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# Assignment Project Exam Help The basic strategy of backtracking is to systematically explore the

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- trying another extension if the extension fails,
- backing up to a smaller partial solution when the options for extending a physical solution are textual to the control of th

Input: A sequence of characters stored in A[1..n].

Output: True if A can be segmented into a sequence of words,

### Assignment Project Exam Help

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 $\begin{array}{c} \text{and } \textit{IsWord(w)} \text{ is True if } \textit{w} \text{ is a word in English, then:} \\ Add \ WeChat \ powcoder \end{array}$ 

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Add True because O'T I De Chat powcoder

BOTH-EARTH-AND-SATURN-SPIN is a valid segmentation of A

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and IsWord(w) is True if w is a word in English, then:

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Output True because

BOTH-EARTH-AND-SATURN-SPIN is a valid segmentation of *A* 

 By the way, BOT·HEART·HANDS·AT·URNS·PIN is another valid segmentation of A.

## Assignment Project Exam Help We consider a process that consumes the input characters in order

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from left to right

### Assignment Project Exam Help We consider a process that consumes the input characters in order

from left to right and produces the output words in order from left

to right:ttps://powcoder.com

### Assignment Project Exam Help We consider a process that consumes the input characters in order

We consider a process that consumes the input characters in order from left to right and produces the output words in order from left to right.

to right ttps://powcoder.com

So, at any point in time of the process we might have the following situation:

BLUE ATEN UMA ENGLAT HORNAGENERISPIN

BLUE STEM UNIT ROBOT HEARTHANDSATURNSPIN

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BLUE STEM UNIT ROBOT HEARTHANDSATURNSPIN

We need to decide on the Pert word Option Exam Help
BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN

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ASSIGNMENT ROBOT HEARTHANDSATURNSPIN

ASSIGNMENT PROBOT HE ARTHANDSATURNSPIN

BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN

BLUE STENSUNITOROBOT CHEART THANDSATURNSPIN

BLUE **STEM** UNIT **ROBOT HEARTHANDSATURNSPIN** We need to decide on the Project Exam Help

BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN BLUE ISTEMS UNIT POBOT HEART THAN PATURNSPIN ROBOT HEART **HANDSATURNSPIN BLUE** Add WeChat powcoder

**HEARTHANDSATURNSPIN BLUE STEM** UNIT **ROBOT** We need to decide on the Pext word Options: ARTHANDSATURNSPIN BLUE ISTEMS UNIT POBOT HEART THAN PATURNSPIN ROBOT HEART **HANDSATURNSPIN BLUE** BLUE STEW UNTEROEGRA THENETEN CARROS ATURNSPIN

**BLUE STEM** UNIT **ROBOT HEARTHANDSATURNSPIN** We need to decide on the Pext word Options X am Help

STEM LINIT ROBOT HE ARTHANDSATURNSPIN BLUE STEMSUNITOROBOT HEART THANPSATURNSPIN ROBOT **HEART HANDSATURNSPIN BLUE** BLUE STEW UNTEROEGRA THENRITY CANOS AFURNSPIN

Which option do we choose?

**HEARTHANDSATURNSPIN** 

ASSIGNMENT PROJECT EXAM Help
BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN

BLUE STEM UNIT ROBOT HEART THANDSATURNSPIN

BLUE STEM UNIT ROBOT HEART HANDSATURNSPIN

BLUE STEW UNTEROE OF THE PRITE CANONS AFTENSPIN

**ROBOT** 

Which option do we choose? We try all of them, using backtracking and letting the recursion fairy do the work.

**BLUE** 

**STEM** 

UNIT

**HEARTHANDSATURNSPIN** 

ASSIGNMENT PROJECT EXAM Help
BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN

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BLUE

**STEM** 

UNIT

If any of the recursive calls returns True, we return True.

**HEARTHANDSATURNSPIN** 

ASSIGNMENT PROJECT Exam Help

BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN

BLUE STEM UNIT ROBOT HEART THANBSATURNSPIN

**ROBOT** 

### BLUE STEND UNTEROEGRACHER CANDS AFTERNSPIN

HEART

Which option do we choose? We try all of them, using backtracking and letting the recursion fairy do the work.

ROBOT

BLUE

**STEM** 

UNIT

If any of the recursive calls returns True, we return True. If none do, we return False.

