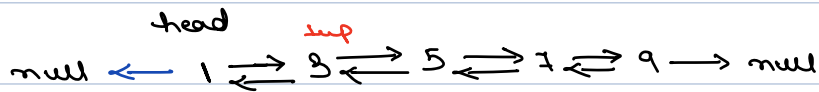


Today's Content :-

- What is doubly linked list?
- How is doubly linked list different from singly linked list?
- 4 coding problems related to doubly linked list

← Doubly LL →

A doubly linked list is a type of data structure used in computer science and programming to store and organize a collection of elements, such as nodes. It is similar to a singly linked list but with an additional feature: each node in a doubly linked list contains pointers or references to both the next and the previous nodes in the list.



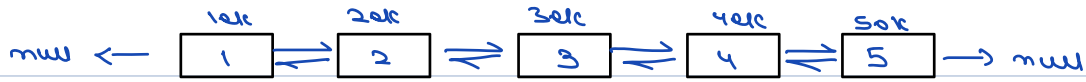
$x = \text{temp.next}$

$y = \text{temp.prev}$

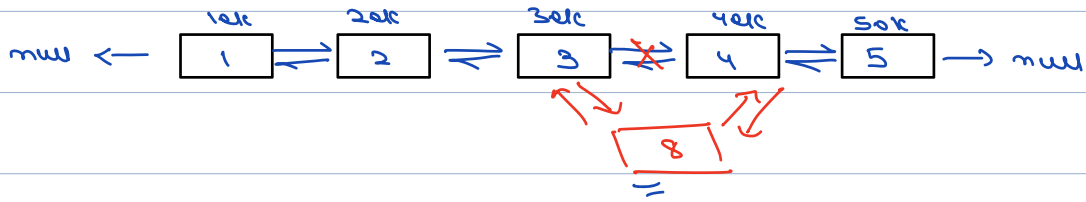
```
class Node {  
    int data;  
    Node next;  
    Node prev;  
}
```

Ques

A doubly linked list is given. A node is to be inserted with data X at position K . The range of K is between 0 and N where N is the length of the doubly Linked list.



Ex $\rightarrow X = 8, K = 3$



data = 8
next = 4th
prev = 3rd

$xm = \text{new Node}(x)$

if (head == null) & return xm

if (x == 0) &

T.C $\rightarrow O(m)$

S.C $\rightarrow O(1)$

xm.next = head;

head.prev = xm;

head = xm;

return head;

3

temp = head;

for (i = 1; i < k; i++);

