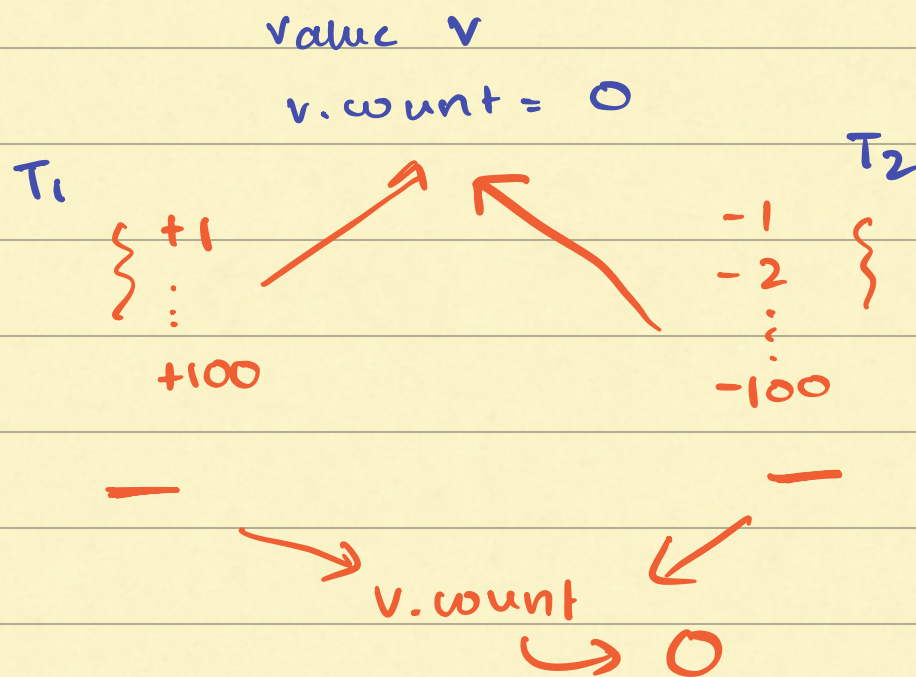


## Agenda

- Why synchronisation error happens!
- Mutex
- synchronised keyword
- synchronised method
- Producer/Consumer



