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COMP3161/9164
Concepts of Programming Languages

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UNSW, Term 3 2020

λ -Calculus

The term language we defined for Higher Order Abstract Syntax is almost a full featured programming language.

Just enrich the syntax slightly:

$$t ::= \text{Symbol} \quad | \quad x \quad (\text{variables})$$

$$| \quad t_1 \ t_2 \quad (\text{application})$$

$$| \quad \lambda x. t \quad (\lambda\text{-abstraction})$$

There is just one rule to evaluate terms, called β reduction.

$$(\lambda x. t) \ u \ \mapsto_{\beta} \ t[x := u]$$

Just as in Haskell, $(\lambda x. t)$ denotes a **function** that, given an argument for x , will return t .

Syntax Concerns

Function application is left-associative:

$$f\ a\ b\ c \quad = \quad ((f\ a)\ b)\ c$$

λ-abstraction extends as far as possible:

$$\lambda a. f\ a\ b \quad = \quad \lambda a. (f\ a\ b)$$

All functions are unary like Haskell. Multiple argument functions are modelled with nested λ-abstractions:

$$\lambda x. \lambda y. x + y$$

β -reduction

β -reduction is a *congruence*:

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$$\frac{t \mapsto_{\beta} t'}{s \ t \mapsto_{\beta} s \ t'} \quad \frac{s \mapsto_{\beta} s'}{s \ t \mapsto_{\beta} s' \ t} \quad \frac{t \mapsto_{\beta} t'}{\lambda x. \ t \mapsto_{\beta} \lambda x. \ t'}$$

This means we can pick any reducible subexpression (called a *redex*) and perform β -reduction.

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Example:

$(\lambda x. \lambda y. f \ (y \ x)) \ 5 \ (\lambda x. x)$

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Example:

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$$(\lambda x. \lambda y. f (y x)) 5 (\lambda x. x) \mapsto_{\beta} (\lambda y. f (y 5)) (\lambda x. x)$$

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$$\begin{aligned}
 (\lambda x. \lambda y. f (y x)) 5 (\lambda x. x) &\mapsto_{\beta} (\lambda y. f (y 5)) (\lambda x. x) \\
 &\mapsto_{\beta} f ((\lambda x. x) 5)
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Confluence

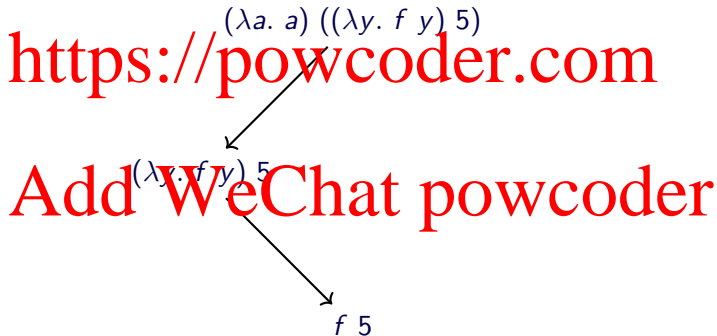
Supposing we arrive via one reduction path to an expression that cannot be reduced further (called a *normal form*), then any other reduction path will result in the same normal form.

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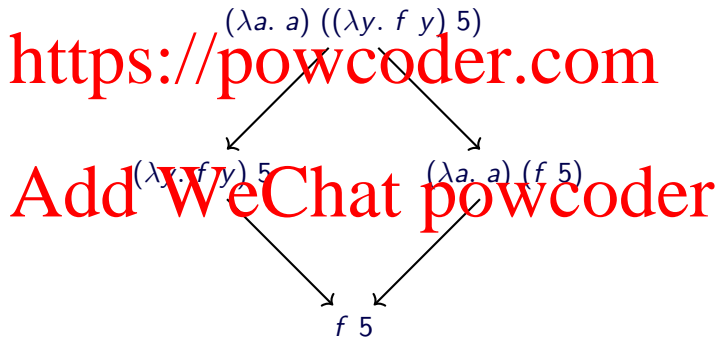
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Equivalence

Confluence means we can define another notion of *equivalence*, which equates more than α -equivalence. Two terms are $\alpha\beta$ -equivalent, written $s \equiv_{\alpha\beta} t$ if they β -reduce to α -equivalent normal forms.

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η

There is also another equation that cannot be proven from β -equivalence alone, called η -reduction:

$\lambda x. f \times \rightarrow_{\eta} f$

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Adding this reduction to the system preserves confluence and uniqueness of normal forms, so we have a notion of $\alpha\beta\eta$ -equivalence also.

Normal Forms

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Does every term in λ -calculus have a normal form?

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Normal Forms

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Does every term in λ -calculus have a normal form?

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 $(\lambda x. x\ x)(\lambda x. x\ x)$

Try to β -reduce this! (the answer is that it doesn't have a normal form)

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Why learn this stuff?

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- λ -calculus is a *Turing-complete* programming language.
- λ -calculus is the foundation for every functional programming language and some non-functional ones.
- λ -calculus is the foundation of *Higher Order Logic* and *Type Theory*, the two main foundations used for mathematics in interactive proof assistants.
- λ -calculus is the smallest example of a usable programming language, so it's good for teaching about programming languages.

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Making λ -Calculus Usable

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In order to demonstrate that λ calculus is actually a usable programming language, we will demonstrate how to encode booleans and natural numbers as λ -terms, along with their operations.

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General Idea

We transform a data type into the type of its *eliminator*. In other words, we make a function that can serve the same purpose as the data type at its use sites.

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Booleans

How do we **use** booleans?

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Booleans

How do we use booleans? To choose between two results!

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Booleans

How do we **use** booleans? **To choose between two results!**

So, a boolean will be a function that, given two arguments, returns the first one if it is true and the second one if it is false:

$\text{TRUE} \equiv \lambda a. \lambda b. a$
 $\text{FALSE} \equiv \lambda a. \lambda b. b$
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How do we write conjunction? to board

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How do we write conjunction? to board

$$\text{AND} \equiv \lambda p. \lambda q. p \ q \ p$$
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Example (Test it out!)

Try β -normalising AND TRUE FALSE .

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What about **OR**?

Natural Numbers

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So, a natural number will be a function that takes a function f and a value x , and applies the function f to x that number of times:

~~ZERO~~ $\equiv \lambda f. \lambda x. x$
~~ONE~~ $\equiv \lambda f. \lambda x. f\ x$
 TWO $\equiv \lambda f. \lambda x. f\ (f\ x)$

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How do we write Δ Sud?

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...

How do we write SUC ?

$\text{SUC} \equiv \lambda n. \lambda f. \lambda x. f\ (n\ f\ x)$

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How do we write **SUC**?

$$\text{SUC} \equiv \lambda n. \lambda f. \lambda x. f\ (n\ f\ x)$$

How do we write **ADD**?

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How do we write **SUC**?

$$\text{SUC} \equiv \lambda n. \lambda f. \lambda x. f\ (n\ f\ x)$$

How do we write **ADD**?

$$\text{ADD} \equiv \lambda m. \lambda n. \lambda f. \lambda x. m\ f\ (n\ f\ x)$$

Natural Numbers

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Natural Number Practice

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Example

Try β -normalising SUC ONE.

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Example

Try writing a different λ -term for defining SUC.

Example

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Try writing a λ -term for defining MULTIPLY.