temp = temp.next

xm.next = temp.next ✓

xm.prev = temp; ✓

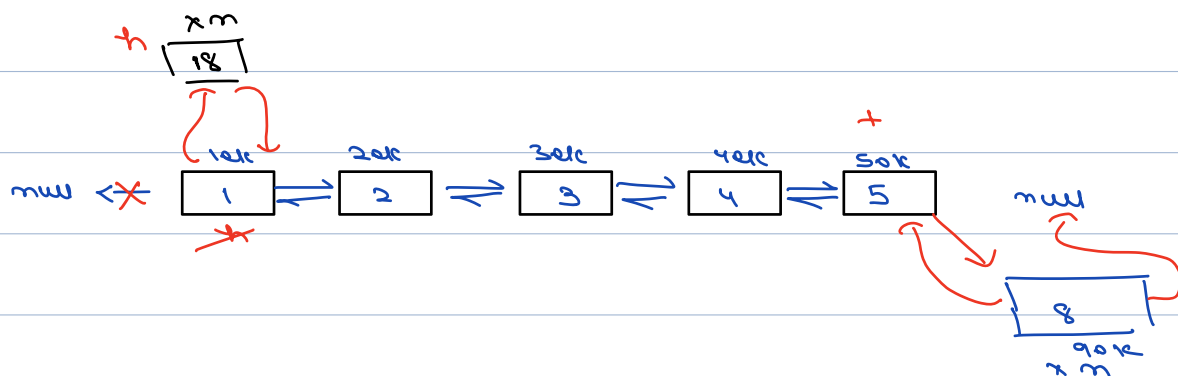
if (temp.next != null) &

3

temp.next.prev = xm; ^{90k}

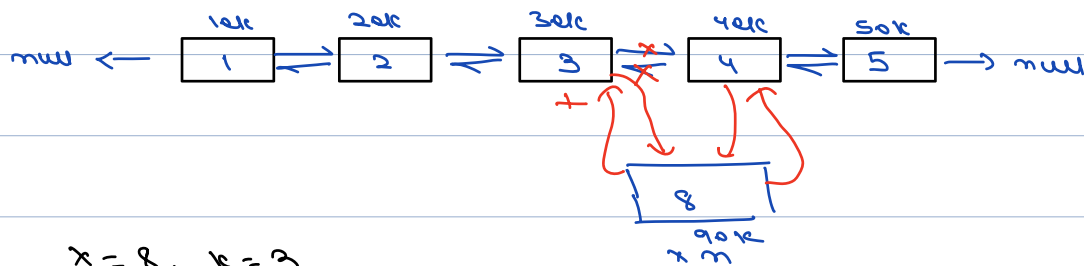
temp.next = xm;

return head;



x = 8, k = 5

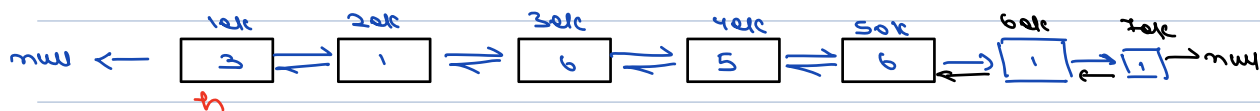
(0 based)



$x = 8, k = 3$
(0 based)

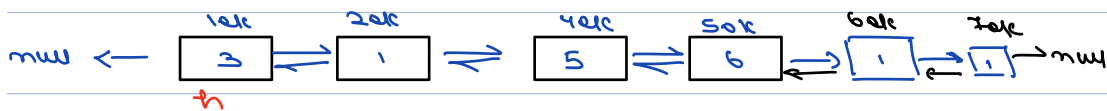
Ques

we have been given a doubly linked list of length N , we have to delete the first occurrence of data X from the given doubly linked list. If element X is not present, don't do anything.

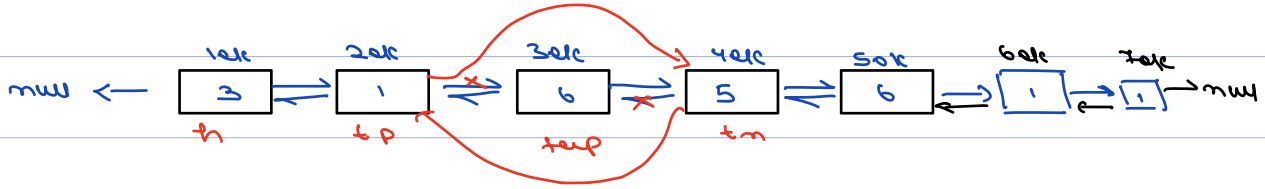


$x = 6$.

Ans.



Soln :- $x = 6$,



temp = head;

11 Searching.

```
while (temp != null) {
```

```
if (heap.data == x) {
```

1 break;

temp - temp. ment.;

13

```
if (temp == null) {
```

return head;

3

```
if (temp.prev == null && temp.next == null) {
```

head = null;

western head;

3

Diagram illustrating a linked list node structure:

```
graph TD; temp --> node; node --> null; node --> head;
```

The diagram shows a node containing the value 6. The node is pointed to by a variable named `temp`. The node's next pointer is set to `null` (indicated by an arrow pointing left to the word `null`). The node is also pointed to by a variable named `head` (indicated by an arrow pointing down to the node).

