

## Laboratory practice No. 2: Big O Notation

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### 3) Practice for final project defense presentation

#### 3.1 Table for exercise number 1:

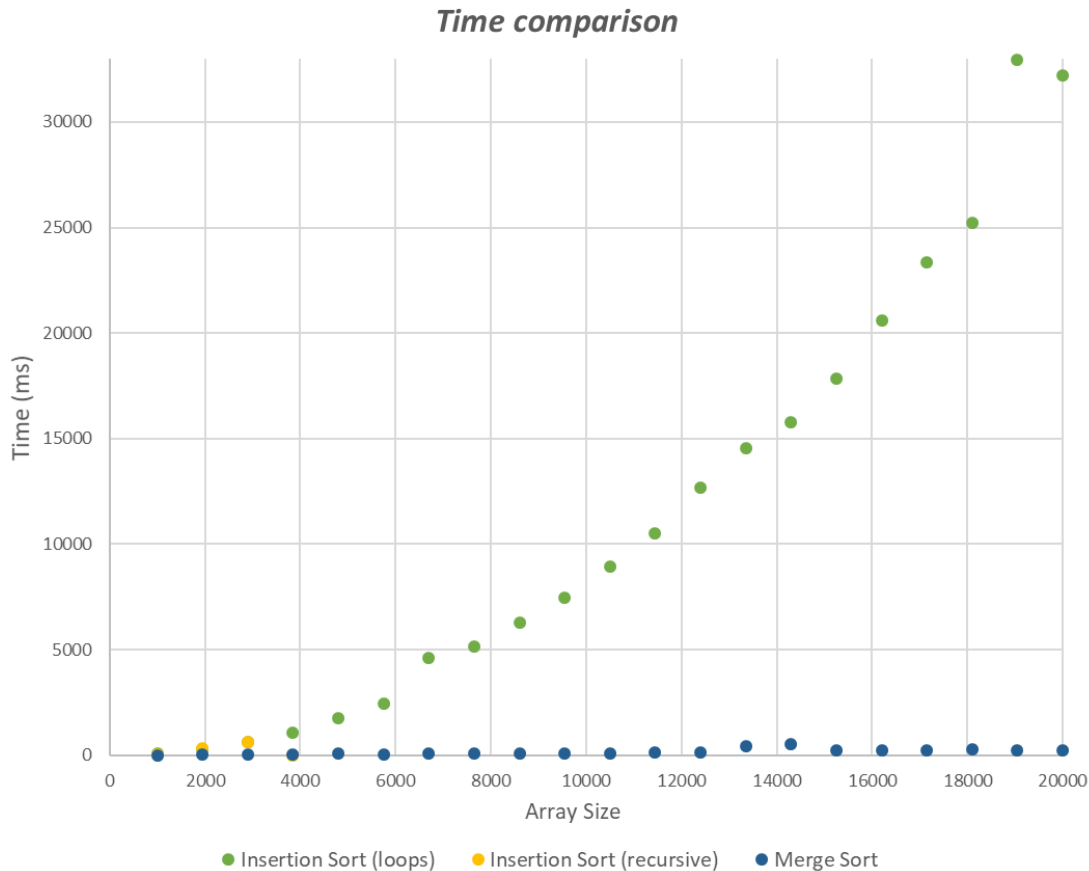
| Insertion Sort |           | Insertion Sort (Recursive) |  | Merge Sort |           |
|----------------|-----------|----------------------------|--|------------|-----------|
| Array size     | Time (ms) | Array size                 | Time (ms)  | Array size | Time (ms) |
| 1000           | 83.10     | 1000                       | 62.41  | 1000       | 10.53     |
| 1950           | 293.46    | 1950                       | 340.75   | 1950       | 34.80     |
| 2900           | 649.24    | 2900                       | 640.26   | 2900       | 40.23     |
| 3850           | 1090.22   | 3850                       | The recursive version of insertion sort doesn't run on our computer for arrays with more than 3932 elements. | 3850       | 64.16     |
| 4800           | 1757.80   | 4800                       |  | 4800       | 68.98     |
| 5750           | 2469.03   | 5750                       |  | 5750       | 63.00     |
| 6700           | 4598.43   | 6700                       |  | 6700       | 66.19     |
| 7650           | 5175.22   | 7650                       |  | 7650       | 79.04     |
| 8600           | 6294.73   | 8600                       |  | 8600       | 84.09     |
| 9550           | 7473.03   | 9550                       |  | 9550       | 102.95    |
| 10500          | 8945.84   | 10500                      |  | 10500      | 112.90    |
| 11450          | 10528.21  | 11450                      |  | 11450      | 126.40    |
| 12400          | 12666.18  | 12400                      |  | 12400      | 136.46    |
| 13350          | 14573.18  | 13350                      |  | 13350      | 432.49    |
| 14300          | 15798.98  | 14300                      |  | 14300      | 529.95    |
| 15250          | 17822.61  | 15250                      |  | 15250      | 229.79    |
| 16200          | 20622.56  | 16200                      |  | 16200      | 217.08    |
| 17150          | 23331.42  | 17150                      |  | 17150      | 221.44    |
| 18100          | 25228.31  | 18100                      |  | 18100      | 261.38    |
| 19050          | 32966.89  | 19050                      |  | 19050      | 227.21    |
| 20000          | 32225.10  | 20000                      |  | 20000      | 251.17    |

