

NetXPTO - LinkPlanner

20 de Março de 2018

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Capítulo 1

Preface

Capítulo 2

Introduction

LinkPlanner is devoted to the simulation of point-to-point links.

Capítulo 3

Simulator Structure

LinkPlanner is a signals open-source simulator. Developed in C++14 language.

The IDE use is the MS Visual Studio, with the Visual Studio 2015 (v140) Platform Toolset, and with a Target Platform Version 8.1. These two parameters must be specified in the solution properties in the Visual Studio. The Platform Toolset v140 uses the Tool Version 14.0 by default.

The major entity is the system.

A system comprises a set of blocks.

The blocks interact with each other through signals.

3.1 System

3.2 Blocks

3.3 Signals

List of available signals:

- Signal

Capítulo 4

Development Cycle

The NetXPTO-LinkPlanner has been developed by several people using git as a version control system. The NetXPTO-LinkPlanner repository is located in the GitHub site <http://github.com/netxpto/linkplanner>. The more updated functional version of the software is in the branch master. Master should be considered a functional beta version of the software. Periodically new releases are delivered from the master branch under the branch name Release<Year><Month><Day>. The integration of the work of all people is performed by Armando Nolasco Pinto in the branch Develop. Each developer has his/her own branch with his/her name.

4.1 Development Language

LinkPlanner is a signals open-source simulator. Developed in C++14 language.

4.2 Development Environment

The IDE use is the MS Visual Studio, with the Visual Studio 2015 (v140) Platform Toolset, and with a Target Platform Version 8.1. These two parameters must be specified in the solution properties in the Visual Studio. The Platform Toolset v140 uses the Tool Version 14.0 by default.

Capítulo 5

Visualizer

visualizer

Capítulo 6

Case Studies

6.1 QPSK Transmitter

2017-08-25, Review, Armando Nolasco Pinto

This system simulates a QPSK transmitter. A schematic representation of this system is shown in figure 6.1.

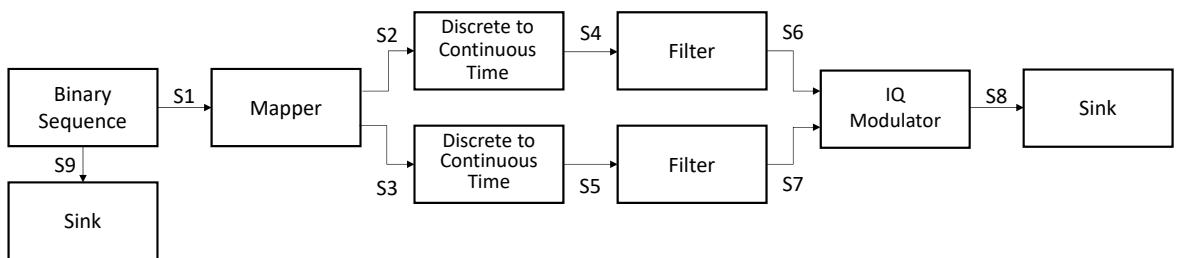


Figura 6.1: QPSK transmitter block diagram.

System Input Parameters

Parameter: *sourceMode*

Description: Specifies the operation mode of the binary source.

Accepted Values: PseudoRandom, Random, DeterministicAppendZeros, DeterministicCyclic.

Parameter: *patternLength*

Description: Specifies the pattern length used by the source in the PseudoRandom mode.

Accepted Values: Integer between 1 and 32.

Parameter: *bitStream*

Description: Specifies the bit stream generated by the source in the DeterministicCyclic and DeterministicAppendZeros mode.

Accepted Values: "XXX..", where X is 0 or 1.

Parameter: *bitPeriod*

Description: Specifies the bit period, i.e. the inverse of the bit-rate.

Accepted Values: Any positive real value.

Parameter: *iqAmplitudes*

Description: Specifies the IQ amplitudes.

Accepted Values: Any four pair of real values, for instance $\{ \{ 1,1 \}, \{ -1,1 \}, \{ -1,-1 \}, \{ 1,-1 \} \}$, the first value correspond to the "00", the second to the "01", the third to the "10" and the forth to the "11".

Parameter: *numberOfBits*

Description: Specifies the number of bits generated by the binary source.

Accepted Values: Any positive integer value.

Parameter: *numberOfSamplesPerSymbol*

Description: Specifies the number of samples per symbol.

Accepted Values: Any positive integer value.

Parameter: *rollOffFactor*

Description: Specifies the roll off factor in the raised-cosine filter.

Accepted Values: A real value between 0 and 1.

Parameter: *impulseResponseTimeLength*

Description: Specifies the impulse response window time width in symbol periods.

Accepted Values: Any positive integer value.

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6.2 BPSK Transmission System

Student Name	:	André Mourato (2018/01/28 - 2018/02/27) Daniel Pereira (2017/09/01 - 2017/11/16)
Goal	:	Estimate the BER in a Binary Phase Shift Keying optical transmission system with additive white Gaussian noise. Comparison with theoretical results.
Directory	:	sdf/bpsk_system

Binary Phase Shift Keying (BPSK) is the simplest form of Phase Shift Keying (PSK), in which binary information is encoded into a two state constellation with the states being separated by a phase shift of π (see Figure 6.2).

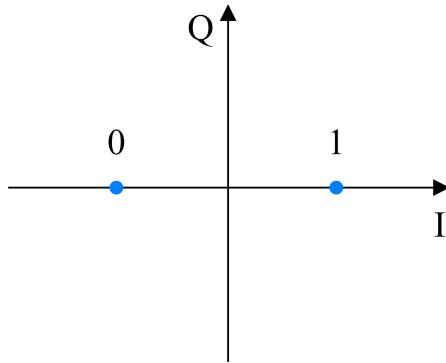


Figura 6.2: BPSK symbol constellation.

White noise is a random signal with equal intensity at all frequencies, having a constant power spectral density. White noise is said to be Gaussian (WGN) if its samples follow a normal distribution with zero mean and a certain variance σ^2 . For WGN its spectral density equals its variance. For the purpose of this work, additive WGN is used to model thermal noise at the receivers.