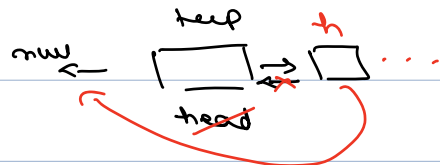
```
else if (temp.prev == null) {
```

temp. next. prev = null;

head = temp.next;

return head;

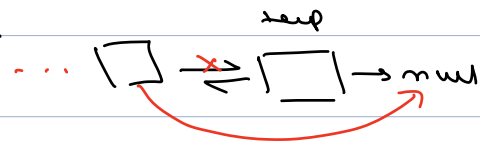
3



```
else if (temp->next == null) {
```

```
    temp->prev->next = null;
```

```
    return head;
```



```
}
```

```
    tp = temp->prev;
```

```
    tn = temp->next;
```

```
    tp->next = tn;
```

```
    tn->prev = tp;
```

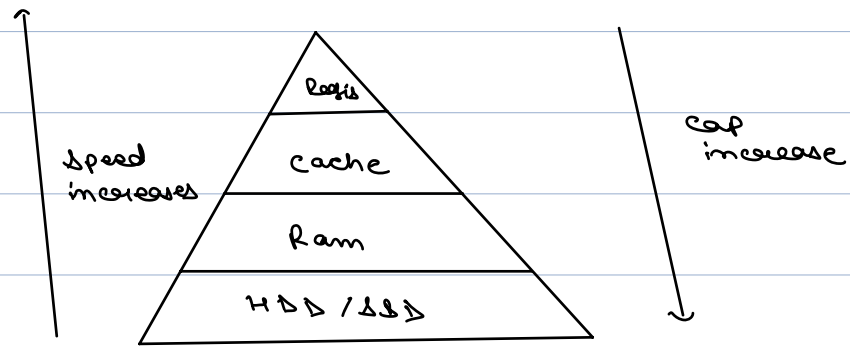
```
    return head;
```

```
}
```

T.C $\rightarrow O(n)$

S.C $\rightarrow O(1)$.

LRU Cache → least recently used.

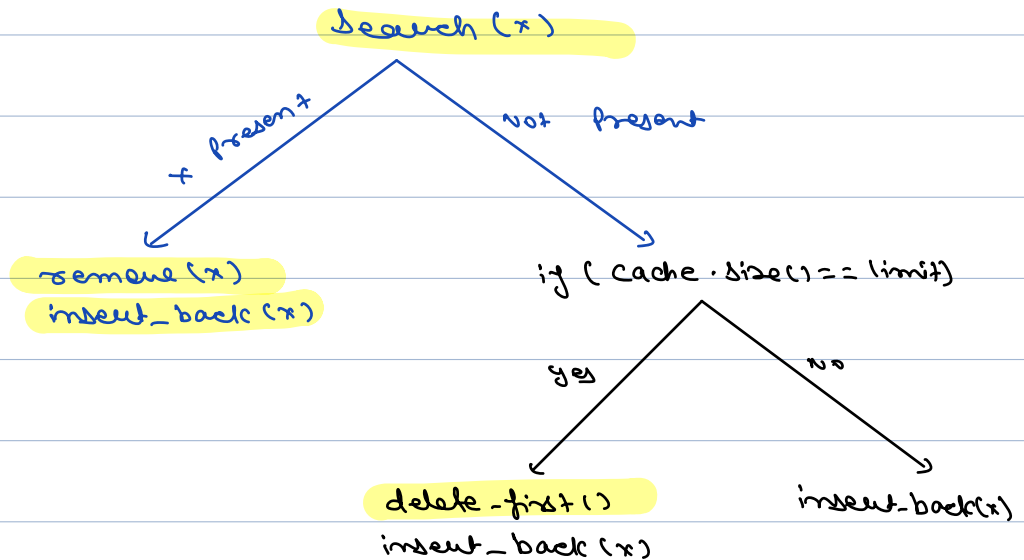


Data 7 3 2 9 6 10 14 2 15 10 8 11 10 19

Cap = 5

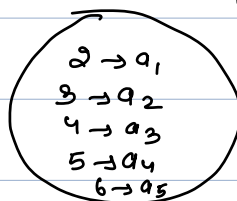
~~7~~ ~~3~~ ~~2~~ ~~9~~ ~~6~~ ~~10~~ ~~14~~ ~~2~~ 15 ~~10~~ 8 11 10 19

flow chart (x)



