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The No-Free-Lunch Theorem

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1 Preliminaries

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2 The No-Free-Lunch Theorem

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Reminder

- If K is event such that $P(K) = p$, $\mathbf{1}_K$ is a random variable

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$$\mathbf{1}_K = \begin{cases} 1 & \text{if } K \text{ takes place} \\ 0 & \text{otherwise.} \end{cases}$$

- If $P(K) = p$, then

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$$\mathbf{1}_K : \begin{pmatrix} 0 & 1 \\ 1-p & p \end{pmatrix}$$

and $E(\mathbf{1}_K) = p$.

- If X is a random variable

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$$X : \begin{pmatrix} x_1 & \cdots & x_n \\ p_1 & \cdots & p_n \end{pmatrix},$$

then $X = \sum_{i=1}^n x_i \mathbf{1}_{X=x_i}$, where

$$\mathbf{1}_{X=x_i} : \begin{pmatrix} 0 & 1 \\ 1-p_i & p_i \end{pmatrix}.$$

First Lemma

Lemma

Let Z be a random variable that takes values in $[0, 1]$ such that $E[Z] = \mu$. Then, for every $a \in (0, 1)$ we have

$$P(Z > 1 - a) \geq \frac{\mu - (1 - a)}{a} \text{ and } P(Z > a) \geq \frac{\mu - a}{1 - a} \geq \mu - a.$$

Proof: The random variable $Y = 1 - Z$ is non-negative with $E(Y) = 1 - E(Z) = 1 - \mu$. By Markov's inequality:

$$P(Z \leq 1 - a) = P(1 - Z \geq a) = P(Y \geq a) \leq \frac{E(Y)}{a} = \frac{1 - \mu}{a}.$$

Therefore,

$$P(Z > 1 - a) \geq 1 - \frac{1 - \mu}{a} = \frac{a + \mu - 1}{a} = \frac{\mu - (1 - a)}{a}.$$

Proof (cont'd)

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By replacing a by $1 - a$ we have:

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$$P(Z > a) \geq \frac{\mu - a}{1 - a} \geq \mu - a.$$

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Second Lemma

Lemma

Let θ be a random variable that ranges in the interval $[0, 1]$ such that $E(\theta) \geq \frac{1}{4}$. We have

$$P\left(\theta > \frac{1}{8}\right) \geq \frac{1}{7}.$$

Proof: From the second inequality of the previous lemma it follows that

$$P(\theta > a) \geq \frac{E(\theta) - a}{1 - a}.$$

By substituting $a = \frac{1}{8}$ we obtain:

$$P\left(\theta > \frac{1}{8}\right) \geq \frac{\frac{1}{4} - \frac{1}{8}}{1 - \frac{1}{8}} = \frac{1}{7}.$$

- A learning task is defined by an unknown probability distribution \mathcal{D} over $\mathcal{X} \times \mathcal{Y}$.

- The goal of the learner is to find (to learn) a hypothesis $h : \mathcal{X} \rightarrow \mathcal{Y}$ such that its risk $L_{\mathcal{D}}(h)$ is sufficiently small.

- The choice of a hypothesis class \mathcal{H} reflects some prior knowledge that the learner has about the task: a belief that a member of \mathcal{H} is a low-error model for the task.

- **Fundamental Question:** There exist a universal learner \mathcal{A} and a training set size m such that for every distribution \mathcal{D} , if \mathcal{A} receives m iid examples from \mathcal{D} , there is a high probability that \mathcal{A} will produce h with a low risk?

- The No-Free-Lunch (NFL) Theorem stipulates that a universal learner (for every distribution) does not exist!
- A learner **fails** if, upon receiving a sequence of iid examples from a distribution \mathcal{D} , its output hypothesis is likely to have a large loss (say, larger than 0.3), whereas for the same distribution there exists another learner that will output a hypothesis with a small loss.
- More precise statment: for **every** binary prediction task and learner, **there exists a distribution \mathcal{D}** for which the learning task fails.
- No learner can succeed on all learning tasks: every learner has tasks on which it fails whereas other learners succeed.

Recall 0/1 Loss Function

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The loss function is the 0/1-loss function ℓ_{0-1} :

$$\ell_{0-1}(h, (x, y)) = \begin{cases} 0 & \text{if } h(x) = y, \\ 1 & \text{if } h(x) \neq y. \end{cases}$$

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The NFL Theorem

For a learning algorithm \mathcal{A} denote by $\mathcal{A}(S)$ the hypothesis returned by the algorithm \mathcal{A} upon receiving the training sequence S .

Theorem

Let \mathcal{A} be any learning algorithm for the task of binary classification with respect to the 0/1-loss function over a domain \mathcal{X} . Let $m < \frac{|\mathcal{X}|}{2}$ be a number representing a training set size.