The purpose of this system is to simulate BPSK transmission in back-to-back configuration with additive WGN at the receiver and to perform an accurate estimation of the BER and validate the estimation using theoretical values.

6.2.1 Theoretical Analysis

The output of the system with added gaussian noise follows a normal distribution, whose first probabilistic moment can be readily obtained by knowledge of the optical power of the received signal and local oscillator,

$$m_i = 2\sqrt{P_L P_S} G_{ele} \cos(\Delta\theta_i), \quad (6.1)$$

where P_L and P_S are the optical powers, in watts, of the local oscillator and signal, respectively, G_{ele} is the gain of the trans-impedance amplifier in the coherent receiver and

$\Delta\theta_i$ is the phase difference between the local oscillator and the signal, for BPSK this takes the values π and 0, in which case (6.1) can be reduced to,

$$m_i = (-1)^{i+1} 2 \sqrt{P_L P_S} G_{ele}, \quad i = 0, 1. \quad (6.2)$$

The second moment is directly chosen by inputting the spectral density of the noise σ^2 , and thus is known *a priori*.

Both probabilist moments being known, the probability distribution of measurement results is given by a simple normal distribution,

$$f(x) = \frac{1}{\sqrt{2\pi}\sigma} e^{-\frac{(x-m_i)^2}{2\sigma^2}}. \quad (6.3)$$

The BER is calculated in the following manner,

$$BER = \frac{1}{2} \int_0^{+\infty} f(x|\Delta\theta = \pi) dx + \frac{1}{2} \int_{-\infty}^0 f(x|\Delta\theta = 0) dx, \quad (6.4)$$

given the symmetry of the system, this can be simplified to,

$$BER = \int_0^{+\infty} f(x|\Delta\theta = \pi) dx = \frac{1}{2} \operatorname{erfc} \left(\frac{-m_0}{\sqrt{2}\sigma} \right) \quad (6.5)$$

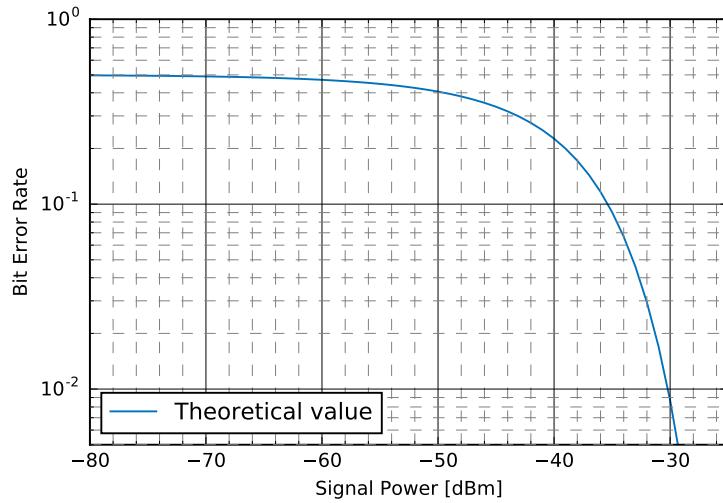


Figura 6.3: Bit Error Rate in function of the signal power in dBm at a constant local oscillator power level of 0 dBm.

6.2.2 Simulation Analysis

A diagram of the system being simulated is presented in the Figure 6.4. A random binary sequence is generated and encoded in an optical signal using BPSK modulation. The decoding of the optical signal is accomplished by an homodyne receiver, which combines the

signal with a local oscillator. The received binary signal is compared with the transmitted binary signal in order to estimate the Bit Error Rate (BER). The simulation is repeated for multiple signal power levels. Each corresponding BER is recorded and plotted against the expectation value.

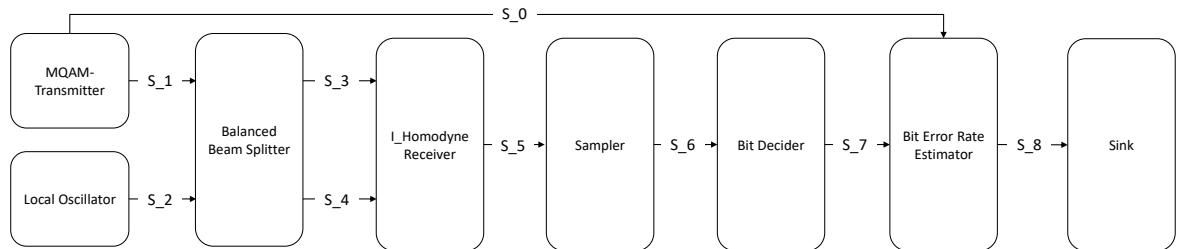


Figura 6.4: Overview of the BPSK system being simulated.

Required files

Header Files		
File	Comments	Status
add_20171116.h		✓
balanced_beam_splitter_20180124.h		✓
binary_source_20180118.h		✓
bit_decider_20170818.h		✓
bit_error_rate_20171810.h		✓
discrete_to_continuous_time_20180118.h		✓
i_homodyne_receiver_20180124.h		✓
ideal_amplifier_20180118.h		✓
iq_modulator_20180130.h		✓
local_oscillator_20180130.h		✓
m_qam_mapper_20180118.h		✓
m_qam_transmitter_20180118.h		✓
netxpto_20180118.h		✓
photodiode_old_20180118.h		✓
pulse_shaper_20180118.h		✓
sampler_20171116.h		✓
sink_20180118.h		✓
super_block_interface_20180118.h		✓
ti_amplifier_20180102.h		✓
white_noise_20180118.h		✓

Source Files		
File	Comments	Status
add_20171116.cpp		✓
balanced_beam_splitter_20180124.cpp		✓
binary_source_20180118.cpp		✓
bit_decider_20170818.cpp		✓
bit_error_rate_20171810.cpp		✓
discrete_to_continuous_time_20180118.cpp		✓
i_homodyne_receiver_20180124.cpp		✓
ideal_amplifier_20180118.cpp		✓
iq_modulator_20180130.cpp		✓
local_oscillator_20180130.cpp		✓
m_qam_mapper_20180118.cpp		✓
m_qam_transmitter_20180118.cpp		✓
netxpto_20180118.cpp		✓
photodiode_old_20180118.cpp		✓
pulse_shaper_20180118.cpp		✓
sampler_20171116.cpp		✓
sink_20180118.cpp		✓
super_block_interface_20180118.cpp		✓
ti_amplifier_20180102.cpp		✓
white_noise_20180118.cpp		✓