<u>Operations :-</u>	<u>Dynamic Array</u>	<u>LH</u>
Search (x)	$O(n)$	$O(n)$
remove (x)	$O(n)$	$O(n)$
insert - back (x)	$O(1)$	$O(1)$ (if you maintain tail)
delete - first (1)	$O(n)$	$O(1)$

a₁ # a₂ # a₃ # a₄ # a₅
 2 → 3 → 4 → 5 → 6
 + +



LH + HashMap

<u>Operations :-</u>	<u>Queue</u>	<u>LH + HashMap</u>
Search (x)	$O(n)$	$O(1)$
remove (x)	$O(n)$	<u>$O(n)$</u>
insert - back (x)	$O(1)$	$O(1)$ (with tail)
delete - first (1)	$O(1)$	$O(1)$

DLH + HashMap

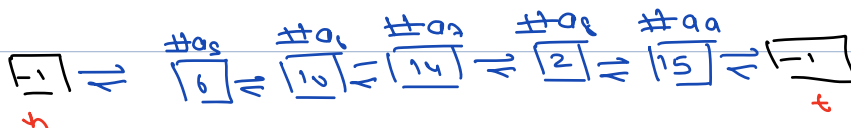
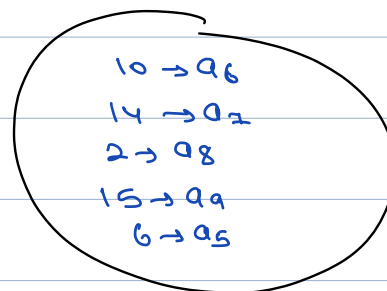
Data 7 3 2 9 6 10 14 2 15 , 19 ,

Cap = 5

~~7~~ ~~3~~ ~~2~~ ~~9~~ 6 10 14 2 15

DLL

Hash Map



Node h = new Node (-1);

Node t = new Node (-1);

h.next = t



t.prev = h

HashMap<int, Node> hm;

- - - LRU (int x, int limit) {

if (hm.search(x)) {

Node t = hm[x];

DeleteNode(t);

```
Node nn = new Node(x);
```

```
insert-back (nn, tail);
```

```
hm[x] = nn;
```

```
}
```

```
else {
```

```
if (hm.size() == limit) {
```

```
Node t = hm.next();
```

```
hm.delete (t.data);
```

```
Delete Node (t);
```

```
}
```

```
Node nn = new Node(x);
```

```
insert-back (nn, tail);
```

```
hm.insert (x, nn);
```

```
}
```

```
void DeleteNode (Node temp) {
```

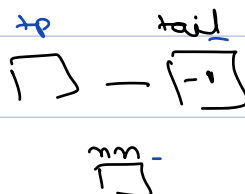
```
Node tp = temp.prev;
```

```
Node tn = temp.next;
```

```
tp.next = tn;
```

```
tn.prev = tp;
```

```
}
```



```
void insert_back (Node nn, Node tail) {
```

```
    Node tp = tail->prev;
```

```
    tp->next = nn;
```

```
    nn->prev = tp;
```

```
    nn->next = tail;
```

```
    tail->prev = nn;
```

```
}
```

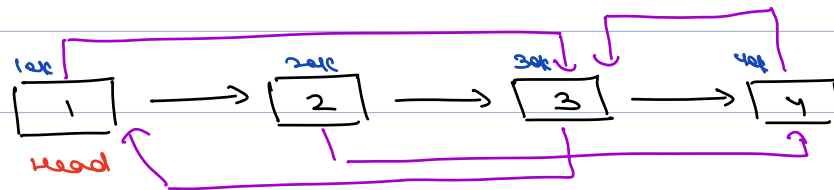
Ques Deep copy of a lh.

class node {

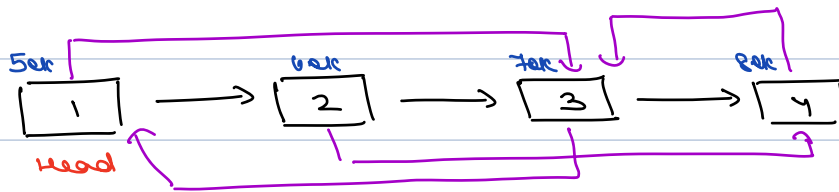
int data;
node next;
node rand;

