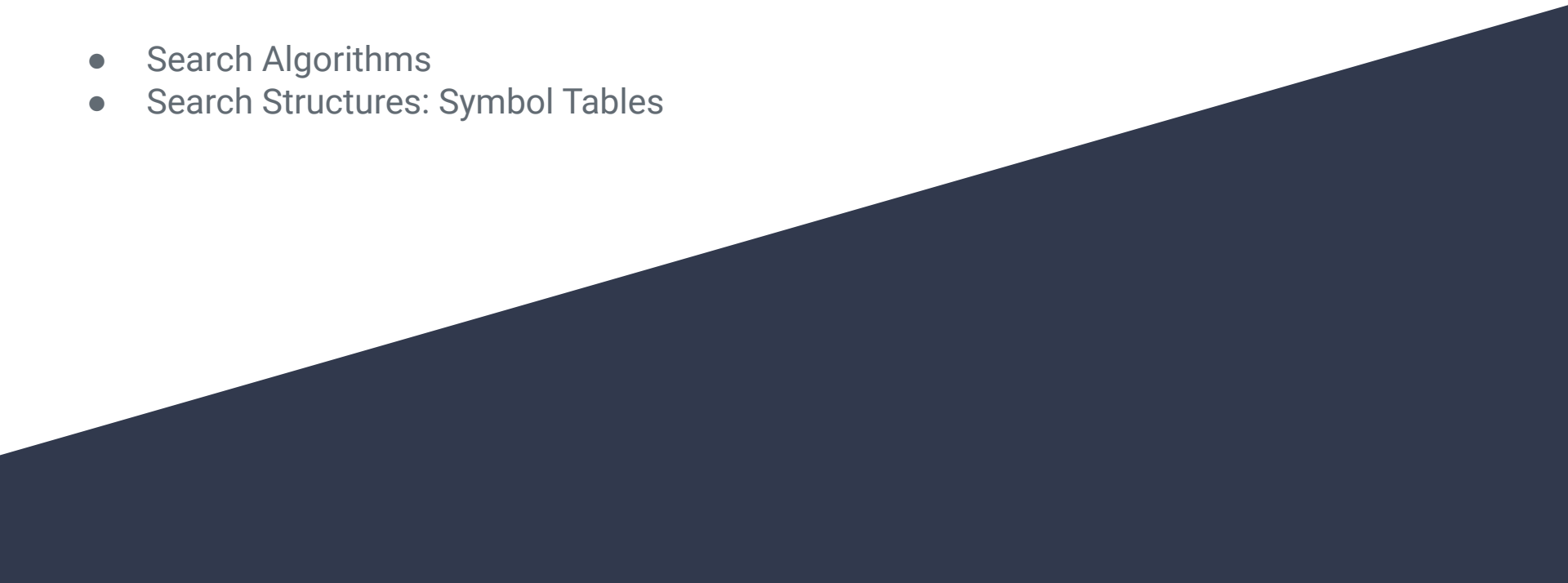


# Searching Basics

- Search Algorithms
  - Search Structures: Symbol Tables
- 
- A large, dark blue, curved shape that starts from the bottom left and extends diagonally upwards towards the right, filling the lower half of the slide.

# Searching

- With all the vast amounts of data out there, it's important to be able to search for specific items efficiently.
- Search algorithms are often closely tied with data structures because how data is stored affects how it can be searched.

# Review of data structures...

We've discussed...

- Stacks
- Queues
- Priority Queues
- Linked Lists

None of these are built for efficient *searching* (i.e. searching for a specific element.)

What about arrays?

# Linear Search on an unsorted array...

4	8	1	9	0	2	3	7	6	3
---	---	---	---	---	---	---	---	---	---

↑  
best:  $O(1)$   
worst:  $O(N)$

# Binary search on a sorted array...

0	1	2	3	3	4	6	7	8	9
---	---	---	---	---	---	---	---	---	---

best:  $O(1)$

worst:  $O(\log N)$

# Symbol Tables

- “...an abstract mechanism where we save information (a *value*) that we can later search for and retrieve by specifying a *key*.” [1]
- Sometimes called *dictionaries* or *indices*.
- A symbol table associates *values* with *keys* (called *key-value pairs*) which get inserted into the table and can later be retrieved using the key.

# Symbol Tables

- The important operations include *insert* (aka *put*) and *search* (aka *get*) but there may be other useful operations as well.
- To implement, we need...
  - an underlying data structure, and
  - definitions for the above operations (and other useful operations).



