

Today's Agenda:-

- Modular Arithmetic Introduction
- Count pairs whose sum mod m is 0
- Introduction to GCD
- Properties of GCD
- Delete One

$A \% B$ \rightarrow remainder when A is divided by B.

$\rightarrow [0 \ B-1]$

$x \% 6 \rightarrow [0 \ 5]$

$30 \% 7 \rightarrow 2$

$40 \% 9 \rightarrow 4$

$5 \% 5 \rightarrow 0$

$1 \% 2 \rightarrow$

$$\begin{array}{r} 2 \overline{) 10} \\ \underline{0} \\ 10 \end{array}$$

$30 \% 7 = 30 - 7 = 23 - 7 = 16 - 7 = 9 - 7 \Rightarrow \underline{2}$

$A \% B \Rightarrow$ keep subtracting B from A,
till $A < B$.

Why do we need mod

$-\infty$

$-\infty$

$\% 10 \rightarrow [0 \text{ to } 9]$

To limit our range.

Divisor) Dividend (Quotient
 remainder

Remainder = Dividend - greatest

multiple of divisor \leq dividend

$$\underline{30} \% 7 \Rightarrow 30 - 28 \Rightarrow \underline{2}$$

$$\underline{61} \% 5 \Rightarrow 61 - 60 \Rightarrow \underline{1}$$

$$\underline{-7} \% 9 \Rightarrow -7 - (-9) \Rightarrow \underline{2}$$

$$\underline{-30} \% 7 \Rightarrow -30 - (-35) \Rightarrow \underline{5}$$

$$\rightarrow -1 + 9$$

$$\rightarrow -2 + 7$$

$$a \% m \rightarrow (a + m) \text{ if } a < 0$$

Rules of Modular arithmetic

$$(0 \text{ } 2m-1) \% m \rightarrow 0m-1$$

$$1 \quad (a+b) \% m \rightarrow (\underbrace{a \% m}_{(0 \text{ } m-1)} + \underbrace{b \% m}_{(0 \text{ } m-1)}) \% m$$

$$a=9, \quad b=8, \quad m=5$$

LHS

$$(9+8) \% 5 \Rightarrow 17 \% 5 \Rightarrow 2$$

RHS

$$\underbrace{a \% m}_{(0 \text{ } m-1)} + \underbrace{b \% m}_{(0 \text{ } m-1)} \% m$$

$$(9+8) \% 5$$

$$\Rightarrow \underline{2}$$

Let's say our

no. range is till 10 only.

$$2) (a * b) \% m \rightarrow (a \% m * b \% m) \% m.$$

$$3) (a + m) \% m \rightarrow \underline{a \% m}$$

$$\downarrow$$

$$(a \% m + m \% m) \% m$$

↓

$$(a \% m) \% m$$

$$4) (a - b) \% m \Rightarrow (a \% m - b \% m + m) \% m.$$

$$\underline{a=10, b=8, m=9}$$

LHS.

$$(10 - 8) \% 9$$

$$\downarrow 2 \% 9 \Rightarrow 2$$

RHS

$$(10 \% 9 - 8 \% 9) \% 9$$

$$(1 - 8) \% 9$$

$$\downarrow$$

$$-7 \% 9$$

$$\downarrow$$

$$-7 + 9 = 2$$

$$5) a^b \% m = (a \% m)^b \% m$$

↓

$$(a * a * a * a * \dots * a) \% m$$

b times

$$\rightarrow (a \% m * a \% m * \dots * a \% m) \% m$$

b times

Ques

$$(37^{10^9} - 1) \cdot 12$$

$$(37^{10^9} \cdot 12 - 1 \cdot 12 + 12) \cdot m$$

$$\left(\underbrace{(37 \cdot 12)}^{10^9} \cdot 12 - 1 + 12 \right) \cdot m$$

$$\downarrow$$
$$(1^{10^9} \cdot 12 - 1 + 12) \cdot 12$$

$$(1 \cdot 12 - 1 + 12) \cdot 12$$

$$(1 - 1 + 12) \cdot 12$$

$$(0 + 12) \cdot 12$$

$$\Rightarrow 12 \cdot 12 \Rightarrow 0$$

Ques

Given n array elements,

find pairs (i, j) s.t.,

$$(arr[i] + arr[j]) \% m = 0, \quad i \neq j,$$

$$arr = \{4, 9, 6, 3, 8, 12\}$$

$$m = 6.$$

Brute Force Approach :-

Two for loops.

