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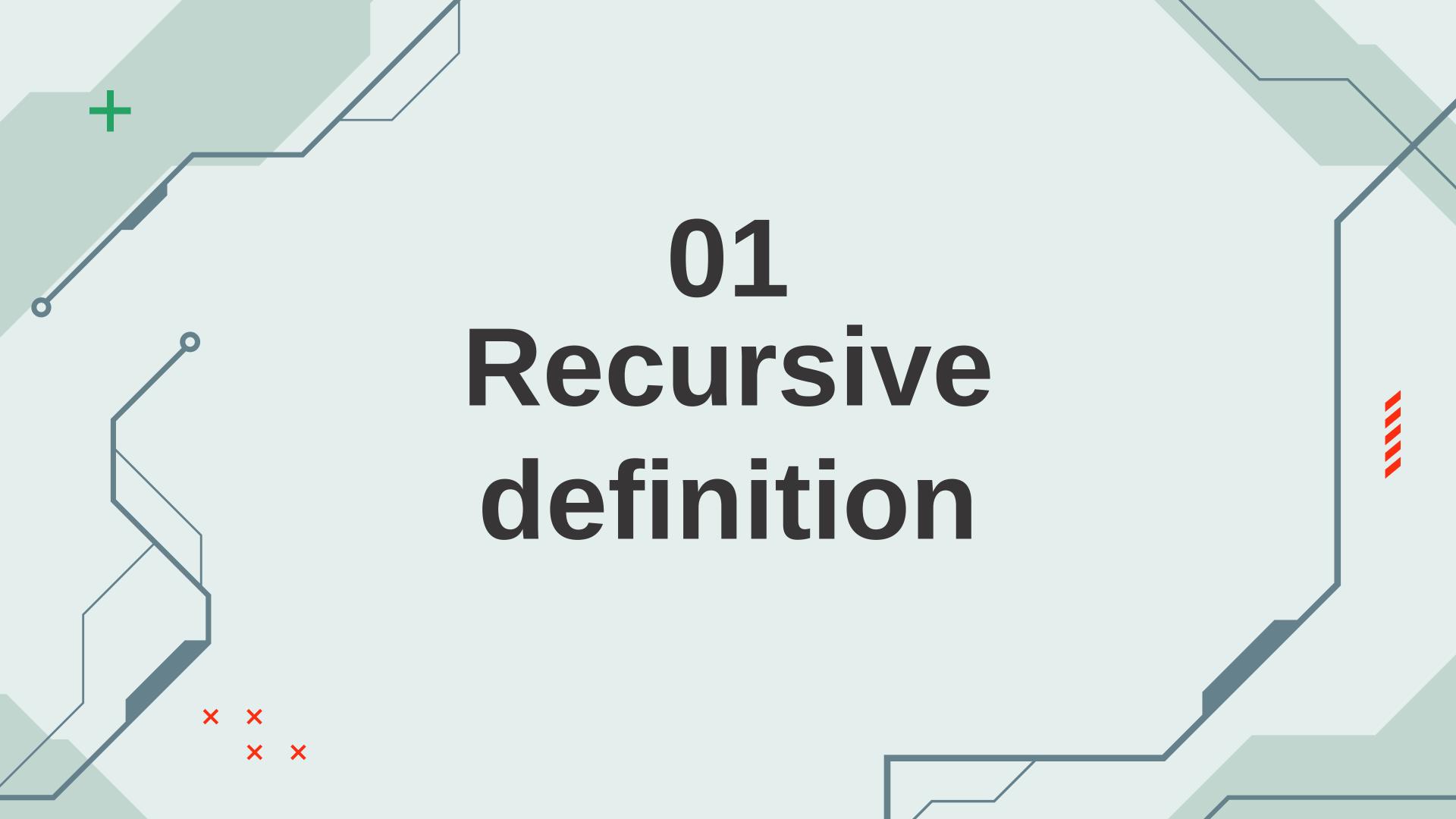
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Subproblem graph of P(14)

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Dynamic
Programming

O4
Compare P(14) with different weights



Recursive Definition of P(C)