# Some Applications

application	purpose of search	key	value
<i>dictionary</i>	find definition	word	definition
<i>book index</i>	find relevant pages	term	list of page numbers
<i>file share</i>	find song to download	name of song	computer ID
<i>account management</i>	process transactions	account number	transaction details
<i>web search</i>	find relevant web pages	keyword	list of page names
<i>compiler</i>	find type and value	variable name	type and value

**Typical symbol-table applications**

```
public class ST<Key, Value>
```

---

```
    ST()
```

*create a symbol table*

```
    void put(Key key, Value val)
```

*put key-value pair into the table  
(remove key from table if value is null)*

```
    Value get(Key key)
```

*value paired with key  
(null if key is absent)*

```
    void delete(Key key)
```

*remove key (and its value) from table*

```
    boolean contains(Key key)
```

*is there a value paired with key?*

```
    boolean isEmpty()
```

*is the table empty?*

```
    int size()
```

*number of key-value pairs in the table*

```
    Iterable<Key> keys()
```

*all the keys in the table*

**API for a generic basic symbol table**

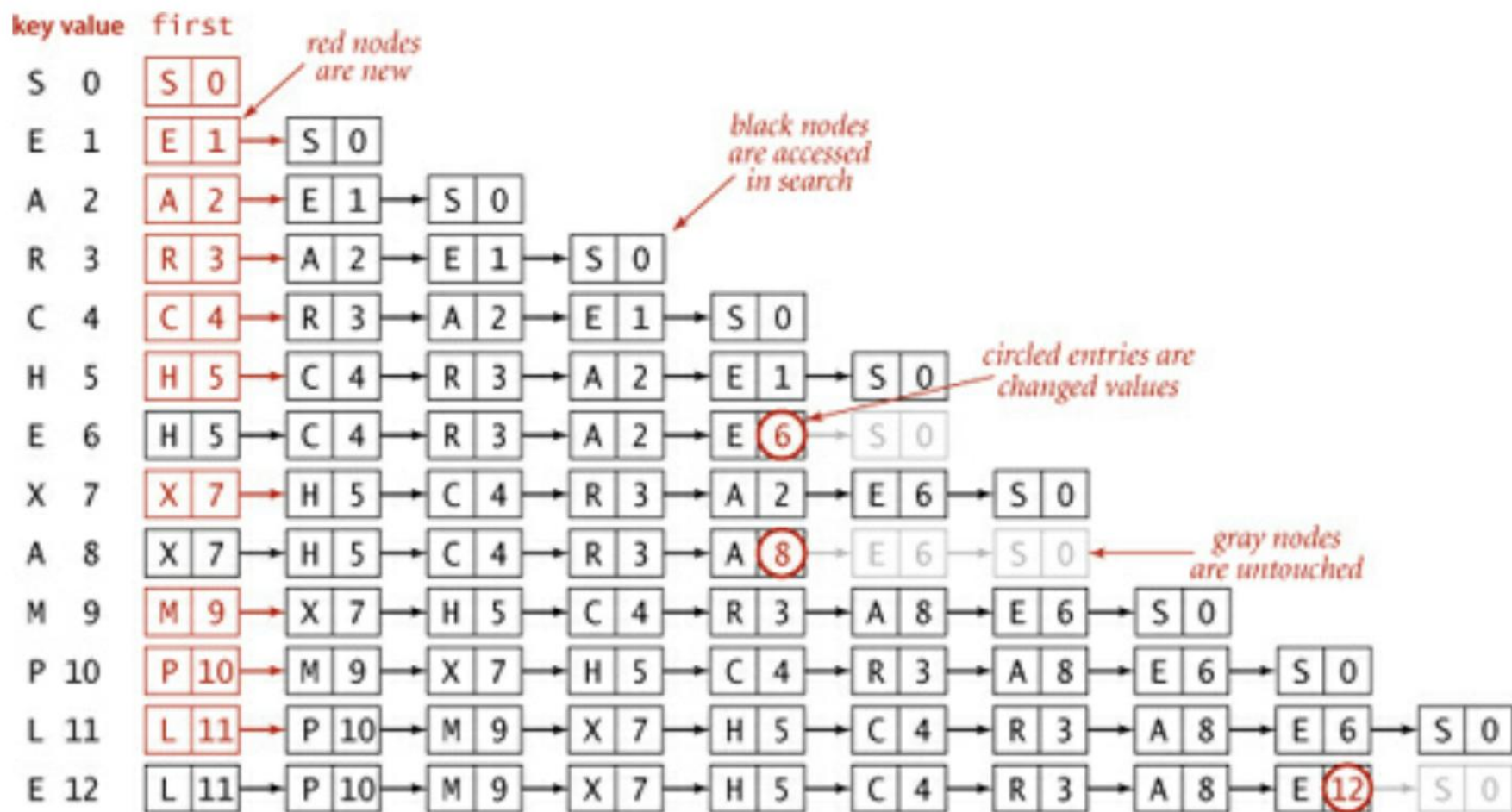
How would you implement  
a symbol table?

