CSE 307 – Midterm Exam 2 Version 1

Question 1

Write the following programs in SML.

a. Mergesort:

```
fun take(L) = if L = nil then nil else hd(L)::skip(tl(L)) and skip(L) = if L=nil then nil else take(tl(L));

fun merge([],M) = M | merge(L,[]) = L | merge(x::xl,y::yl) = if (x:int)

fun sort(L) = if L=[] then [] else if tl(L)=[] then L else merge(sort(take(L)),sort(skip(L)));

c. Reverse a list. (2 points)
```

```
fun reverse(L) = if L=[] then []
else reverse(tl(L)) @ [hd(L)];

Alternative solution:
fun reverseHelper(L,Acc) = if L=[] then Acc
else reverseHelper(tl(L),hd(L)::Acc);
fun reverse(L) = reverseHelper(L,[]);
```

d. Find out whether a list is a palindrome. A palindrome can be read forward or backward; e.g. [r,a,c,e,c,a,r]. (4 points)

```
fun palindrome(L) = L=reverse(L);

palindrome(["r","a","c","e","c","a","r"]);

% Alternative solution
fun last(L) = if L=[] then ""
else if tl(L)=[] then hd(L)
else last(tl(L));

fun removeLast(L) = if L=[] then []
else if tl(L)=[] then []
else hd(L)::removeLast(tl(L));

removeLast([2,3,4,5]);

fun palindrome2(L) = if L=[] then true else if tl(L)=[] then true else if hd(L)=last(L) then palindrome2(removeLast(tl(L))) else false;
```

a. Describe in English the language defined by the regular expression. a*(b a*b a*)*

(2 points)

The set of all strings of as and bs containing an even number of bs.

b. Write an unambiguous context-free grammar that generates the language above.

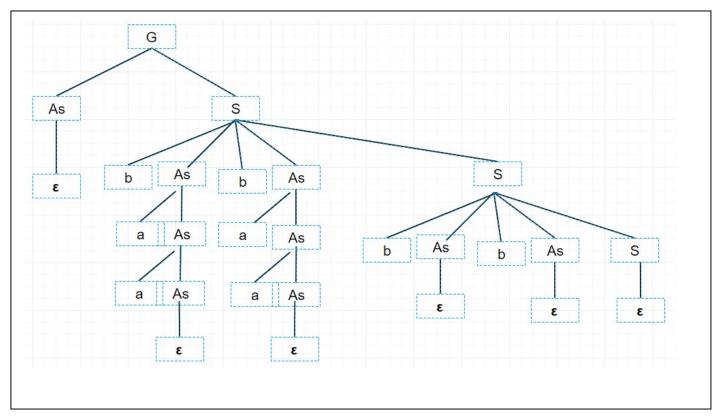
(3 points)

```
G -> As S
S -> b As b As S
| ε
As -> a As
| ε
```

c. Using your grammar derive the parse tree for the following string.

(3 points)

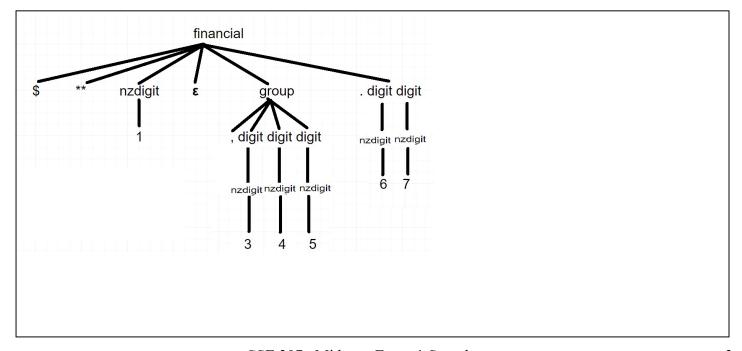
"baabaabb".