C: capacity i = 0, 1, ..., n-1

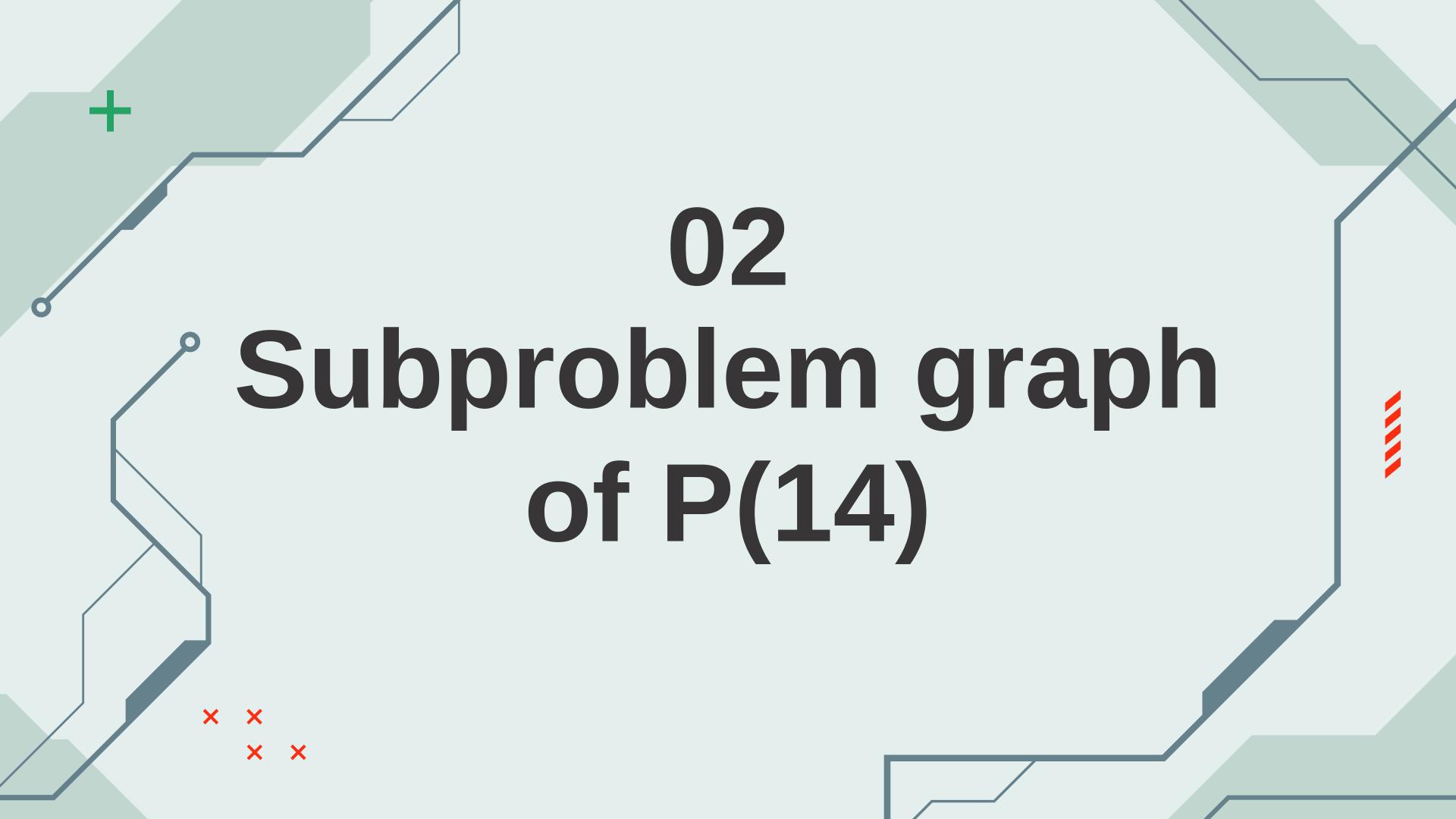
