Searching and Sorting Part One

Outline for Today

- Gauss's Sum
 - A famous, ubiquitous sum.
- Sorting Algorithms
 - How quickly can we get things in order?
- Inventing an Algorithm
 - Building better sorts with big-O.

Recap from Last Time

```
double averageOf(const Vector<int>& vec) {
  double total = 0.0;
  for (int i = 0; i < vec.size(); i++) {</pre>
      total += vec[i];
  return total / vec.size();
```

Assume any individual statement takes one unit of time to execute. If the input Vector has *n* elements, how many time units will this code take to run?

```
double averageOf(const Vector<int>& vec) {
1 double total = 0.0;
                       n+1
  for (int i = 0; i < vec.size(); i++) {</pre>
      total += vec[i];
  return total / vec.size(); 1
```

Assume any individual statement takes one unit of time to execute. If the input Vector has *n* elements, how many time units will this code take to run?

```
double averageOf(const Vector<int>& vec) {
1 double total = 0.0;
                       n+1
  for (int i = 0; i < vec.size(); i++) {</pre>
      total += vec[i];
  return total / vec.size(); 1
```

One possible answer: 3n + 4.

```
double averageOf(const Vector<int>& vec) {
1 double total = 0.0;
                       n+1
  for (int i = 0; i < vec.size(); i++) {</pre>
      total += vec[i];
  return total / vec.size(); 1
```

One possible answer: 3n + 4. More useful answer: O(n).

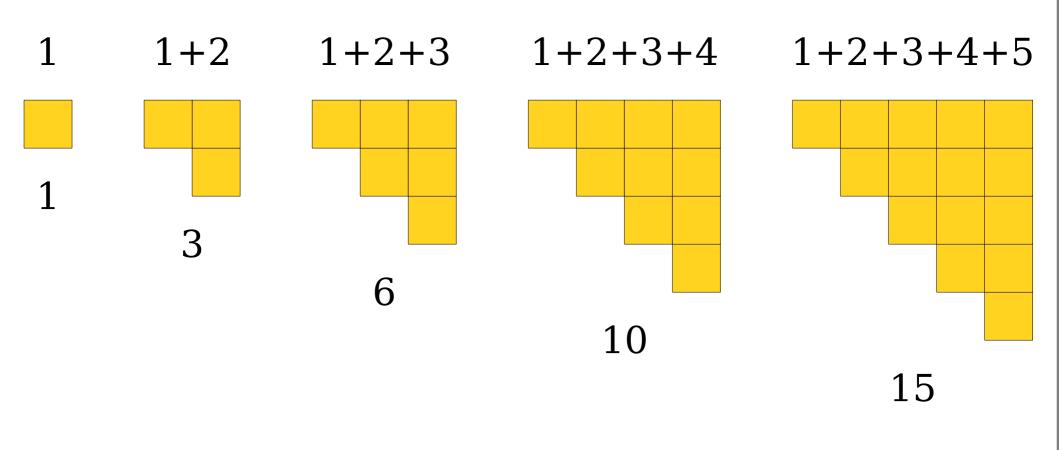
```
void printStars(int n) {
    for (int i = 0; i < n; i++) {
        for (int j = 0; j < n; j++) {
            cout << '*' << endl;
        }
    }
}</pre>
```

Work Done: $O(n^2)$.

New Stuff!

Gauss's Sum

Gauss's Sum



How does the sum 1 + 2 + 3 + ... + n scale as n increases?

n	1+2++ <i>n</i>
10	55
20	210
30	465
40	820
50	1275
60	1830
70	2485
80	3240
90	4095
100	5050

Which best describes the rate at which the quantity 1 + 2 + ... + n grows as a function of n?

A. O(n)

B. $O(n^2)$

C. $O(n^3)$

Formulate a hypothesis!

How does the sum 1 + 2 + 3 + ... + n scale as n increases?

n	1+2++n
10	55
20	210
30	465
40	820
50	1275
60	1830
70	2485
80	3240
90	4095
100	5050

Which best describes the rate at which the quantity 1 + 2 + ... + n grows as a function of n?

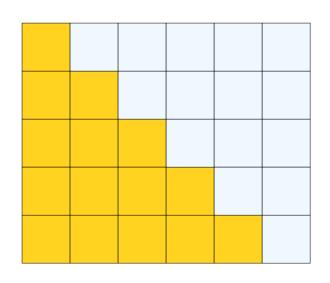
A. O(n)

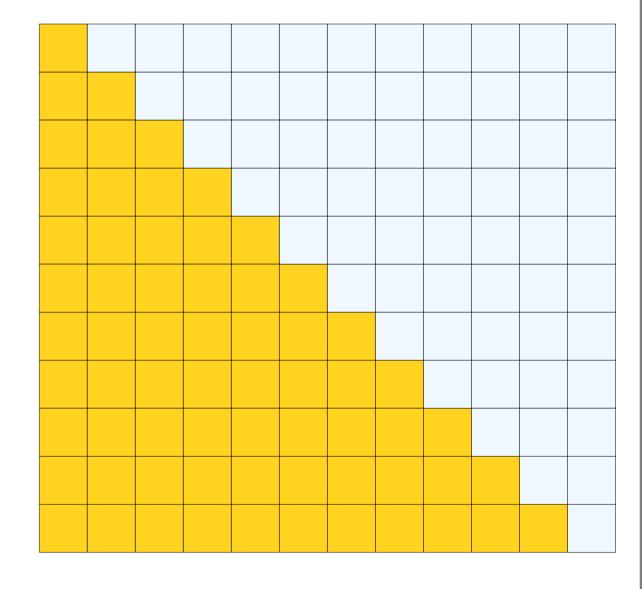
B. $O(n^2)$

C. $O(n^3)$

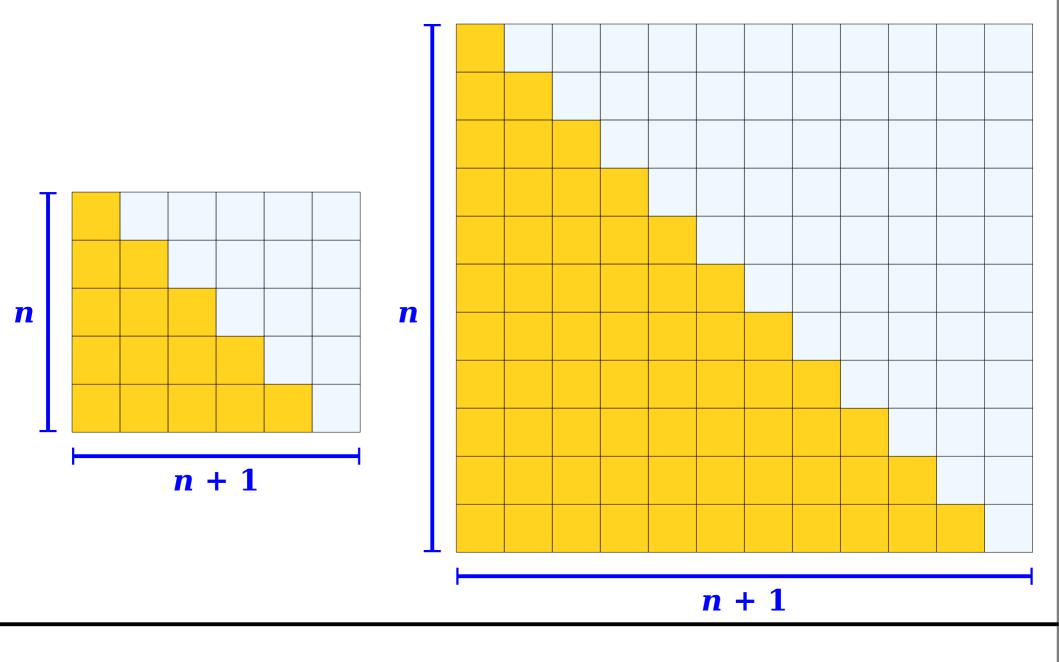
Discuss with your neighbor!

How does the sum 1 + 2 + 3 + ... + n scale as n increases?

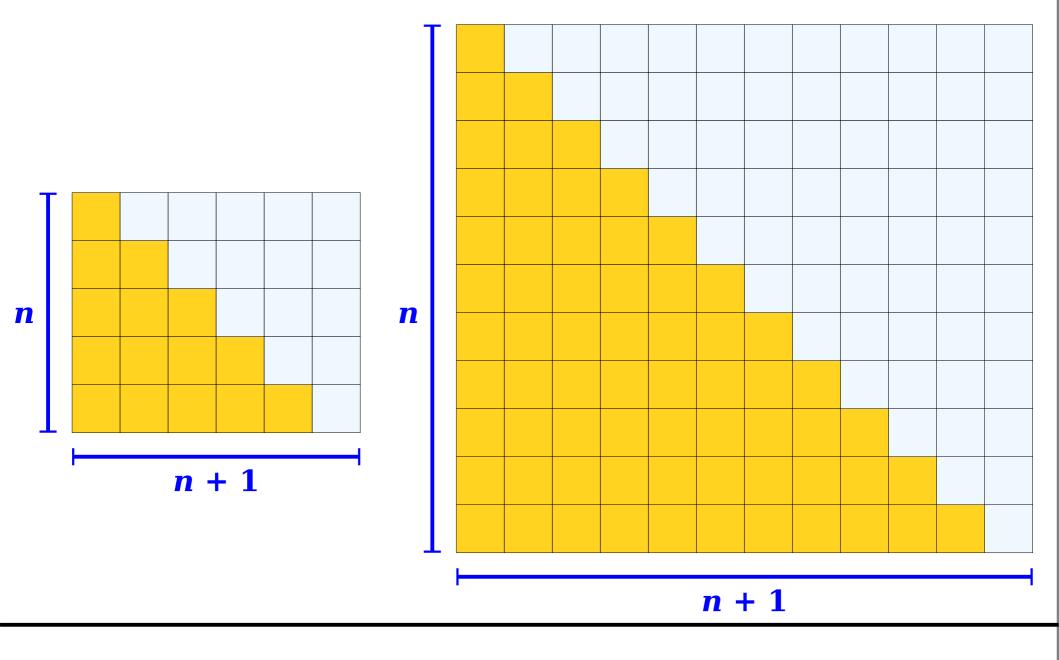




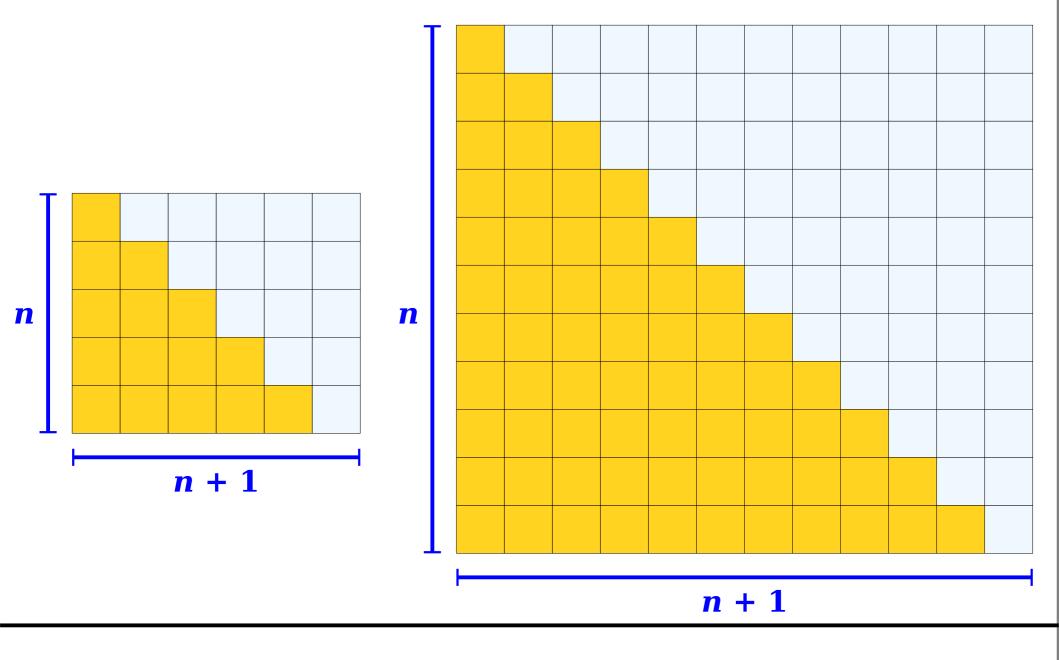
How does the sum 1 + 2 + 3 + ... + n scale as n increases?



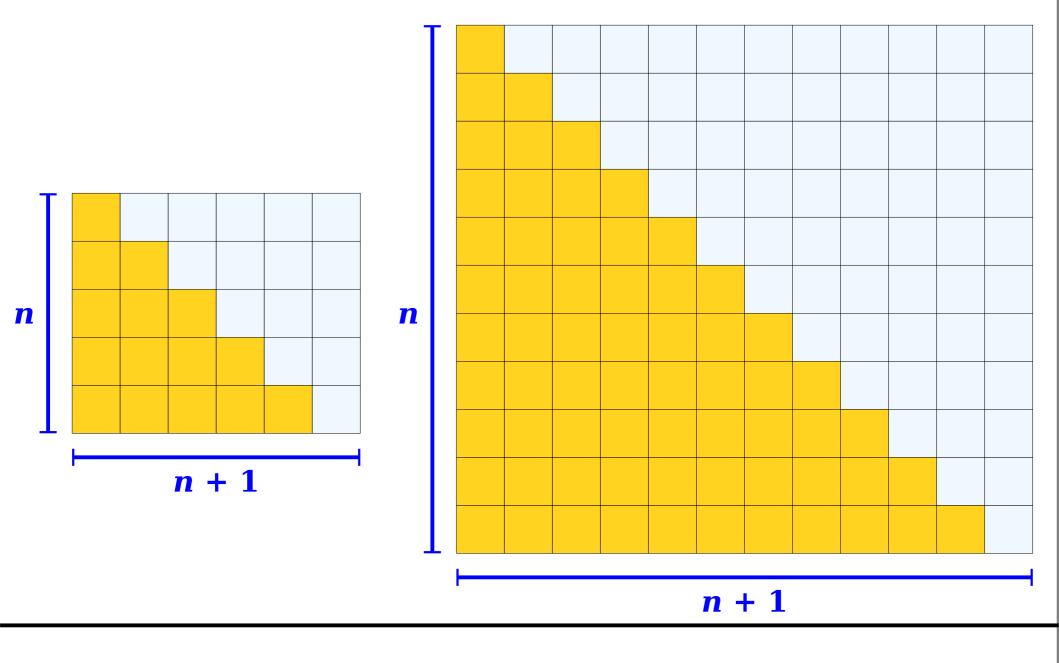
Each figure has area $n(n + 1) = n^2 + n$.



Half that area is the gold figure, which is 1 + 2 + 3 + ... + n.



Big-O ignores leading coefficients and low order-terms.



So
$$1 + 2 + 3 + ... + n = O(n^2)$$
.

Sorting Algorithms

What is sorting?



Time	Auto	Athlete	Nationality	Date	Venue
4:37.0		Anne Smith	United Kingdom	3 June 1967 ^[8]	London
4:36.8		Maria Gommers	Netherlands	14 June 1969 ^[8]	Leicester
4:35.3		Ellen Tittel	West Germany	20 August 1971 ^[8]	Sittard
4:29.5		Paola Pigni	■ Italy	8 August 1973 ^[8]	Viareggio
4:23.8		Natalia Mărășescu	Romania	21 May 1977 ^[8]	Bucharest
4:22.1	4:22.09	Natalia Mărășescu	Romania	27 January 1979 ^[8]	Auckland
4:21.7	4:21.68	Mary Decker	United States	26 January 1980 ^[8]	Auckland
	4:20.89	Lyudmila Veselkova	Soviet Union	12 September 1981 ^[8]	Bologna
	4:18.08	Mary Decker-Tabb	United States	9 July 1982 ^[8]	Paris
	4:17.44	Maricica Puică	Romania	9 September 1982 ^[8]	Rieti
	4:16.71	Mary Decker-Slaney	United States	21 August 1985 ^[8]	Zürich
	4:15.61	Paula Ivan	Romania	10 July 1989 ^[8]	Nice
	4:12.56	Svetlana Masterkova	Russia	14 August 1996 ^[8]	Zürich
	4:12.33	Sifan Hassan	Netherlands	12 July 2019	Monaco

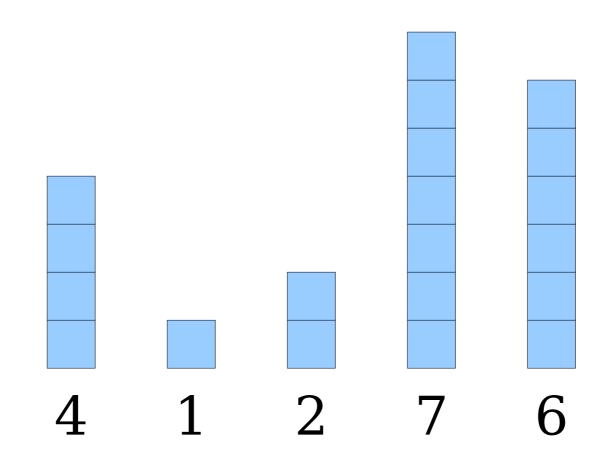
Problem: Given a list of data points, sort those data points into ascending / descending order by some quantity.

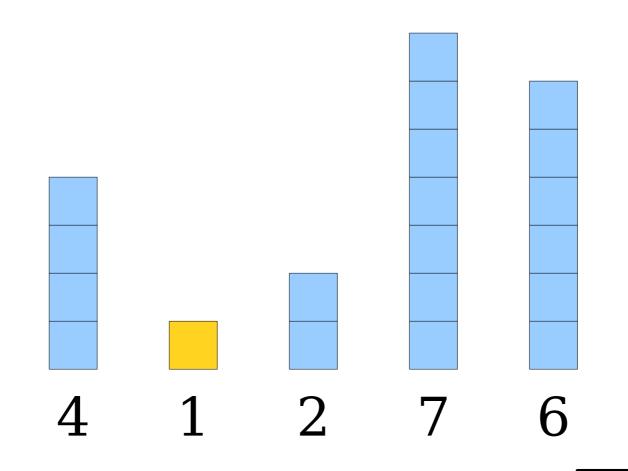
Suppose we want to rearrange a sequence to put elements into ascending order.

