



PennState

# CMPSC 311 - Introduction to Systems Programming

## Bit/Byte Operations

Professors

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(Slides are mostly by Professor Patrick McDaniel  
and Professor Abutalib Aghayev)



Hobbit

Hobbyte

# Number Systems



- All base- $X$  systems have the following characteristic:

Assume a base  $b$  and digits  $P = \{p_k, p_{k-1}, p_{k-2}, \dots, p_1, p_0\}$

$$value = \sum_{i=0}^k b^i * p_i$$

where  $\forall p_i \in P, p_i = [0, b - 1]$

Base-10 (decimal)

- Digits: 0,1,2,3,4,5,6,7,8,9
- Place values:  $10^0$ ,  $10^1$ ,  $10^2$ , ...
- Example:  $123 = 3 \times 1 + 2 \times 10 + 1 \times 100$   
= 123

Base-2(binary)

- Digits: 0,1
- Place values:  $2^0$ ,  $2^1$ ,  $2^2$ , ...
- Example:  $1011 = 1 \times 1 + 1 \times 2 + 0 \times 4 + 1 \times 8$   
= 11

Base-16 (hexadecimal)

- Digits: 0,1,2,3,4,5,6,7,8,9,A,B,C,D,E,F
- Place values:  $16^0$ ,  $16^1$ ,  $16^2$ , ...
- Example:  $0xAFC = C \times 1 + F \times 16 + A \times 256$   
= 2812

# Converting Decimal to Binary



- while  $n \neq 0$ :
  - next binary digit (from right to left) =  $n \% 2$
  - $n = n / 2$

n	remainder	digit #
235	1	0
117	1	1
58	0	2
29	1	3
14	0	4
7	1	5
3	1	6
1	1	7

235 (base-10) = 11101011 (base-2)

# Converting Decimal and Binary to Hex



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- Converting decimal to hexadecimal
  - while  $n \neq 0$ :
    - next hex digit (from right to left) =  $n \% 16$
    - $n = n / 16$

n	remainder	digit #
235	B (11)	0
14	E (14)	1

235 (base-10) = EB (base-16)

- Converting binary to hexadecimal
  - group binary digits to into groups of 4 bits (nibbles) starting from right
  - Convert each nibble to hexadecimal digit

235 (base-10) = 11101011 (base-2) = 1110 1011 = EB

- Be familiar with hex notation: 0xDEADBEEF has 8 nibbles, 4 bytes (or octets)
- Hexadecimal to binary: just convert each nibble to binary

# Conversion Summary



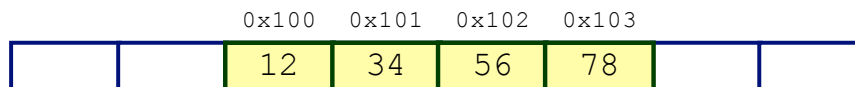
- Binary to decimal
  - sum of binary digits times powers of 2
- Hexadecimal to decimal
  - sum of hex digits times power of 16
- Decimal to binary
  - division method
- Decimal to hex
  - division method
  - or, convert to binary and use binary to hex method below
- Binary to hex
  - group binary digits to nibbles, convert nibbles to hex digits
- Hex to binary
  - convert each hex digit to binary

# Byte Ordering Example

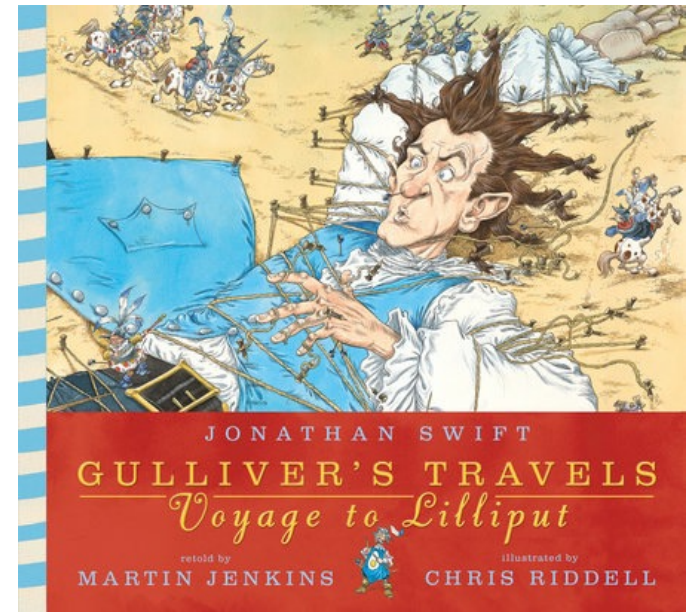
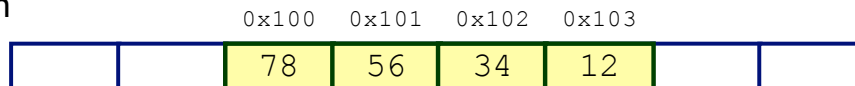


