathbf{x} = \mathbf{z}, \mathbf{x} \in \mathbb{R}^n\}$$

- In general,  $q(\mathbf{z})$  is a set
- If the Law of One Price holds:

$$\mathbf{S}\mathbf{x}_1 = \mathbf{S}\mathbf{x}_2 \Rightarrow \mathbf{p}^\top \mathbf{x}_1 = \mathbf{p}^\top \mathbf{x}_2 \Rightarrow q(\mathbf{z}) \text{ has a unique value}$$

## Payoff pricing function

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- When  $q(\mathbf{z})$  is unique for all  $\mathbf{z} \in \Delta_1$ , we call  $q$  a *payoff pricing function*.

- $q(\mathbf{z})$  unique  $\Leftrightarrow$  Law of One Price holds  $\Leftrightarrow \mathbf{p}^\top = \mathbf{q}^\top \mathbf{S}$
- Putting both together for a portfolio  $\mathbf{x}$  achieving payoff  $\mathbf{z}$ :

$$q(\mathbf{z}) = \mathbf{p}^\top \mathbf{x} = \mathbf{q}^\top \mathbf{S} \mathbf{x} = \mathbf{q}^\top \mathbf{z}$$

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### Theorem

*The Law of One Price holds in a market if and only if the payoff pricing function is linear.*

## Example in a complete market

Example

- $p = \begin{bmatrix} 1 \\ 2 \end{bmatrix}$ ,  $S = \begin{bmatrix} 1 & 1 \\ 0 & 1 \end{bmatrix}$
- Check LOP holds:  $S^T x = 0 \Rightarrow x = 0 \Rightarrow p^T x = 0$
- $q^T = \begin{bmatrix} 1 & 1 \end{bmatrix}$  satisfies  $q^T S = p^T$  so  $q(z) = q^T z$
- If we want to achieve payoff  $z = \begin{bmatrix} 3 \\ 0 \end{bmatrix}$  then we need to buy a portfolio worth

$$q^T z = \begin{bmatrix} 1 & 1 \end{bmatrix} \begin{bmatrix} 3 \\ 0 \end{bmatrix} = 3$$

## Example in an incomplete market

### Example

- $p = 1$ ,  $S = \begin{bmatrix} 1 \\ 1 \end{bmatrix}$
- Check LOP holds:  $S^T x = 0 \Rightarrow x = 0 \Rightarrow p^T x = 0$
- $q_1 = \begin{bmatrix} 1 & 0 \end{bmatrix}$  or  $q_2 = \begin{bmatrix} 0 & 1 \end{bmatrix}$  both satisfy  $q^T S = p^T$
- We cannot achieve a payoff  $z = \begin{bmatrix} 3 \\ 0 \end{bmatrix}$  at any price
- To achieve a payoff  $z = \begin{bmatrix} 2 \\ 2 \end{bmatrix}$ , we must pay

$$q_1^T z = \begin{bmatrix} 1 & 0 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \end{bmatrix} = 2 = \begin{bmatrix} 0 & 1 \end{bmatrix} \begin{bmatrix} 2 \\ 2 \end{bmatrix} = q_2^T z$$

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Arbitrage  
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# Strong arbitrage

## Definition

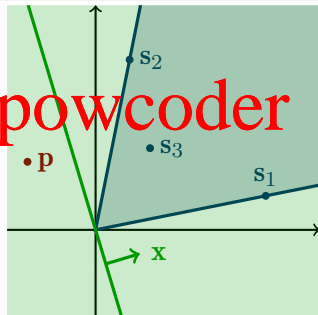
In a market with price vector  $p$  and payoff matrix  $S$ , a strong arbitrage is a portfolio  $x$  that verifies:

- $p^T x < 0$  (guaranteed gain at time 0)
- $Sx \geq 0$  (guaranteed no losses at time 1)

- $Sx \geq 0 \Leftrightarrow x^T s_i \geq 0 \forall i \in [m]$

- $p^T x < 0 \Leftrightarrow x^T p < 0$

- Equivalently,  
 $p \notin \text{cone}(s_1, \dots, s_m)$



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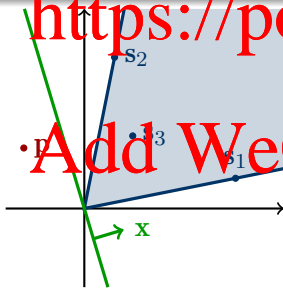