System Input Parameters

This system takes into account the following input parameters:

System Input Parameters		
Parameter	Default Value	Comments
numberOfBitsReceived	-1	
numberOfBitsGenerated	1000	
samplesPerSymbol	16	
pLength	5	
bitPeriod	20×10^{-12}	
rollOffFactor	0.3	
signalOutputPower_dBm	-20	
localOscillatorPower_dBm	0	
localOscillatorPhase	0	
iqAmplitudesValues	$\{ \{-1, 0\}, \{1, 0\} \}$	
transferMatrix	$\{ \{\frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{-1}{\sqrt{2}}\} \}$	
responsivity	1	
amplification	10^6	
electricalNoiseAmplitude	$5 \times 10^{-4}\sqrt{2}$	
samplesToSkip	$8 \times \text{samplesPerSymbol}$	
bufferLength	20	
shotNoise	false	

Inputs

This system takes no inputs.

Outputs

This system outputs the following objects:

System Output Signals	
Signal	Associated File
Initial Binary String (S_0)	S0.sgn
Optical Signal with coded Binary String (S_1)	S1.sgn
Local Oscillator Optical Signal (S_2)	S2.sgn
Beam Splitter Outputs (S_3, S_4)	S3.sgn & S4.sgn
Homodyne Receiver Electrical Output (S_5)	S5.sgn
Sampler Output (S_6)	S6.sgn
Decoded Binary String (S_7)	S7.sgn
BER Result String (S_8)	S8.sgn
Report	Associated File
Bit Error Rate Report	BER.txt

Bit Error Rate - Simulation Results

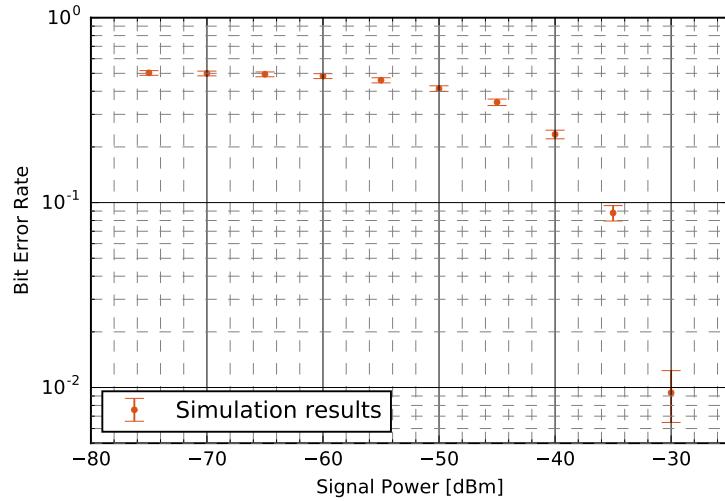


Figura 6.5: Bit Error Rate in function of the signal power in dBm at a constant local oscillator power level of 0 dBm.

6.2.3 Comparative Analysis

The following results show the dependence of the error rate with the signal power assuming a constant Local Oscillator power of 0 dBm, the signal power was evaluated at levels between -70 and -25 dBm, in steps of 5 dBm between each. The simulation results are presented in orange with the computed lower and upper bounds, while the expected value, obtained from (6.5), is presented as a full blue line. A close agreement is observed between the simulation results and the expected value. The noise spectral density was set at $5 \times 10^{-4}\sqrt{2} \text{ V}^2$ [1].

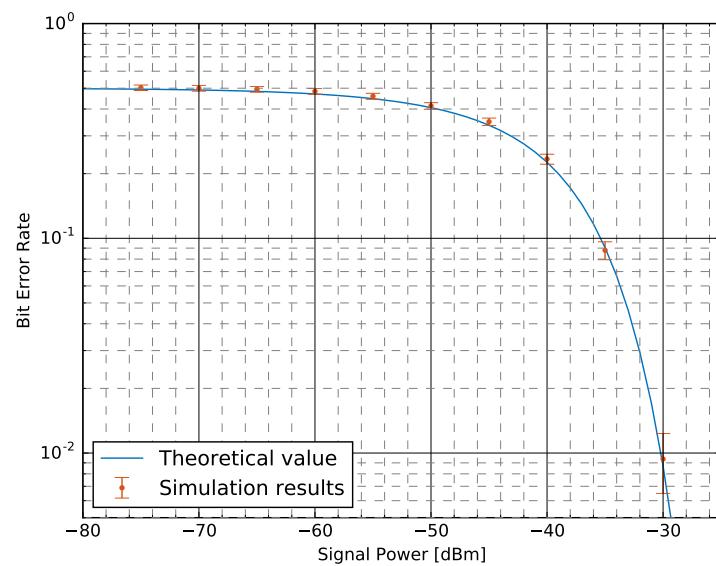


Figura 6.6: Bit Error Rate in function of the signal power in dBm at a constant local oscillator power level of 0 dBm. Theoretical values are presented as a full blue line while the simulated results are presented as a errorbar plot in orange, with the upper and lower bound computed in accordance with the method described in 7.4

References

- [1] Thorlabs. *Thorlabs Balance Amplified Photodetectors: PDB4xx Series Operation Manual*, 2014.

6.3 M-QAM Transmission System

Student Name	:	Ana Luisa Carvalho
Goal	:	M-QAM system implementation with BER measurement and comparison with theoretical values.
Directory	:	sdf/m_qam_system

The goal of this project is to simulate a Quadrature Amplitude Modulation transmission system with M points in the constellation diagram (M-QAM) and to perform a Bit Error Rate (BER) measurement that can be compared with theoretical values.

M-QAM systems can encode $\log_2 M$ bits per symbol which means they can transmit higher data rates keeping the same bandwidth when compared, for example, to PSK systems. However, because the states are closer together, these systems can be more susceptible to noise.

The Bit Error Rate (BER) is a measurement of how a bit stream is altered by a transmission system due to noise (among other factors). To study this effect we introduced Additive White Gaussian Noise (AWGN) to model thermal noise at the receiver.