1. Critical Section → any part of code dealing with shared memory or data is called CS.

① `cout ("Hello");`

② `v.count += 1;`

③ `cout ("v = " + v);`

Critical Section

① `cout ("Hello");`

② `v.count -= 1;`

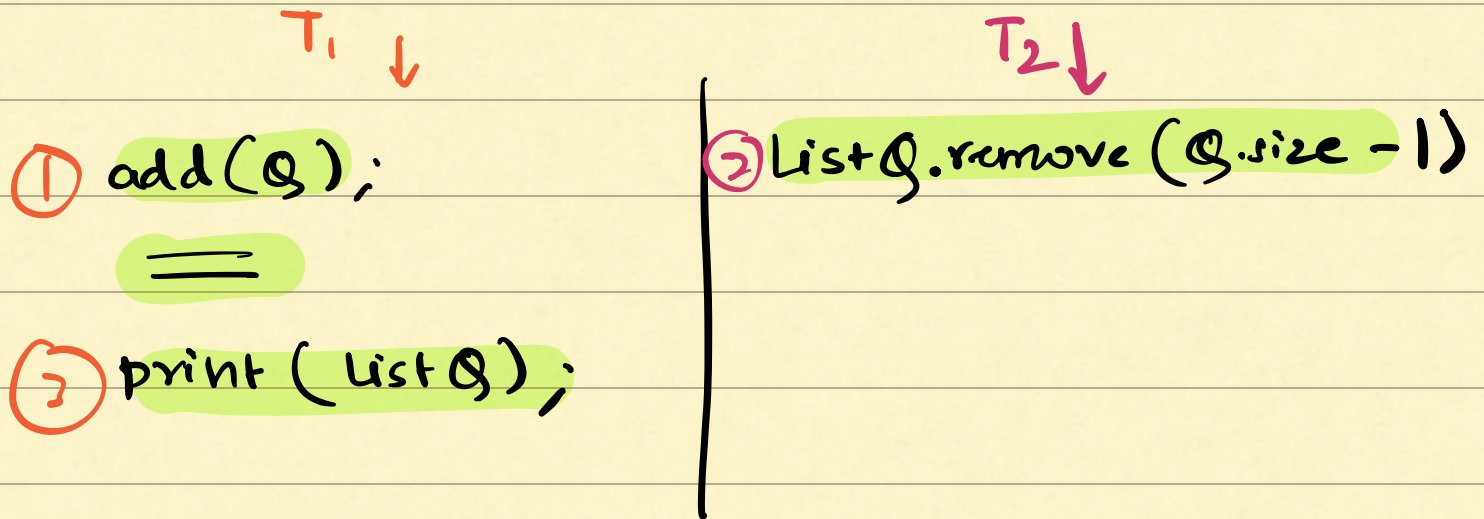
③ `cout ("Hello2");`

④ `v.count *= 1;`

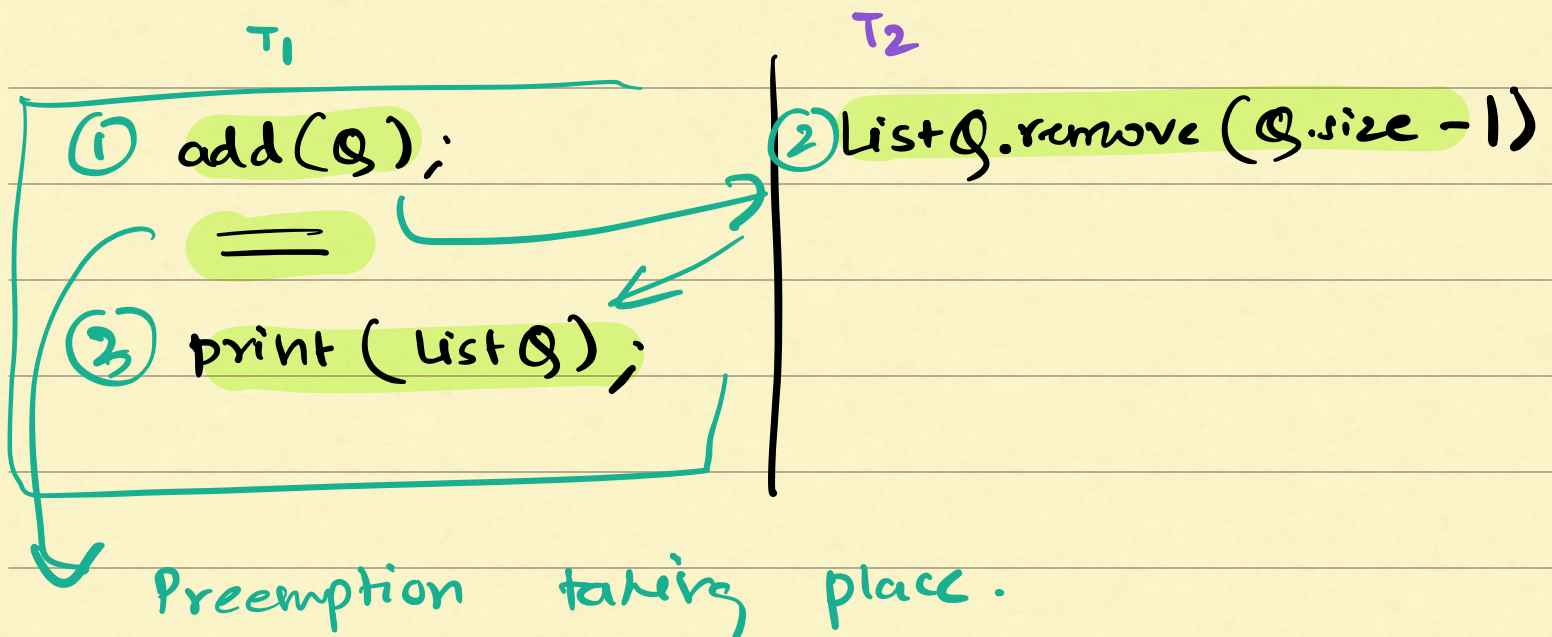
⑤ `cout ("Hello3");`

Critical Section

2. Race Condition → when two threads enter the same critical section at the same time.



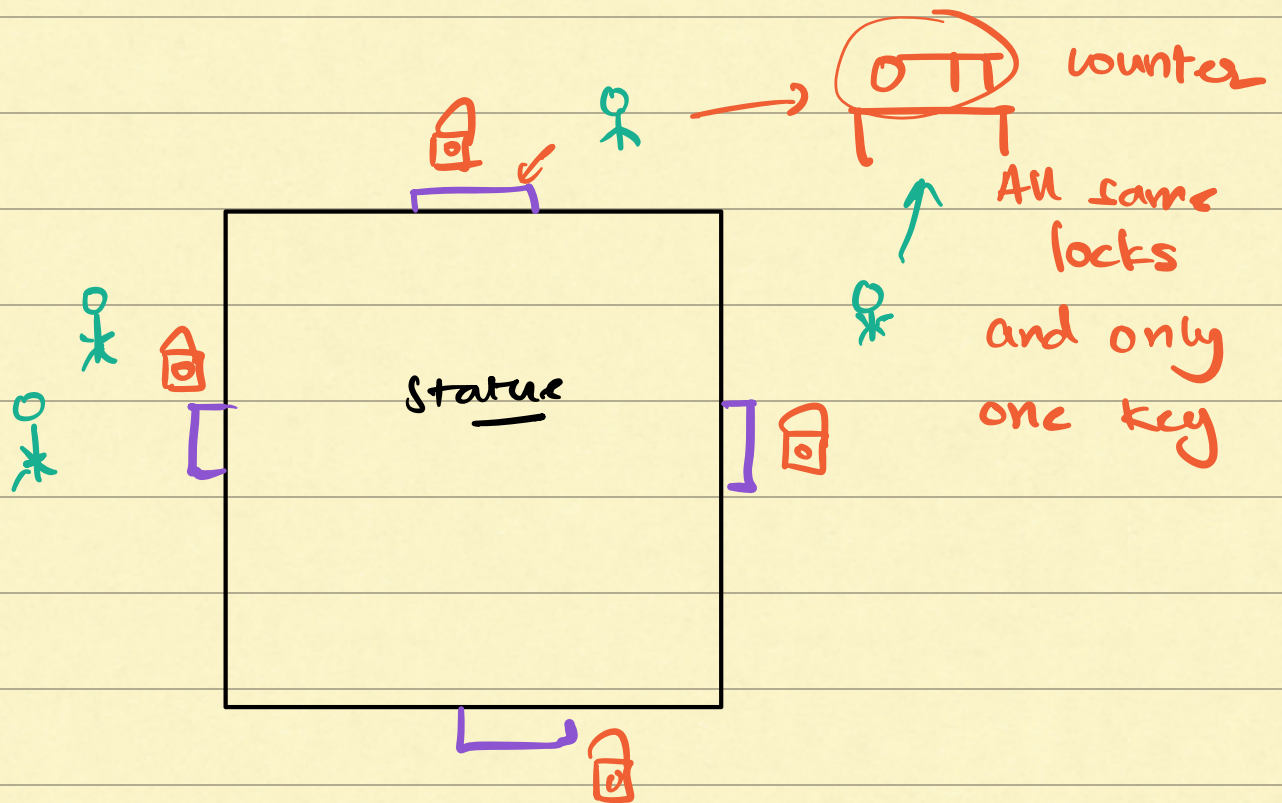
3. Preemptive ness: Ending a task before completion



CONTEXT SWITCHING



Mutex → Mutual Exclusion



Adder

$T_1$  → Lock  
 lock. take Lock C);  $(T_1 - 1)$   
 $v. count += i;$

lock. release Lock C);

=====

Lock

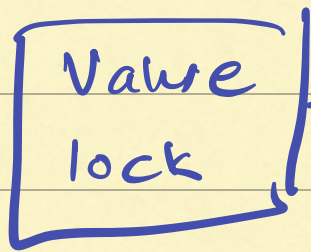
Subtractor

$T_2$  → Lock  
 lock. take Lock C);  
 $v. count -= i;$   
 lock. release Lock C);

=====

→  $(T_1 - 2)$

# Synchronised keyword



Synchronised block.

## Synchronised methods :

Adder {

synchronised add (int i) { lock.lock();

} lock.unlock();