What are some strategies we could use?

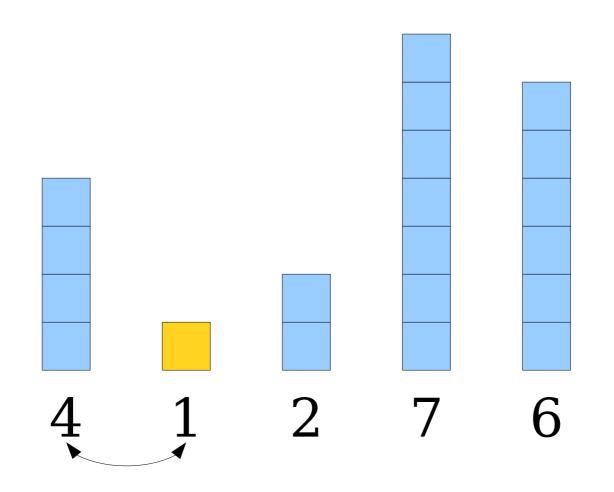
How do those strategies compare?

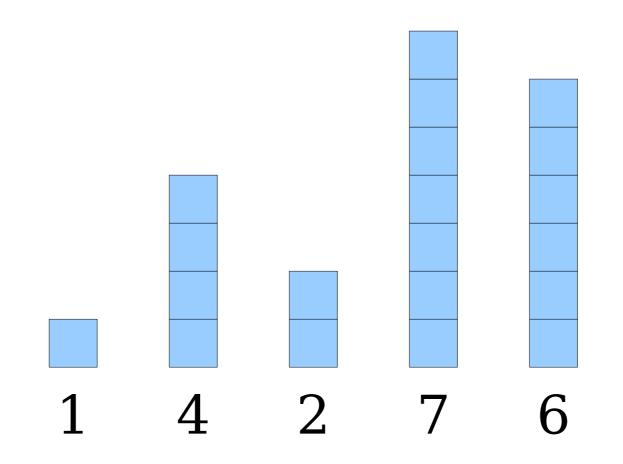
Is there a "best" strategy?

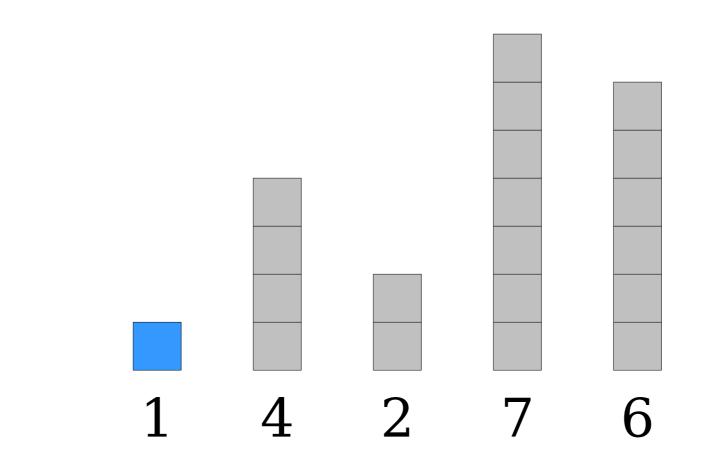




The smallest element should go in front.

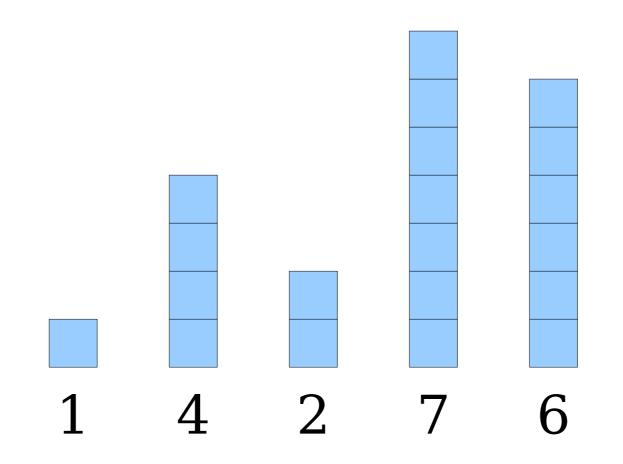


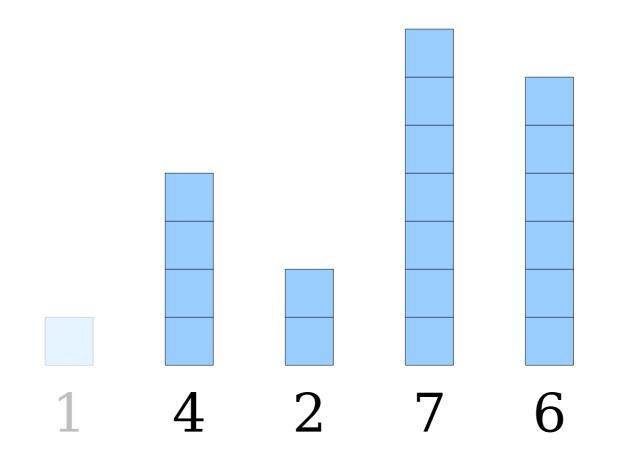


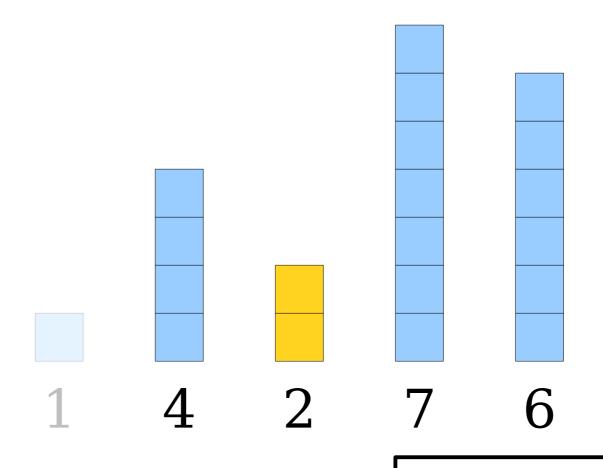


This element is in the right place now.

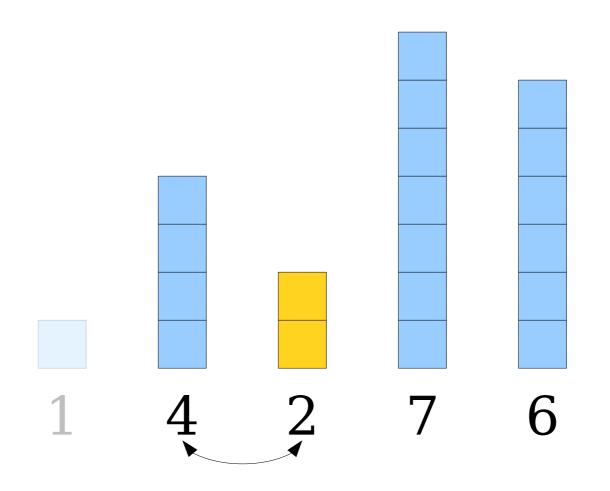
The remaining elements are in no particular order.

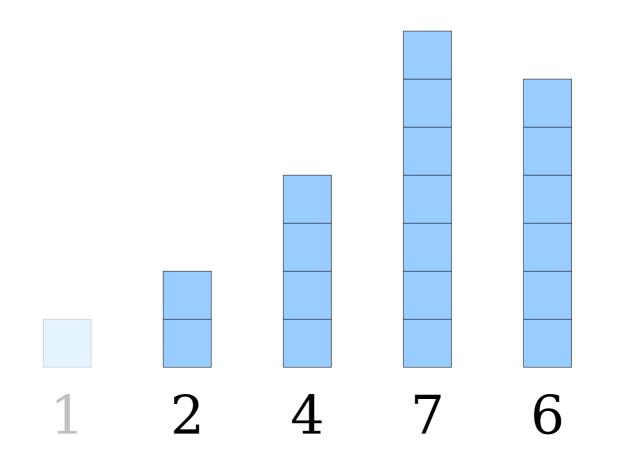


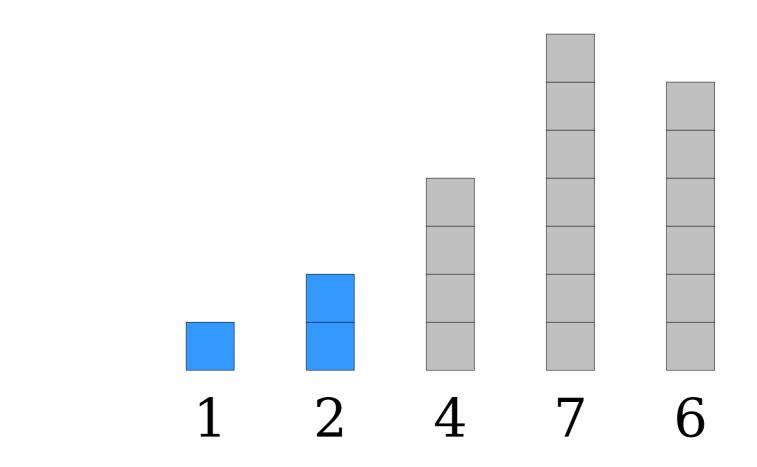




The smallest element of the remaining elements goes at the front of the remaining elements.

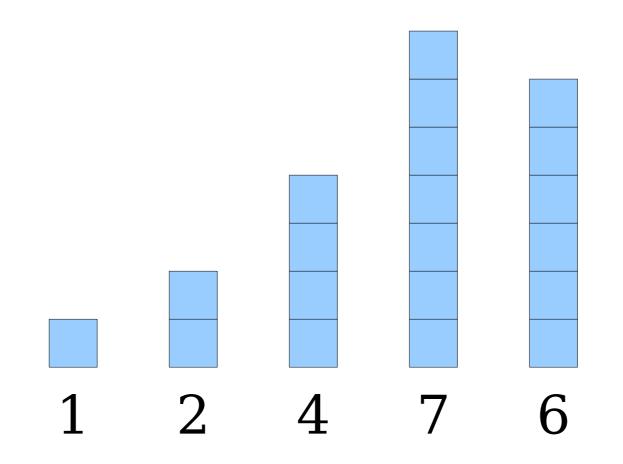


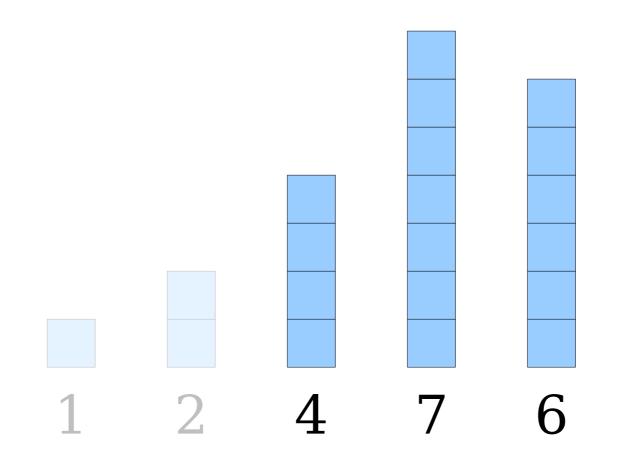


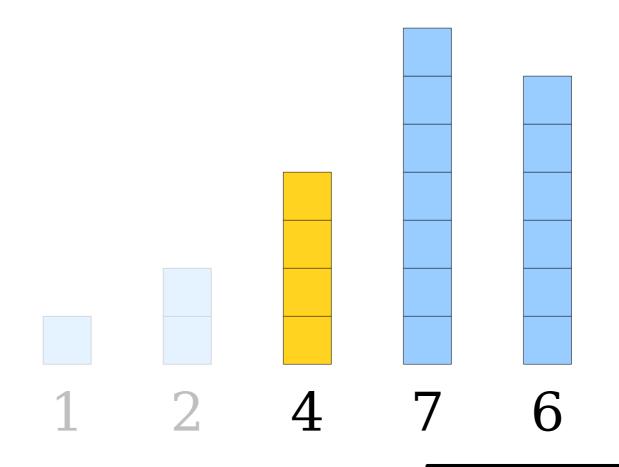


These elements are in the right place now.

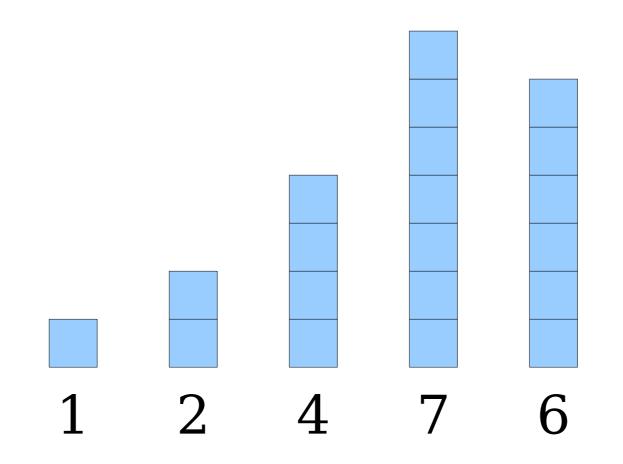
The remaining elements are in no particular order.

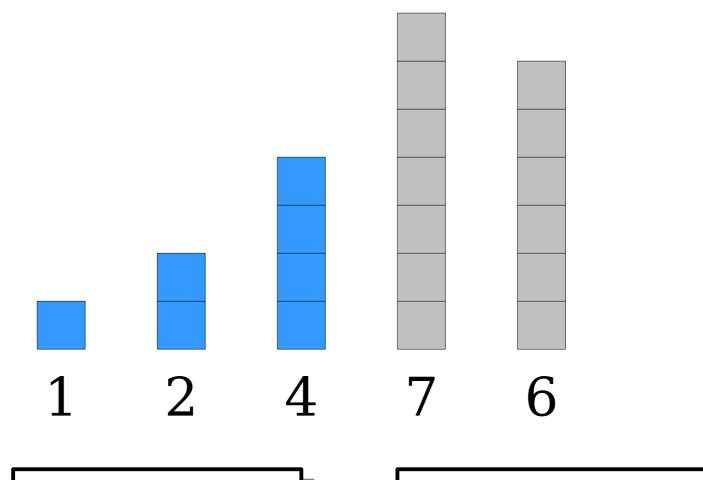






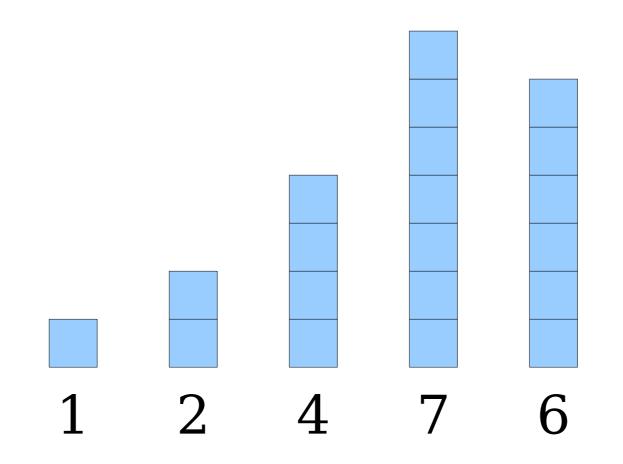
The smallest of these remaining elements goes at the front of the remaining elements.

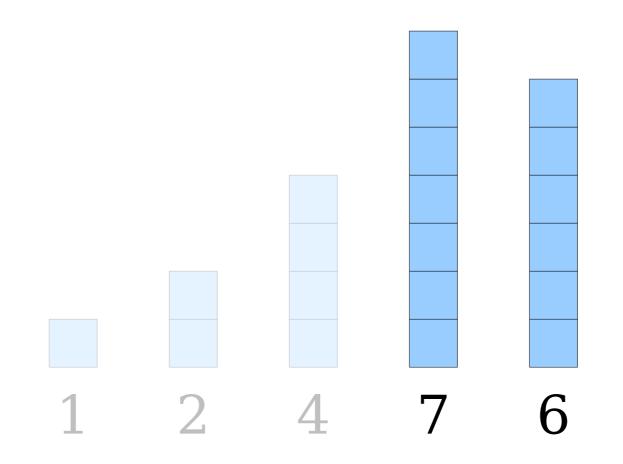


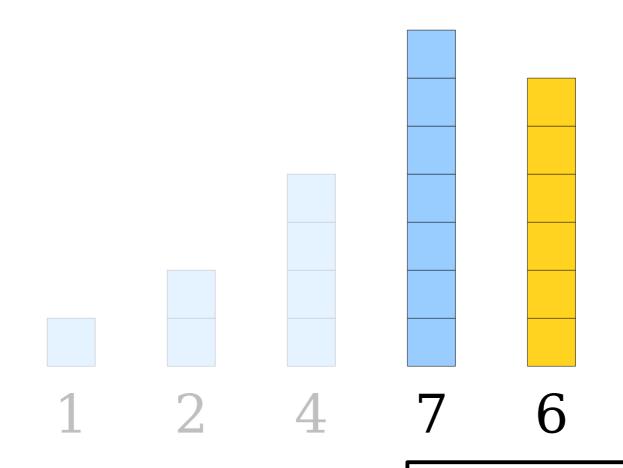


These elements are in the right place now.