For $M = 4$ the M-QAM system reduces to a Quadrature Phase Shift Keying system (QPSK) system that uses four equispaced points in the constellation diagram (see figure 6.7).

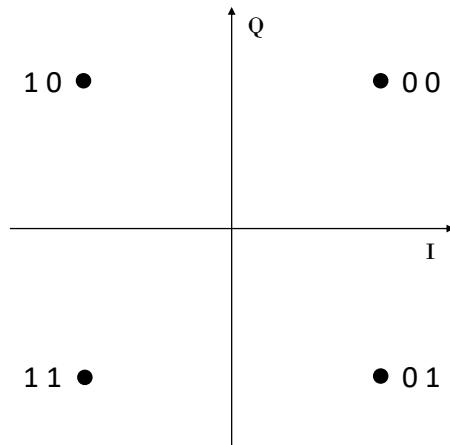


Figura 6.7: 4-QAM constellation points.

6.3.1 Theoretical Analysis

M-QAM is a modulation scheme that takes advantage of two carriers (usually sinusoidal waves) with a phase difference of $\frac{\pi}{2}$. The resultant output consists of a signal with both amplitude and phase variations. The two carriers, referred to as I (In-phase) and Q (Quadrature), can be represented as

$$I = A \cos(\phi) \quad (6.6)$$

$$Q = A \sin(\phi) \quad (6.7)$$

which means that any sinusoidal wave can be decomposed in its I and Q components:

$$A \cos(\omega t + \phi) = A (\cos(\omega t) \cos(\phi) - \sin(\omega t) \sin(\phi)) \quad (6.8)$$

$$= I \cos(\omega t) - Q \sin(\omega t), \quad (6.9)$$

where we have used the expression for the cosine of a sum and the definitions of I and Q.

The probability of symbol error, P_s , in coherent M-PSK demodulation with AWGN is given by

$$P_s = 2 Q \left(\sqrt{2 \log_2 M \left(\frac{E_b}{n_0} \right) \sin^2 \frac{\pi}{M}} \right) \quad (6.10)$$

where E_b is the energy of one bit, n_0 is the noise power and the function Q is defined as

$$Q(x) = \frac{1}{2} \operatorname{erfc} \left(\frac{x}{\sqrt{2}} \right). \quad (6.11)$$

The probability of bit errors, P_b is related to P_s by

$$P_b = \frac{1}{\log_2 M} P_s. \quad (6.12)$$

For QPSK we get, using $M = 4$ in equations 6.11 and 6.12,

$$P_b = \frac{1}{2} \operatorname{erfc} \left(\sqrt{\frac{2 E_b}{n_0}} \right). \quad (6.13)$$

This function is plotted in figure 6.8 for $n_0 = 10^{-6}$.

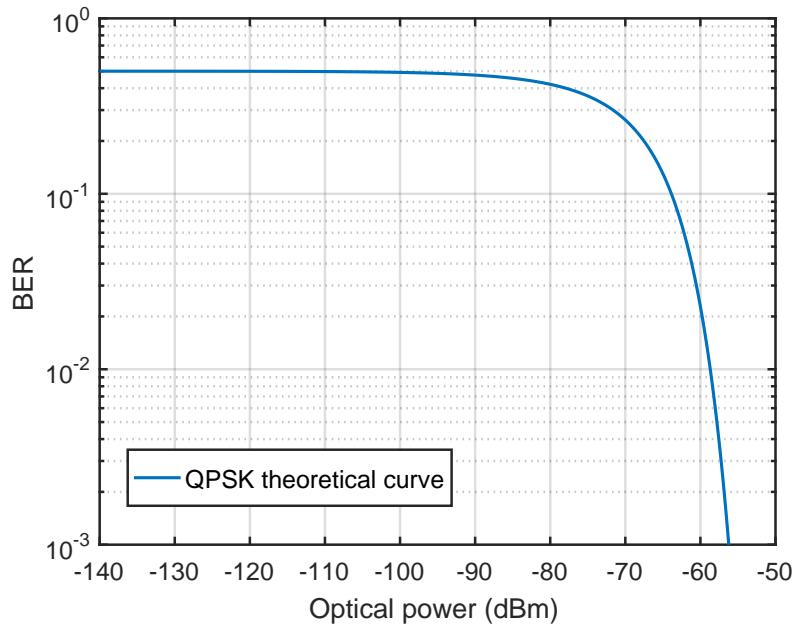


Figura 6.8: QPSK theoretical BER values as a function of the output optical power in dBm.

6.3.2 Simulation Analysis

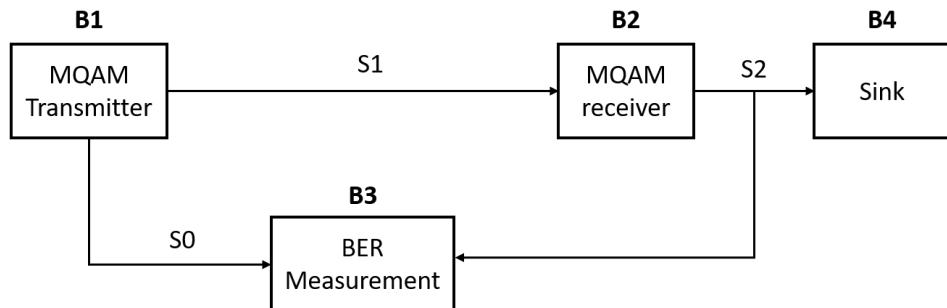


Figura 6.9: Schematic representation of the MQAM system.

The M-QAM transmission system is a complex block of code that simulates the modulation, transmission and demodulation of an optical signal using M-QAM modulation. It is composed of four blocks: a transmitter, a receiver, a sink and a block that performs a Bit Error Rate (BER) measurement. The schematic representation of the system is presented in figure 6.9.

Current state: The system currently being implemented is a QPSK system ($M=4$).

Future work: Extend this block to include other values of M .