**HANDSATURNSPIN** 

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We start with:

SCIENT PROJECT EXAMELY PROBLEMS TO THE PROJECT OF THE PROJECT OF THE PROJECT OF THE PROBLEMS TO THE PROBLEMS T

We start with:

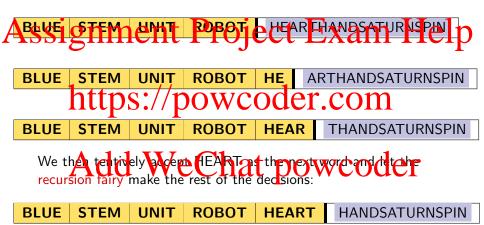
### Assignment Rosoje otar Etwasnik Holp

BLUE STEM UNIT ROBOT HE ARTHANDSATURNSPIN ARTHANDSATURNSPIN

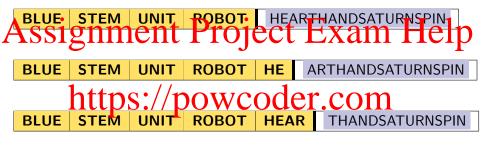
We then tentively accept HEAR as the next word and let the recursion fairy make the rest of the decisions:

BLUE ATEN UM COEDE THE DRIVE CARRISPIN

We start with:



We start with:

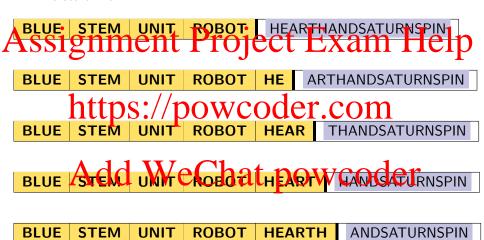


### BLUE ASTELO UM EKOEDA THERETY CANCELE TRNSPIN

We then tentively accept HEARTH as the next word and let the recursion fairy make the rest of the decisions:

BLUE STEM UNIT ROBOT HEARTH ANDSATURNSPIN

We start with:



If the recursion fairy reports success on any of these, we report success.

We start with:



If, on the other hand, the recursion fairy does not report success then we report failure.



#### Backtracking algorithm strategy, simplified



BLUES TEMS NITED BOT DE ARTHANDSATURNSPIN

So the recursive process can be simplified by ignoring the past:

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#### Backtracking algorithm strategy, simplified

# Assignment effettojeeve texaith rulelp BLUE STEM UNIT ROBOT HEARTHANDSATURNSPIN

### BLUES TENS: NUTRO BOT DE ARTHANDSATURNSPIN

So the recursive process can be simplified by ignoring the past:

### Add Weetharspoweder

Backtracking stategy: Select the first output word, and recursively segment the rest of the input string.

A recursive algorithm needs a base case. https://powcoder.com

A recursive algorithm needs a base case. https://powcoder.com
We stop when the input sequence is of length 0.

A recursive algorithm needs a base case.

\*\*This://powcoder.com\*\*

We stop when the input sequence is of length 0. What is returned then? Success or Failure?

```
Splittable(A[1 .. n]):

if n = 0

retting So n powcoder.com

if IsWord(A[1 .. i]) and Splittable(A[i + 1 .. n])

return True

return True

return True
```

Passing arrays as input parameters to recursive functions is

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Passing arrays as input parameters to recursive functions is inefficient.

It is better to treat the original input array as a global variable and reformulate the problem and the algorithm in terms of array indices.

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SSIGNMENT Project Exam Help
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https://powcoder.com
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For the string segmentation problem, the argument of any recursive call is always a suffix A[i..n] of the original input array A[1..n].

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So if we teache Wurd ay Alant ap gold Grane Can reformulate our recursive problem as follows:

Given an index i, find a segmentation of the suffix A[i..n].

Note that function IsWord(i, j) now takes two indices of array A.

For any sequence S,

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#### Longest Increasing Subsequence

For any sequence S, a subsequence of S is another sequence obtained from S by deleting zero or more elements, without changing the order of the remaining elements.

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## Longest Increasing Subsequence

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Problem: Given a sequence of integers S find the Longest Increasing Space of the Communication of the Communicatio

#### Longest Increasing Subsequence

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changing the order of the remaining elements.

ASSIGNMENT Project Exam Help

Problem: Given a sequence of integers S find the Longest Increasing Subsquencial WEOGER.COM

Input: Integer array A[1..n].

Output: Longest possible sequence of indices

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Example: See above.