Checking all pairs.

$$T.C \rightarrow O(n^2).$$

idea 2

$$(a+b) \% m = \underline{0}$$

↓

$$(a \% m + b \% m) \% m = \underline{0}$$

$$[0 \dots m-1] \quad [0 \dots m-1]$$

↓

↓

$$[1 + (m-1)] \% m$$

$$[2 + (m-2)] \% m$$

$$[3 + (m-3)] \% m$$

$$[0 + 0] \% m$$

$$\text{arr}[12] = \{ 6, 7, 5, 11, 19, 20, 9, 15, 14, 13, 12, 23 \}$$

$m = 5$

↓

$$\text{arr}[12] \% 5 = \{ 1, 2, 0, 1, 4, 0, 4, 0, 4, 3, 2, 3 \}$$

remainder freq

0	→	3
1	→	2
2	→	2
3	→	2
4	→	3

$$3 \times \frac{(3-1)}{2} \Rightarrow$$

$$\frac{n(n-1)}{2}$$

$$n C_2$$

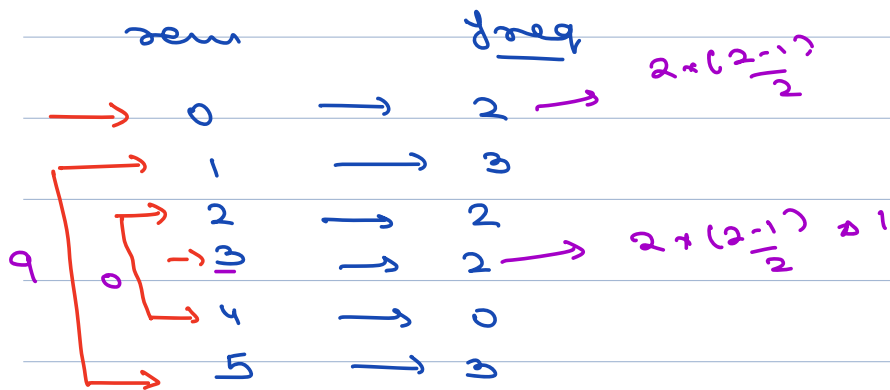
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no. of ways to
pick 2 items
out of n

arr[12] :- { 0 1 2 3 4 5 6 7 8 9 10 11
 { 6, 7, 5, 11, 19, 20, 9, 15, 14, 13, 12, 23 }

M = 6

arr[12] % 6 :- { 0 1 2 3 4 5 6 7 8 9 10 11
 { 0, 1, 5, 5, 1, 2, 3, 3, 2, 1, 0, 5 }



given arr[] & M,

hashmap <int, int> hmap;

→ Insert all elements in hmap.

↳ arr[i] % M.

C = 0;

Pairs of 0

x = hmap[i], // freq of 0

C = C + $\frac{x * (x-1)}{2}$

if (m % 2 == 0) {

x = hm[$\frac{m}{2}$]

c = c + x * ($\frac{x-1}{2}$)

}

m = 9

for (i = ¹~~1~~, i < ($\frac{m+1}{2}$); i++) {

c = c + hm[i] * hm[m-i];
hm[2] * hm[3]

}

T.C $\rightarrow O(N + M)$

S.C \rightarrow

Case - 1

N = 100

M = 10,

S.C $\rightarrow O(\underline{M})$.

Case - 2

N = 10

M = 100

S.C $\rightarrow O(\underline{N})$.

S.C $\rightarrow O(\min(N, M))$

Break 8:34am - 8:44am.