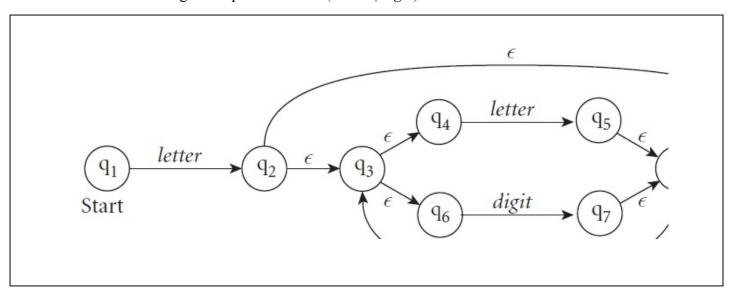
Ouestion 3

- a. Write a grammar to recognize financial expressions as defined below. (4 pts) Financial quantities in American notation have:
 - a leading dollar sign (\$),
 - an optional string of leading asterisks (*—used on checks to discourage fraud),
 - a string of decimal digits, and an <u>optional</u> fractional part consisting of a decimal point (.) and <u>exactly</u> two decimal digits.
 - o the string of digits to the left of the decimal point may consist of a single zero (0). Otherwise $i\underline{t}$ must not start with a zero.
 - o If there are more than three digits to the left of the decimal point, groups of three (counting from the right) must be separated by commas (,).

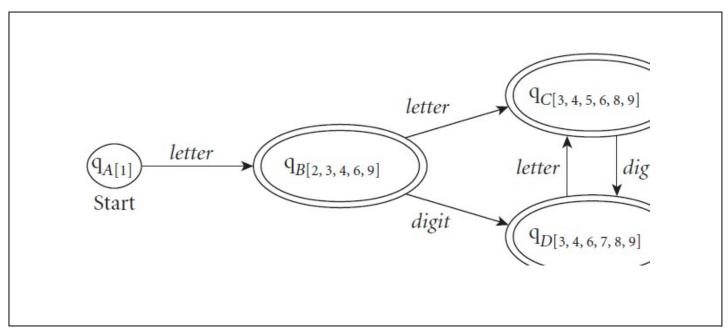
b. Use your grammar to derive the parse tree for the string "\$**2,345.67". (4 pts)



a. Build the NFA for the regular expression letter (letter | digit)*.



b. Convert the NFA you constructed to create an equivalent DFA.



c. Consider the following grammar. Describe in English the language the grammar generates. (4 points)

The grammar generates all strings of a's and b's (terminated by an end marker), in which there are more a's than b's.

a. Consider the following pseudocode:

```
x : integer — global
procedure set x(n : integer)
       x := n
procedure print x()
       write integer(x)
procedure first()
       set x(1)
       print x()
procedure second()
       x:integer
       set x(2)
       print x()
set x(0)
first()
print x()
second()
print x()
```

```
What does this program print if the language uses static scoping? (2 pts)
```

With static scoping it prints 1 1 2 2. The

What does it print with dynamic scoping? (2 pts)

With dynamic scoping it prints 1 1 2 1.

Why?
$$(1 pt)$$

The difference lies in whether set x sees the global x or the x declared in second when it is called from second.

```
b. Consider the following pseudocode:
```

```
int x = 5
int y = 6
procedure add() {
       x := x * y
       print x
       print y
procedure second(procedure P) {
        int x := 3
        call P()
       print x
       print y
procedure first() {
       int y = 4
        call second(add)
       print x
       print y
call first()
print x
print y
```

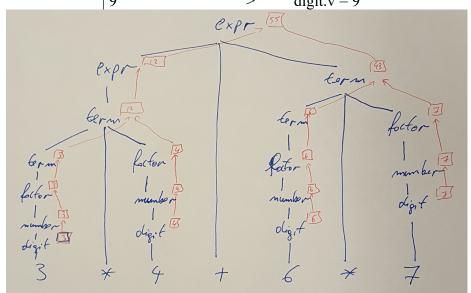
```
What does this program print if the language uses static scoping? (2 pts)
```

