



# Cache Replacement Based on Reuse-Distance Prediction

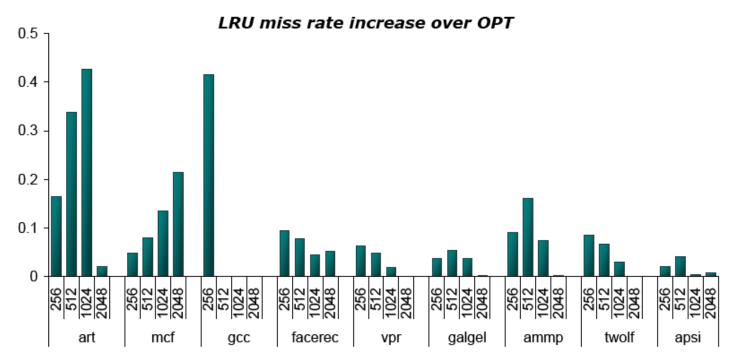
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#### Introduction

- We need caches
  - The memory wall rises more and more

- Contributions:
  - Instruction based reuse distance prediction
  - Potential for cache level optimizations
  - Case study: a replacement algorithm for second level caches

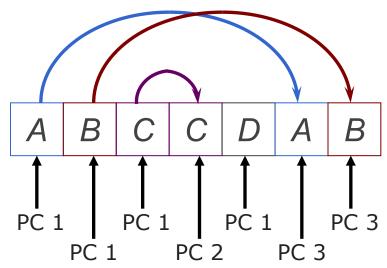
## LRU vs Belady's Optimal Replacement



- LRU: Inefficient for L2 caches
- Reasons:
  - L1 filtering
  - Highly associative caches

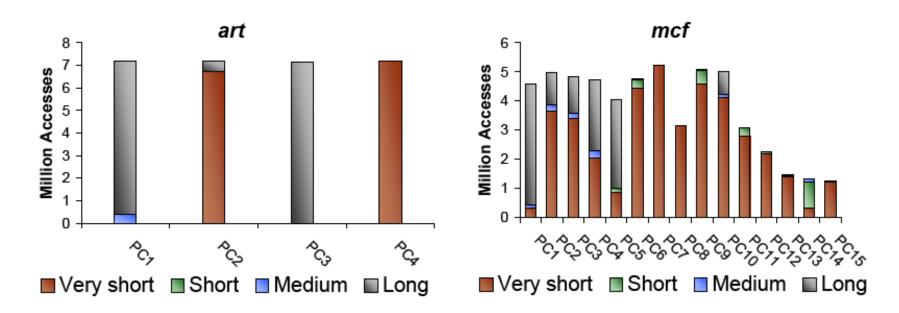
#### Can we see the future?

- Yes, via prediction
- Memory behavior → repeating patterns
- Our Motivation:
  - Strong correlation between instructions (PC) and moment of next access (reuse distance)



Reuse Distance of A = 4 Reuse Distance of B = 4 Reuse Distance of C = 0

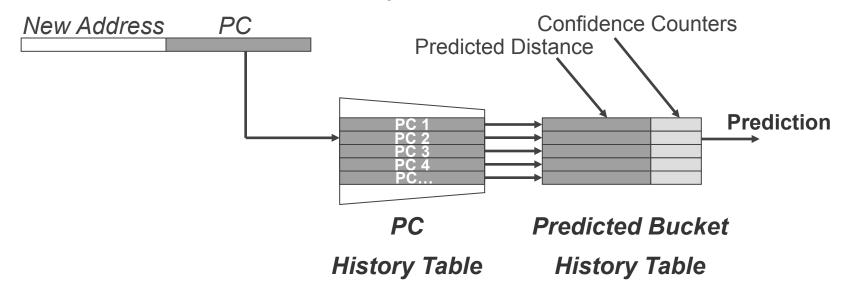
#### Reuse distances are predictable



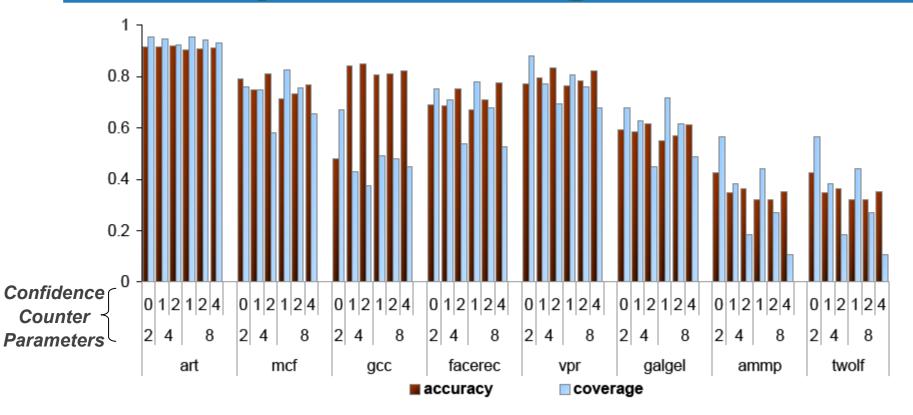
- Significant PCs:
  - Are few → Easily stored
  - Have predictable reuse distances