Tabela 6.1: Main system files

System blocks	Source file	Header file	Status
Main	m_qam_system_sdf.cpp	—	✓
M-QAM transmitter	m_qam_transmitter.cpp	m_qam_transmitter.h	✓
M-QAM receiver	homodyne_receiver.cpp	homodyne_receiver.h	
Sink	sink.cpp	sink.h	✓
BER estimator	bit_error_rate.cpp	bit_error_rate.h	

Functional description

A complete description of the M-QAM transmitter and M-QAM homodyne receiver blocks can be found in the *Library* chapter of this document as well as a detailed description of the independent blocks that compose these blocks.

The M-QAM transmitter generates one or two optical signals by encoding a binary string using M-QAM modulation. It also outputs a binary signal that is used to perform the BER measurement.

The M-QAM homodyne receiver accepts one input optical signal and outputs a binary signal. It performs the M-QAM demodulation of the input signal by combining the optical signal with a local oscillator.

The demodulated optical signal is compared to the one produced by the transmitter in order to estimate the Bit Error Rate (BER).

The files corresponding to each of the system's blocks are summarized in table 6.1. Along with the library and corresponding source files these allow for the full operation of the M-QAM system.

Required Files

The required header and source files needed to run this system are summarized in table 6.2.

Input Parameters

The system accepts several input parameters that can be defined by the user. These are described in table 6.3.

Simulation results

In this section we show the eye diagrams for the S1 signals for two different values of the output optical power.

6.3.3 Comparative Analysis

In this section we show the simulation results and compared them with the theoretical predictions. Figures 6.11 shows the variation of the BER with the power of the signal, using

Tabela 6.2: Required files

Header file	Source file	Description	Status
add.h	add.cpp	Adds two signals.	✓
binary_source.h	binary_source.cpp	Produces a binary sequence.	✓
bit_error_rate.h	bit_error_rate.cpp	Computes the BER and writes it to a text file.	✓
discrete_to_continuous_time.h	discrete_to_continuous_time.cpp	Converts a signal from discrete in time to continuous in time.	✓
homodyne_receiver.h	m_qam_homodyne_receiver.cpp		
ideal_amplifier.h	ideal_amplifier.cpp	Amplifies the signal.	✓
iq_modulator.h	iq_modulator.cpp	Divides the signal in its quadrature and in phase components	✓
local_oscillator.h	local_oscillator.cpp		
m_qam_mapper.h	m_qam_mapper.cpp	Maps the signal using the defined constellation	✓
m_qam_transmitter.h	m_qam_transmitter.cpp		✓
netxpto.h	netxpto.cpp	General class that contains definition from signals and buffers.	✓
optical_hybrid.h	optical_hybrid.cpp	Implements an optical hybrid.	✓
photodiode_old.h	photodiode_old.cpp	Pair of photodiodes and current subtraction.	✓
pulse_shaper.h	pulse_shaper.cpp	Electrical filter.	✓
sampler_20171119.h	sampler_20171119.cpp	Samples the signal.	✓
sink.h	sink.cpp	Deletes signal.	✓
super_block_interface.h	super_block_interface.cpp		✓
white_noise.h	white_noise.cpp	Generates white gaussian noise.	✓

4000 bits and a pseudorandom binary sequence with pattern length 2^7 . To produce this plots we considered a noise amplitude of 10^{-6} and an amplification of 10^3 .

Tabela 6.3: Input parameters

Parameter	Type	Description
numberOfBitsGenerated	t_integer	Determines the number of bits to be generated by the binary source
samplesPerSymbol	t_integer	Number of samples per symbol
prbsPatternLength	int	Determines the length of the pseudorandom sequence pattern (used only when the binary source is operated in <i>PseudoRandom</i> mode)
bitPeriod	t_real	Temporal interval occupied by one bit
rollOffFactor	t_real	Parameter of the raised cosine filter
signalOutputPower_dBm	t_real	Determines the power of the output optical signal in dBm
numberOfBitsReceived	int	Determines when the simulation should stop. If -1 then it only stops when there is no more bits to be sent
iqAmplitudeValues	vector<t_iqValues>	Determines the constellation used to encode the signal in IQ space
symbolPeriod	double	Given by bitPeriod / samplesPerSymbol
localOscillatorPower_dBm	t_real	Power of the local oscillator
responsivity	t_real	Responsivity of the photodiodes (1 corresponds to having all optical power transformed into electrical current)
amplification	t_real	Amplification provided by the ideal amplifier
noiseAmplitude	t_real	Amplitude of the white noise
samplesToSkip	t_integer	Number of samples to be skipped by the <i>sampler</i> block
confidence	t_real	Determines the confidence limits for the BER estimation
midReportSize	t_integer	
bufferLength	t_integer	Corresponds to the number of samples that can be processed in each run of the system

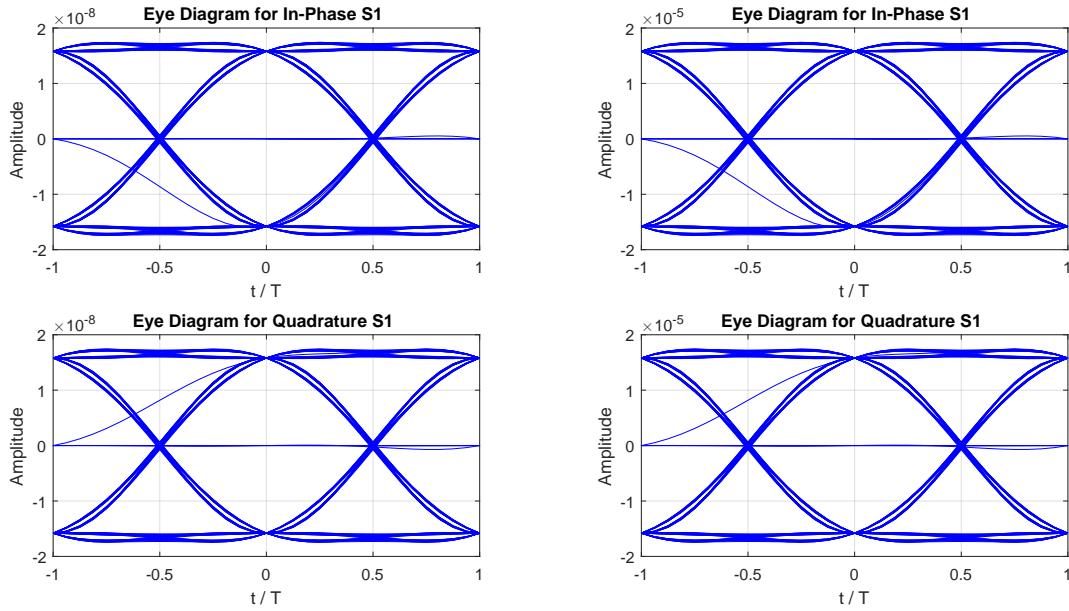


Figura 6.10: Eye diagrams for the S1 bandpass signal with an output optical power of -120dBm (left) and -60dBm (right). Note differente scales on y axis.