What does it print with dynamic scoping? (2 pts)

Why?
$$(1 pt)$$

Add attributed rules for the following CFG grammar and show the annotated parse tree for "3 * 4 + 6 * 7":

```
expr -> expr + term
                                        expr1.v = expr2.v + term.v
        term
                                        expr.v = term.v
                                        term1.v = term2.v * factor.v
term -> term * factor
        | factor
                                        term.v = factor.v
factor -> ( expr )
                                        factor.v = expr.v
        number
                                        factor.v = number.v
                                        number1.v = number2.v * 10 + digit.v
number -> number digit
        digit
                                        number.v = digit.v
                               >
digit \rightarrow 0
                                        digit.v = 0
                                        digit.v = 1
                                        digit.v = 2
         2
         3
                                        digit.v = 3
         4
                                        digit.v = 4
                                        digit.v = 5
         5
         6
                                        digit.v = 6
                                        digit.v = 7
         7
                                       \overline{\text{digit.v}} = 8
         8
         9
                               >
                                        digit.v = 9
```



```
tokens = (
    'NAME','NUMBER',
    'PLUS', 'MINUS', 'TIMES', 'DIVIDE', 'EQUALS',
    'LPAREN', 'RPAREN',
    )
# Tokens
t PLUS = r' + '
t MINUS = \mathbf{r'}-'
t TIMES = r' \ *'
t DIVIDE = r'/'
t EQUALS = r'='
t LPAREN = r' \setminus ('
t RPAREN = r' \setminus '
          = r'[a-zA-Z][a-zA-Z0-9]*'
t NAME
def t NUMBER(t):
    r' \setminus d+'
    try:
        t.value = int(t.value)
    except ValueError:
        print("Integer value too large %d", t.value)
        t.value = 0
    return t
# Ignored characters
t ignore = " \t"
def t newline(t):
    r' \setminus n + '
    t.lexer.lineno += t.value.count("\n")
def t error(t):
    print("Illegal character '%s'" % t.value[0])
    t.lexer.skip(1)
# Build the lexer
import ply.lex as lex
lexer = lex.lex()
lexer.input("1+2")
while True:
    tok = lexer.token()
    if not tok:
        break
    print(tok)
lexer.input("abc=123")
while True:
```

```
tok = lexer.token()
    if not tok:
        break
    print (tok)
# Parsing rules
precedence = (
    ('left', 'PLUS', 'MINUS'),
    ('left','TIMES','DIVIDE'),
    ('right', 'UMINUS'),
# dictionary of names
names = { }
def p statement assign(t):
    'statement : NAME EQUALS expression'
    names[t[1]] = t[3]
def p_statement_expr(t):
    'statement : expression'
    print(t[1])
def p expression binop(t):
    '''expression : expression PLUS expression
                  | expression MINUS expression
                  | expression TIMES expression
                  | expression DIVIDE expression'''
    if t[2] == '+' : t[0] = t[1] + t[3]
    elif t[2] == '-': t[0] = t[1] - t[3]
    elif t[2] == '*': t[0] = t[1] * t[3]
    elif t[2] == '/': t[0] = t[1] / t[3]
def p expression uminus(t):
    'expression : MINUS expression %prec UMINUS'
    t[0] = -t[2]
def p expression group(t):
    'expression : LPAREN expression RPAREN'
    t[0] = t[2]
def p expression number(t):
    'expression : NUMBER'
    t[0] = t[1]
def p expression name(t):
    'expression : NAME'
    try:
        t[0] = names[t[1]]
    except LookupError:
        print("Undefined name '%s'" % t[1])
        t[0] = 0
def p error(t):
```

```
print("Syntax error at '%s'" % t.value)

import ply.yacc as yacc
yacc.yacc()

while 1:
    try:
        s = input('calc > ')  # Use raw_input on Python 2
    except EOFError:
        break
    yacc.parse(s)
```