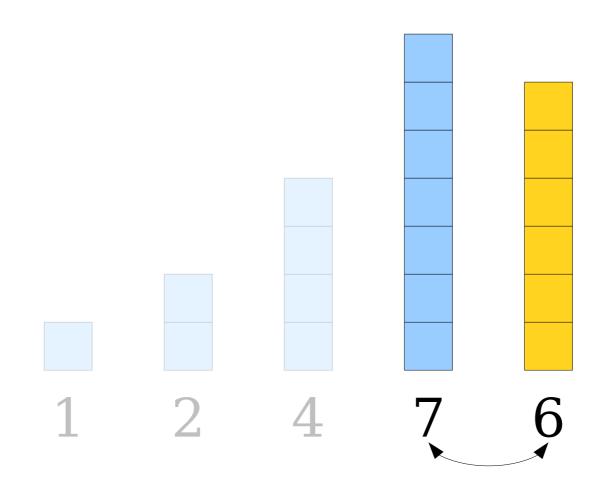
The remaining elements are in no particular order.

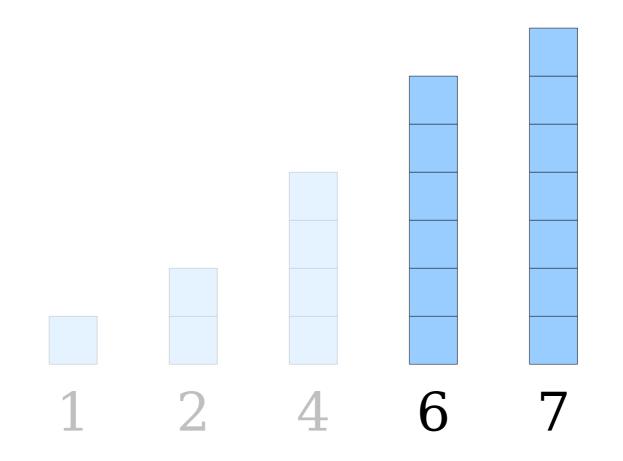


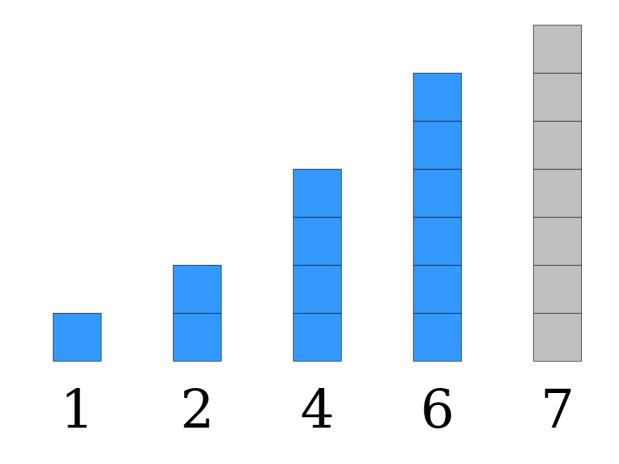




The smallest of these elements needs to go at the front of this group of elements.

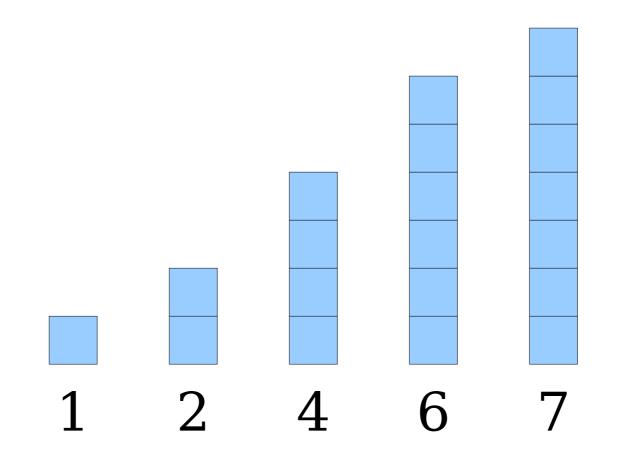


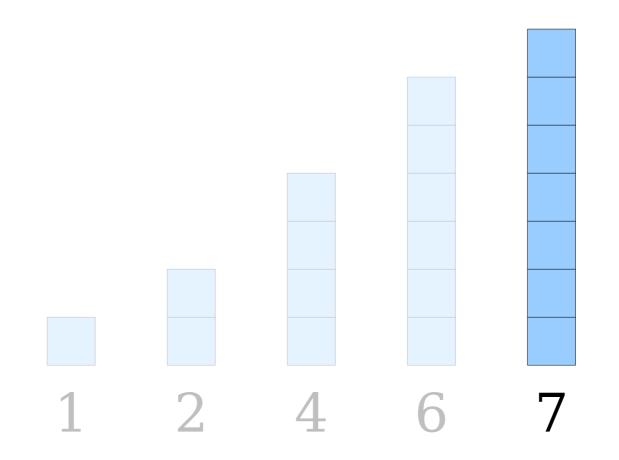


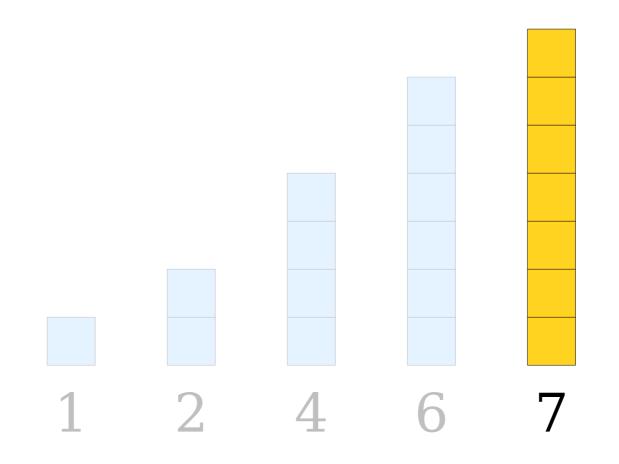


These elements are in the right place now.

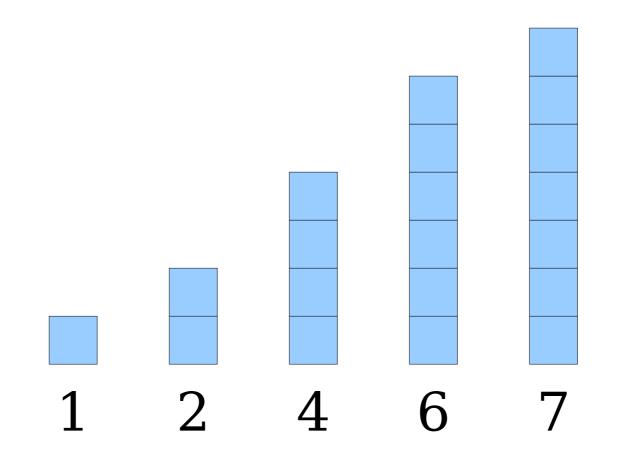
The remaining elements are in no particular order.

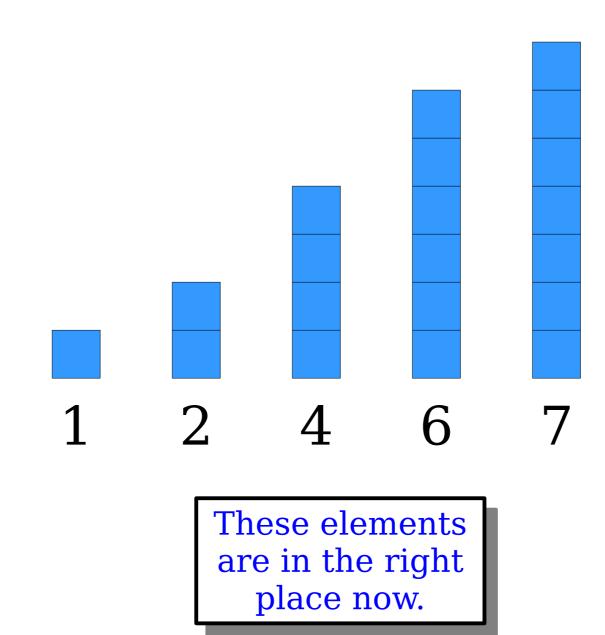






The smallest element from this group needs to go at the front of the group.





Selection Sort

- Find the smallest element and move it to the first position.
- Find the smallest element of what's left and move it to the second position.
- Find the smallest element of what's left and move it to the third position.
- Find the smallest element of what's left and move it to the fourth position.
- (etc.)

```
/**
 * Sorts the specified vector using the selection sort algorithm.
void selectionSort(Vector<int>& elems) {
  for (int index = 0; index < elems.size(); index++) {</pre>
    int smallestIndex = indexOfSmallest(elems, index);
    swap(elems[index], elems[smallestIndex]);
/**
 * Given a vector and a starting point, returns the index of the
 * smallest element in that vector at or after the starting point.
int indexOfSmallest(const Vector<int>& elems, int startPoint) {
  int smallestIndex = startPoint;
  for (int i = startPoint + 1; i < elems.size(); i++) {</pre>
    if (elems[i] < elems[smallestIndex]) {</pre>
      smallestIndex = i;
  return smallestIndex;
```

{ 46, 69, 20, 16, 09, 10, 29, 90, 67, 18, 53, 20, 38, 20, 46 }

 $\{46, 69, 20, 16, 09, 10, 29, 90, 67, 18, 53, 20, 38, 20, 46\}$

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 2 requires us to scan n-2 elements.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 2 requires us to scan n-2 elements.

. . .

Number of elements scanned:

$$n + (n-1) + (n-2) + ... + 2 + 1$$
.

Finding the element that goes in position 0 requires us to scan all n elements.

Finding the element that goes in position 1 requires us to scan n-1 elements.

Finding the element that goes in position 2 requires us to scan n-2 elements.

. . .

Number of elements scanned:

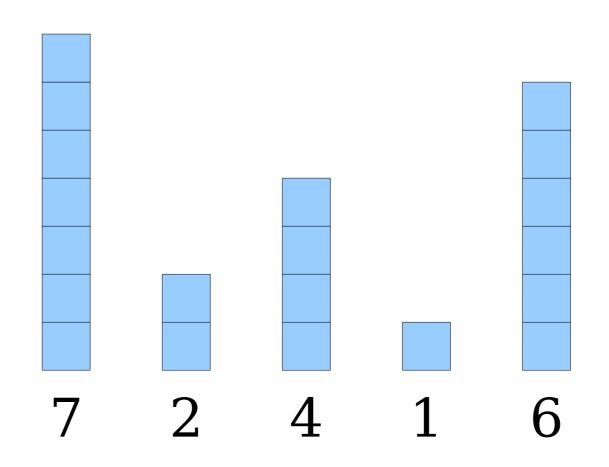
 $O(n^2)$

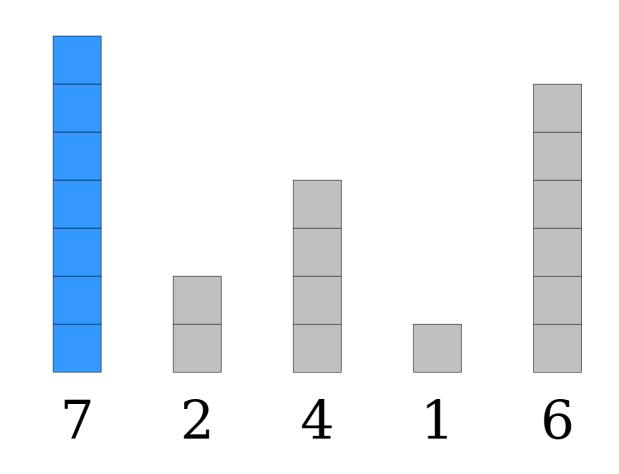
Our theory predicts that the runtime of selection sort is $O(n^2)$.

Does that match what we see in practice?

What should we expect to see when we look at a runtime plot?

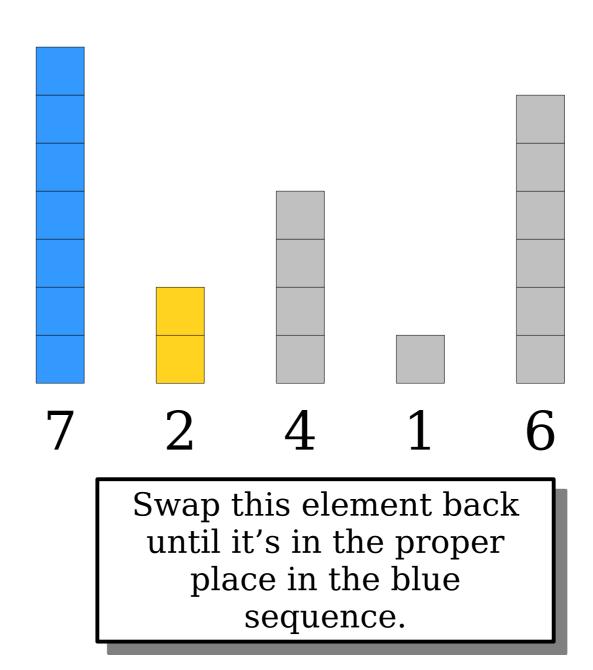
Another Sorting Algorithm

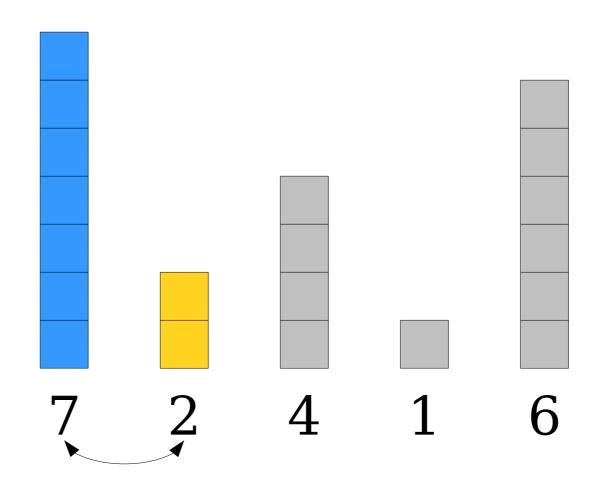


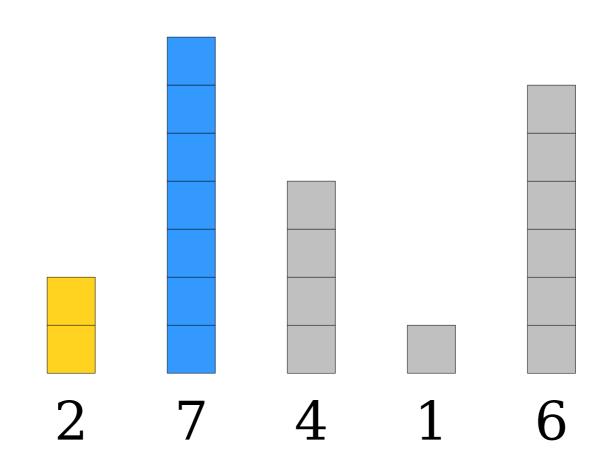


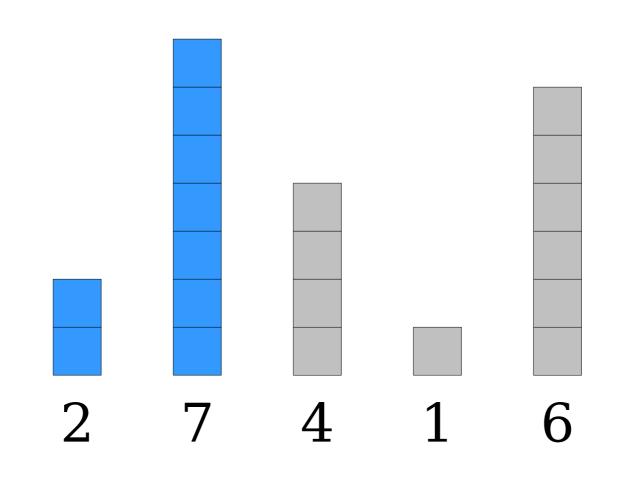
This sequence in blue, taken in isolation, is in sorted order.

This sequence in gray is in no particular order.