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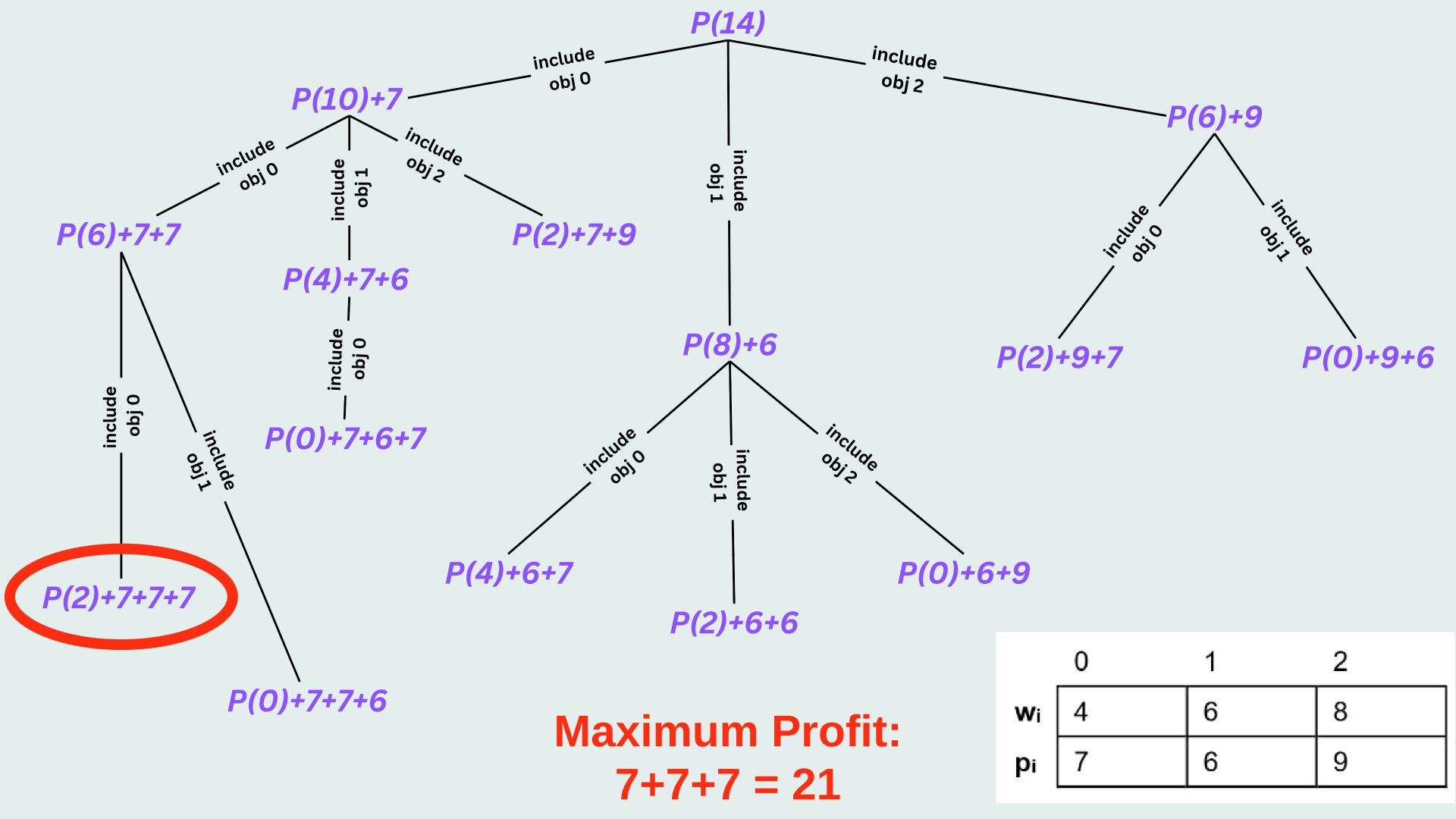
Base Case:

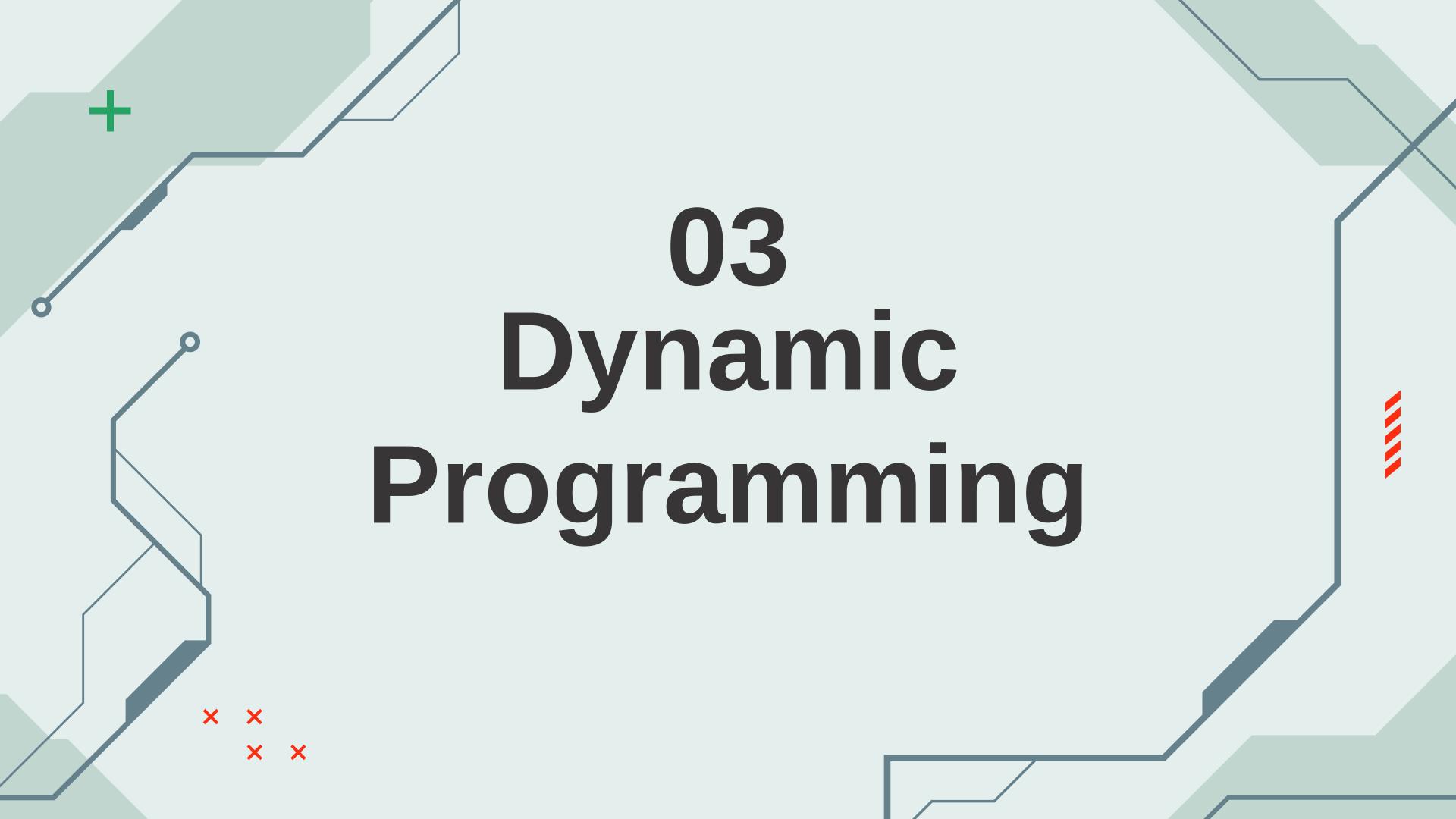
P(C)=0 when C=0 or n=0

P(C) = max(
$$P(C)$$
, $P(C-w_i) + p_i$)

ith item is excluded ith item is included







Dynamic Programming with Bottom-Up Approach:

- 1.2D DP Array: dp[i][j] stores the maximum profit for capacity j using items up to index i.
- 2. **Initialize DP Array**: Each row represents an item and each column represents a capacity from 0 to C.
- 3. Outer Loop:
- iterates over each capacity j from 1 to C.
- 4.Inner Loop:
- iterates over each item i.
- Excluding the item: If we exclude the item, dp[i][capacity] is set to dp[i 1] [capacity] if i > 0, otherwise it's 0 (base case).
- Including the item: If weights[i] <= capacity, the code calculates the profit by including the item, allowing reuse by taking dp[i][capacity - weights[i]] + profits[i]
- 5. Final Solution:
- dp[n-1][C] holds the maximum profit for capacity C using all items.

