

Languages and Algorithms for Artificial Intelligence (Third Module)

Introduction

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ALMA MATER STUDIORUM
UNIVERSITÀ DI BOLOGNA



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Section 1

The Logistics

Organisation

- ▶ **Webpage**

- ▶ At IOL: <http://iol.unibo.it/>

- ▶ **Email**

- ▶ ugo.dallago@unibo.it

- ▶ **Office Hours**

- ▶ *Where:* see <http://www.cs.unibo.it/~dallago>
 - ▶ *When:* there is no fixed office hours, just write me an email and we will fix an appointment.

- ▶ **Teaching Assistant**

- ▶ Francesco Gavazzo, francesco.gavazzo2@unibo.it
 - ▶ Please contact both of us if you need help, this way your request will be taken into account ASAP.

Structure and Schedule

- ▶ The *schedule* of this module's lectures will vary along the coming weeks.
 - ▶ Please keep an eye to to the online schedule.
- ▶ **Content:** the module is divided into two parts:
 1. The very basics of the theory of *computational complexity*, with implications for AI.
 2. Some rudiments of *computational learning theory*.
- ▶ **Requirements:**
 - ▶ We assume all students know what a *program* is and what an *algorithm* is.
 - ▶ We also assume familiarity with the basics of *probability theory*, *discrete mathematics* (natural numbers, induction, graphs, and the like) and *mathematical logic* (propositional logic and its semantics).
- ▶ All lectures will be given by Ugo Dal Lago.
- ▶ Some of the exercise sessions will be given by Francesco Gavazzo

Textbooks and Exams

► Textbooks

The module will be as self contained as possible, but there are plenty of nice books from which you can learn more than what we will do, if you are interested.

- Sanjeev Arora and Boaz Barak. *Computational Complexity: A Modern Approach*. Cambridge University Press. 2007.
- Christos Papadimitriou. *Computational Complexity*. Addison-Wesley. 1994.
- Shai Shalev-Shwartz and Shai Ben-David. *Understanding Machine Learning: from Theory to Algorithms* Cambridge University Press. 2014.
- Michael Kearns and Umesh Vazirani. *An Introduction to Computational Learning Theory* The MIT Press. 1994.

► Exams

- You need to pass a *written exam*.
- As you know, you will get just *one* mark for the whole of this course, and the three modules have the same weight.

Section 2

What? Why?

Is There any Limit on what AI can do?

- ▶ **AI and ML are Everywhere.**

- ▶ The great impact to our society of technologies derived from Artificial Intelligence and Machine Learning is in front of us.
- ▶ Learning the way tools from AI and ML can be employed in a variety of application scenarios (vision, social-media, recommendation systems, etc.) is of course important for you and your career.
- ▶ A number of courses in the Master Degree you are enrolled in are specifically about the above.

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- ▶ **This is *not* What this Module is About!**

- ▶ The question we will try to ask you is rather the following:
is there any intrinsic limit to what we can do with AI and ML, or to the accuracy or efficiency of algorithms solving a give problem?

A Paradigm Shift

**“Algorithm X solves Problem Y Efficiently
and with Great Accuracy”**

A Paradigm Shift

“Problem Y Cannot be Solved by any
Algorithm X ”.



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“Algorithm X solves Problem Y Efficiently
and with Great Accuracy”



“Any Algorithm Solving Y is bound to be
Inefficient, or not Accurate”.

The Value of Negative Results

A Quote from Thomas A. Edison

“Negative results are just what I want. They’re just as valuable to me as positive results. I can never find the thing that does the job best until I find the ones that don’t.”

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- ▶ If we know that the task we want to accomplish is intrinsically *hard*, and we know *in what* its hardness resides, we can:
 1. Direct our work towards methodologies *specifically designed* for hard tasks.
 2. *Avoid* spending our time trying to design solutions that cannot exist.
 3. Inform the stakeholders *in advance* that the solution we are going to design will likely be inefficient.
 4. Decide that the task is *too difficult*, and that it cannot be accomplished the way we want.
 5. Decide that it is better to switch to a *relaxed* version of the task, which instead admits reasonable solutions.

Example Problems - I

Natural Number Multiplication

Suppose you want to write a **Python** program which multiplies two positive integer numbers, which can both be expressed as n -digits numbers, but which are represented as lists rather than scalars. Suppose that the operations at your disposal are just the basic arithmetic operations *on digits*, so that you cannot just multiply the two numbers. How should you write your program?

Example Problems - II

Maximizing the Number of Invitees

Suppose you are going to marry your boyfriend (or girlfriend) soon, and you want to decide the list of invitees. Since your soon-to-be father-in-law will pay all bill, you have all the interest in keeping the list of invitees as large as possible. But actually, it would be impossible to invite all the people in the set $L = \{a_1, \dots, a_m\}$ of all candidate invitees, since some of them are kind of incompatible, and cannot be invited together, i.e. there is another set $I = \{(b_1, c_1), \dots, (b_n, c_n)\} \subseteq L \times L$ whose elements are those candidate invitees pairs which are incompatible. Would it be possible to maximize the length of the list of invitees?

Example Problems - III

Learning Programs with Finite Inputs and Outputs

Suppose your ML teacher asks you to write a learning algorithm capable of reconstructing a program by querying it as a black box, without looking at the code. You can assume that the program takes in input a list of booleans of a fixed (rather big) length n , and produces in output one boolean value. How would you proceed? On which input would you query the program? Would it be possible to query the program on significantly less inputs than the 2^n possible ones?

The Module's Structure

1. Introduction to Computational Complexity

- ▶ Historical, Conceptual, and Mathematical Preliminaries
- ▶ The Computational Model, and Uncomputability
- ▶ The Class P
- ▶ The Class NP and NP-Completeness

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2. Some Rudiments of Computational Learning Theory

- ▶ From Computation to Learning
- ▶ Positive and Negative Results about the Learning Process
- ▶ The VC Dimension

Thank You!

Questions?