# Strong arbitrage and the Farkas lemma

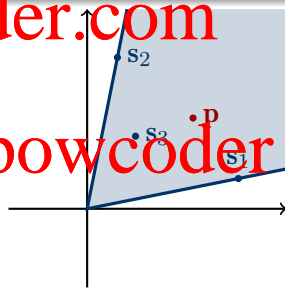
## Theorem (Farkas Lemma)

Let  $S \in \mathbb{R}^{n \times m}$  and  $p \in \mathbb{R}^n$ . Exactly one of the following is true:

1. There exists some  $x$  such that  $Sx \geq 0$  and  $p^T x < 0$
2. There exists some  $q \geq 0$  such that  $S^T q = p$



**Strong arbitrage:**  
 $p \notin \text{cone}(s_1, \dots, s_m)$



**No strong arbitrage:**  
 $p \in \text{cone}(s_1, \dots, s_m)$

# Fundamental theorem of asset pricing

- Farkas' lemma gives us two alternatives
  - Either the market has a strong arbitrage
  - Or there exists  $\mathbf{q} \geq 0$  such that  $\mathbf{S}^T \mathbf{q} = \mathbf{p}$

- Any portfolio  $\mathbf{x}$  yielding payoff  $\mathbf{S}\mathbf{x} = \mathbf{z}$  will have price

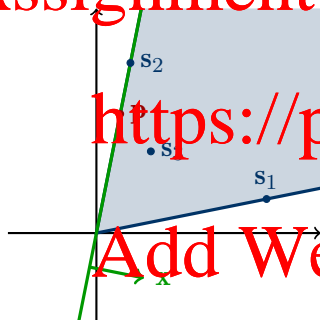
$$\mathbf{p}^T \mathbf{x} = \mathbf{q}^T \mathbf{S}\mathbf{x} = \mathbf{q}^T \mathbf{z}$$

- We have a linear pricing rule again (so the LOP holds!) but this time it is non-negative

Theorem (Fundamental theorem of asset pricing)

*A market is strong-arbitrage-free if and only if there exists a linear and non-negative payoff pricing function.*

## An edge case



- $p \in \text{cone}(s_1, \dots, s_m) \Leftrightarrow$  no strong arbitrage
- Consider the portfolio  $x$ 
  - $p^T x = 0$  (no cost at time 0)
  - $x^T s_i \geq 0 \ \forall i \in [m] \Leftrightarrow Sx \geq 0$  (no losses at time 1)
  - $s^T x > 0$  (at least one scenario has positive payoff)
- Portfolio  $x$  is called a *weak arbitrage*

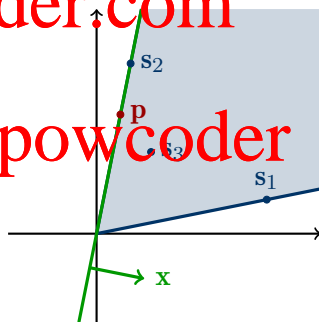
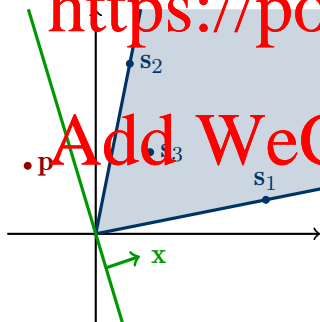
# Strong vs. weak arbitrage

## Strong arbitrage

- $p^T x < 0$
- $Sx \geq 0$

## Weak arbitrage

- $p^T x = 0$
- $Sx \geq 0$
- $\exists i \in [m], s_i^T x > 0$



## Ruling out weak arbitrage

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How to rule out weak arbitrage?

- $\mathbf{p} \in \text{cone}(\mathbf{s}_1, \dots, \mathbf{s}_m)$  rules out strong arbitrage
- To rule out weak arbitrage,  $\mathbf{p}$  cannot be on the boundary of the cone:  
$$\mathbf{p} \in \text{int}(\text{cone}(\mathbf{s}_1, \dots, \mathbf{s}_m))$$

- Equivalently, there exists  $\mathbf{q} > 0$  such that  $\mathbf{S}^T \mathbf{q} = \mathbf{p}$
- We need a *stronger* fundamental theorem of asset pricing to rule out *weak* arbitrage
- Ironical, I know!