Dynamic Programming with Bottom-Up Approach:

```
37
     def max_profit(C, weights, profits):
          n = len(weights)
38
39
         # Initialize a 2D DP array where rows are items and columns are capacities
40
         dp = [0] * (C + 1) for _ in range(n)]
41
42
         # Fill the DP table
43
          for capacity in range(1, C + 1):
              for i in range(n):
44
45
                  # Option 1: Exclude the item
                  dp[i][capacity] = dp[i - 1][capacity] if i > 0 else 0
46
47
                  # Option 2: Include the item (if capacity allows)
48
                  if weights[i] <= capacity:</pre>
49
                      dp[i][capacity] = max(dp[i][capacity], dp[i][capacity - weights[i]] + profits[i])
50
51
52
         # The maximum profit with the given capacity and all items is in dp[n-1][C]
53
          return dp[n - 1][C]
```

The code allows item reuse by repeatedly updating dp[i][capacity] with dp[i][capacity-weights[i]] + profits[i] whenever an item fits.

Time Complexity

Outer Loop (Capacities):

The outer loop iterates over each capacity from 1 to C, so it runs C times.

Inner Loop (Items):

For each capacity, the inner loop iterates over all n items.

In Each Iteration:

Within the inner loop, the program checks two options (excluding and including the item) and performs a constant-time update of the dp array.

Overall Time Complexity:

Since the outer loop runs C times and the inner loop runs n times for each capacity, the total time complexity is: **O(n×C)**

Space Complexity

DP Array:

The 2D dp array has n rows (one for each item) and C + 1 columns (one for each capacity from 0 to C), resulting in a space complexity of: $O(n \times C)$

vs Recursive Solution (Without Memoization)

- try each item at every capacity level, making recursive calls for each item that can fit within the remaining capacity.
- The recursion continues until the remaining capacity is exhausted (or until there are no items left to consider).

Time Complexity

For each capacity C, we have n choices (one for each item) if it fits. For each choice, the function calls itself recursively with a smaller capacity (C - weight[i]).

Branching Factor: **At each step, the function can branch out n times (one for each item), and it can potentially reach down to a depth of C** (if each call reduces capacity by 1 in the worst case). Worst-Case Time Complexity: **O(n^C), which is exponential**.

Space Complexity

Recursive Call Stack: The depth of the recursive calls depends on the capacity C. worst case: each call reduces the capacity by the smallest weight, leading to a maximum depth of C. -> O(C)

Code for running the DP algo and printing part 4 results

```
def max_profit(C, weights, profits):
         n = len(weights)
         # Initialize a 2D DP array where rows are items and columns are capacities
         dp = [0] * (C + 1) for _ in range(n)
         # Fill the DP table
          for capacity in range(1, C + 1):
              for i in range(n):
 8
                 # Option 1: Exclude the item
                 dp[i][capacity] = dp[i - 1][capacity] if i > 0 else 0
10
11
                 # Option 2: Include the item (if capacity allows)
12
                 if weights[i] <= capacity:</pre>
13
                      dp[i][capacity] = max(dp[i][capacity], dp[i][capacity - weights[i]] + profits[i])
14
15
16
         # Print the DP table with labels
         print("DP Table (Items x Capacity):")
17
                   ", " ".join([f"C={j:2}" for j in range(C + 1)]))
         print("
18
                     ", "-" * (5 * (C + 1)))
         print("
19
20
21
          for i in range(n):
              print(f"Item {i} |", " ".join(f"{dp[i][capacity]:4}" for capacity in range(C + 1)))
22
23
         # The maximum profit with the given capacity and all items is in dp[n-1][C]
24
          return dp[n - 1][C]
25
26
     # Test with the data
27
     C = 14
     weights = [5, 6, 8]
     profits = [7, 6, 9]
     max_profit = max_profit(C, weights, profits)
     print("\nMaximum profit for P(14) with weights [5, 6, 8] and profits [7, 6, 9] is:", max_profit)
```