## Details:

- No duplicate keys (i.e. Each key can only have one value, but that value could be a list.)
- If you try to insert  $(k, v)$ , and  $k$  is already associated with  $w$ , then  $(k, w)$  is updated.

# Idea 1: Use a linked list.

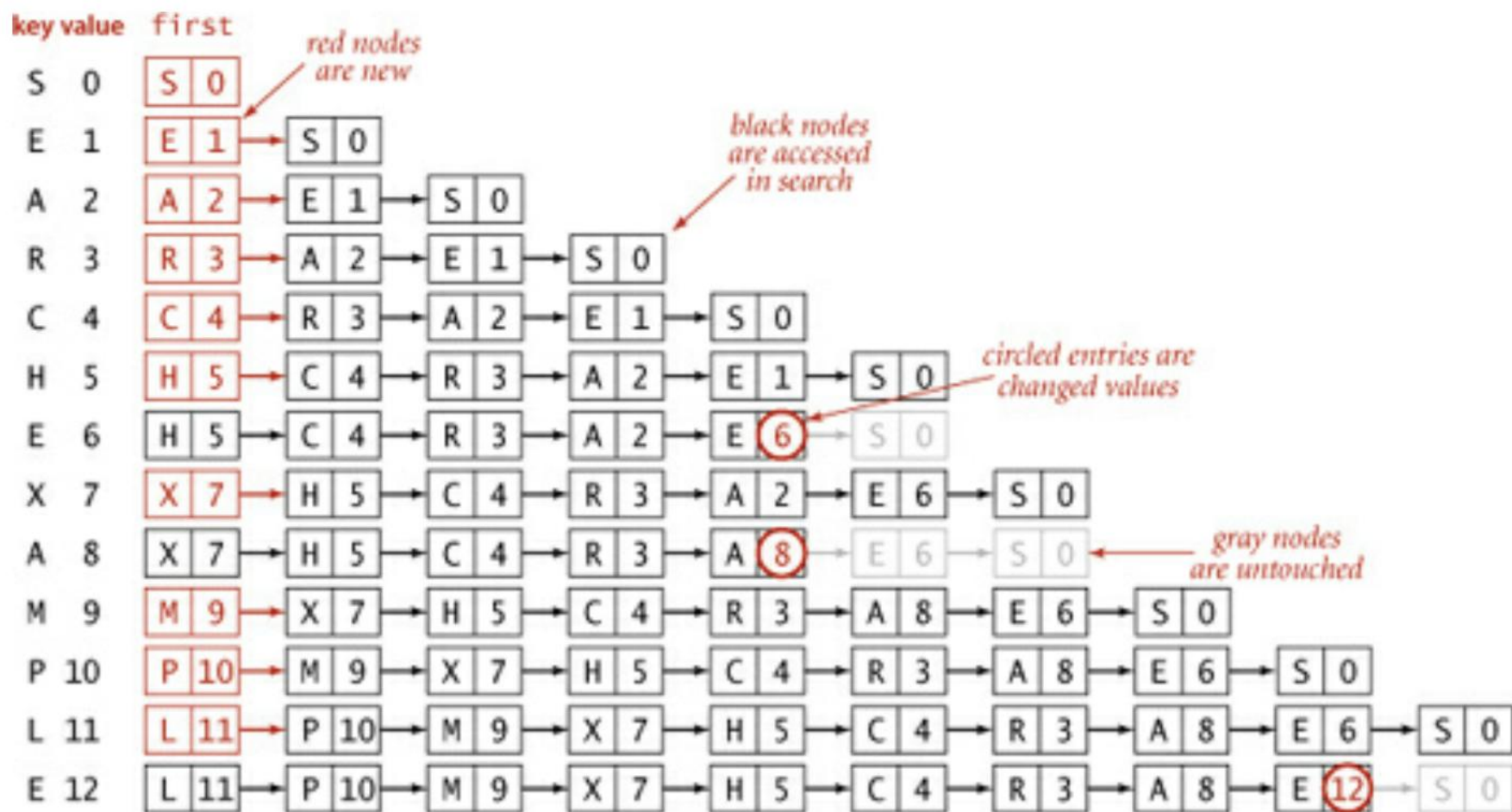




Trace of linked-list ST implementation for standard indexing client

What is the worst-case runtime for *put* and *get* if you are using a linked list and the size is  $N$ ?

$$O(N)$$



Trace of linked-list ST implementation for standard indexing client



Idea 2: Use an array (or a pair of arrays).