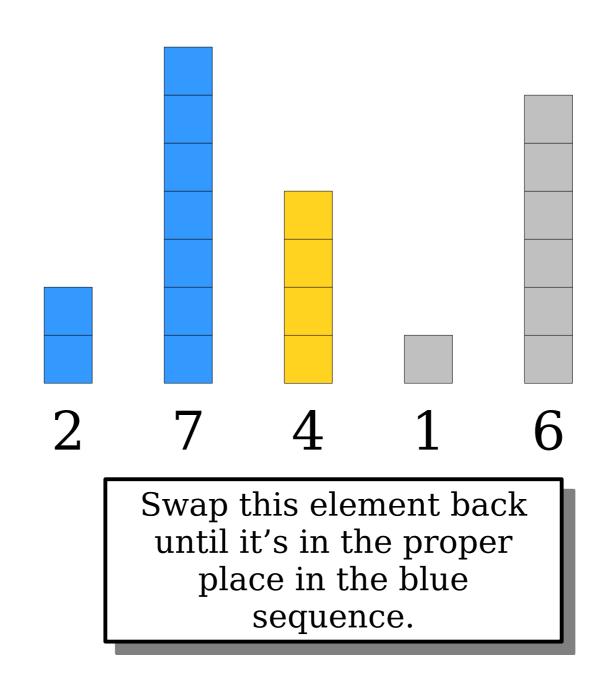


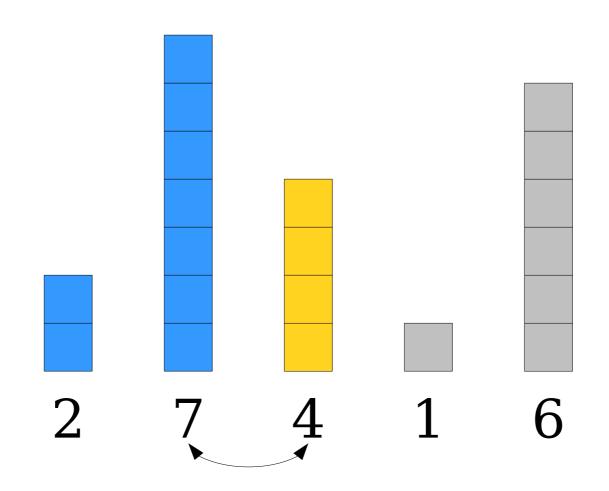


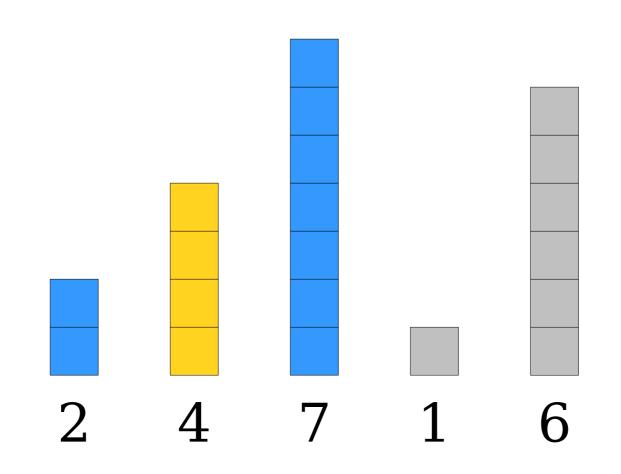


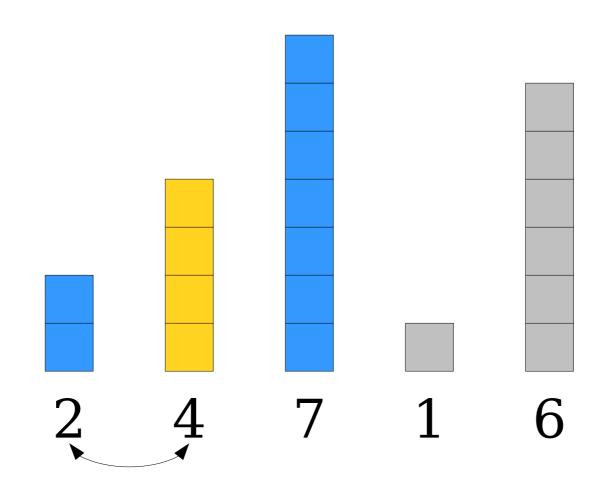
This sequence in blue, taken in isolation, is in sorted order.

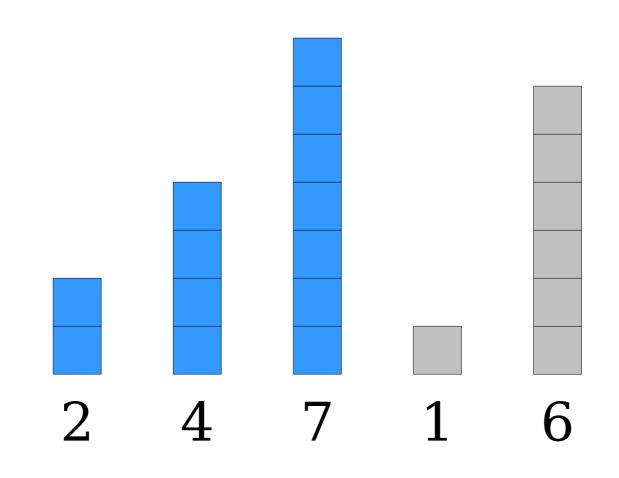
This sequence in gray is in no particular order.





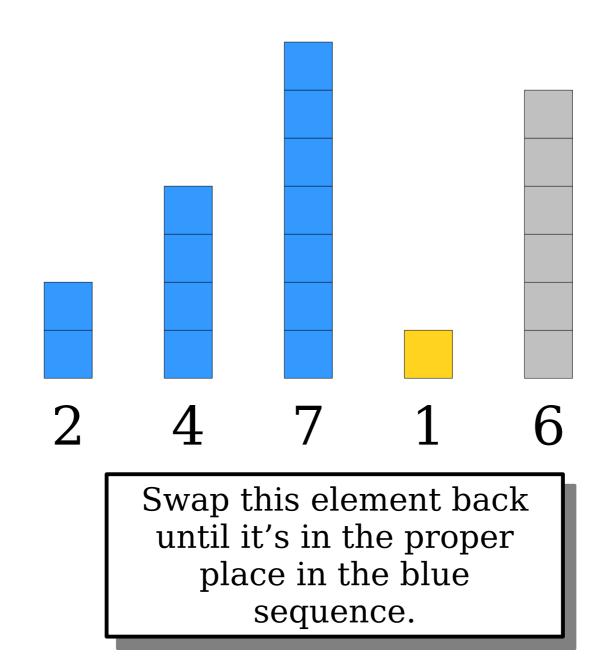


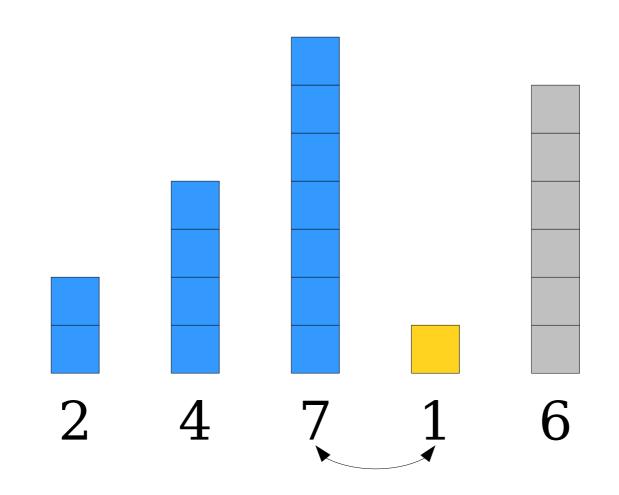


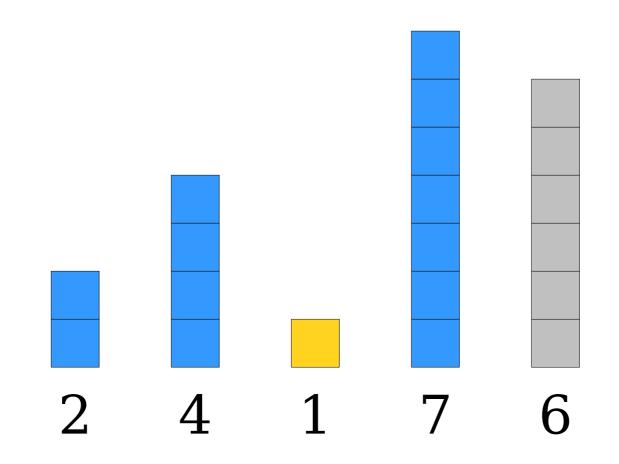


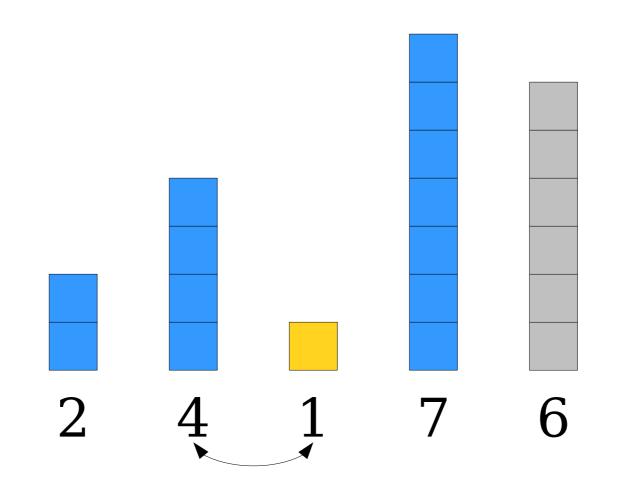
This sequence in blue, taken in isolation, is in sorted order.

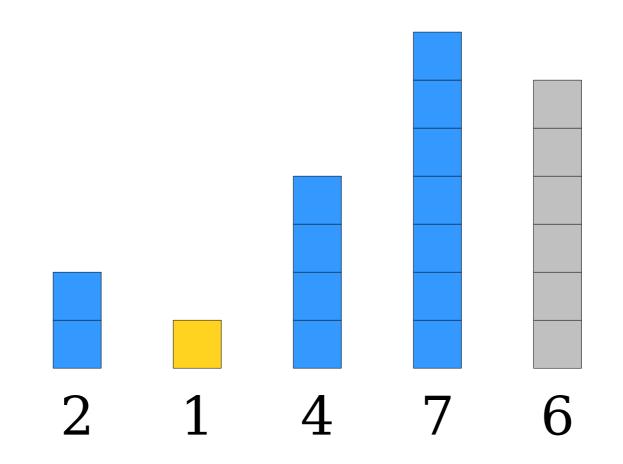
This sequence in gray is in no particular order.

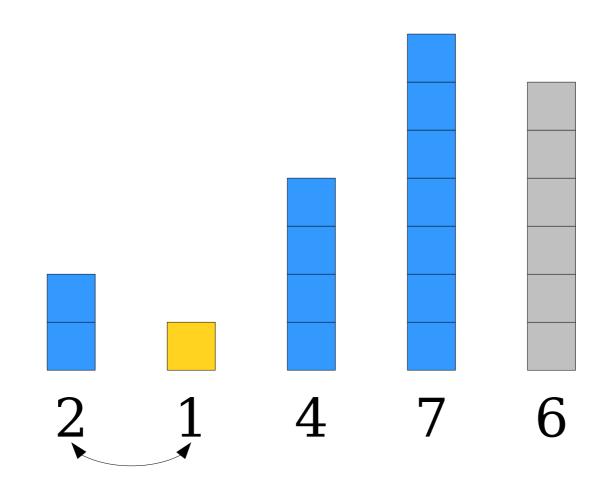


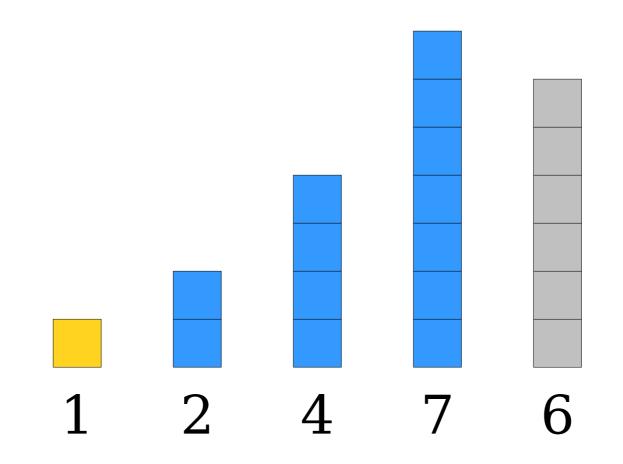


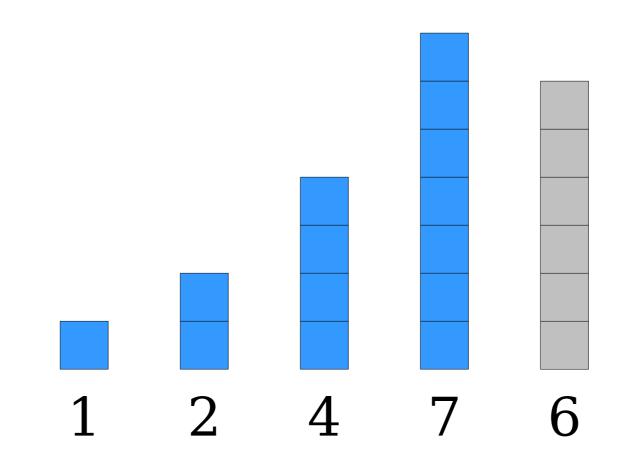






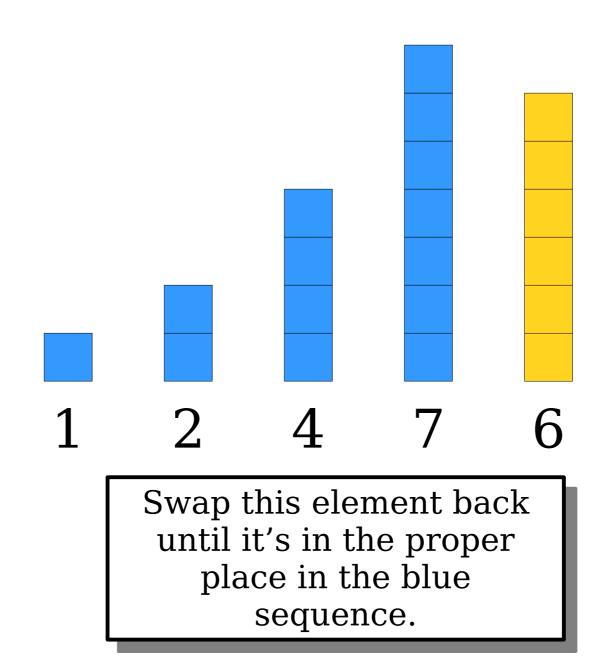


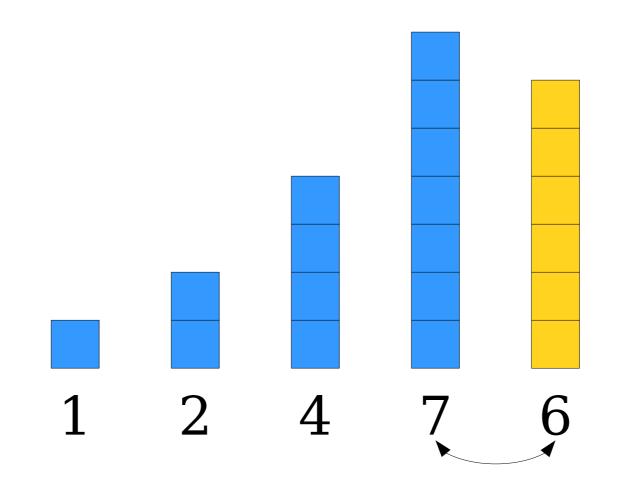


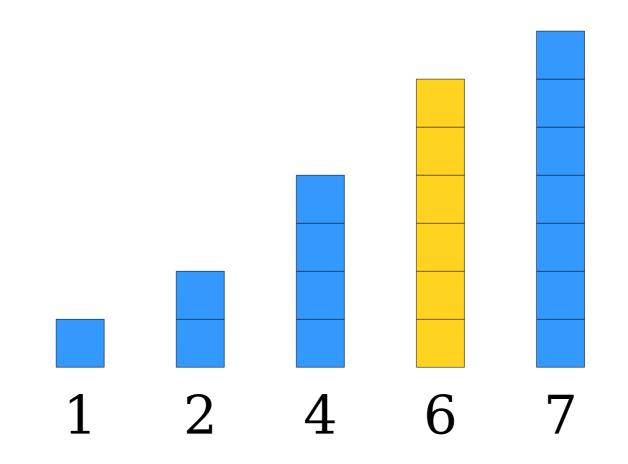


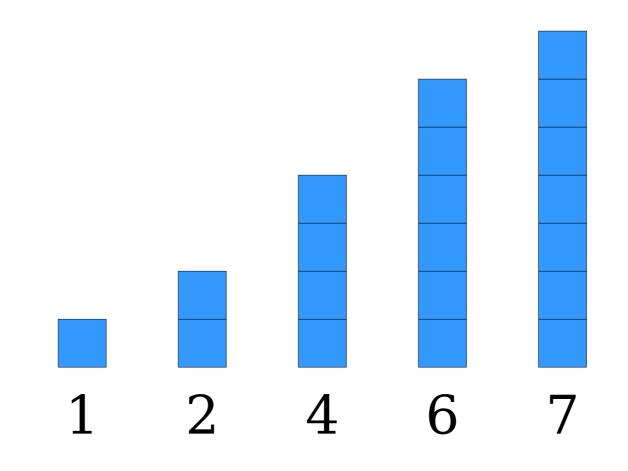
This sequence in blue, taken in isolation, is in sorted order.

This sequence in gray is in no particular order.









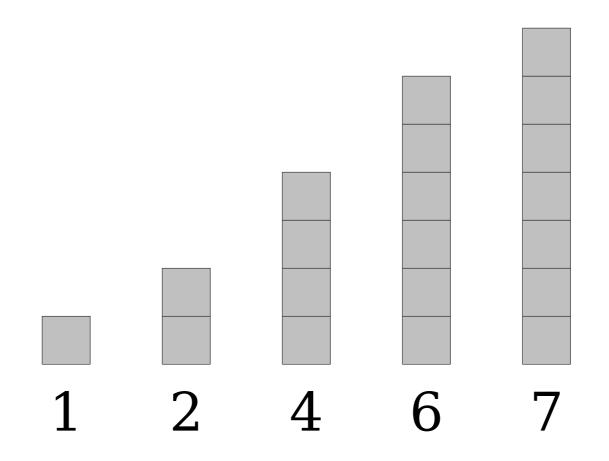
This sequence in blue, taken in isolation, is in sorted order.

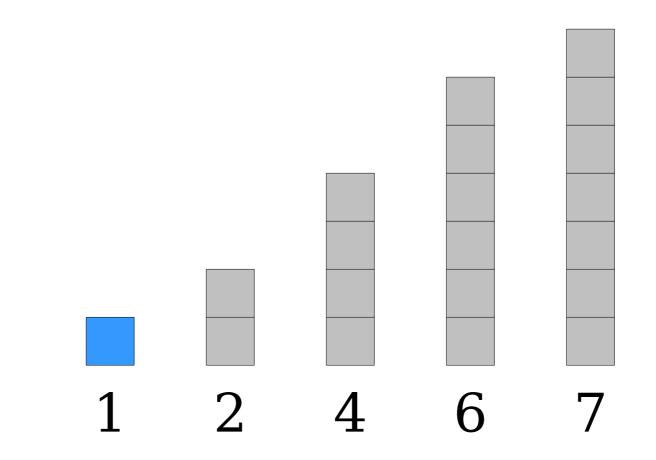
There are no more gray elements, so the sequence is sorted!