← gcd →

Maths

↳ Greatest Common divisor
↳ hcf → Highest Common factor

$$\gcd(a, b) = \underline{x},$$

↳ x is the greatest no.,

$$a \div x = 0$$

$$b \div x = 0$$

$$\gcd(15, 25) \Rightarrow 5$$

↓ ↓
1 1
3 5
5 25
15

$$\gcd(10, -25) \Rightarrow 5$$

↓ ↓
1 -25
2 -5
5 -1
10 5
25

$$\gcd(12, 30) \Rightarrow 6$$

↓ ↓
1 1
2 2
3 3
4 4
5 5
6 6
12 15
30

$$\gcd(10, 8) \Rightarrow 2$$

↓ ↓
1 1
2 2
3 3
4 4
5 5
6 6
7 7
8 8

$$\gcd(0, -10) = 10.$$

↓ ↓
1 1
0 -10
1 -5
2 -2
3 3
4 4
5 5
6 6
7 7
8 8
9 9
10

$$\gcd(0, 0) \Rightarrow 10$$

↓ ↓
1 1
0 -10
1 -5
2 -2
3 3
4 4
5 5
6 6
7 7
8 8
9 9
10

$$\gcd(-2, -3) \Rightarrow 1.$$

↓ ↓
1 1
-2 -3
1 1
2 2
3 3
4 4
5 5
6 6
7 7
8 8
9 9
10

$$\gcd(0, -5) \Rightarrow \underline{5}, \quad \gcd(0, 5) = \underline{5}.$$

\downarrow
 0
 \vdots
 2
 \vdots
 ∞

\downarrow
 -5
 \vdots
 -1
 \vdots
 5

$$\gcd(0, x) = \underline{|x|}, \quad (\underline{x \neq 0})$$

Properties of \gcd ,

$$1) \gcd(a, b) = \gcd(b, a)$$

$$2) \gcd(0, x) = \underline{|x|}.$$

$$3) \gcd(A, B, C) = \gcd(A, \gcd(B, C)) = \\ \gcd(B, \gcd(A, C)) = \\ \gcd(C, \gcd(A, B))$$

Special Property :-

$$A, B \geq 0, \quad (A \geq B)$$

$$\gcd(a, b) = \gcd(a - b, \underline{b}) \quad \checkmark$$

$$\gcd(10, 5) = \gcd(5, 5) = \gcd(0, 5)$$

Proof

$$\gcd(a, b) = \gcd(a - b, b) \quad \checkmark$$

$$\checkmark \gcd(a, b) = d$$

$$a \cdot \checkmark d = 0, \quad b \cdot \checkmark d = 0,$$

$$a = d + k_1, \quad b = d + k_2$$

$$\Rightarrow (a-b) = d(k_1 - k_2)$$

$$\checkmark (a-b) \cdot \checkmark d = 0,$$

$\gcd(a, b) = d$, & m is also
a factor of a & b .

$$d \geq m$$

$$\gcd(a-b, b) = m$$

$$(a-b) \cdot \checkmark m = 0, \quad \& \quad b \cdot \checkmark m = 0,$$

$$a-b = k_3 m, \quad \& \quad b = k_4 m$$

$$\swarrow + \searrow$$

$$a = m(k_3 + k_4)$$

$$a \cdot \checkmark m = 0 \quad \checkmark$$

$\gcd(a-b, b) = m$ & d is a
factor of $(a-b)$ & b .

$$m \geq d$$

$$m = d$$

$$\gcd(a, b) = \gcd(a-b, b).$$

$$\gcd(23, 5) = \gcd(18, 5) = \gcd(13, 5) = \gcd(8, 5)$$

↓

$$\gcd(3, 5)$$

$$\gcd(a, b) = \gcd(a-b, b)$$

↓

$$\gcd(a-2b, b)$$

↓

$$\gcd(a-3b, b)$$

⋮

$$\gcd(a-yb, b)$$

$$\Rightarrow \gcd(a \cdot \checkmark b, b)$$

Property :-

$$\gcd(a, b) = \gcd(a \div b, b) \quad \checkmark$$

$$\gcd(24, 16) = \gcd(8, 16) = \gcd(8, 16) = \gcd(8, 16)$$

- - - - -
∞

$$\gcd(a, b) = \gcd(b, a \% b) \quad \underline{\underline{\underline{\quad}}}$$

$$\gcd(\underline{24}, 16) = \gcd(16, 8) = \gcd(8, 0)$$

⌈
8

$$\gcd(14, \underline{24}) = \gcd(\underline{24}, 14) = \gcd(14, \underline{10})$$

$$= \gcd(10, 4)$$

$$= \gcd(4, 2)$$

$$= \gcd(\underline{2}, 0) \Rightarrow \underline{\underline{2 \text{ only}}}$$

$$\gcd(|a|, |b|);$$

```
int gcd(a, b) {  
    if (b == 0) return a;  
    return gcd(b, a % b);  
}
```

$$T.C \rightarrow O(\log(\max(a, b)))$$

Ques) Given arr \rightarrow , delete one element, i.e.,
gcd of remaining elements become
max.