#### **Instruction based Predictor**

- Predicts the reuse distance of an access based on the PC which initiated the access
- Tracks the accessed line
- Upon reuse calculates the reuse distance
- Updates the predictor entry associated with the PC



#### **Accuracy and Coverage**

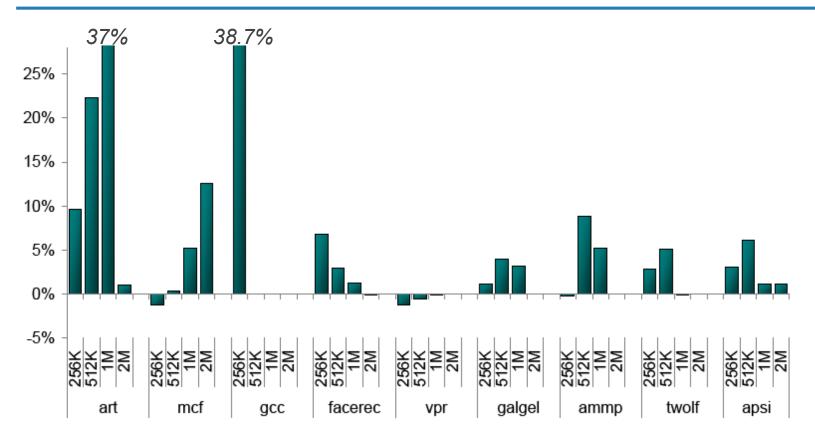


- Not bad, but also not perfect
  - The concept works
  - Even better implementations are possible

#### A case study: replacement policy

- Not all accesses are predictable
  - We cannot implement a perfect optimal algorithm
- Hybrid algorithm:
  - Replace the line used farthest in the future (OPT) when you have enough info
  - Replace the line used farthest in the past (LRU) if you don't.
- More reliable predictor → policy closer to OPT

#### Miss rate reduction

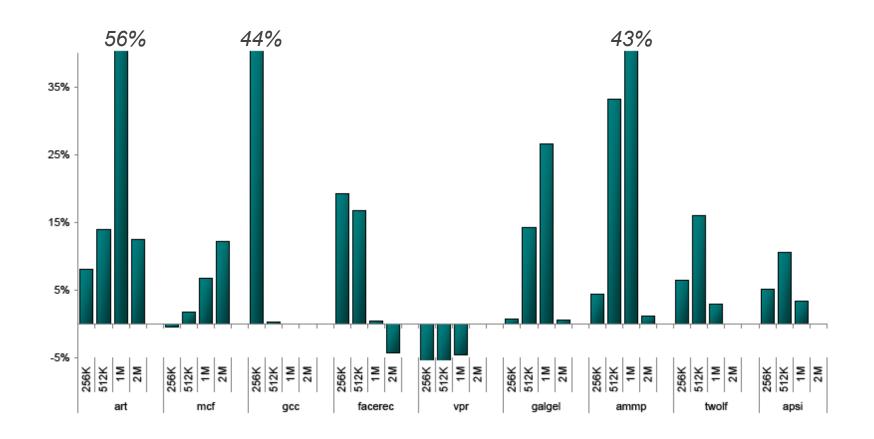


- Significant improvement
- Benchmarks with low accuracy are not greatly affected

#### **Conclusions and Future Work**

- Instruction based reuse distance predictor
  - Proof of concept
  - Low overhead (< 2%)</li>
  - Good accuracy
- Replacement policy for L2 caches
  - Speed ups in many benchmarks
  - 60% in art
  - 43% in ammp
  - Slowdown in vpr
- Future Work:
  - Improve the effectiveness of the predictor
  - More cache level optimizations

### Speedup (IPC)



# Dead Block Prediction vs Reuse Distance Prediction

#### **Reuse Distance Prediction**

- Predicts when the line is going to be accessed again
- Fine grained prediction
- Can identify which line is less needed
  - Very large RDs indicate dead lines
- More information can be extracted
- Proof of concept predictor

#### **Dead Block Prediction**

- Predicts when the line is going to be evicted
- Coarse grained prediction
- Can identify only dead lines
  - Not which is less needed

- Less information is stored
- Sophisticated predictor