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## ESTRUCTURA DE DATOS 1

### Código ST0245

#### 3.2 Chart for exercise number 1:



**3.3** After the previous test, it is easily seen that merge sort method is more useful than insertion sort. The former doesn't work on bigger problems, and the only alternative is to use a different version with loops, which takes more time than the normal type.

**3.4** The use of Insertion Sort in videogames with millions of elements wouldn't be optimal. For starters, the recursive algorithm won't work with bigger numbers; likewise, the algorithms that uses loops takes too much time when running with bigger arrays. A better solution for this would be to use a method such as Merge Sort.

**3.5** Insertion sort is only a better algorithm when it runs with arrays that are already (or almost) sorted. The reason is that this process doesn't do any operation in these cases; on the other hand, merge sort always divides and groups the array, even when it is sorted.

**3.6** MaxSpan's objective is to find how many elements are between two other elements which have the same value. The one in the left is going to be called as "the leftmost element" and the one in the right is going to be called "the rightmost element", so we are going to called "Span" the number of elements between the leftmost element and the rightmost element, and we want to know which one is the largest Span of any given array.

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## ESTRUCTURA DE DATOS 1

### Código ST0245

For example, if the given array is an array with just one value, the span will be 1, another thing we must consider is that we need to include in our count also the rightmost and the leftmost element.

Now the algorithm works like this: we created a variable called “m” which is going to store the Span, then we make a nested loop, the outside loop is for the leftmost element and the inside loop is for the rightmost element, in both loops the condition is that the element has to be less than the array length, next we introduce a conditional if to compare the equality between the leftmost element and the rightmost element and also the rightmost element minus the leftmost element plus 1 have to be more than our variable m, in case that the previous conditional sentence be true, the variable m will be equal to the following operation : rightmost element – leftmost element + 1

This program will be continue doing the previous mentioned steps until the element in the leftmost position be more or equal to the length of the array and then the program will return the variable m, which will be the largest Span and this is the answer to the problem.

### 3.7 Complexity for online exercises:

| EXERCISE         | LEVEL | COMPLEXITY(T(N)) | COMPLEXITY(BIG(O)) |
|------------------|-------|------------------|--------------------|
| 1. CountEvens    | 2     | $T(n) = n + C$   | $O(n)$             |
| 2. Sum13         | 2     | $T(r) = r + C$   | $O(r)$             |
| 3. Only14        | 2     | $T(n) = n$       | $O(n)$             |
| 4. is Everywhere | 2     | $T(n) = n + C$   | $O(n)$             |
| 5. Fizz Buzz     | 2     | $T(k) = k + C$   | $O(k)$             |
| 1. maxSpan       | 3     | $T(m) = m^2 + C$ | $O(m^2)$           |
| 2. Fix34         | 3     | $T(t) = t^2 + C$ | $O(t^2)$           |
| 3. can Balance   | 3     | $T(n) = 2n + c$  | $O(n)$             |
| 4. squareUp      | 3     | $T(s) = s^2 + C$ | $O(s^2)$           |
| 5. maxMirror     | 3     | $T(l) = l^2 + C$ | $O(l^2)$           |

### 3.8

| VARIABLE | DEFINITION  |
|----------|---|
| n        | Array for which the algorithm will count the even numbers, or also Split it |
| r        | Sum of the array.   |
| m        | The largest span between two numbers.                                       |
| k        | Difference between the end and the start of the array.                      |
| l        | Length of the array.  |
| s        | Multiplication of the array parameter by itself.                            |
| t        | Temporal variable which stored the position of the array plus one.          |

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#### 4) Practice for midterms

4.1 b

4.2 b

4.3 b

4.4 b

4.5 d

4.6 a

4.7

4.7.1  $T(n) = T(n-1) + c$

4.7.2  $O(n) = n$

4.8 a

4.9 d

4.11 c

4.12 b

4.13 a

4.14 a

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