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# Fundamental theorem of asset pricing

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### Theorem (Weak form)

*A market is strong-arbitrage-free if and only if there exists a linear and non-negative payoff pricing function.*

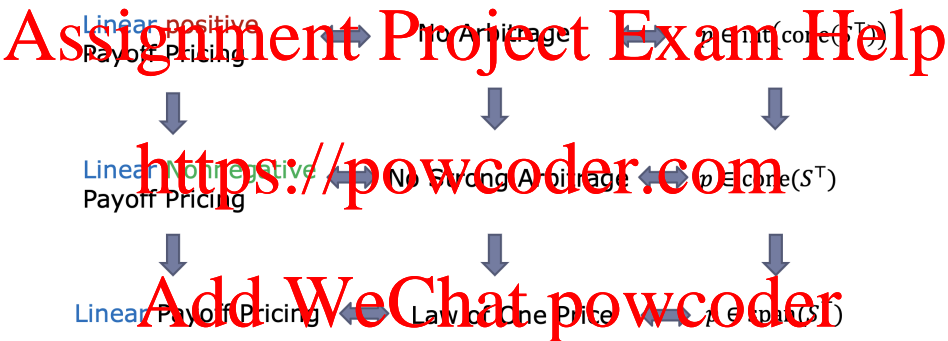
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### Theorem (Strong form)

*A market is arbitrage-free if and only if there exists a linear and positive payoff pricing function.*

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## Payoff pricing summary



Why do we care so much about linear pricing rules?

Because it gives us a way to price new assets (e.g. derivatives!)

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Risk-neutral probabilities

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# Risk-neutral probabilities

- A market is arbitrage-free if and only if there exists a positive linear pricing function  $\mathbf{q} > 0$ , i.e.

$$\mathbf{p} = \mathbf{S}^\top \mathbf{q}$$

- For each asset  $j$

$$p_j = \mathbf{S}_j^\top \mathbf{q} = \sum_{i=1}^n q_i S_{ij}$$

- Let  $\delta = \sum_{i=1}^n q_i$

$$p_j = \sum_{i=1}^n \underbrace{\frac{q_i}{\delta}}_{\tilde{q}_i \in (0,1]} \delta S_{ij} = \sum_{i=1}^n \tilde{q}_i \delta S_{ij} = \mathbb{E}_{\tilde{\mathbf{q}}}[\delta S_j],$$

where we notice that  $\tilde{\mathbf{q}}$  defines a valid probability distribution, which we call a *risk-neutral probability distribution*

# The meaning of $\delta$

- In most markets, there is at least one *risk-free asset* (e.g. US Treasury bonds; assume our first asset is risk-free with interest rate  $r$ )

- Price is  $p_1$

- Future price is  $S_1 = \begin{bmatrix} (1+r)p_1 \\ (1+r)p_1 \end{bmatrix}$  for all scenarios

- Arbitrage-free market:

$$p_1 = \sum_{i=1}^m \tilde{q}_i \delta S_{i1} = \sum_{i=1}^m \tilde{q}_i \delta (1+r) p_1 = \delta (1+r) p_1 \Rightarrow \delta = \frac{1}{1+r}$$

- $\delta$  is a discount rate! Risk-neutral pricing is thus:  $p_j = \mathbb{E}_{\tilde{q}} \left[ \frac{S_j}{1+r} \right]$

## Arbitrage-free and complete markets

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- Arbitrage-free means there exists a positive linear pricing function, equivalently a risk-neutral probability distribution
- Complete means  $S$  has rank  $m \leq n$
- If  $S$  has rank  $m$ , then so does  $S^T$ , and  $S^T q = p$  has a unique solution

Theorem (Second fundamental theorem of asset pricing)

*An arbitrage-free market is complete if and only if the risk-neutral probability distribution is unique.*

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Derivative pricing  
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# Derivatives