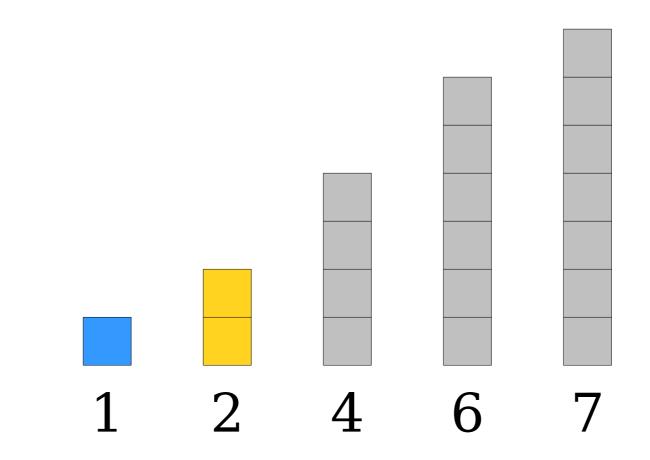
Insertion Sort

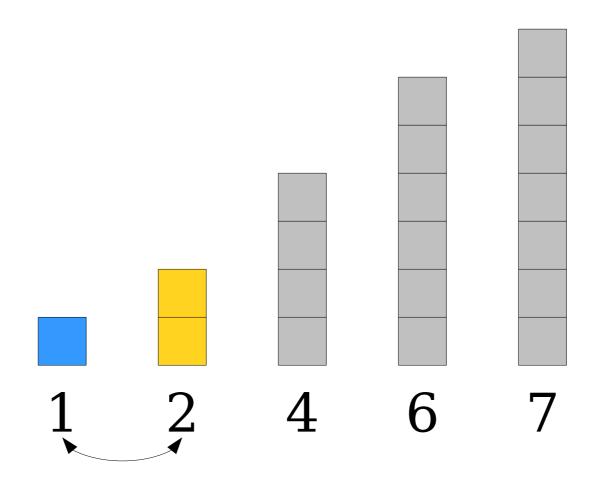
- Repeatedly *insert* an element into a sorted sequence at the front of the array.
- To insert an element, swap it backwards until either
 - (1) it's at least as big as the element before it in the sequence, or
 - (2) it's at the front of the array.

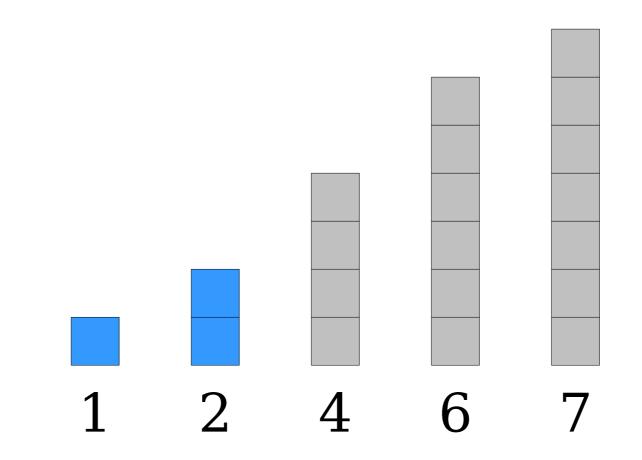
```
/**
 * Sorts the specified vector using insertion sort.
 * @param v The vector to sort.
void insertionSort(Vector<int>& v) {
  for (int i = 0; i < v.size(); i++) {</pre>
    /* Scan backwards until either (1) there is no
     * preceding element or the preceding element is
     * no bigger than us.
    for (int j = i - 1; j >= 0; j--) {
      if (v[j] <= v[j + 1]) break;</pre>
      /* Swap this element back one step. */
      swap(v[j], v[j + 1]);
```