void subtract ( ) {

}

synchronised void multiply ( ) {

Adder a = new

Adder b = new

T <sub>1</sub>	T <sub>2</sub>
a.add	a.add

✓

✗ wait

T <sub>1</sub>	T <sub>2</sub>
a.add	a.multiply

✓

✗ wait

T <sub>1</sub>	T <sub>2</sub>
a.add	a.subtract

✓

✓

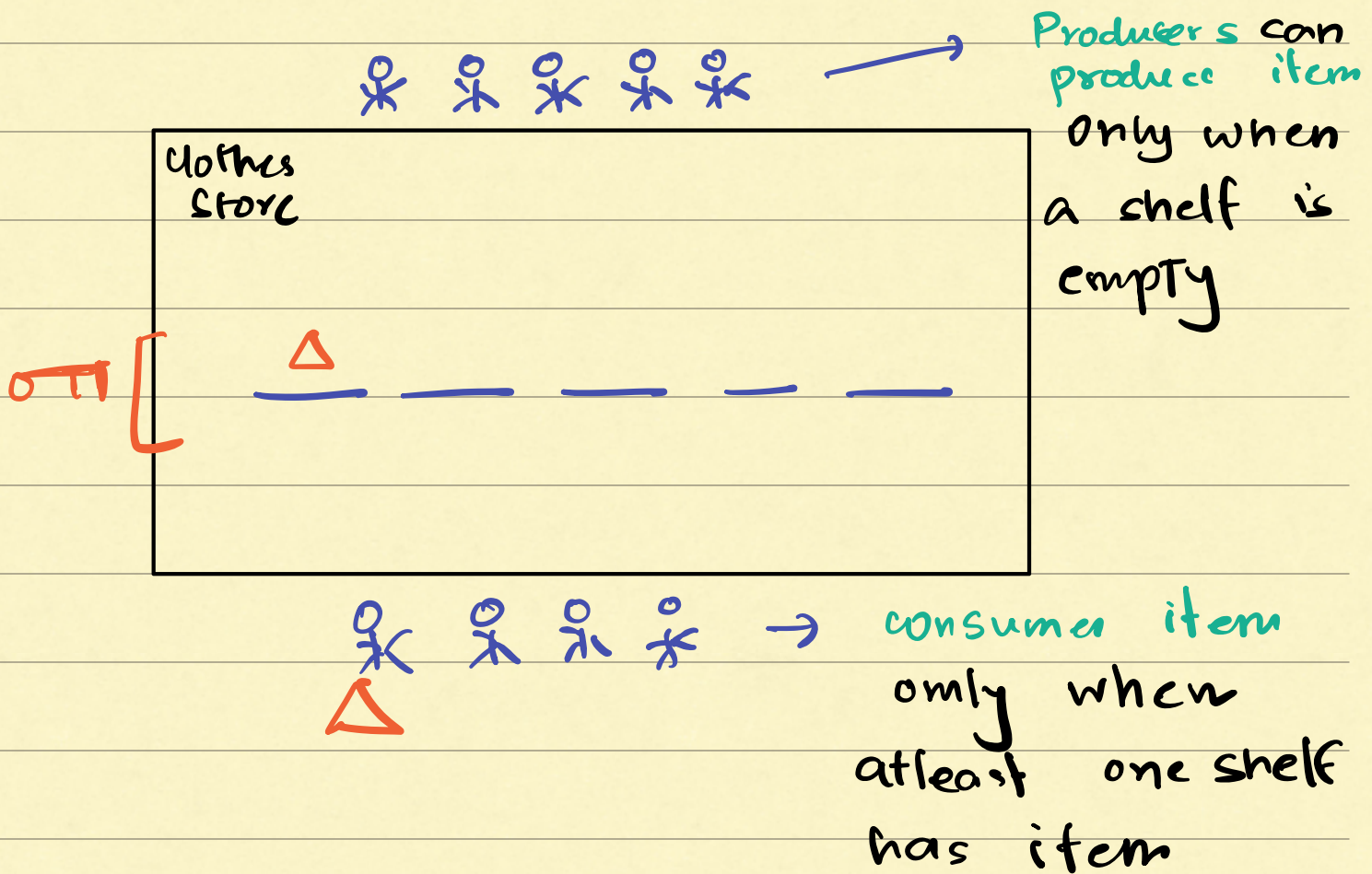
T <sub>1</sub>	T <sub>2</sub>
----------------	----------------

a.add b.add



## Producer/consumer

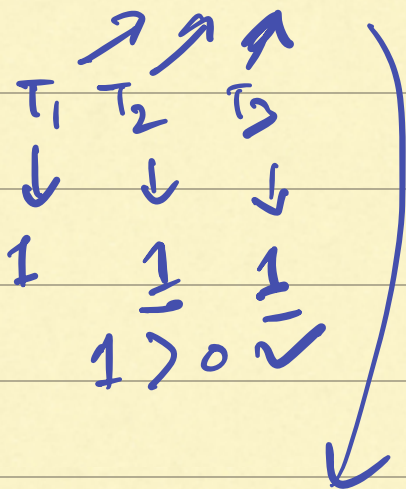
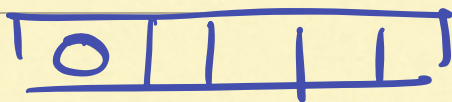
↳ is restricting 1 thread to  
work on a CS always beneficial



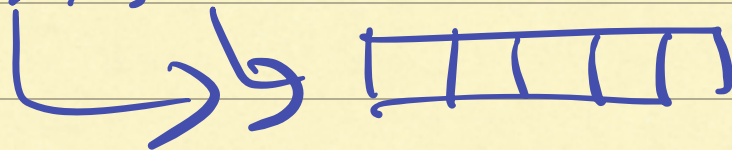


## Producer

$T_1, T_2, T_3$  store size = 3  
↓ ↓ ↓  
if ( store.size() < store.maxsize )  
{  
    store.add(new Object());  
}  
  
store size → 6



$T_1, T_2, T_3 \rightarrow T_1$  removes successfully.



## Consumer

$T_1, T_2, T_3$  store size = 4  
↓ ↓ ↓  
if ( stor.size() > 0 ) {  
    stor.remove();  
}  
  
 $T_2, T_3$  will throw outofindex error.

synchronization issues

↳ Semaphores