- Let's say we have a new payoff vector  $\mathbf{z}$  (also called a contingency claim)
- If  $\mathbf{z}$  is a deterministic function of the existing  $n$  asset payoffs, it is called a *derivative security*:

$$\mathbf{z} = f(\mathbf{S}_1, \dots, \mathbf{S}_n)$$

## Examples

Derivative	Outcome (at time 1)	Payoff
Forward contract	Must buy asset $j$ at price $K$	$\mathbf{S}_j - K$
Call option	Can buy asset $j$ at price $K$	$\max(\mathbf{S}_j - K, 0)$
Put option	Can sell asset $j$ at price $K$	$\max(K - \mathbf{S}_j, 0)$

# Put-call parity

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$$z_{\text{call}} - z_{\text{put}} = \max(S_j - K, 0) - \max(K - S_j, 0)$$

$$= S_j - K$$

$$= z_{\text{forward}}$$

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In a complete and arbitrage-free market

$$q(z_{\text{call}} - z_{\text{put}}) = q(z_{\text{call}}) - q(z_{\text{put}}) \quad (\text{unique linear pricing})$$

$$= q(z_{\text{forward}})$$

$$= \mathbb{E}_{\tilde{q}} \left[ \frac{S_j - K}{1 + r} \right] \quad (\text{risk-neutral probability})$$

$$= p_j - \frac{K}{1 + r}$$

# Arbitrage-free derivative pricing

How to price derivatives?

Given a derivative with payoff/contingency claim  $\mathbf{z}$ , how to determine the current price  $q(\mathbf{z})$ ?

Basic idea

- Start with a market  $M_1 = (\mathbf{p}, \mathbf{S})$  with  $n$  assets
- Create a new market  $M_2 = \left( \begin{bmatrix} \mathbf{p} \\ p_{n+1} \end{bmatrix}, \begin{bmatrix} \mathbf{S} & \mathbf{z} \end{bmatrix} \right)$  with  $n+1$  assets (the new contingency claim is the  $n+1$ -th asset)
- Does there always exist an arbitrage-free price  $p_{n+1}$  for  $\mathbf{z}$ ?
- Is the arbitrage-free price of  $\mathbf{z}$  unique?

# Arbitrage-free derivative pricing

## Theorem

Suppose the derivative has non-negative payoff. If  $M_1$  is arbitrage-free and complete, there exists a unique price of the derivative such that  $M_2$  is also arbitrage-free and complete. Moreover,  $M_1$  and  $M_2$  have the same risk-neutral probability distribution.

- If  $M_1$  is complete, then  $M_2$  is also complete (adding an asset/column cannot decrease the column space)
- $M_1$  is arbitrage-free and complete, so there is a unique risk-neutral probability distribution
- We can use this distribution to price  $\mathbf{z}$ , i.e.

$$q(\mathbf{z}) = \mathbb{E}_{\tilde{q}} \left[ \frac{z}{(1+r)} \right]$$



## Example

### Example

Consider the following complete arbitrage-free market.

$$\mathbf{p} = \begin{bmatrix} 1 \\ 50 \end{bmatrix}, \mathbf{S} = \begin{bmatrix} 1 & 60 \\ 1 & 20 \end{bmatrix}$$

What is the price of a put option with strike price  $K = 40$ ?

- First work out the risk-neutral probabilities (notice asset 1 is risk-free with 0 interest):

$$50 = q \cdot 60 + (1 - q) \cdot 20 \Rightarrow q = \frac{3}{4}$$

- The risk-neutral probabilities are  $\begin{bmatrix} \frac{3}{4} & \frac{1}{4} \end{bmatrix}$
- The payoff is  $\mathbf{z} = \begin{bmatrix} 0 \\ 20 \end{bmatrix}$ , so  $q(\mathbf{z}) = \frac{3}{4} \cdot 0 + \frac{1}{4} \cdot 20 = 5$ .

## Arbitrage-free pricing in incomplete markets

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- As long as the market is arbitrage-free, we can always find arbitrage-free prices for derivatives
- If the market is not complete, the price may not be unique
- This is the topic of Problem 3 on HW4!

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## Summary

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- Farkas' lemma underpins fundamental theorems in asset pricing
- Markets that obey the law of one price have linear pricing rules
- Complete markets have a unique pricing rule
- Arbitrage-free markets have positive pricing rules

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