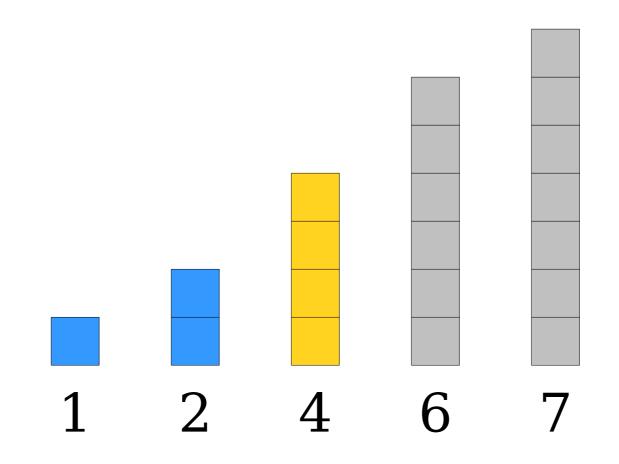


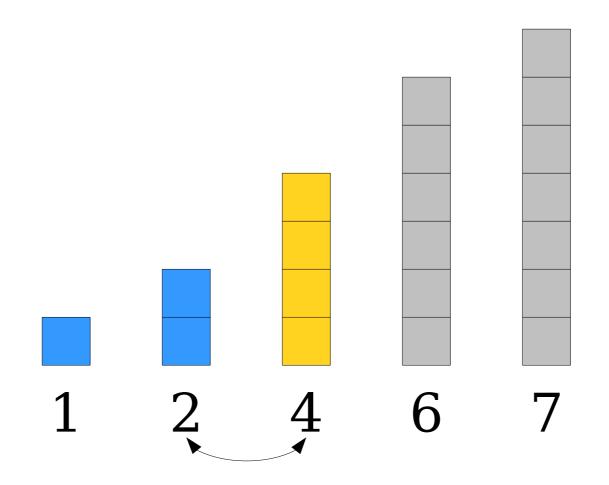


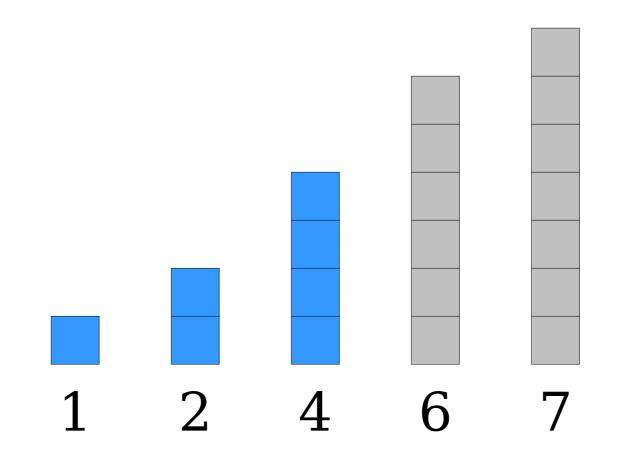


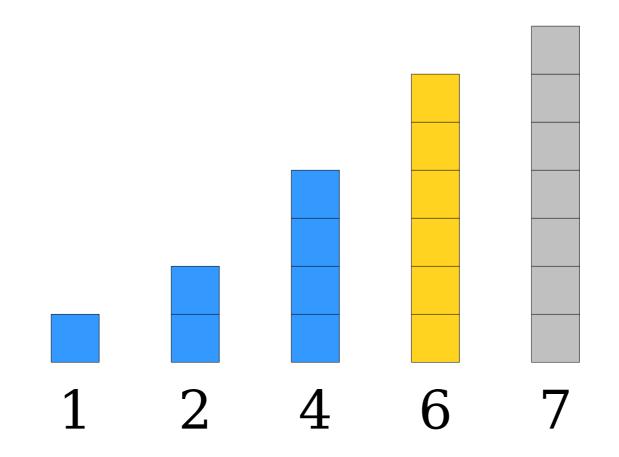


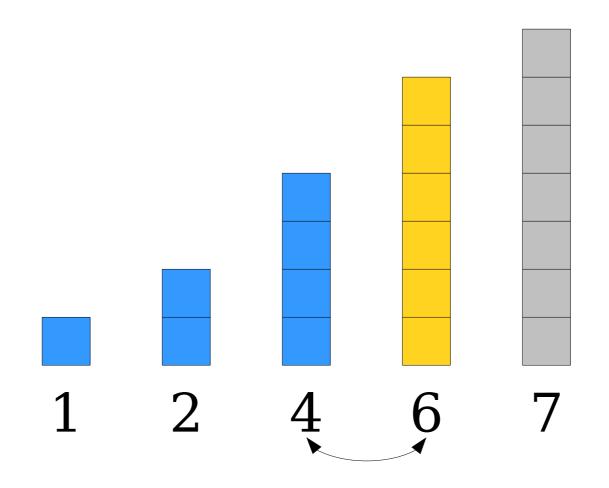


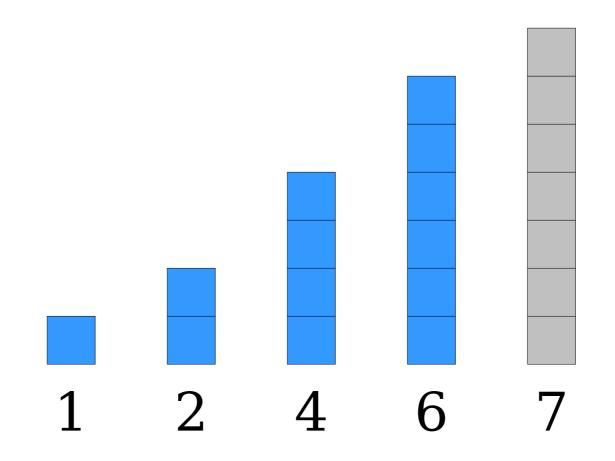


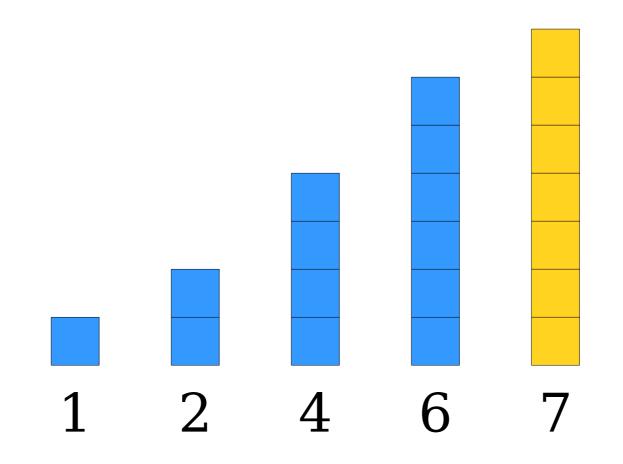


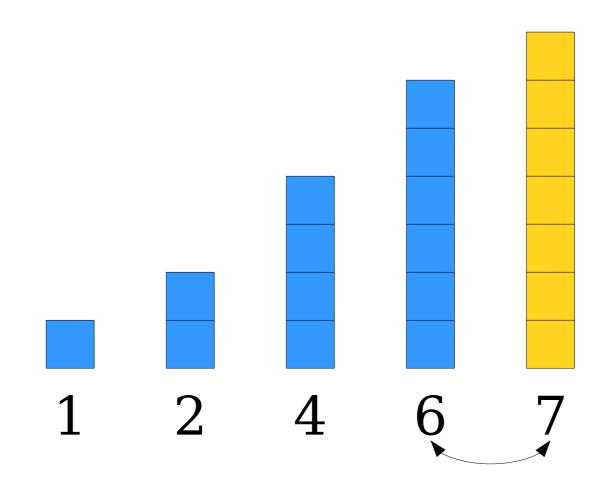


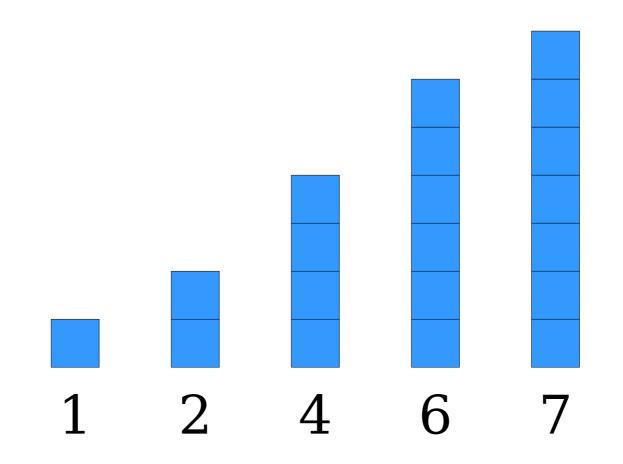


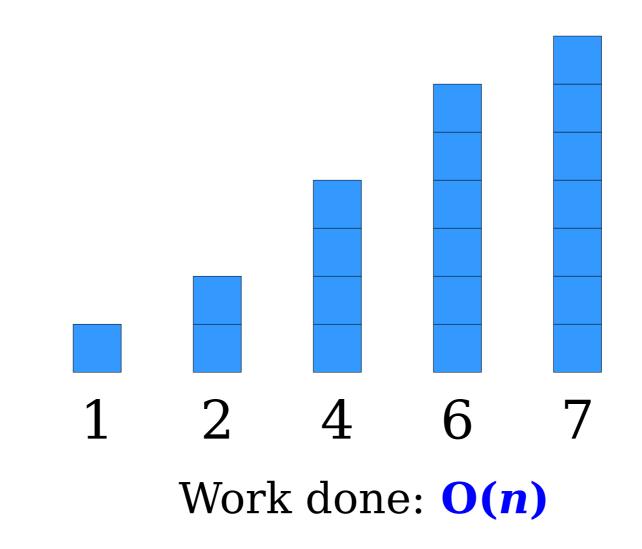


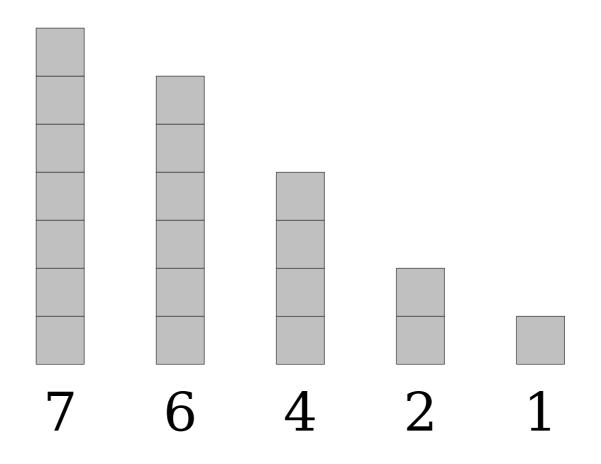


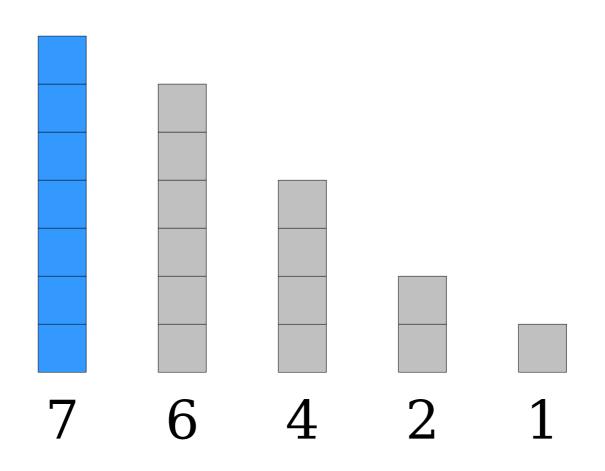


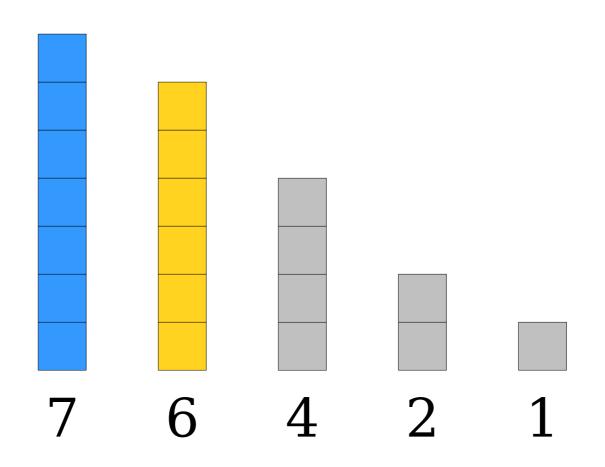


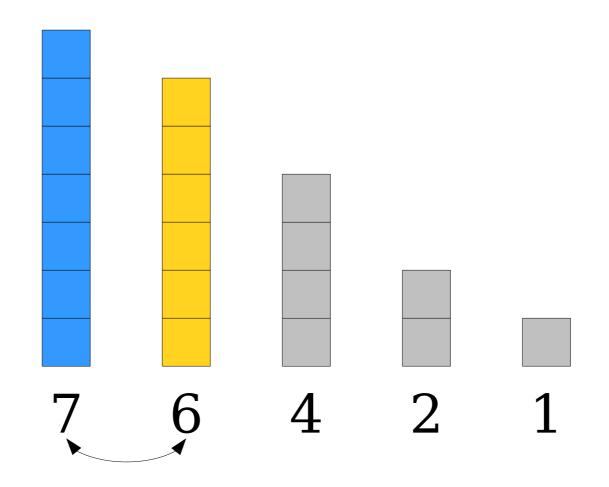


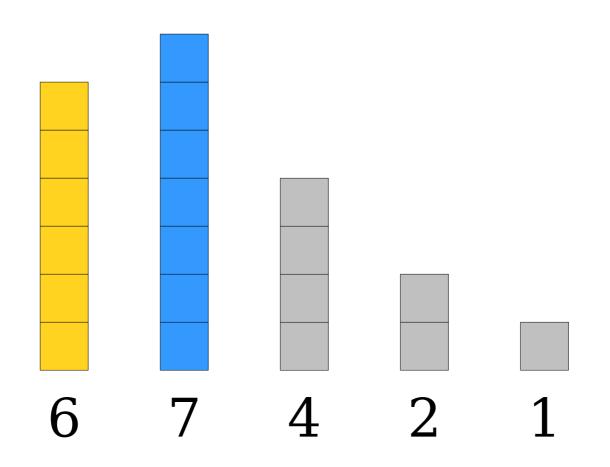


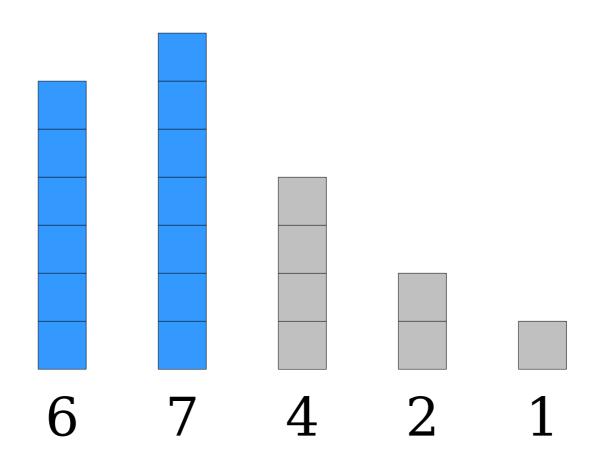


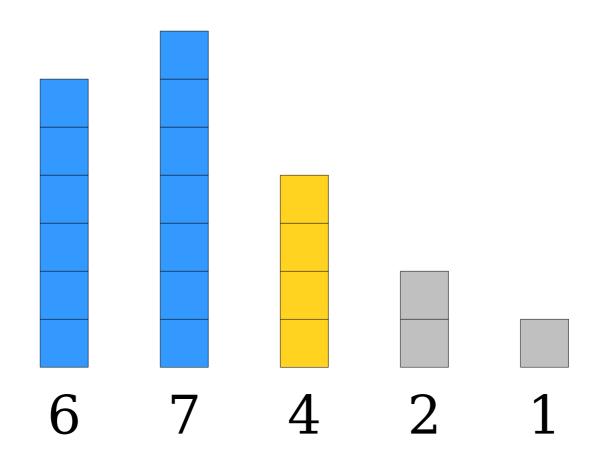


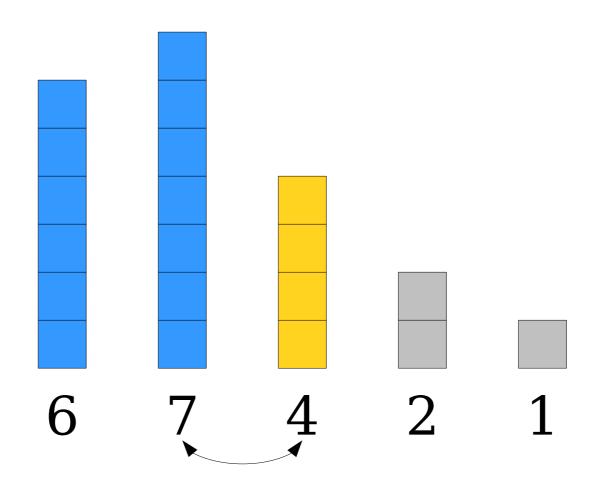


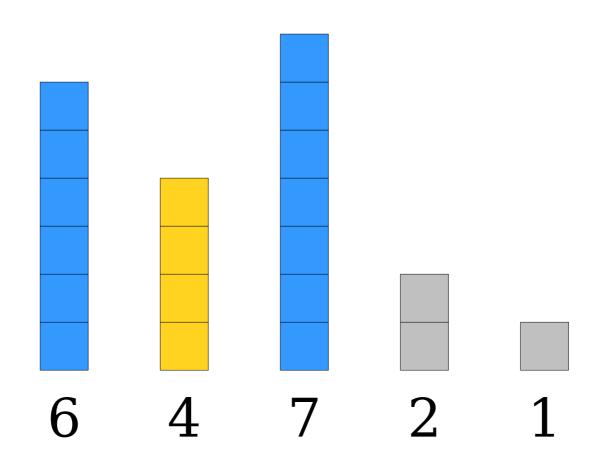


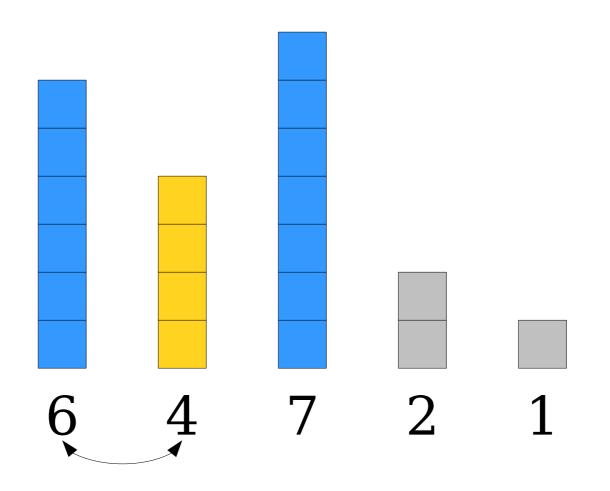


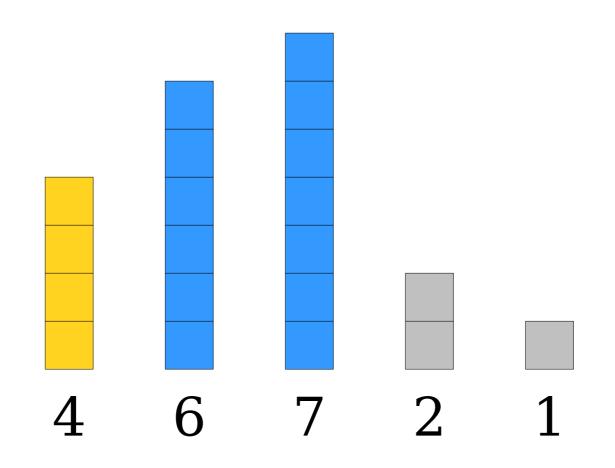


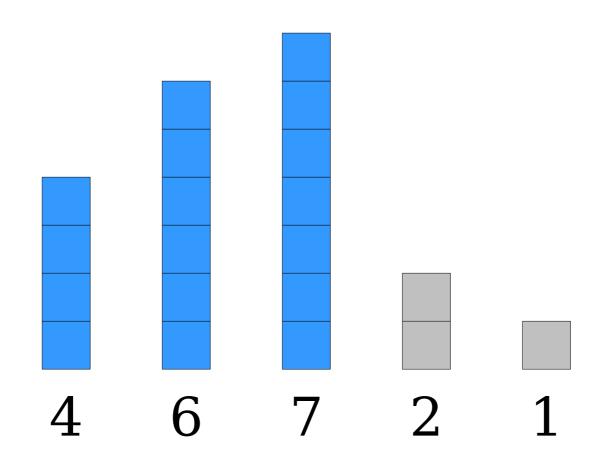


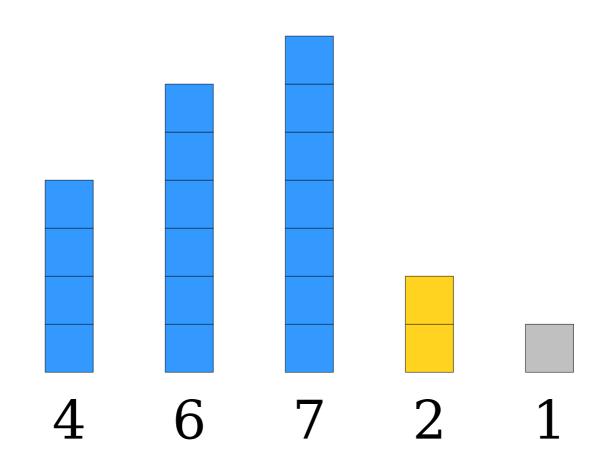


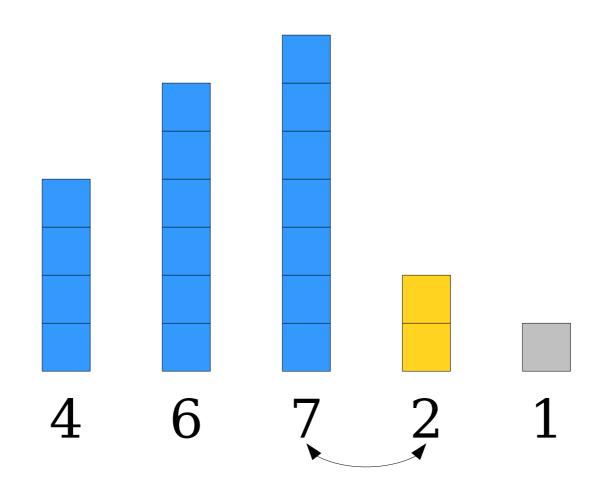


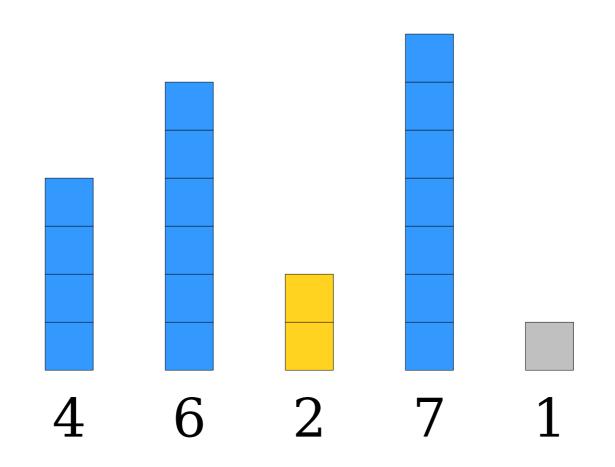


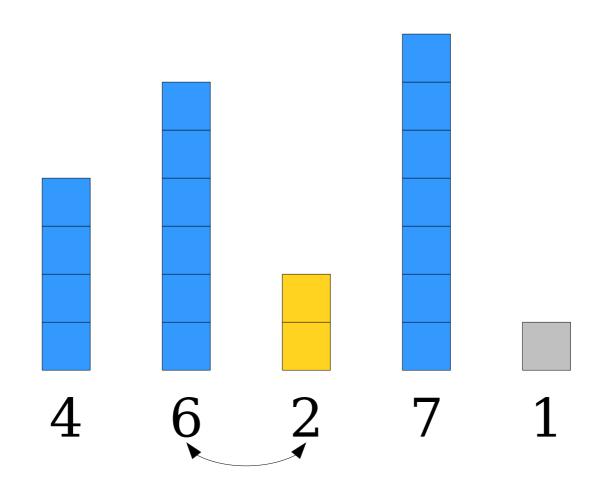


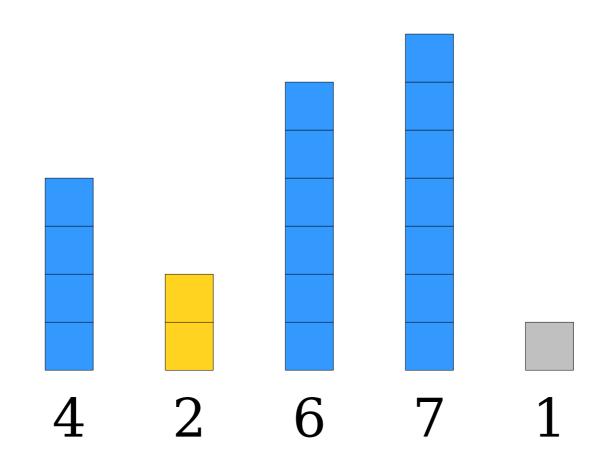


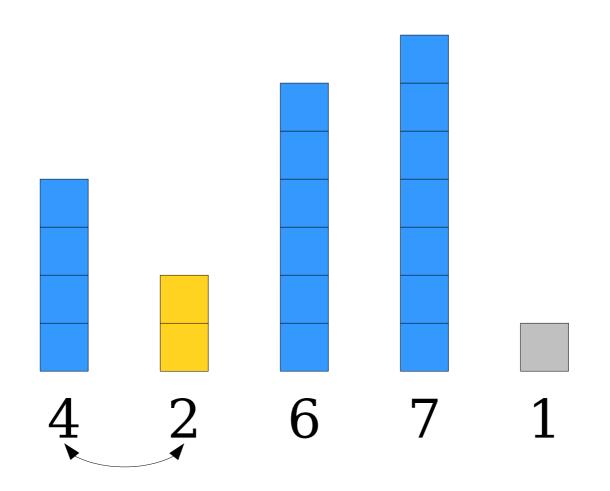


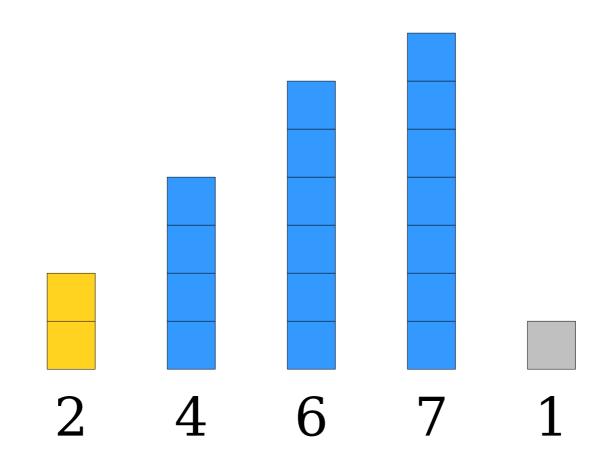


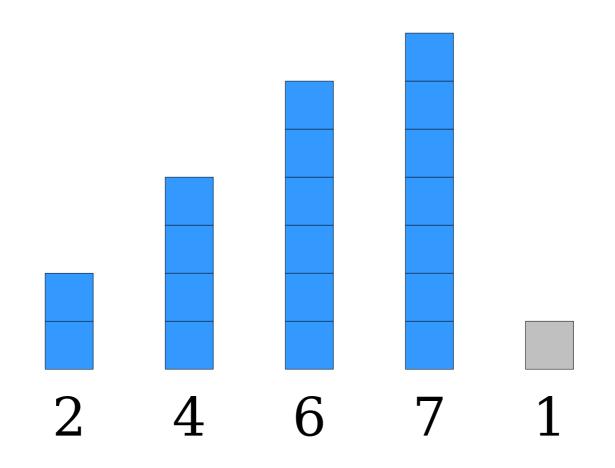


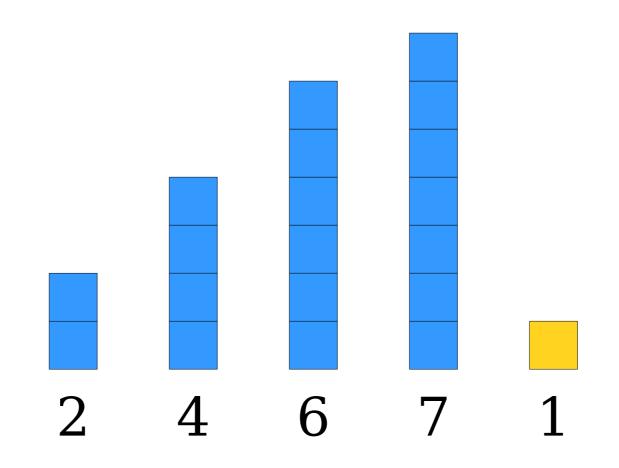


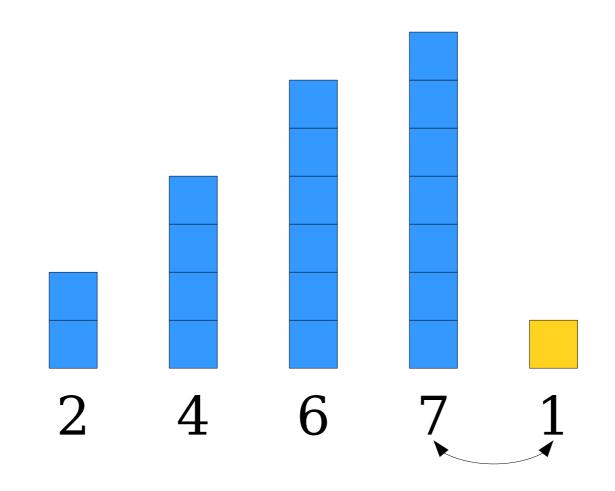


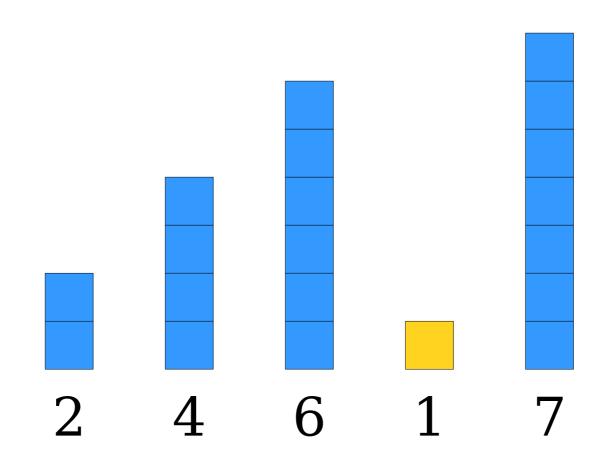


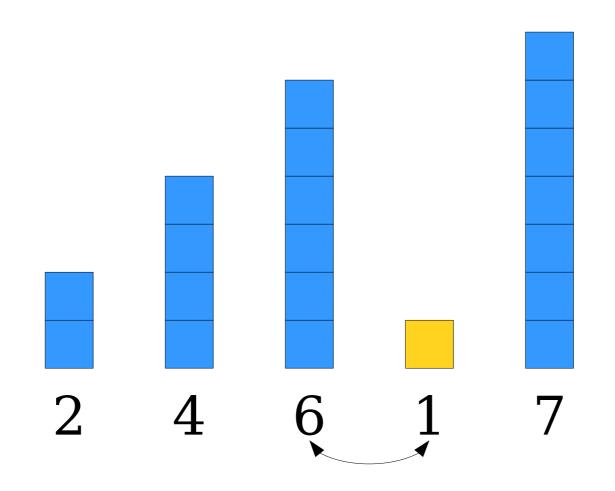


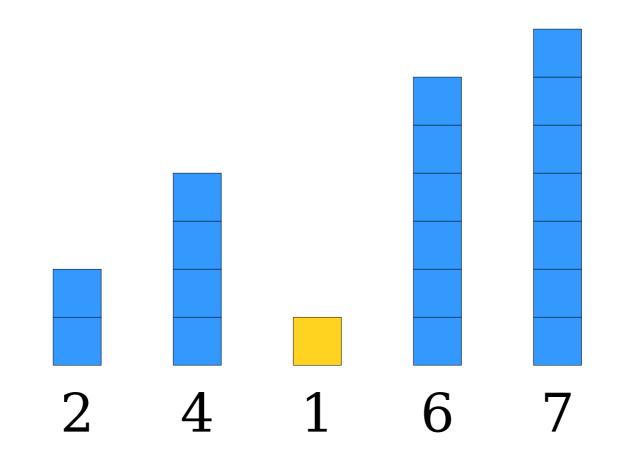


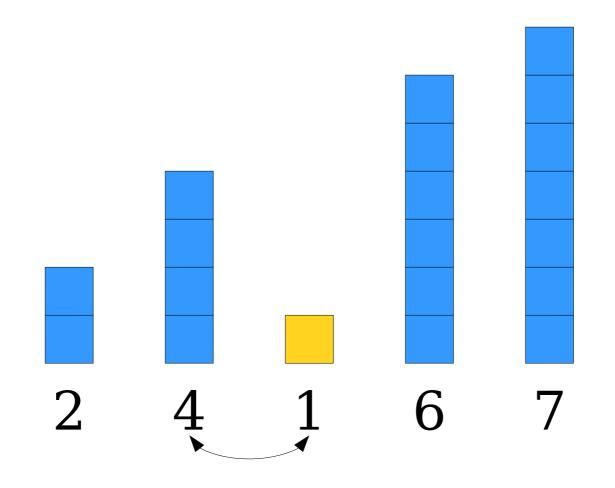


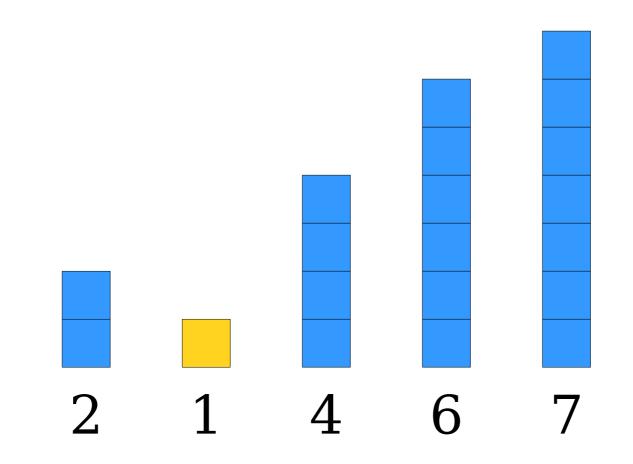


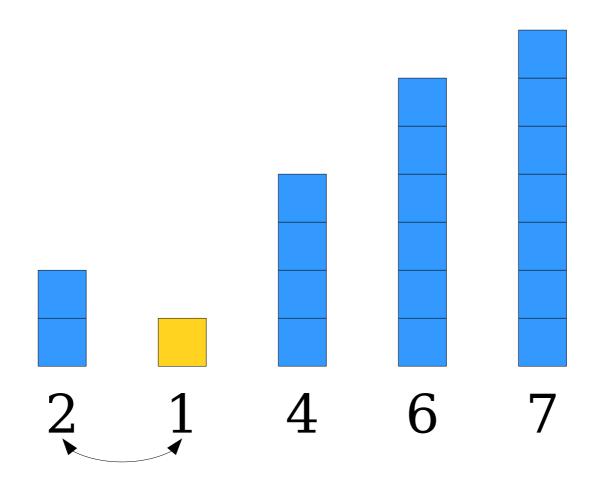


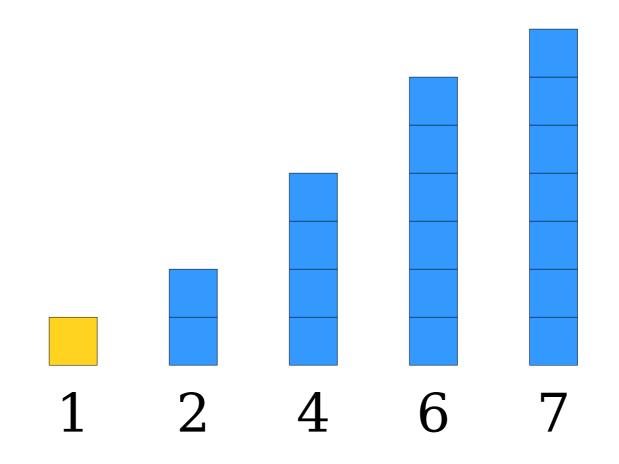


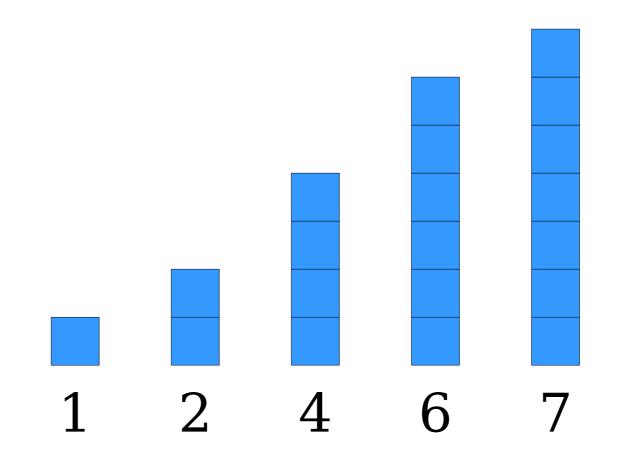