04 Results using Dynamic Programming

0	1	2
4	6	8
7	6	9

DP Tal		e (Item C= 0 C=			The state of the s	1 C-	5 C- (,	8 C-	a c-	10 C-	11 C-1	12 C-1	3 C-1	14
						4 (- 			/ C- 		9 C					
Item	0	0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
Item :	1	0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
Item :	2	j 0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
		•														
Maxim	um	profit	for F	(14)	with	weigh	ts [4	, 6,	8] ar	nd pro	ofits	[7, 6	5, 9]	is: 2	21	

Maximum Profit = 21

0	1	2
4	6	8
7	6	9

DP Tabl	e (Ite C= 0 C				4 C=	5 C=	6 C=	7 C=	8 C=	9 C=1	L0 C=:	11 C=1	L2 C=1	L3 C=:	14
Item 0 Item 1 Item 2	0 0 0	 0 0 0	 0 0 0	0 0 0	7 7 7	 7 7 7	7 7 7 7	 7 7 7	 14 14 14	 14 14 14	 14 14 14	14 14 14 14	21 21 21 21	21 21 21 21	 21 21 21
Maximum	profi	t for I	P(14)	with v	yeigh	ts [4	, 6,_	8] ar	nd pro	ofits	[7, 6	5, 9]	is: 2	21	

Maximum Profit = 21

at C=4, item 0 will be placed in the bag bag = [0]

0	1	2
4	6	8
7	6	9

DP Ta							o Nex				Market silve		PROPERTY COLUMN		estature Parcela A		
		C= 0	C=	1 C=	2 C=	3 C=	4 C=	5 C=	6 C=	7 C=	8 C=	9 C=:	10 C=	11 C=:	12 C=:	13 C=:	L4
Item	0	 	0 0		 0	 0	 7	7	 7	7	14	 14 14	 14	 14	 21	21	 21
Item	1	İ	0	0	0	0	7	7 7 7	7	7	14	14	14	14	21	21	21
Item	2	İ	0	0	0	0	7	7	7	/7/	14	14	14	14	21	21	21
Maxim	um	prof	fit	for P	(14)	with	weigh	nts [4	6,	81 ar	nd pr	ofits	[7, (6, 9]	is: 2	21	

Maximum Profit = 21

at C=8, another item 0
will be placed in the bag
bag = [0,0]

0	1	2
4	6	8
7	9	9

DP Ta						The state of the s	4 C=	5 C= 6	C=	7 C=	8 C=	9 C=:	10 C=:	11 C=	12 C=:	13 C=:	14
Item Item		•	 0 a	0 0		 0 0				 7 7		 14 14	14 14		 21 21	 21 21	 21 21
Item		1	0					7				14	14		21	21	21
Maxim	num	prof	it fo	or P	(14)	with	weigh	nts [4,	6,	8] aı	nd pr	ofits	[7,	6, 9]	is: 2	21	

рi

Maximum Profit = 21

at C=10,

if we consider including item 1

bag = [0,0]: profit = 7+7=14

bag = [0,1]: profit = 7+6=13<14

0	1	2
4	6	8
7	6	9

```
DP Table (Items x Capacity):
       C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14
                                        7 7 14
7 7 14
7 7
                                                                          21
                                                                     14
                                                          14
                                                               14
                                                                               21
Item 0
                                                                                    21
                                                                               21
Item 1
                                                          14
                                                               14
                                                                     14
                                                                                    21
                                                          14
                                                               14
                                                                     14
                                                                                    21
Item 2
Maximum profit for P(14) with weights [4, 6, 8] and profits [7, 6, 9]/ is: 21
```

Maximum Profit = 21

at C=12, another item 0 will be placed in the bag bag = [0,0,0]

Wi

рi

0	1	2
4	6	8
7	6	9

```
DP Table (Items x Capacity):
       C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14
Item 0
                                                   14
                                                         14
                                                              14
                                                                   14
                                                                        21
                                                                             21
                                                                                  21
                              7 7 7 7 14
7 7 7 7 14
                                                                                  21
                                                        14
                                                              14
                                                                   14
                                                                        21
                                                                             21
Item 1
                                                                        21
                                                         14
                                                              14
                                                                   14
                                                                             21
                                                                                  21
Item 2
Maximum profit for P(14) with weights [4, 6, 8] and profits [7, 6, 9]/is: 21
```