We use the same approach as we used for Subset Sum:

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We decide, for each index in order from left to right, whether or Assignation the respect Exam Help

https://powcoder.com

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 $\underset{\text{Do we include } A[7] = \frac{1}{2}?}{\text{https://powcoder.com}}$ 

We use the same approach as we used for Subset Sum:

We decide, for each index in order from left to right, whether or the left to right, whether or left to right to right.

https://powcoder.com
Do we include A[7] = 2? No, we can't.

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2 6 1 4 2 4 2 9 5 3 5 7 8 3

https://powcoder.com
Do we include A[8] = 9? We can, but should we?

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Instead of trying to be smart, we apply the backtracking strategy:

• We for the decisions.

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- Warfict Centalive Circlude Cand Detrews Cooking Imake the rest of the decisions.
- We then tentatively exclude 9, and let the recursion fairy make the rest of the decisions.

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https://powcoder.com
Do we include A[8] = 9? We can, but should we?

Instead of trying to be smart, we apply the backtracking strategy:

- We first Unitarities Circlade Cand Dethew Cookie Imake the rest of the decisions.
- We then tentatively exclude 9, and let the recursion fairy make the rest of the decisions.

Whichever choice leads to a longer increasing subsequence is the right one.

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Assignment Project Exam Help
In other words, what do we need to remember about our past
decisions?

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The distribution of the last is the partial continumber selected so far:

Assignment Project Exam Help In other words, what do we need to remember about our past decisions?

The day this compared to the selected so far:

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The problem can be restated as:

Given an integer *prev* and array A[1..n], find the Longest Increasing Subsequence of A in which every element is larger than *prev*.

```
LISbigger(prev, A[1 .. n]):

if hether s.//powcoder.com

else if A[1] \le prev

return LISbigger(prev, A[2 .. n])

else dd WeChat powcoder

skip \(
-\ LISbigger(prev, A[2 .. n]) + 1

return max{skip, take}
```

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We will restate the algorithm in terms of array indices, assuming that the array A[1..n]/is a global variable the array array

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The integer prev is typically an array element A[i], and the remaining array is always a suffix A[j..n] of the original input array.

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The integer pray is typically an array element A[i] and the

The integer prev is typically an array element A[i], and the remaining array is always a suffix A[j..n] of the original input array.

## The padd best better powcoder

Given two indices i and j, where i < j, find the Longest Increasing Subsequence of A[j..n] in which every element is larger than A[i].

```
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return 0

else if A[i] > A[j]

return Stieger(i, j + 1)

take delisting for the provider of the
```

# Assignment Project Exam Help return 0 else if A[i] > A[j] netups bigg powcoder.com else skip LISbigger(i, j + 1) take delse return 0 place to the coder. take delse coder. return 0 place to the coder. place to the coder. place to the coder.

What is the initial call on A[1..n]? There is not index i initially...

```
LIS(A[1, . n])://powcoder.com
return LISbigger(0, 1)
```

The running time for a successful search in a binary search tree is proportional to the depth of the target node.

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• The control case ting in top to with the tree.

The running time for a successful search in a binary search tree is proportional to the depth of the target node.

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- The control case terms time is proportion with the tree.
- To minimize the worst-case search time, the tree should be balanced.

It is often more important to minimize the total cost of several searches

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It is often more important to minimize the total cost of several searches

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If x is a more frequent search target than y, we can save time by building a tree where the depth of x is smaller than the depth of y

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• Even if that means increasing the overall depth of the tree.

It is often more important to minimize the total cost of several searches

## Assignment Project Exam Help



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If x is a more frequent search target than y, we can save time by building a tree where the depth of x is smaller than the depth of y

- Even if that means increasing the overall depth of the tree.
- A perfectly balanced tree is not the best choice if some items are significantly more popular than others.

This state of affairs suggests the following problem:

Given a sorted array of keys A[1,i] and an array of corresponding p search tree to store the p keys and that minimizes the total search time, assuming that there will be exactly f[i] searches for each key A[i].

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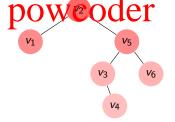
It is helpful D write a good recursive definition of the function we are trying to optimize.

This state of affairs suggests the following problem:

Given a sorted array of keys A[1.4] and an array of corresponding access frequencies I[1.n], build the binary search tree to store the keys and that minimizes the total search time, assuming that there will be exactly f[i] searches for each key A[i].

It is helpful por recursive definition of the function we are trying to optimize.

Given a Archer Wee With at poweoder n nodes, let  $v_1, v_2, \ldots, v_n$  be the nodes of T, indexed in sorted order, so that each node *v<sub>i</sub>* stores the corresponding key A[i].



