

# MATH3411 INFORMATION, CODES & CIPHERS

## Test 1      2016 S2      SOLUTIONS

### Version A

Multiple Choice: d, a, e, c, d, b, e, b, e, c

1. (d): One or two errors.
2. (a): Test the words with  $x_1 = 1, x_2 = 0$  to see whether they are codewords.
3. (e): There are  $3^4 = 81$  codewords: the 81 linear combinations of the rows of  $G$ .
4. (c): There is just one codeword starting with 1021, namely the sum of rows 1, 3 (twice) and 4 of  $G$ .
5. (d): By the Sphere-Packing Theorem with  $t = 1$ ,  $2^k \times (1 + 7) \leq 2^7$ , so  $k \leq 4$ , and indeed, the binary Hamming  $(7, 4)$  code has these parametres.
6. (b):
7. (e): None since the message 11 cannot be decoded unambiguously.
8. (b): The Kraft-McMillan number  $K = \sum r^{-r}$  must be at most 1 for UD codes. Testing values of  $r = 2, 3, \dots$  gives us that  $r = 3$  is the minimum radix that satisfies this. (You could also draw a decision tree.)
9. (e): Draw a decision tree in the standard way.
10. (c): One dummy symbol is needed.
11. (a) The Kraft-McMillan number is

$$K = \frac{1}{2} + \frac{1}{2^2} + \frac{1}{2^3} + \frac{3}{2^4} = \frac{17}{16} > 1$$

so there is no UD-code.

- (b) We find that  $s_1 s_1 \mapsto 0$ ,  $s_1 s_2 \mapsto 11$ ,  $s_2 s_1 \mapsto 100$ ,  $s_2 s_2 \mapsto 101$ .  
The average length per original source symbol is  
 $\frac{1}{2} \left( \frac{5}{25} + \frac{9}{25} + 1 \right) = \frac{39}{50}$  by Knuth's Lemma.

## Version B

Multiple choice: b, d/e, c, a, b, b, e, c, b, c

1. (b):
2. (d/e): None of the three words containing 10 (in correct order) are codewords. However, you could argue that we could be allowed to specify a different order for the message digits, in which case, (d) is valid.
3. (c): There are  $2^4 = 16$  codewords: the 16 linear combinations of the rows of  $G$ .
4. (a): There is just one codeword starting with 1011, namely the sum of the first and last rows  $G$ .
5. (b): By the Sphere-Packing Theorem with  $t = 2$ ,  $2^k \times (1 + 8 + 36) \leq 2^8$ , so  $k \leq 2$ , and indeed, the code with basis  $\{11111000, 00000111\}$  has these parameters.
6. (b):
7. (e):
8. (c): The Kraft-McMillan number  $K = \sum \frac{1}{r^{\ell_i}}$  must be at most 1 for UD codes. Testing values of  $r = 2, 3, \dots$  gives us that  $r = 4$  is the minimum radix that satisfies this. (You could also draw a decision tree.)
9. (b): The codewords are 00, 01, 100, 101, 1100, 1101.
10. (c): One dummy symbol is needed.
11. (a) The Kraft-McMillan number is

$$K = \frac{2}{3} + \frac{3}{3^2} + \frac{1}{3^3} = \frac{28}{27} > 1$$

so there is no UD-code.

- (b) We find that  $s_1 s_1 \mapsto 0$ ,  $s_1 s_2 \mapsto 11$ ,  $s_2 s_1 \mapsto 100$ ,  $s_2 s_2 \mapsto 101$ .

The average length per original source symbol is

$$\frac{1}{2} \left( \frac{6}{36} + \frac{11}{36} + 1 \right) = \frac{53}{72} \text{ by Knuth's Lemma.}$$