There exists a distribution \mathcal{D} over $\mathcal{X} \times \{0, 1\}$ such that:

- there exists a function $f: \mathcal{X} \rightarrow \{0, 1\}$ with $L_{\mathcal{D}}(f) = 0$;
- with probability at least $1/7$ over the choice of a sample S of size m there exists a hypothesis $h = \mathcal{A}(S)$ such that we have $L_{\mathcal{D}}(h) \geq 1/8$.

Interpretation of the NFL Theorem

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For every learner, there is a task for which it fails, even though the task can be successfully learned by another learner.

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Proof

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Let C be a subset of \mathcal{X} of size $2m$; this set exist because we assume that $|\mathcal{C}| > 2m$.

Intuition of the proof: any algorithm that observes only half of the instances of C has no information of what should be the labels of the other half. Therefore, there exists a target function f which would contradict the labels that $\mathcal{A}(S)$ predicts on the unobserved instances of C .

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Note that:

- There are $T = 2^{2^m}$ possible functions from C to $\{0, 1\}$: f_1, \dots, f_T .

- The set $C \times \{0, 1\}$ consists of the pairs

$$C \times \{0, 1\} = \{(x_1, 0), (x_1, 1), \dots, (x_{2^m}, 0), (x_{2^m}, 1)\}$$

For each f_i let \mathcal{D}_i be the distribution over $C \times \{0, 1\}$ given by

$$\mathcal{D}_i(\{(x, y)\}) = \begin{cases} \frac{1}{|C|} & \text{if } y = f_i(x) \\ 0 & \text{otherwise} \end{cases}$$

The probability to choose a pair (x, y) is $\frac{1}{|C|}$ if y is the true label according to f_i and 0, otherwise (if $y \neq f_i(x)$). Clearly $L_{\mathcal{D}_i}(f_i) = 0$.

Intuition

Let $m = 3$, $C = \{x_1, x_2, x_3, x_4, x_5, x_6\}$.

Suppose that

$$\begin{aligned} f_i(x_1) &= 1 & f_i(x_2) &= 0 & f_i(x_3) &= 1 \\ f_i(x_4) &= 1 & f_i(x_5) &= 1 & f_i(x_6) &= 0, \end{aligned}$$

The distribution \mathcal{D}_i is:

$(x_1, 0)$	$(x_2, 0)$	$(x_3, 0)$	$(x_4, 0)$	$(x_5, 0)$	$(x_6, 0)$
0	$\frac{1}{6}$	0	0	0	$\frac{1}{6}$
$(x_1, 1)$	$(x_2, 1)$	$(x_3, 1)$	$(x_4, 1)$	$(x_5, 1)$	$(x_6, 1)$
$\frac{1}{6}$	0	$\frac{1}{6}$	$\frac{1}{6}$	$\frac{1}{6}$	0

Clearly, we have:

$$L_{\mathcal{D}_i}(f_i) = P(\{(x, y) \mid f_i(x) \neq y\}) = 0.$$

Claim (*):

For every algorithm \mathcal{A} that receives a training set of m examples from $C \times \{0, 1\}$ and returns a function $\mathcal{A}(S) : C \rightarrow \{0, 1\}$ we have:

$$\max_{1 \leq i \leq |C|} E_{S \sim \mathcal{D}_i} (L_{\mathcal{D}_i}(\mathcal{A}(S))) \geq \frac{1}{4}$$

This means that for every \mathcal{A}' that receives a training set of m examples from $\mathcal{X} \times \{0, 1\}$ there exists $f : \mathcal{X} \rightarrow \{0, 1\}$ and a distribution \mathcal{D} over $\mathcal{X} \times \{0, 1\}$ such that $L_{\mathcal{D}}(f) = 0$ and $E_{S \sim \mathcal{D}} (L_{\mathcal{D}}(\mathcal{A}'(S))) \geq \frac{1}{4}$

Note that the index j refers to samples while i refers to hypotheses.

- There are $k = (2m)^m$ possible sequences (samples)

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S_1, \dots, S_k
of m examples from \mathcal{X} .

- If $S_j = (x_1, \dots, x_m)$, the sequence labeled by a function f_i is denoted by

$S_j^i = ((x_1, f_i(x_1)), \dots, (x_m, f_i(x_m)))$.

- If the distribution is \mathcal{D}_i , then the possible training sets that \mathcal{A} can receive are S_1^i, \dots, S_k^i and all these training sets have the same probability of being sampled. Therefore, the expected error of the sample S is:

$$E_{S \sim \mathcal{D}^m}(L_{\mathcal{D}_i}(\mathcal{A}(S))) = \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)).$$

Recall that there are $T = 2^{2^m}$ possible functions from C to $\{0, 1\}$: f_1, \dots, f_T , so $1 \leq i \leq T$, where i the superscript of S_j^i reflecting the labeling function f_i .