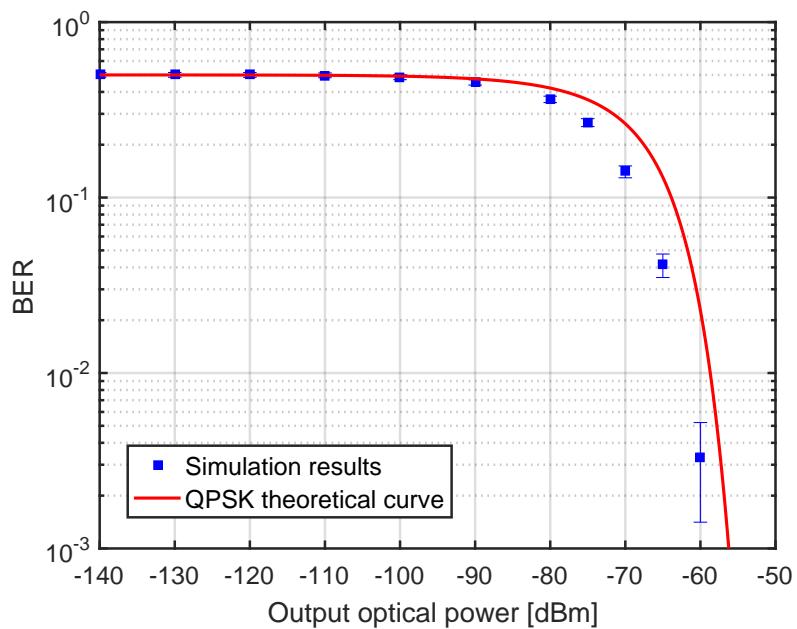


Figura 6.11: Simulation result for a random binary sequence with 4000 bits, a noise amplitude of 10^{-6} and an amplification of 10^3 .

6.4 Quantum Random Number Generator

Students Name : Mariana Ramos (12/01/2018 -)
Goal : sdf/quantum_random_number_generator.

True random numbers are indispensable in the field of cryptography to guarantee the security of the communication protocols [?]. There are two approaches for random number generation: the pseudorandom generation which are based on an algorithm implemented on a computer, and the physical random generators which consist in measuring some physical observable with random behaviour. Since classical physics description is deterministic, all classical processes are in principle predictable. Therefore, a true random number generator must be based on a quantum process.

In this chapter, it is presented the theoretical analysis and simulation of a quantum random generator by measuring the polarization of single photons with a polarizing beam splitter.

6.4.1 Theoretical Analysis

Nowadays, the only known way to generate truly random numbers is by building a physical source by using quantum mechanical decisions, since the occurrence of each individual result of such a quantum mechanical decision is truly random, ie it is inestimable or unknowable [?]. One of the optical processes available as a source of randomness is the splitting of a polarized single photon beam.

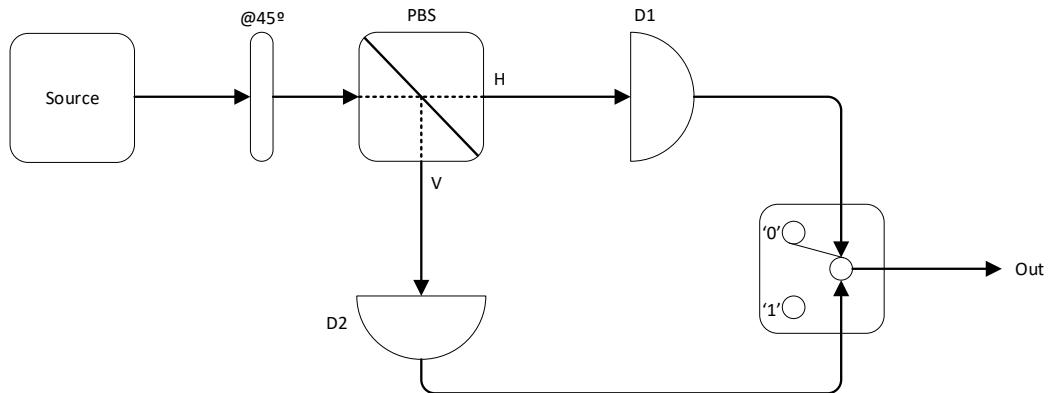


Figura 6.12: Source of randomness with a polarization beam splitter PBS where the incoming light is linearly polarized at 45° with respect to the PBS. Figure adapted from [?].

The principle of operation of the random generator is shown in figure 6.12. Each individual photon coming from the source is linearly polarized at 45° and has equal probability of found in the horizontal polarization (H) or in the vertical polarization (V) output of the PBS. Quantum theory estimates for both cases the individual choices are truly

random and independent one from each other. This way, the detection of the photons in each output of the polarization beam splitter is done with single photon detectors and combining the detection pulse in a switch, which has two possible states: "0" or "1". When the detector **D1** fires, the switch is flipped to state "0" and does not move until a detection event in detector **B2** occurs and it does not move until a detection occurs in detector **D1**. In the case of some detections occur in a row in the same detector, only the first detection clicks and the following detections leave the switch unaltered.