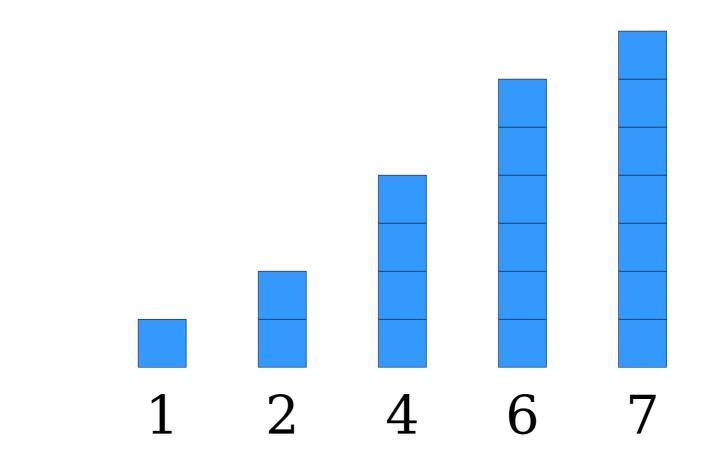








How Fast is Insertion Sort?



Work Done:
$$1 + 2 + 3 + ... + n-1$$

= $O(n^2)$

Three Analyses

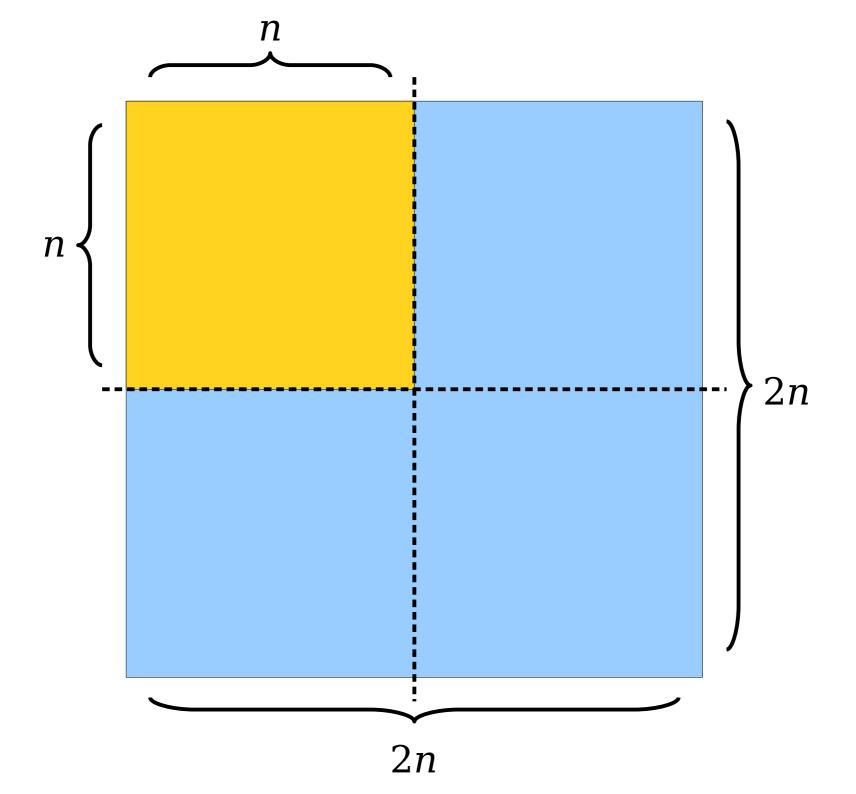
- Worst-Case Analysis
 - What's the worst possible runtime for the algorithm?
 - Useful for "sleeping well at night."
- Best-Case Analysis
 - What's the best possible runtime for the algorithm?
 - Useful to see if the algorithm performs well in some cases.
- Average-Case Analysis
 - What's the *average* runtime for the algorithm?
 - Far beyond the scope of this class; take CS109, CS161, or CS265 for more information!

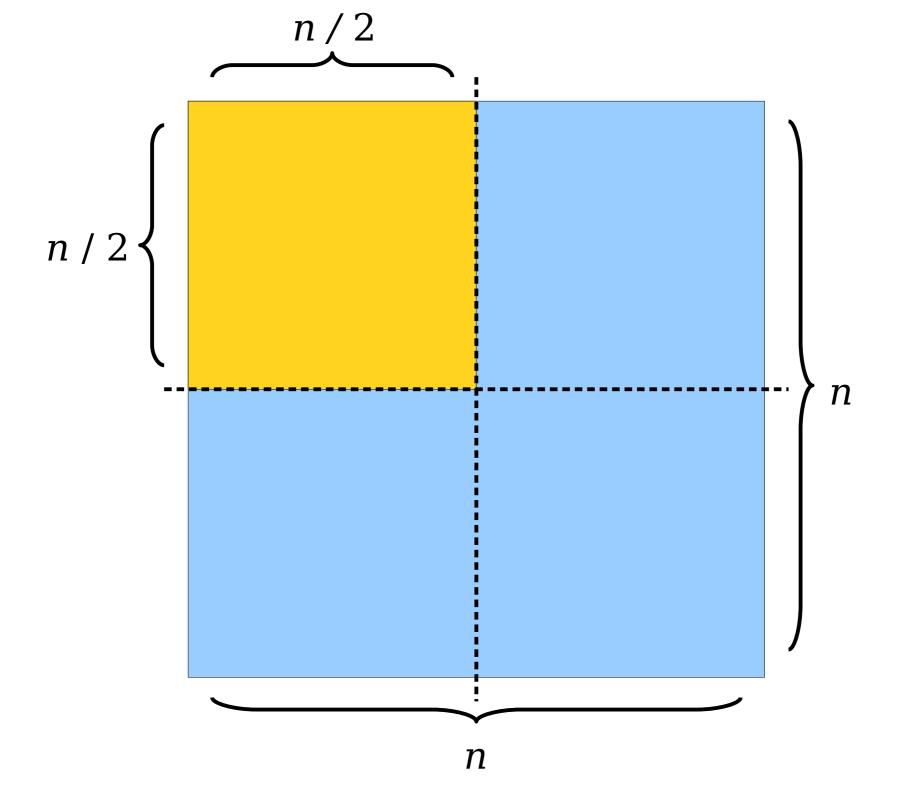
The Complexity of Insertion Sort

- In the best case (the array is sorted), insertion takes time O(n).
- In the worst case (the array is reverse-sorted), insertion sort takes time $O(n^2)$.
- **Fun fact:** Insertion sorting an array of (uniformly) random values takes, on average, $O(n^2)$ time.
 - Curious why? Come talk to me after class!

How do selection sort and insertion sort compare against one another?

Building a Better Sorting Algorithm

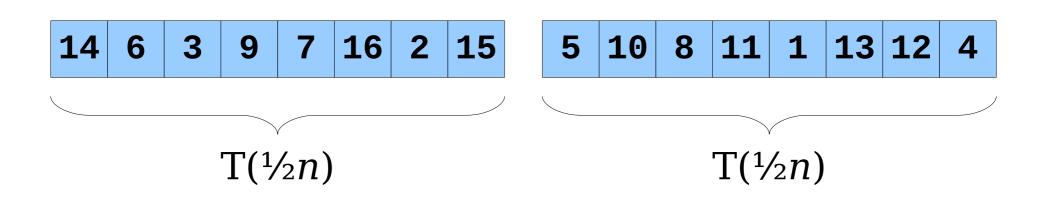




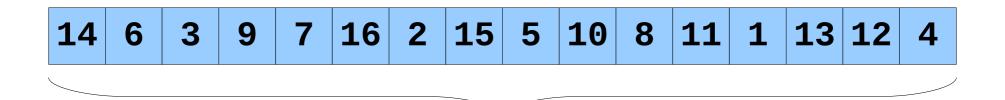
Thinking About $O(n^2)$



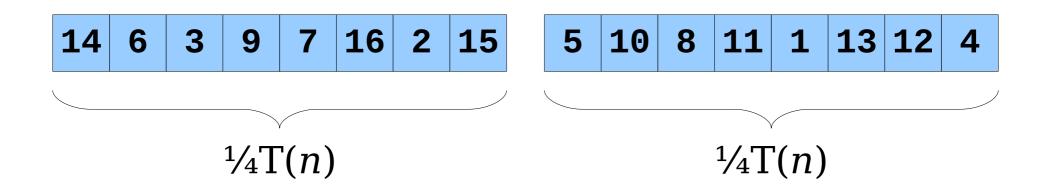
T(n)



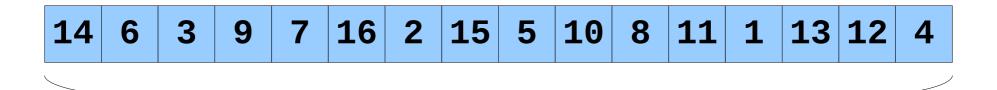
Thinking About $O(n^2)$

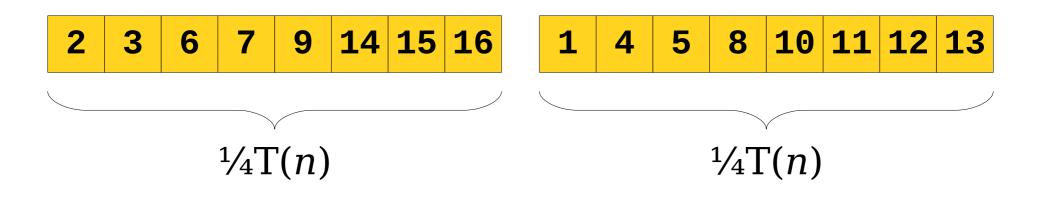


T(n)



Thinking About $O(n^2)$

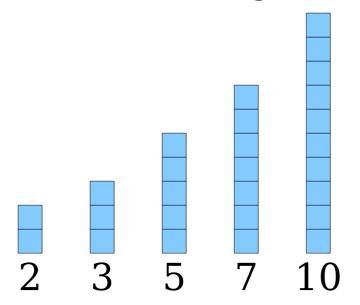


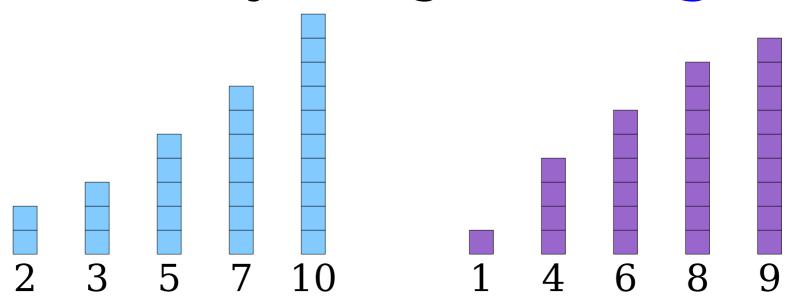


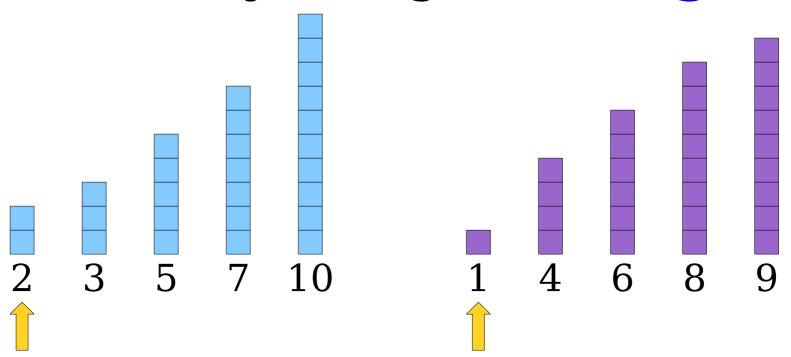
$$2 \cdot \frac{1}{4}T(n) = \frac{1}{2}T(n)$$

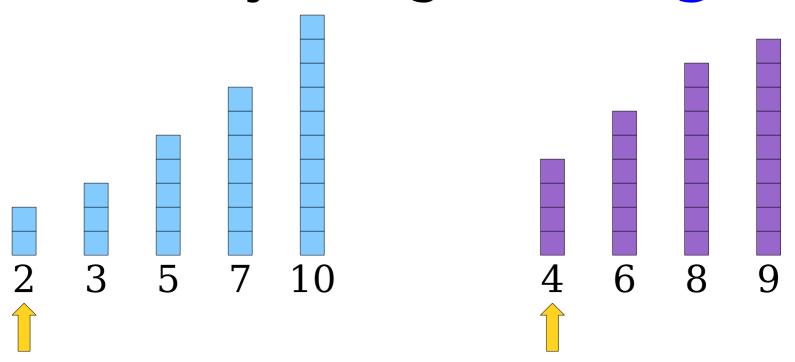
With an $O(n^2)$ -time sorting algorithm, it takes twice as long to sort the whole array as it does to split the array in half and sort each half.