#### Optimization function

The total cost of performing all the binary searches is given by the following expression:

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where  $d_T(v_i)$  is the depth of  $v_i$  in T (where depth of root is 1). https://powcoder.com

#### Optimization function

The total cost of performing all the binary searches is given by the following expression:

# Assignment, Project Exam Help

where  $d_T(v_i)$  is the depth of  $v_i$  in T (where depth of root is 1). Let  $v_r$ , where  $1 \le r \le p$  be the root of r and let  $\ell(T)$  and r(T) be the left and right subtrees of  $v_r$ .

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r(T)

 $\ell(T)$ 



Then:

$$\begin{array}{l} \mathbf{Assignment}^{\textit{Cost}(T, f[1..n])} = \sum\limits_{i=1}^{n} f[i] \times d_{T}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=r+1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=r+1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=r+1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=r+1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=r+1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{r(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) \\ = \sum\limits_{i=1}^{n} f[i] \times d_{\ell(T)}(v_{i}) + \sum\limits_{i=1}^{n} f[i]$$

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r(T)



$$\underset{\sum}{\textbf{Assignment Project Exam}} \underbrace{\textbf{Help}}_{\text{Nost}(T), f[1..r-1]) + \textit{Cost}(r(T), f[r+1\dot{n}])}$$

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$$\underbrace{ \text{Assignment Project Exam}}_{\sum\limits_{i=1}^{r} f[i] + \textit{Cost}(\ell(T), f[1...r-1]) + \textit{Cost}(r(T), f[r+1\dot{n}]) }_{\text{Help}}$$

Let vhttps://phtw/2000 demton constrate

$$\underset{i=1}{\textbf{Assignment Project Exam Help}} \\ \sum_{i=1}^{Cost} f[i] + Cost(\ell(T), f[1...r-1]) + Cost(r(T), f[r+1n])$$

Let v the left subtree  $\ell(T_{opt})$  must be the optimal search tree for

• the left subtree  $\ell(\mathcal{T}_{\text{opt}})$  must be the optimal search tree for the keys A[1..r-1] and access frequencies f[1..r-1] Add WeChat powcoder

$$\underset{i=1}{\text{Assignment Project Exam Help}} \\ \sum_{i=1}^{\text{Cost}(\mathcal{T}, f[1..n])} \bar{\mathbb{E}}_{\text{nt Project Exam Help}} \\ + \sum_{i=1}^{\text{Cost}(\mathcal{T}, f[1..n-1])} + Cost(r(\mathcal{T}), f[r+1n]) \\ + Cos$$

Let v the left subtree  $\ell(T_{opt})$  must be the optimal search tree for

- the keys A[1..r-1] and access frequencies f[1..r-1]
- the regulative recent between the keys A[r+1..n] and access frequencies f[r+1..n].

$$\underbrace{ \text{Assign}^{Cost}(\mathcal{T}, f[1..n])}_{\sum_{i=1}^{r} f[i] + Cost(\ell(\mathcal{T}), f[1..r-1]) + Cost(r(\mathcal{T}), f[r+1n]) }_{\text{Help}}$$

Let v the left subtree  $\ell(T_{opt})$  must be the optimal search tree for

- the keys A[1..r-1] and access frequencies f[1..r-1]
- the resolution of the the the two the for the keys A[r+1..n] and access frequencies f[r+1..n].

So, if we choose the correct key to store at the root, the recursion fairy will construct the rest of the optimal tree.

$$\underset{\sum_{i=1}}{\textbf{Assignment Project Exam}} \underbrace{\textbf{Exam Help}}_{\text{Cost}(\ell(T), f[1...r-1]) + \textit{Cost}(r(T), f[r+1\dot{n}])}$$

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So how do we choose the correct key?



$$\underset{i=1}{\text{Assignment Project Exam Help}} \\ \sum_{i=1}^{\text{Cost}(\mathcal{T}, f[1..n])} \bar{\mathbb{E}}_{\text{nt Project Exam Help}} \\ + \sum_{i=1}^{\text{Cost}(\mathcal{T}, f[1..n-1])} + Cost(r(\mathcal{T}), f[r+1n]) \\ + Cos$$

# Let v the left subtree $\ell(T_{opt})$ must be the optimal search tree for

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So, if we choose the correct key to store at the root, the recursion fairy will construct the rest of the optimal tree.

So how do we choose the correct key? Try each one!