$$\text{arr}[] = \{ \overset{0}{24} \quad \overset{1}{16} \quad \overset{2}{18} \quad \overset{3}{30} \quad \overset{4}{15} \}$$

gcd

$$\{ \overset{0}{\cancel{24}} \quad \overset{1}{16} \quad \overset{2}{18} \quad \overset{3}{30} \quad \overset{4}{15} \}$$

1

$$\text{arr}[] = \{ \overset{0}{24} \quad \overset{1}{\cancel{16}} \quad \overset{2}{18} \quad \overset{3}{30} \quad \overset{4}{15} \}$$

3

$$\text{arr}[] = \{ \overset{0}{24} \quad \overset{1}{16} \quad \overset{2}{\cancel{18}} \quad \overset{3}{30} \quad \overset{4}{15} \}$$

1

$$\text{arr}[] = \{ \overset{0}{24} \quad \overset{1}{16} \quad \overset{2}{18} \quad \overset{3}{\cancel{30}} \quad \overset{4}{15} \}$$

1

$$\text{arr}[] = \{ \overset{0}{24} \quad \overset{1}{16} \quad \overset{2}{18} \quad \overset{3}{30} \quad \overset{4}{\cancel{15}} \}$$

2

Brute force

$$n \times (n \times \log \max(a_i))$$

delete an $a[i]$ element, calculate gcd of remaining elements & get overall max.

ideq 2 :-

	0	1	2	3	4
	24	16	18	30	15
pfgcd \rightarrow	24	8	2	2	1
sfgcd \rightarrow	1	1	3	15	15

```
int deleteOne (int[] arr, int n) {
```

```
    pfgcd[n];  
    sfgcd[n]; } todo.
```

```
    ans = Map(sfgcd[1], pfgcd[n-2]);
```

```
    for (i=1; i < n-1; i++) {
```

```
        // deleting ith element.
```

```
        left = pfgcd[i-1];
```

```
        right = sfgcd[i+1];
```

```
        val = gcd(left, right);
```

```
        ans = Map(ans, val);
```

$\begin{array}{|c} \hline \\ \hline \end{array}$
 $\begin{array}{|c} \hline \\ \hline \end{array}$

 return ans;

$$T.C \Rightarrow O(n \log \max(a, b))$$

$$S.C \Rightarrow O(1)$$

// $a > b$

$$\gcd(a, b) = \gcd(a \cdot b, b)$$

Case-1

$$b < \frac{a}{2}$$

$$\underbrace{a \cdot b}_{0 \text{ to } b-1} < b < \frac{a}{2}$$

$$\underline{a \cdot b < \frac{a}{2}}$$

Case-2

$$b = \frac{a}{2}$$

$$\underbrace{a \cdot b}_{0 \text{ to } b-1} < b = \frac{a}{2}$$

$$\underline{a \cdot b < \frac{a}{2}}$$

Case-3

$$\underline{b > \frac{a}{2}}$$

$$\rightarrow 2b > a$$

$$\rightarrow 2b - a > 0$$

// (-) both sides

$$a - 2b < 0$$

adding a on both sides

$$2a - 2b < a$$

$$(a - b) < \frac{a}{2}$$

$$\boxed{a \cdot b = a - b}$$

$$a - 2b$$

$$a - 3b$$

$$\boxed{a \cdot b < \frac{a}{2}}$$

$$a$$

$$-30 \div 7 = -2 \text{ r } 6$$

(sign follows the rule of sign of remainder will be same as sign of dividend)

$$7 \overline{) 30} \begin{array}{r} 28 \\ \underline{2} \end{array}$$

$$-30 \div 7 = -30 - (-35) = 5$$

$$\approx 10^9, \text{ Prime no.}$$

$$(37.2)^{103} \approx 10^3$$

$$\downarrow$$

$$10^3$$