- How should bytes within a multi-byte word be ordered in memory?
  - Big Endian: Least significant byte has *highest* address (SPARC)
  - Little Endian: Least significant byte has *lowest* address (x86)
- Example: `int x = 305419896`
  - Variable `x` has 4-byte representation `0x12345678`
  - Address given by `&x` is `0x100`

Big Endian



Little Endian



# Examining Data Representations



- Code to find endianness of the architecture
  - Casting pointer to `uint8_t *` creates byte array

```
#include <stdio.h>
#include <stdint.h>

void show_bytes(uint8_t *start, int len) {
    for (int i = 0; i < len; ++i)
        printf("%p\t0x%.2x\n", start+i, start[i]);
    printf("\n");
}

int main(void) {
    int a = 305419896;
    printf("a lives at address %p\n\n", &a);
    show_bytes((uint8_t *)&a, sizeof(int));
}
```

Result (Linux on x86):

a lives at address 0x7ffebf803174

0x7ffebf803174	0x78
0x7ffebf803175	0x56
0x7ffebf803176	0x34
0x7ffebf803177	0x12

printf directives

%p	Print pointer
%x	Print hexadecimal

# Boolean Algebra



- Developed by George Boole in 19th Century
  - Algebraic representation of logic based on “True” (as 1) and “False” (as 0)

## And

$A \& B = 1$  when both  $A=1$  and  $B=1$

$\&$	0	1
0	0	0
1	0	1

## Or

$A | B = 1$  when either  $A=1$  or  $B=1$

$ $	0	1
0	0	1
1	1	1

## Not

$\sim A = 1$  when  $A=0$

$\sim$	
0	1
1	0

## Exclusive-Or (Xor)

$A \wedge B = 1$  when either  $A=1$  or  $B=1$ , but not both

$\wedge$	0	1
0	0	1
1	1	0



# Bit-Level Operations in C



- Operations  $\&$ ,  $|$ ,  $\sim$ ,  $\wedge$  Available in C
  - Apply to any “integral” data type
    - long, int, short, char, unsigned
  - View arguments as bit vectors
  - Arguments applied bit-wise
- Examples (Char data type)
  - $\sim 0x41 \rightarrow 0xBE$ 
    - $\sim 01000001_2 \rightarrow 10111110_2$
  - $\sim 0x00 \rightarrow 0xFF$ 
    - $\sim 00000000_2 \rightarrow 11111111_2$
  - $0x69 \& 0x55 \rightarrow 0x41$ 
    - $01101001_2 \& 01010101_2 \rightarrow 01000001_2$
  - $0x69 | 0x55 \rightarrow 0x7D$ 
    - $01101001_2 | 01010101_2 \rightarrow 01111101_2$



# Contrast: Logic Operations in C



- Contrast to Logical Operators (`&&`, `||`, `!`)
  - View 0 as “False”
  - Anything nonzero as “True”
  - Always return 0 or 1
  - Early termination

# Representing & Manipulating Sets PennState

- Representation

- Width  $w$  bit vector represents subsets of  $\{0, \dots, w-1\}$  for set “A”
- $a_j=1$  if  $j \in A$

01010101	{ 0, 2, 4, 6 }
76543210	

01101001	{ 0, 3, 5, 6 }
76543210	

## Operations On Sets:

& Intersection	01000001	{ 0, 6 }
Union	01111101	{ 0, 2, 3, 4, 5, 6 }
^ Symmetric difference	00111100	{ 2, 3, 4, 5 }
~ Complement	10101010	{ 1, 3, 5, 7 }

# Representing Signed Numbers



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b3	b2	b1	b0	Unsigned	One's complement	Two's complement
0	0	0	0	0	0	0
0	0	0	1	1	1	1
0	0	1	0	2	2	2
0	0	1	1	3	3	3
0	1	0	0	4	4	4
0	1	0	1	5	5	5
0	1	1	0	6	6	6
0	1	1	1	7	7	$7 = 2^{n-1} - 1$
1	0	0	0	8	-7	$-8 = -2^{n-1}$
1	0	0	1	9	-6	-7
1	0	1	0	10	-5	-6
1	0	1	1	11	-4	-5
1	1	0	0	12	-3	-4
1	1	0	1	13	-2	-3
1	1	1	0	14	-1	-2
1	1	1	1	15	-0	-1

- Computing one's complement negative representation:
  - Complement the positive number
- Computing two's complement negative representation:
  - Complement the positive number and add 1
- Most architectures use two's complement
- Given n-bit signed integer:
  - Positive range:  $0 - 2^{n-1} - 1$
  - Negative range:  $-2^{n-1}$

type	bytes (32-bit)	bytes (64-bit)	32-bit range	printf
char	1	1	[0, 255]	%c
short int	2	2	[-32768, 32767]	%hd
unsigned short int	2	2	[0, 65535]	%hu
int	4	4	[-2147483648, 2147483647]	%d
unsigned int	4	4	[0, 4294967295]	%u
long int	4	8	[-2147483648, 2147483647]	%ld
long long int	8	8	[-9223372036854775808, 9223372036854775807]	%lld
float	4	4	[-3.4e+38, 3.4e+38]	%f