Maximum Profit = 21

at C=12,

if we consider including item 2

bag = [0,0,0]: profit = 7+7+7=21

bag = [0,2]: profit = 7+9=15<21

0	1	2
4	6	8
7	6	9

DP Tab																
	C= (0 C= 	1 C=	2 C= 	3 C=	4 C= 	5 C= (6 C= 	7 C= 	8 C= 	9 C=: 	10 C=: 	11 C=: 	12 C=: 	L3 C=:	14
Item 0		0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
Item 1	İ	0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
Item 2	Ì	0	0	0	0	7	7	7	7	14	14	14	14	21	21	21
Maximur	n pro	ofit	for F	P(14)	with	weigh	ts [4	, 6,	8] ar	nd pro	ofits	[7, (6 , 9]	is:	21	

Maximum Profit = 21-

Wi

рi

0	1	2
5	6	8
7	6	9

Maximum Profit = 16

Wi

рi

0	1	2
5	6	8
7	6	9

```
DP Table (Items x Capacity):

C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14

Item 0 | 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 1 | 0 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 2 | 0 0 0 0 0 0 7 7 7 9 9 14 14 14 16 16

Maximum profit for P(14) with weights [5, 6, 8] and profits [7, 6, 9] is: 16
```

Maximum Profit = 16

at C=5, item 0 will be placed in the bag bag = [0]

Wi

рi

0	1	2
5	6	8
7	6	9

```
DP Table (Items x Capacity):

C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14

Item 0 | 0 0 0 0 0 7 7 7 7 7 7 14 14 14 14 14

Item 1 | 0 0 0 0 0 7 7 7 7 7 7 14 14 14 14 14

Item 2 | 0 0 0 0 0 7 7 7 7 7 9 9 14 14 14 16 16

Maximum profit for P(14) with weights [5, 6, 8] and profits [7, 6, 9] is: 16
```

at C=8, only item 0
is in the bag
bag = [0]

Maximum Profit = 16

at C=8, consider up till item 2

Wi

рi

0	1	2
5	6	8
7	6	9

```
DP Table (Items x Capacity):

C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14

Item 0 | 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 1 | 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 2 | 0 0 0 0 0 7 7 7 7 9 9 14 14 14 16 16

Maximum profit for P(14) with weights [5, 6, 8] and profits [7, 6, 9] is: 16
```

Maximum Profit = 16

at C=10, another item 0 is added to the bag bag = [0, 0] --> 7+7=14

at C=10, consider **up till item 2**

Wi

рi

0	1	2
5	6	8
7	6	9

Maximum Profit = 16

at C=11, consider **item 1**

W

рi

0	1	2
5	6	8
7	6	9

Maximum Profit = 16

at C=13, another item 0 is added to the bag bag = [0, 0] --> 7+7=14 at C=13, consider **up till item 2**

Wi

рi

0	1	2
5	6	8
7	6	9

```
DP Table (Items x Capacity):

C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14

Item 0 | 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 1 | 0 0 0 0 0 7 7 7 7 7 14 14 14 14 14

Item 2 | 0 0 0 0 0 7 7 7 7 9 9 14 14 14 16 16

Maximum profit for P(14) with weights [5, 6, 8] and profits [7, 6, 9] is: 16
```

Maximum Profit = 16

Conclusion

Dynamic Programming improved the overall efficiency

time complexity (DP): O(Cn) vs time complexity(recursive function): O(noc) space complexity (DP): O(Cn) using 2D array

to further improve the programming, we can implement the 1D array, we can improve the **space complexity to O(C)** but it decreases the readability.

qusing 2D array

```
DP Table (Items x Capacity):

C= 0 C= 1 C= 2 C= 3 C= 4 C= 5 C= 6 C= 7 C= 8 C= 9 C=10 C=11 C=12 C=13 C=14

Item 0 | 0 0 0 0 7 7 7 7 14 14 14 14 21 21 21

Item 1 | 0 0 0 0 7 7 7 7 14 14 14 14 14 21 21 21

Item 2 | 0 0 0 0 7 7 7 7 14 14 14 14 14 21 21 21

Maximum profit for P(14) with weights [4, 6, 8] and profits [7, 6, 9] is: 21
```