6.4.2 Simulation Analysis

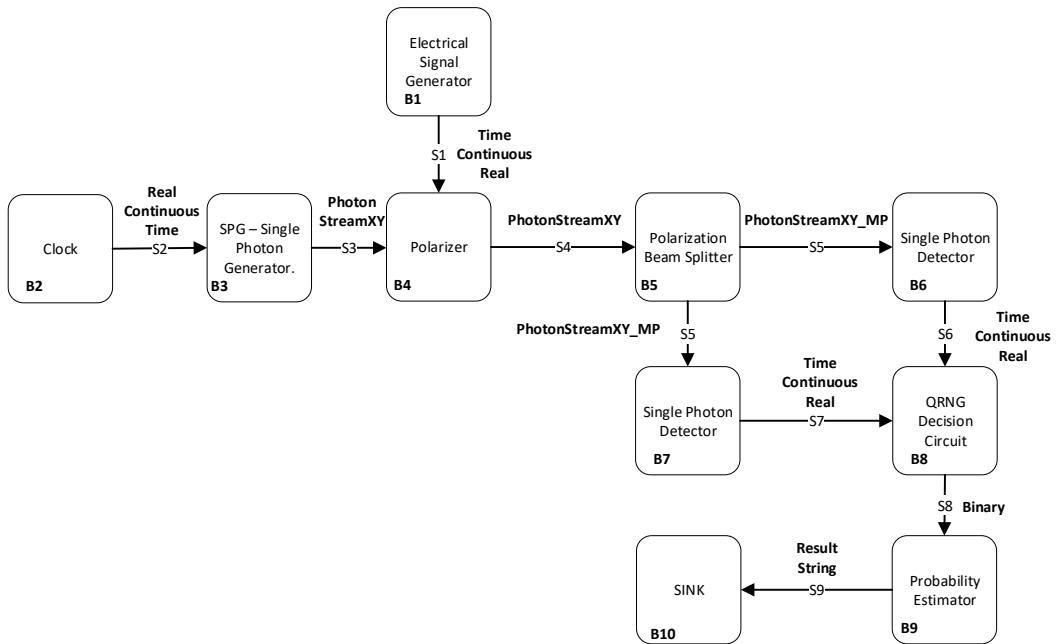


Figura 6.13: Block diagram of the simulation of a Quantum Random Generator.

The simulation diagram of the setup described in the previous section is presented in figure 6.13. The linear polarizer has an input control signal (S1) which allows to change the rotation angle. Nevertheless, the only purpose is to generate a time and amplitude continuous real signal with the value of the rotation angle in degrees.

In addition, the photons are generated by single photon source block at a rate defined by the clock rate.

At the end of the simulation there is a circuit decision block which will outputs a binary signal with value "0" if the detector at the end of the horizontal path clicks or "1" if the detector at the end of the vertical path clicks.

In table 6.4 are presented the input parameters of the system.

Tabela 6.4: System Input Parameters

Parameter	Default Value
RateOfPhotons	1e6
NumberOfSamplesPerSymbol	16
PolarizerAngle	45.0

In table 6.5 are presented the system signals to implement the simulation presented in figure 6.13.

Tabela 6.5: System Signals

Signal name	Signal type
S1	TimeContinuousAmplitudeContinuousReal
S2	TimeContinuousAmplitudeContinuousReal
S3	PhotonStreamXY
S4	PhotonStreamXY
S5	PhotonStreamXYMP
S6	TimeContinuousAmplitudeContinuousReal
S7	TimeContinuousAmplitudeContinuousReal
S8	Binary
S9	Binary

Table 6.6 presents the header files used to implement the simulation as well as the specific parameters that should be set in each block. Finally, table 6.7 presents the source files.

Tabela 6.6: Header Files

File name	Description	Status
netxpto.h		✓
electrical_signal_generator_20180124.h	setFunction(), setGain()	✓
clock_20171219.h	ClockPeriod(1 / RateOfPhotons)	✓
polarization_beam_splitter_20180109.h		✓
polarizer_20180113.h		✓
single_photon_detector_20180111.h	setPath(0), setPath(1)	✓
single_photon_source_20171218.h		✓
probability_estimator_20180124.h		✓
sink.h		✓
qrng_decision_circuit.h		✓

Tabela 6.7: Source Files

File name	Description	Status
netxpto.cpp		✓
electrical_signal_generator_20180124.cpp		✓
clock_20171219.cpp		✓
polarization_beam_splitter_20180109.cpp		✓
polarizer_20180113.cpp		✓
single_photon_detector_20180111.cpp		✓
single_photon_source_20171218.cpp		✓
probability_estimator_20180124.cpp		✓
sink.cpp		✓
qrng_decision_circuit.cpp		✓
qrng_sdf.cpp		✓

In theory, considering the results space Ω associated with a random experience and A an event such that $P(A) = p \in]0, 1[$. Lets $X : \Omega \rightarrow \mathbb{R}$ such that

$$\begin{aligned} X(\omega) &= 1 & \text{if } \omega \in A \\ X(\omega) &= 0 & \text{if } \omega \in \bar{A} \end{aligned} \tag{6.14}$$

This way,

$$\begin{aligned} P(X = 1) &= P(A) = p \\ P(X = 0) &= P(\bar{A}) = 1 - p \end{aligned} \tag{6.15}$$

X follows the Bernoulli law with parameter p , $X \sim \mathbf{B}(p)$, being the expected value of the Bernoulli random value $E(X) = p$ and the variance $\text{VAR}(X) = p(1 - p)$ [?].

Lets calculate the margin error for N of samples in order to obtain X inside a specific confidence interval, which in this case we assume 99%.

We will use *z-score* (in this case 2.576, since a confidence interval of 99% was chosen) to calculate the expected error margin,

$$E = z_{\alpha/2} \frac{\sigma}{\sqrt{N}},$$

where, E is the error margin, $z_{\alpha/2}$ is the *z-score* for a specific confidence interval, $\sigma = \sqrt{\text{VAR}(X)} = \sqrt{\hat{p}(1 - \hat{p})}$ is the standard deviation and N the number of samples. Lets assume, for an angle of 45° , a number of samples $N = 1 \times 10^6$ and the expected probability of reach each detector of $\hat{p} = 0.5$. We have an error margin of $E = 1.288 \times 10^{-3}$, which is acceptable. This way, the simulation will be performed for $N = 1 \times 10^6$ samples for different