# Beware of integer overflows



- Spot the bug in the following:

```
for (uint8_t n = 10; n >=0; --n) {  
    printf("do I ever terminate?\n");  
}
```



- What does the following print?

```
int x = 0xffffffff;  
printf("%d\n", x);
```

---

Extra, Extra - Read All About It: Nearly All Binary Searches and Mergesorts are Broken

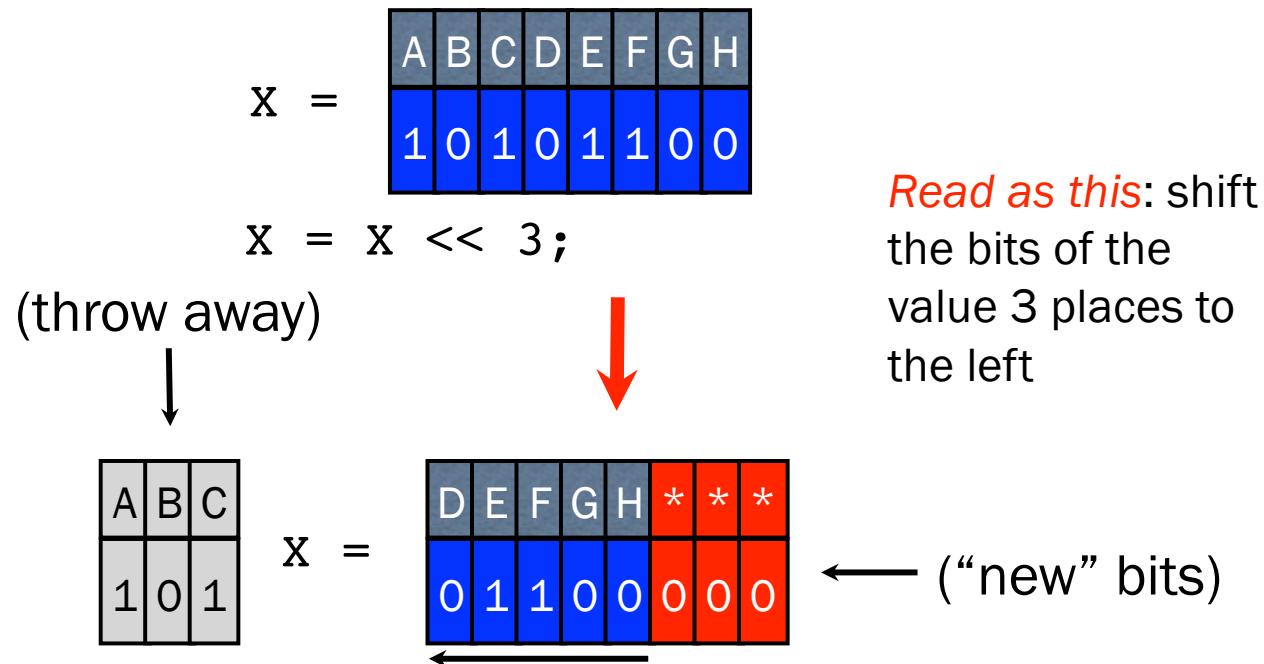
Friday, June 2, 2006

Posted by Joshua Bloch, Software Engineer

# Shift Operations



- A shift operator ( $\ll$  or  $\gg$ ) moves bits to the right or left, throwing away bits and adding bits as necessary



# Putting it all together



- Suppose you want to place multiple values in the same 32-bit integer
  - Value **a** in least significant byte
  - Value **b** in 2nd byte
  - Value **c** in 3rd byte
  - Value **d** in 4th byte

Bits	31-24	23-16	8-15	0-7
Values	d	c	b	a

# Using bit operations ...



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```
uint32_t pack_bytes(uint32_t a, uint32_t b, uint32_t c, uint32_t d) {  
  
    // Setup some local values  
    uint32_t retval = 0x0, tempa, tempb, tempc, tempd;  
  
    tempa = a&0xff; // Make sure you are only getting the bottom 8 bits  
    tempb = (b&0xff) << 8; // Shift value to the second byte  
    tempc = (c&0xff) << 16; // Shift value to the third byte  
    tempd = (d&0xff) << 24; // Shift value to the top byte  
    retval = tempa|tempb|tempc|tempd; // Now combine all of the values  
  
    // Print out all of the values  
    printf("A: 0x%08x\n", tempa);  
    printf("B: 0x%08x\n", tempb);  
    printf("C: 0x%08x\n", tempc);  
    printf("D: 0x%08x\n", tempd);  
  
    // Return the computed value  
    return retval;  
}  
  
...  
  
printf("Packed bytes : 0x%08x\n", pack_bytes(0x111, 0x222, 0x333, 0x444));
```

```
A: 0x00000011  
B: 0x00002200  
C: 0x00330000  
D: 0x44000000  
Packed bytes : 0x44332211
```