}

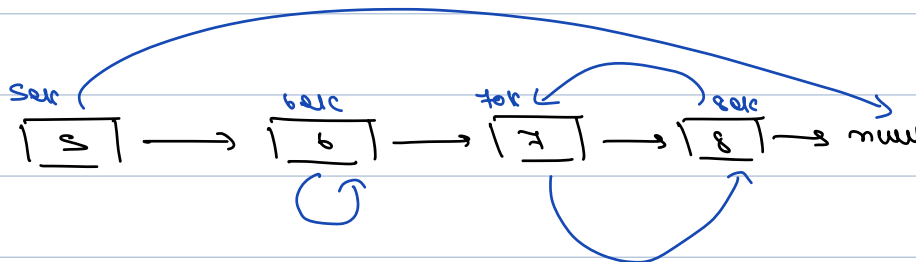
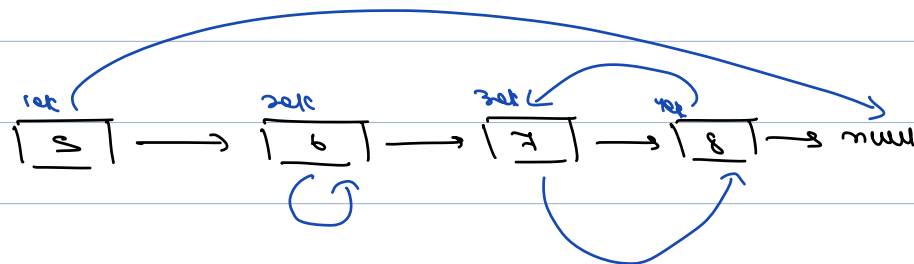
I/p



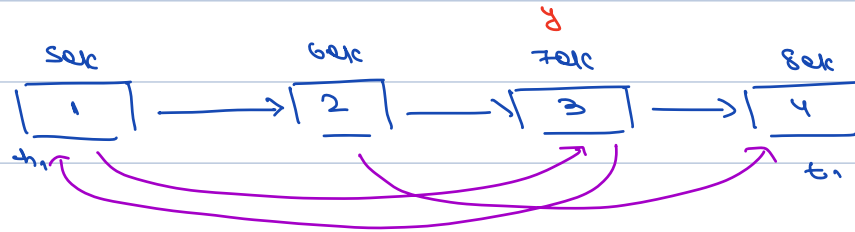
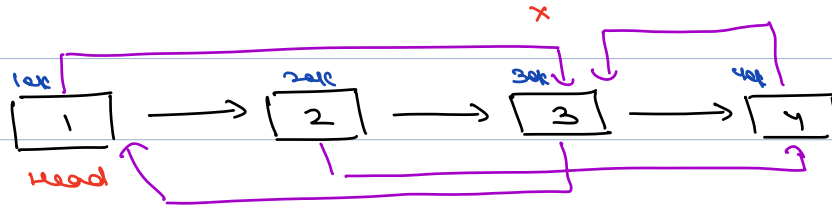
O/p



I/p



idea



while (x != null) {

1st
temp = x.val; 1st
y.val = hm.get(temp)
2nd

x = x.next;
y = y.next;