We have:

$$\max_{1 \leq i \leq T} \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

$$\geq \frac{1}{T} \sum_{i=1}^T \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

$$= \frac{1}{k} \sum_{j=1}^k \frac{1}{T} \sum_{i=1}^T L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

$$\geq \min_{1 \leq j \leq k} \frac{1}{T} \sum_{i=1}^T L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)).$$

Fix some j and let $S_j = \{x_1, \dots, x_m\}$. Let v_1, \dots, v_p be the examples in C that do not appear in S_j . Clearly $p \geq m$. Therefore, for each $h : C \rightarrow \{0, 1\}$ and every i we have:

$$L_{D_i}(h) = \frac{1}{2m} \sum_{x \in C} \mathbf{1}_{h(x) \neq f_i(x)} \geq \frac{1}{2m} \sum_{r=1}^p \mathbf{1}_{h(v_r) \neq f_i(v_r)} \\ \geq \frac{1}{2p} \sum_{r=1}^p \mathbf{1}_{h(v_r) \neq f_i(v_r)}.$$

Hence,

$$\frac{1}{T} \sum_{i=1}^T L_{D_i}(A(S_j^i)) \geq \frac{1}{T} \sum_{i=1}^T \frac{1}{2p} \sum_{r=1}^p \mathbf{1}_{A(S_j^i)(v_r) \neq f_i(v_r)} \\ = \frac{1}{2p} \sum_{r=1}^p \frac{1}{T} \sum_{i=1}^T \mathbf{1}_{A(S_j^i)(v_r) \neq f_i(v_r)} \\ \geq \frac{1}{2} \min_{1 \leq t \leq p} \frac{1}{T} \sum_{i=1}^T \mathbf{1}_{A(S_j^i)(v_r) \neq f_t(v_r)}.$$

Let v_r be an example in C that does not appear in a sample S_j . We can partition all functions in $\{f_1, \dots, f_T\}$ into $T/2$ disjoint sets $\{f_i, f_{i'}\}$ such that we have

$$f_i(c) \neq f_{i'}(c) \text{ if and only if } c = v_r.$$

Since for a set $\{f_i, f_{i'}\}$ we must have $S_j^i = S_j^{i'}$, it follows that

$$1_{A(S_j^i)(v_r) \neq f_i(v_r)} + 1_{A(S_j^{i'})(v_r) \neq f_{i'}(v_r)} = 1,$$

which implies

$$\frac{1}{T} \sum_{i=1}^T 1_{A(S_j^i)(v_r) \neq f_i(v_r)} = \frac{1}{2}.$$

Since

$$\frac{1}{T} \sum_{i=1}^T L_D(A(S_j^i)) \geq \frac{1}{2} \min_{1 \leq p \leq n} \frac{1}{T} \sum_{i=1}^T L_D(S_j^i)(v_i) \neq f(v_i)$$

and

$$\frac{1}{T} \sum_{i=1}^T L_D(S_j^i)(v_i) \neq f(v_i) \leq \frac{1}{2},$$

we have

$$\frac{1}{T} \sum_{i=1}^T L_D(A(S_j^i)) \geq \frac{1}{4}.$$

Thus Assignment Project Exam Help

$$\max_{1 \leq i \leq T} \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)) \geq \min_{1 \leq j \leq k} \frac{1}{T} \sum_{i=1}^T L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

implies

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$$\max_{1 \leq i \leq T} \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)) \geq \frac{1}{4}.$$

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We combined

$$\frac{1}{T} \sum_{i=1}^T L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)) \geq \frac{1}{2} \min_{1 \leq t \leq p} \frac{1}{T} \sum_{i=1}^T \mathbf{1}_{\mathcal{A}(S_j^i)(v_r) \neq f_i(v_r)}$$

$$\max_{1 \leq i \leq T} \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i)) \geq \min_{1 \leq j \leq k} \frac{1}{T} \sum_{i=1}^T L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

$$E_{S \sim \mathcal{D}^m}(L_{\mathcal{D}_i}(\mathcal{A}(S))) = \frac{1}{k} \sum_{j=1}^k L_{\mathcal{D}_i}(\mathcal{A}(S_j^i))$$

$$\frac{1}{T} \sum_{i=1}^T \mathbf{1}_{\mathcal{A}(S_j^i)(v_r) \neq f_i(v_r)} = \frac{1}{2}$$

to obtain:

$$\max_{1 \leq i \leq T} E_{S \sim \mathcal{D}_i^m}(L_{\mathcal{D}_i}(\mathcal{A}(S))) \geq \frac{1}{4}.$$

Thus, the Claim (*) is justified.

This means that for every algorithm A' that receives a training set of n examples from $\mathcal{X} \times \{0, 1\}$ there exists a function $f: \mathcal{X} \rightarrow \{0, 1\}$ and a distribution \mathcal{D} over $\mathcal{X} \times \{0, 1\}$ such that $L_{\mathcal{D}}(f) = 0$ and

$$\mathbb{E}_{S \sim \mathcal{D}^n} (L_{\mathcal{D}}(A'(S))) \geq \frac{1}{4}.$$

By the second Lemma this implies:

$$\mathbb{P} \left(L_{\mathcal{D}}(A'(S)) \geq \frac{1}{8} \right) \geq \frac{1}{7}.$$