

QUANTUM COMPUTATION AND QUANTUM INFORMATION: THE QUANTUM FOURIER TRANSFORM

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We consider the linear map in \mathbb{C}^N which acts on the computational basis as

$$|j\rangle \mapsto \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{\frac{2i\pi jk}{N}} |k\rangle$$

Let A be the matrix of the transformation in the computational basis.

$$\forall (k, l) \in \llbracket 0, N-1 \rrbracket^2, \quad a_{kl} = \frac{1}{\sqrt{N}} e^{\frac{2i\pi kl}{N}}$$

The adjoint matrix A^\dagger is then

$$\begin{aligned} \forall (k, l) \in \llbracket 0, N-1 \rrbracket^2, \quad b_{kl} &= a_{lk}^* \\ &= \frac{1}{\sqrt{N}} e^{-\frac{2i\pi kl}{N}} \end{aligned}$$

We compute the coefficient k, l of the product AA^\dagger :

$$\begin{aligned} \forall (k, l) \in \llbracket 0, N-1 \rrbracket^2, \quad c_{kl} &= \sum_{j=0}^{N-1} a_{kj} b_{jl} \\ &= \frac{1}{N} \sum_{j=0}^{N-1} e^{\frac{2i\pi j}{N} (k-l)} \\ &= \frac{1}{N} \sum_{j=0}^{N-1} (e^{\frac{2i\pi}{N} (k-l)})^j \\ &= \begin{cases} \frac{1}{N} \frac{1 - (e^{\frac{2i\pi}{N} (k-l)})^N}{1 - e^{\frac{2i\pi}{N} (k-l)}} = 0 & \text{if } e^{\frac{2i\pi}{N} (k-l)} \neq 1, \\ 1 & \text{if } e^{\frac{2i\pi}{N} (k-l)} = 1. \end{cases} \\ &= \begin{cases} 0 & \text{if } k \neq l, \\ 1 & \text{if } k = l. \end{cases} \\ &= \delta_{kl} \end{aligned}$$

which shows that $AA^\dagger = A^\dagger A = I$ i.e. A is unitary.

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Here the dimension of the state space is $N = 2^n$. The Fourier transform of the n qubit state $|00 \dots 0\rangle$ is

$$A|0\rangle = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} |k\rangle$$

we can write k in binary $k_{n-1} \dots k_1 k_0$

$$A|0\rangle = \frac{1}{2^{n/2}} \sum_{k_0, k_1, \dots, k_{n-1}=0}^1 |k_{n-1} \dots k_1 k_0\rangle$$

or in product representation,

$$= \frac{1}{2^{n/2}} \underbrace{(|0\rangle + |1\rangle)(|0\rangle + |1\rangle) \dots (|0\rangle + |1\rangle)}_{n \text{ qubits}}$$

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Let $N = 2^n$ and $Y = (y_k)_{k \in \llbracket 0, N-1 \rrbracket}$ be the classical fourier transform of $X = (x_k)_{k \in \llbracket 0, N-1 \rrbracket}$.

$$\forall k \in \llbracket 0, N-1 \rrbracket, \quad y_k = \sum_{j=0}^{N-1} e^{\frac{2i\pi k j}{2^n}} x_j$$

The factor $\frac{1}{\sqrt{N}}$ is omitted for clarity. We can write j in binary $j_{n-1} \dots j_1 j_0$

$$\begin{aligned} y_k &= \sum_{j_0, j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-1}j_{n-1} + \dots + 2j_1 + j_0)}{2^n}} x_j \\ &= \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-1}j_{n-1} + \dots + 2j_1)}{2^n}} x_{j_{n-1} \dots j_1 0} + \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-1}j_{n-1} + \dots + 2j_1 + 1)}{2^n}} x_{j_{n-1} \dots j_1 1} \\ &= \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-1}j_{n-1} + \dots + 2j_1)}{2^n}} x_{j_{n-1} \dots j_1 0} + e^{\frac{2i\pi k}{2^n}} \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-1}j_{n-1} + \dots + 2j_1)}{2^n}} x_{j_{n-1} \dots j_1 1} \\ &= \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-2}j_{n-1} + \dots + j_1)}{2^{n-1}}} x_{j_{n-1} \dots j_1 0} + e^{\frac{2i\pi k}{2^n}} \sum_{j_1, \dots, j_{n-1}=0}^1 e^{\frac{2i\pi k(2^{n-2}j_{n-1} + \dots + j_1)}{2^{n-1}}} x_{j_{n-1} \dots j_1 1} \end{aligned}$$

We see the first sum is the k^{th} coefficient of the FT of the sequence $(x_{2k})_{k \in \llbracket 0, N/2-1 \rrbracket}$ and the second is the k^{th} coefficient of the FT of $(x_{2k+1})_{k \in \llbracket 0, N/2-1 \rrbracket}$. This shows that to compute FT of sequence of length N , we have to compute 2 FT of sequence of length $\frac{N}{2}$ and do $2N$ complex additions/multiplications. The complexity of the operation $T(N)$ follows the recurrence:

$$T(N) = 2T\left(\frac{N}{2}\right) + 2N$$

We can use the Master theorem ¹:

Theorem. Let $a \geq 1$ and $b > 1$ be constants, let $f(n)$ be a function, and let $T(n)$ be defined on the non negative integers by the recurrence

$$T(n) = aT\left(\frac{n}{b}\right) + f(n)$$

where we interpret $\frac{n}{b}$ to mean either $\lfloor \frac{n}{b} \rfloor$ or $\lceil \frac{n}{b} \rceil$. Then $T(n)$ has the following asymptotic bounds:

- (1) If $f(n) = O(n^{\log_b a - \epsilon})$ for some constant $\epsilon > 0$, then $T(n) = \Theta(n^{\log_b a})$.
- (2) If $f(n) = \Theta(n^{\log_b a})$, then $T(n) = \Theta(n^{\log_b a} \log n)$.
- (3) If $f(n) = \Omega(n^{\log_b a + \epsilon})$ for some constant $\epsilon > 0$, and if $af(\frac{n}{b}) \leq cf(n)$ for some constant $c < 1$ and n sufficiently large, then $T(n) = \Theta(f(n))$.

Here we are in the second case of the theorem, so $T(N) = \Theta(N \log(N)) = \Theta(n2^n)$.

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The inverse Fourier Transform

$$|j\rangle \mapsto \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} e^{-\frac{2i\pi j k}{N}} |k\rangle$$

is the adjoint of the Fourier Transform. The quantum circuit of figure 1 is obtained from the FT's circuit, replacing each R_k gate by its adjoint

$$R_k^\dagger = \begin{bmatrix} 1 & 0 \\ 0 & e^{-\frac{2i\pi}{2^k}} \end{bmatrix}$$

¹Thomas H. Cormen and Charles E. Leiserson : *Introduction to algorithms*, MIT Press (2009)



FIGURE 1. Quantum circuit for IFT.



FIGURE 2. Sequence of controlled U.

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In figure 2, the t qubits of the first register are prepared with $|j\rangle = |j_{t-1} \dots j_1 j_0\rangle$, the second register is prepared with some state $|u\rangle$. After the first controlled-U operation, the state is $|j\rangle |U^{j_0 2^0} u\rangle$. After the second controlled-U, the state is $|j\rangle |U^{j_1 2^1} U^{j_0 2^0} u\rangle = |j\rangle |U^{j_0 2^0 + j_1 2^1} u\rangle$ and so on. The final state is $|j\rangle |U^{j_0 2^0 + j_1 2^1 + \dots + j_{t-1} 2^{t-1}} u\rangle = |j\rangle |U^j u\rangle$.

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By linearity, the phase estimation algorithm takes input $|0\rangle |\sum_{u \in A} c_u |u\rangle\rangle$, where A is some orthonormal basis of eigenstates of U , to output $\sum_{u \in A} c_u |\widetilde{\varphi}_u\rangle |u\rangle$, where $\widetilde{\varphi}_u$ is an estimation of the phase of the eigenvalue associated with eigenstate u . If we fix $u_0 \in A$ beforehand, the probability to measure $\widetilde{\varphi}_{u_0}$ when measuring the first register in the computational basis is

$$\begin{aligned}
 & \left(\sum_{u \in A} c_u^* \langle \widetilde{\varphi}_u | \langle u | \right) P_{\widetilde{\varphi}_{u_0}} \otimes I \left(\sum_{u \in A} c_u |\widetilde{\varphi}_u\rangle |u\rangle \right) = \left(\sum_{u \in A} c_u^* \langle \widetilde{\varphi}_u | \langle u | \right) \left(\sum_{\substack{u \in A \\ \widetilde{\varphi}_u = \widetilde{\varphi}_{u_0}}} c_u |\widetilde{\varphi}_u\rangle |u\rangle \right) \\
 & = \left(\sum_{u \in A} c_u^* \langle \widetilde{\varphi}_u | \langle u | \right) \left(\sum_{\substack{u \in A \\ \widetilde{\varphi}_u = \widetilde{\varphi}_{u_0}}} c_u |\widetilde{\varphi}_{u_0}\rangle |u\rangle \right) \\
 & = \sum_{\substack{v \in A \\ u \in A \\ \widetilde{\varphi}_u = \widetilde{\varphi}_{u_0}}} c_v^* c_u \langle \widetilde{\varphi}_v | \widetilde{\varphi}_u \rangle \langle v | u \rangle \\
 & = \sum_{\substack{v \in A \\ u \in A \\ \widetilde{\varphi}_u = \widetilde{\varphi}_{u_0}}} c_v^* c_u \langle \widetilde{\varphi}_v | \widetilde{\varphi}_u \rangle \delta_{vu} \\
 & = \sum_{\substack{u \in A \\ \widetilde{\varphi}_u = \widetilde{\varphi}_{u_0}}} |c_u|^2 \\
 & \geq |c_{u_0}|^2
 \end{aligned}$$

I is the identity operator of whatever state space U operates on, while $P_{\widetilde{\varphi_{u_0}}}$ is the orthonormal projector onto the space generated by the vector $|\widetilde{\varphi_{u_0}}\rangle$ of the computational basis. Besides, following the analysis of the book, $\widetilde{\varphi_{u_0}}$ is an approximation to φ_{u_0} to an accuracy 2^{-n} with probability at least $1 - \epsilon$ if we make use of $t = n + \lceil \log(2 + \frac{1}{2\epsilon}) \rceil$ bits in the first register. We conclude we get the desired approximation of φ_{u_0} at the end of the phase estimation algorithm with probability at least $|c_{u_0}|^2(1 - \epsilon)$.