Lab 4 实验报告

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实验进度

我完成了格式化输入函数,进程通信,信号量相关系统调用,进程同步问题

实验结果

格式化输入函数

进程通信

```
Father Process: 2020, 0
Child Process: 2020, 1000
Father Process: 3020, 1000
Father Process: 2020, 3020
Child Process: 4020, 1000
Father Process: 4020, 1000
Father Process: 5020, 4020
Child Process: 5020, 1000
-
```

信号量相关系统调用

```
Father Process: Semaphore Initializing.
Father Process: Sleeping.
Child Process: Semaphore Waiting.
Child Process: In Critical Area.
Child Process: Semaphore Waiting.
Child Process: In Critical Area.
Child Process: Semaphore Waiting.
Father Process: Semaphore Posting.
Father Process: Sleeping.
Child Process: In Critical Area.
Child Process: Semaphore Waiting.
Father Process: Semaphore Waiting.
Father Process: Semaphore Posting.
Father Process: Sleeping.
Child Process: In Critical Area.
Child Process: Semaphore Destroying.
Father Process: Semaphore Posting.
Father Process: Semaphore Posting.
Father Process: Semaphore Posting.
Father Process: Semaphore Destroying.
Father Process: Semaphore Destroying.
```

进程同步问题

生产者消费者问题

哲学家就餐问题

读者写者问题

```
Writer 7: write
Writer 1: write
Writer 4: write
Writer 3: write
Reader 2: read, total 1 reader
Reader 6: read, total 2 reader
Reader 8: read, total 2 reader
Reader 6: read, total 2 reader
Reader 6: read, total 3 reader
Reader 6: read, total 3 reader
Reader 6: read, total 1 reader
Reader 7: read, total 2 reader
Reader 8: read, total 3 reader
Reader 5: read, total 3 reader
Reader 5: read, total 4 reader
Writer 7: write
Writer 1: write
Writer 4: write
Writer 3: write
Writer 3: write
Reader 6: read, total 1 reader
Reader 6: read, total 2 reader
Reader 5: read, total 3 reader
Reader 6: read, total 3 reader
Reader 5: read, total 3 reader
Reader 6: read, total 3 reader
```

实验修改的代码

所有标记了TODO in lab4的方法

实验中遇到的问题

不知道怎么实现随机函数

因为随机函数跟时间可能是相关的,所以在irqHandle中直接用inbyte把TIME_PORT的数据写到了eax 里面,发现确实是具有随机性的数据

读者写者问题不知道怎么在多个进程之间同步数据

读者需要同步当前正在读的人数,本来以为是可以不用管道的,但是想了一下好像只有管道可以解决这个问题。

生产者消费者问题中的PV顺序有影响吗?

有,我们把代码改写成这样

```
BoundedBuffer::Deposit(c){
 2
        emptyBuffers->P();
 3
        mutex->P();
 4
        Add c to the buffer;
 5
        mutex->V();
 6
        fullBuffers->V();
 7
 8
    BoundedBuffer::Remove(c){
 9
        mutex->P();
10
        fullBuffers->P();
        Remove c from buffer;
11
12
        emptyBuffers->V();
13
        mutex->V();
```

这样,如果消费者先执行,并且在第10行被挂起,然后进程切换到生产者,然后在第三行被挂起,这样 就会产生死锁,因为生产者需要mutex资源,而消费者需要fullBuffers资源,形成了一个资源等待链

有没有更好的方式处理哲学家就餐问题?

暂时没想到