HashMap <Node, Node>

1st → 1st

2nd → 2nd

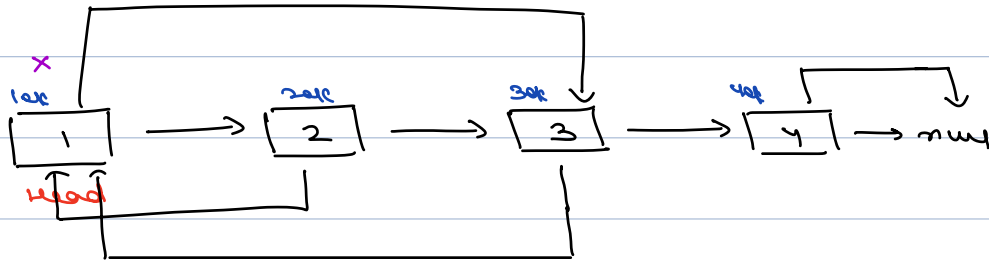
3rd → 3rd

4th → 4th

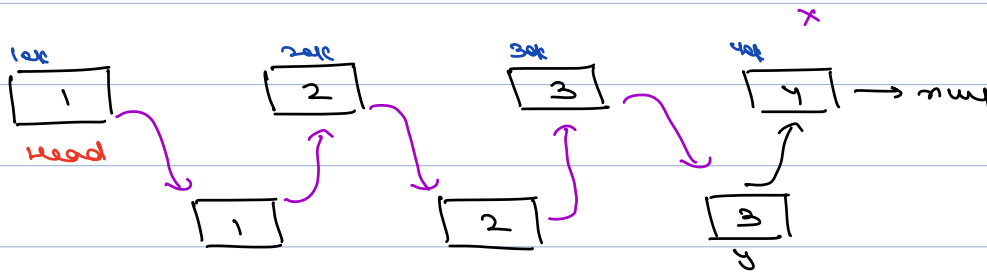
T.C → O(n)

S.C → O(n)

Constraint :- Do it in Constant Space :-



1)

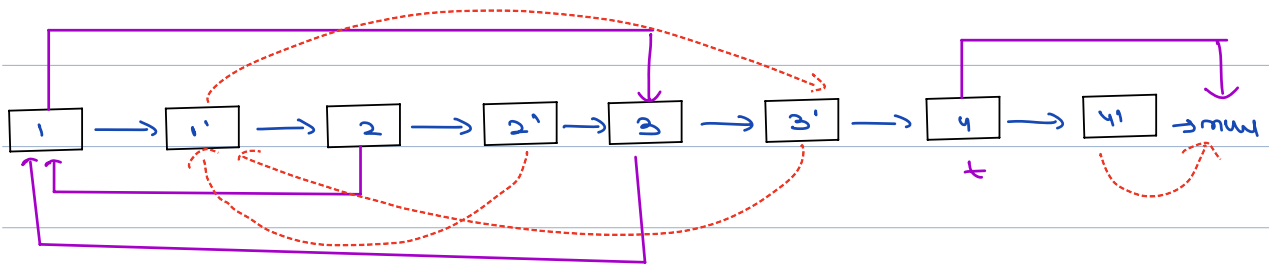


$y = \text{new Node}(x.\text{data})$ ✓

$y.\text{next} = x.\text{next};$ ✓

$x.\text{next} = y$ ✓

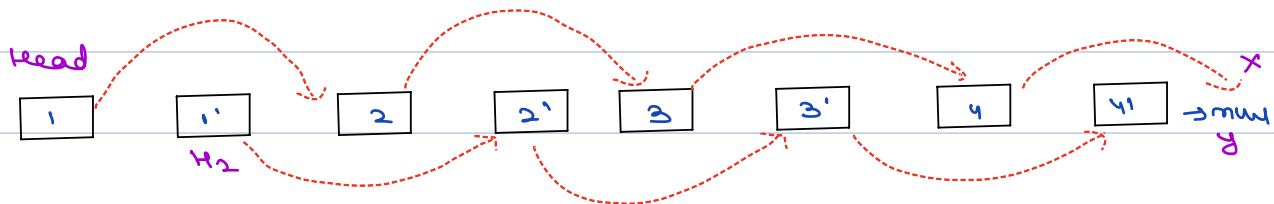
$x = x.\text{next}.\text{next}$



temp = head;

```
while (temp != null) {
    if (temp.random == null) { temp.next.random = null; }
    else { temp.next.random = temp.random.next; }
    temp = temp.next.next;
}
```

Step 3 detach



x2 = head.next

x = head;

y = head.next

while (x != null) {

x.next = x.next.next

if (y.next != null) {
y.next = y.next.next;

x = x.next;

y = y.next;

return x2;