What are some questions to consider with this approach?

# Questions to consider regarding Idea 2...

- How will you manage the size when an array has a fixed capacity?
- What if you have millions of key-value pairs?
- How is the data organized?
- How are keys associated with a location in the array?

keys[]vals[]

key	value	0	1	2	3	4	5	6	7	8	9	N	0	1	2	3	4	5	6	7	8	9
S	0	S										1	0									
E	1	E	S									2	1	0								
A	2	A	E	S								3	2	1	0							
R	3	A	E	R	S							4	2	1	3	0						
C	4	A	C	E	R	S						5	2	4	1	3	0					
H	5	A	C	E	H	R	S					6	2	4	1	5	3	0				
E	6	A	C	E	H	R	S					6	2	4	6	5	3	0				
X	7	A	C	E	H	R	S	X				7	2	4	6	5	3	0	7			
A	8	A	C	E	H	R	S	X				7	8	4	6	5	3	0	7			
M	9	A	C	E	H	M	R	S	X			8	8	4	6	5	9	3	0	7		
P	10	A	C	E	H	M	P	R	S	X		9	8	4	6	5	9	10	3	0	7	
L	11	A	C	E	H	L	M	P	R	S	X	10	8	4	6	5	11	9	10	3	0	7
E	12	A	C	E	H	L	M	P	R	S	X	10	8	4	12	5	11	9	10	3	0	7
		A	C	E	H	L	M	P	R	S	X		8	4	12	5	11	9	10	3	0	7

entries in red were inserted

entries in gray did not move

entries in black moved to the right

circled entries are changed values

Trace of ordered-array ST implementation for standard indexing client

## Ordered Symbol Table:

Keeps the Keys in order

How would this affect the runtime for the basic operations?

How would you implement this?

public class <b>ST</b> <Key extends Comparable<Key>, Value>		
	<b>ST()</b>	<i>create an ordered symbol table</i>
	<b>void put</b> (Key key, Value val)	<i>put key-value pair into the table (remove key from table if value is null)</i>
	<b>Value get</b> (Key key)	<i>value paired with key (null if key is absent)</i>
	<b>void delete</b> (Key key)	<i>remove key (and its value) from table</i>
	<b>boolean contains</b> (Key key)	<i>is there a value paired with key?</i>
	<b>boolean isEmpty</b> ()	<i>is the table empty?</i>
	<b>int size</b> ()	<i>number of key-value pairs</i>
	<b>Key min</b> ()	<i>smallest key</i>
	<b>Key max</b> ()	<i>largest key</i>
	<b>Key floor</b> (Key key)	<i>largest key less than or equal to key</i>
	<b>Key ceiling</b> (Key key)	<i>smallest key greater than or equal to key</i>
	<b>int rank</b> (Key key)	<i>number of keys less than key</i>
	<b>Key select</b> (int k)	<i>key of rank k</i>
	<b>void deleteMin</b> ()	<i>delete smallest key</i>
	<b>void deleteMax</b> ()	<i>delete largest key</i>
	<b>int size</b> (Key lo, Key hi)	<i>number of keys in [lo..hi]</i>
	<b>Iterable&lt;Key&gt; keys</b> (Key lo, Key hi)	<i>keys in [lo..hi], in sorted order</i>
	<b>Iterable&lt;Key&gt; keys</b> ()	<i>all keys in the table, in sorted order</i>

API for a generic ordered symbol table

# Analysis Summary

method      order of growth  
of running time

put()	$N$
get()	$\log N$
delete()	$N$
contains()	$\log N$
size()	1
min()	1
max()	1
floor()	$\log N$
ceiling()	$\log N$
rank()	$\log N$
select()	1
deleteMin()	$N$
deleteMax()	1

algorithm (data structure)	worst-case cost (after $N$ inserts)		average-case cost (after $N$ random inserts)		efficiently support ordered operations?
	search	insert	search hit	insert	
<i>sequential search (unordered linked list)</i>	$N$	$N$	$N/2$	$N$	no
<i>binary search (ordered array)</i>	$\lg N$	$2N$	$\lg N$	$N$	yes

Cost summary for basic symbol-table implementations

Which one is BETTER?  
Could we improve upon it?

BinarySearchST costs

# References

[1] *Algorithms, Fourth Edition*; Robert Sedgewick and Kevin Wayne (and associated slides)