Can we exploit this?

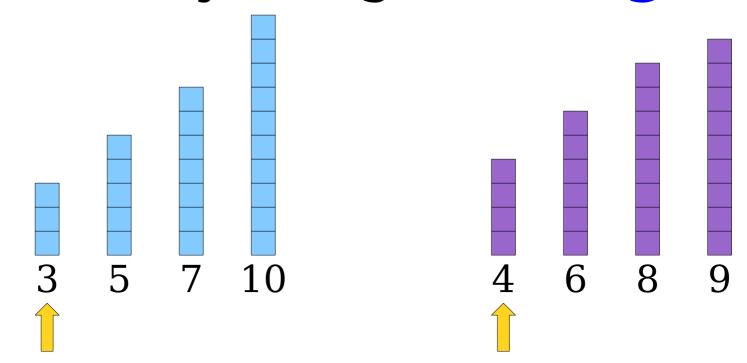




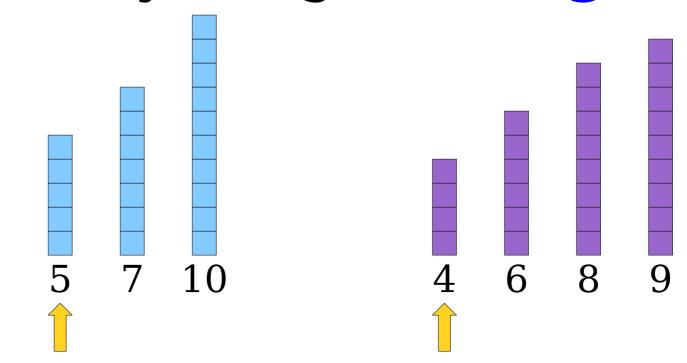


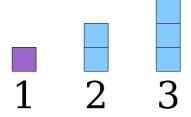


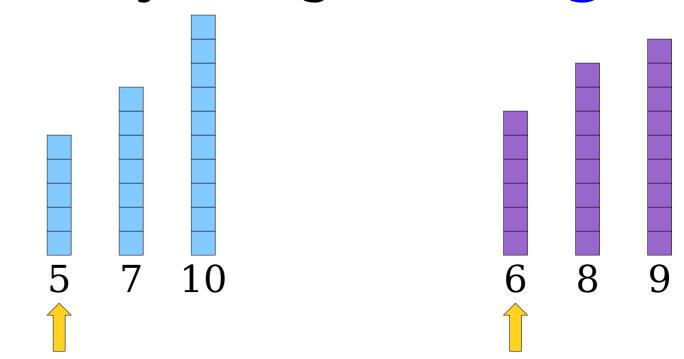


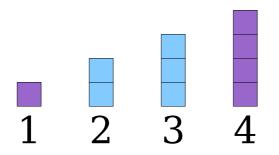


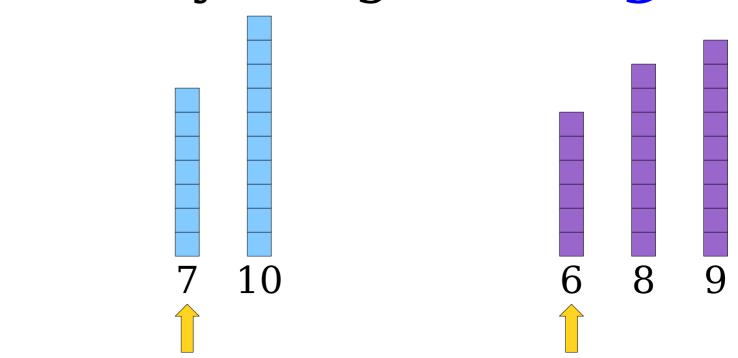


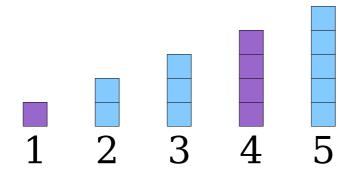


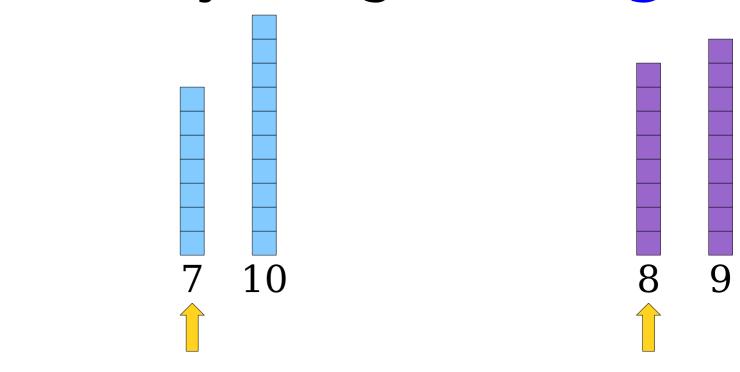


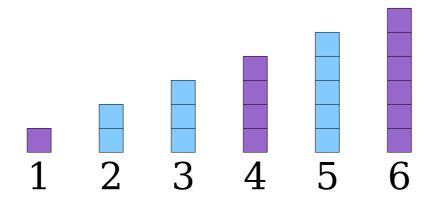


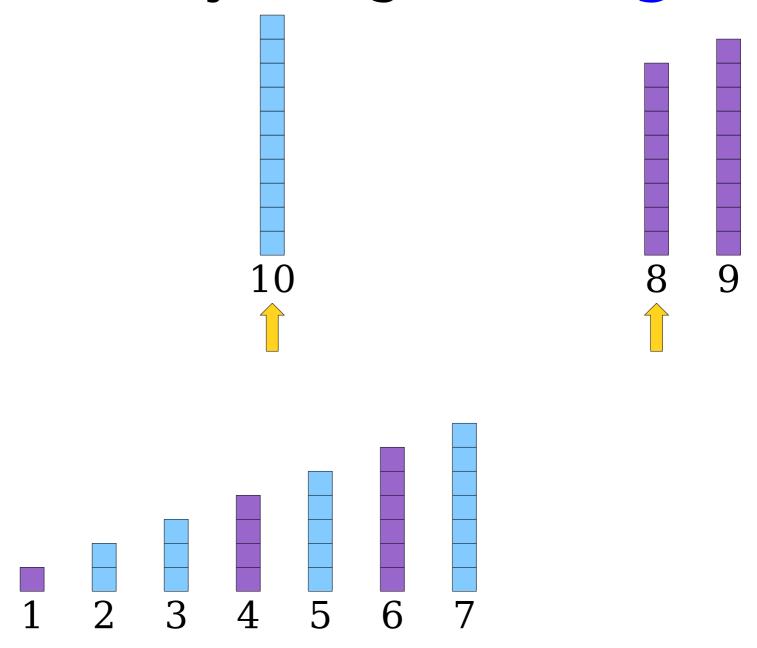


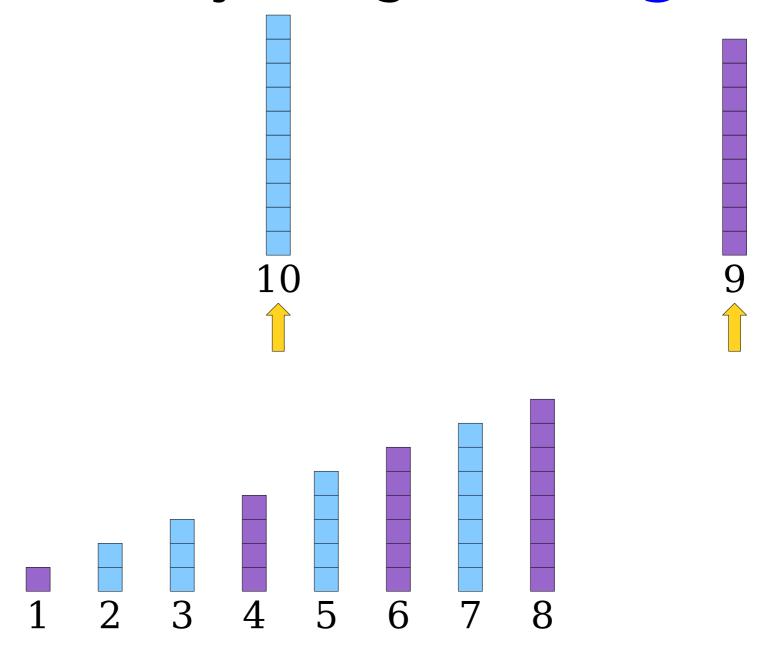


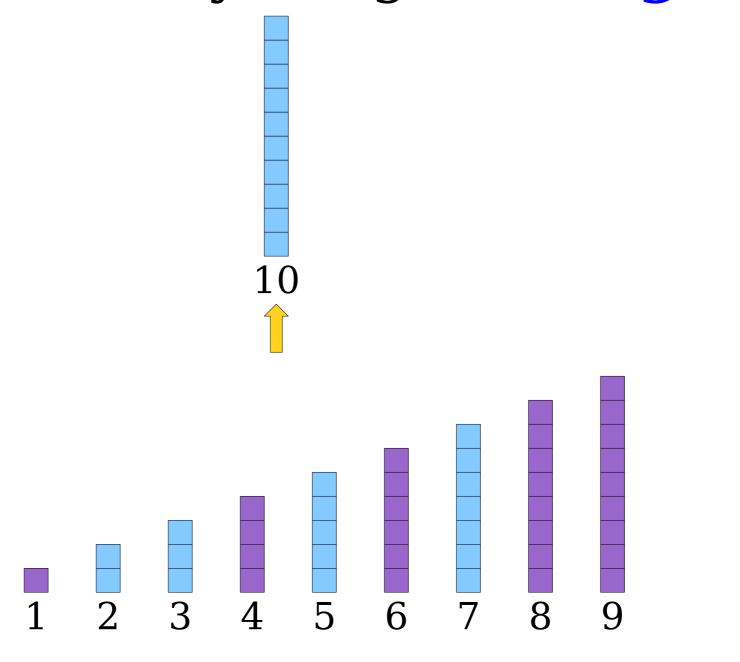


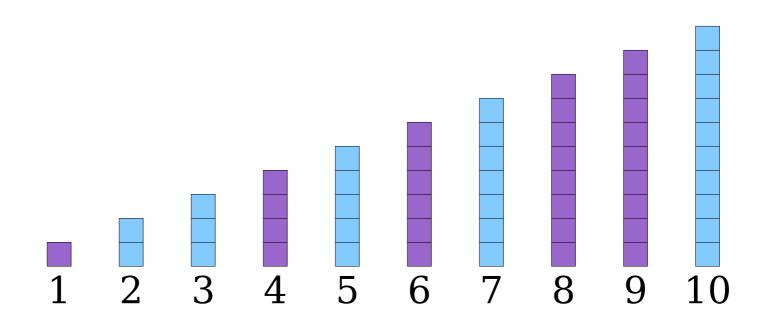






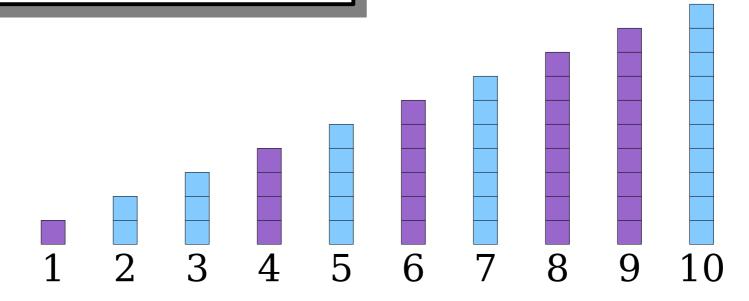




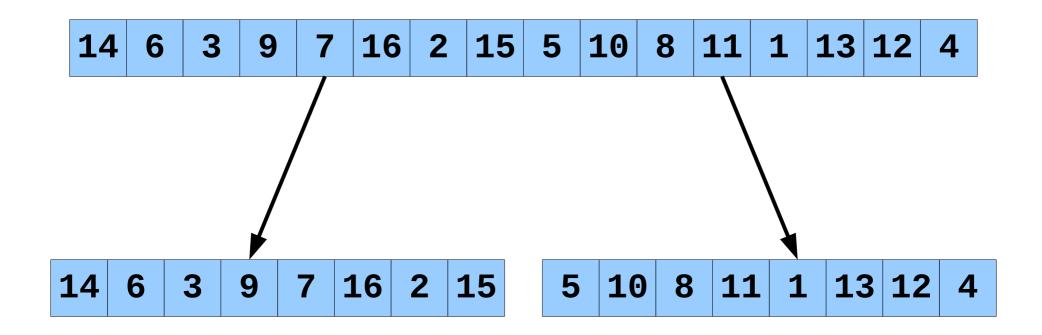


Each step makes a single comparison and reduces the number of elements by one.

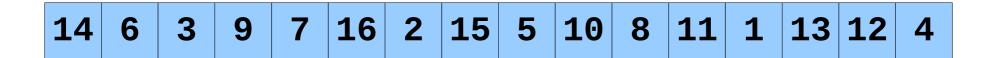
If there are n total elements, this algorithm runs in time O(n).



- The *merge* algorithm takes in two sorted lists and combines them into a single sorted list.
- While both lists are nonempty, compare their first elements. Remove the smaller element and append it to the output.
- Once one list is empty, add all elements from the other list to the output.
- We'll leave the code for this as an Exercise to the Reader.



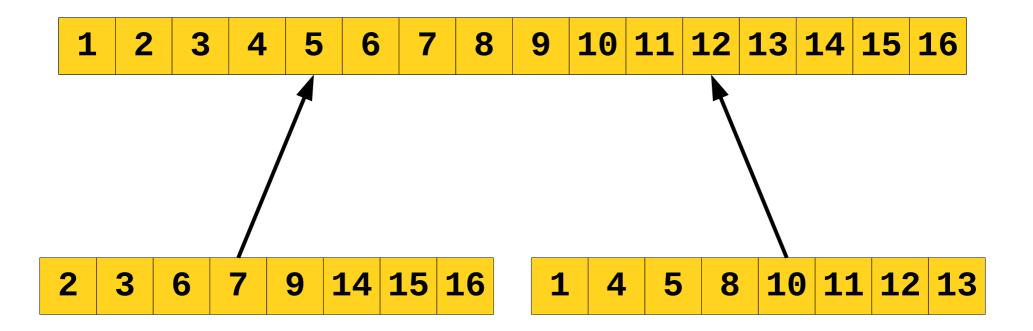
1. Split the input in half.





5 **10** 8 **11** 1 **13** 12 4

- 1. Split the input in half.
- 2. Insertion sort each half.



- 1. Split the input in half.
- 2. Insertion sort each half.
- 3. Merge the halves back together.

```
void splitSort(Vector<int>& v) {
    /* Split the vector in half */
    int half = v.size() / 2;
    Vector<int> left = v.subList(0, half);
    Vector<int> right = v.subList(half);
    /* Sort each half. */
    insertionSort(left);
    insertionSort(right);
    /* Merge them back together. */
    v = merge(left, right);
```

```
void splitSort(Vector<int>& v) {
    /* Split the vector in half */
                                                    Takes O(n) time,
    int half = v.size() / 2;
                                                    since we copy all
    Vector<int> left = v.subList(0, half);
                                                     n elements into
    Vector<int> right = v.subList(half);
                                                      new Vectors.
    /* Sort each half. */
    insertionSort(left);
    insertionSort(right);
    /* Merge them back together. */
    v = merge(left, right);
```

```
void splitSort(Vector<int>& v) {
    /* Split the vector in half */
                                                      Takes O(n) time,
    int half = v.size() / 2;
                                                      since we copy all
    Vector<int> left = v.subList(0, half);
                                                      n elements into
    Vector<int> right = v.subList(half);
                                                        new Vectors.
    /* Sort each half. */
                                    Takes O(n^2) time, but
    insertionSort(left);
                                    about half as much as
    insertionSort(right);
                                     what we did before.
    /* Merge them back together. */
    v = merge(left, right);
```

```
void splitSort(Vector<int>& v) {
    /* Split the vector in half */
                                                      Takes O(n) time,
    int half = v.size() / 2;
                                                      since we copy all
    Vector<int> left = v.subList(0, half);
                                                       n elements into
    Vector<int> right = v.subList(half);
                                                        new Vectors.
    /* Sort each half. */
                                     Takes O(n^2) time, but
    insertionSort(left);
                                    about half as much as
    insertionSort(right);
                                     what we did before.
    /* Merge them back together. */
    v = merge(left, right);
           Takes O(n)
             time.
```

```
void splitSort(Vector<int>& v) {
    /* Split the vector in half */
                                                      Takes O(n) time,
    int half = v.size() / 2;
                                                      since we copy all
    Vector<int> left = v.subList(0, half);
                                                      n elements into
    Vector<int> right = v.subList(half);
                                                       new Vectors.
    /* Sort each half. */
                                    Takes O(n^2) time, but
    insertionSort(left);
                                    about half as much as
    insertionSort(right);
                                     what we did before.
    /* Merge them back together. */
    v = merge(left, right);
                                              Prediction: This
                                            should still take time
           Takes O(n)
```

time.

Prediction: This should still take time $O(n^2)$, but be about twice as fast as insertion sort.

Next Time

- Mergesort
 - A beautiful, elegant sorting algorithm.
- Analyzing Mergesort
 - An unusual runtime analysis.
- Hybrid Sorting Algorithms
 - Improving on mergesort.
- Binary Search
 - Finding things fast!