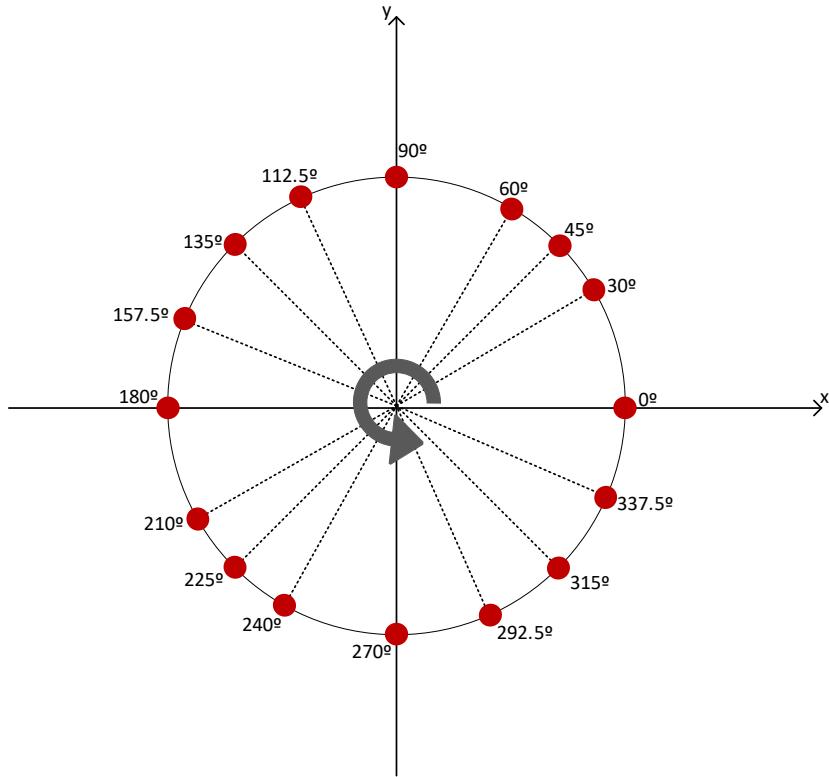


Figura 6.14: Angles used to perform the qrng simulation for $N = 1 \times 10^8$ samples.

angles of polarization shown in figure 6.14 with different error margin's values since the expected probability changes depending on the polarization angle.

For a quantum random number generator with equal probability of obtain a "0" or "1" the polarizer must be set at 45° . This way, we have 50% possibilities to obtain a "0" and 50% of possibilities to obtain a "1". This theoretical value meets the value obtained from the simulation when it is performed for the number of samples mentioned above.

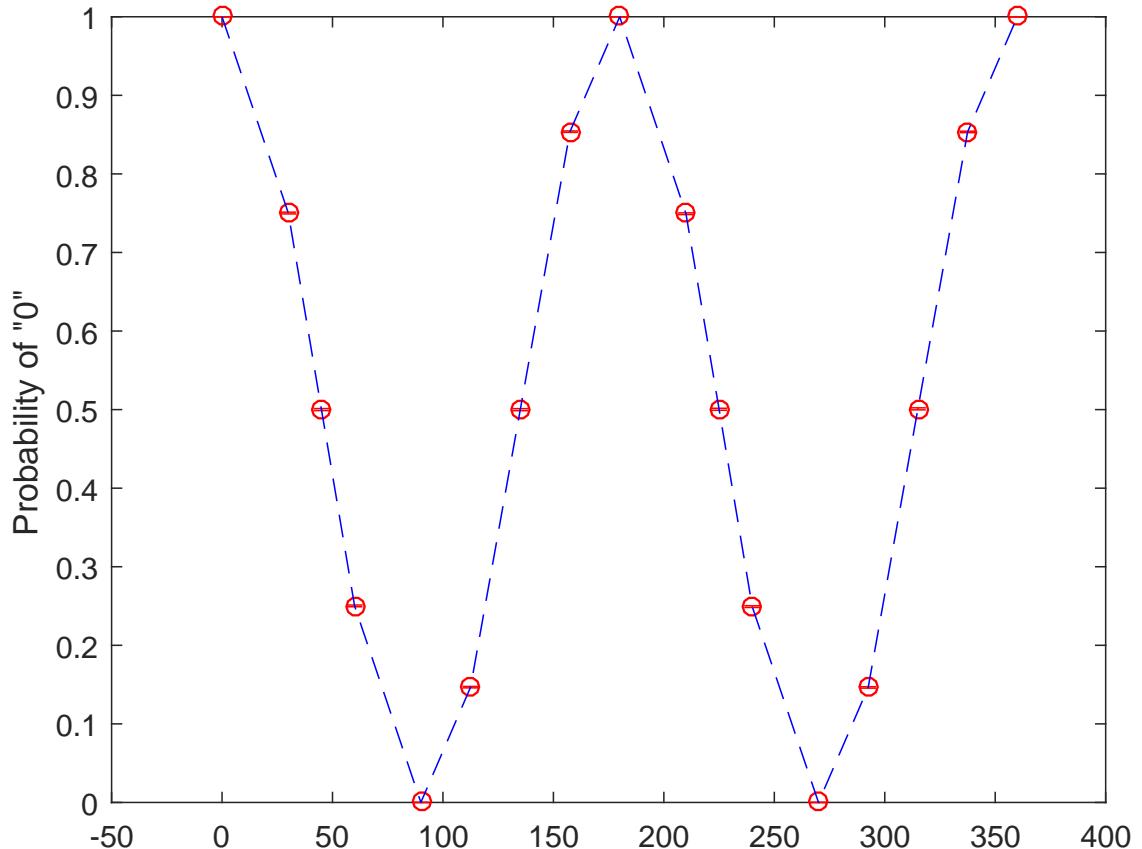


Figura 6.15: Probability of outputs a number "0" depending on the polarization angle.

Figure 6.15 shows the probability of a single photon reaches the detector placed on Horizontal axis depending on the polarization angle of the photon, and this way the output number is "0". On the other hand, figure 6.16 shows the probability of a single photon reaches the detector placed on Vertical component of the polarization beam splitter, and this way the output number is "1". As we can see in the figures the two detectors have complementary probabilities, i.e the summation of both values must be equals to 1.

One can see that "Probability of 1" behaves almost like a sine function and "Probability of 0" behaves almost like a cosine function with a variable angle.

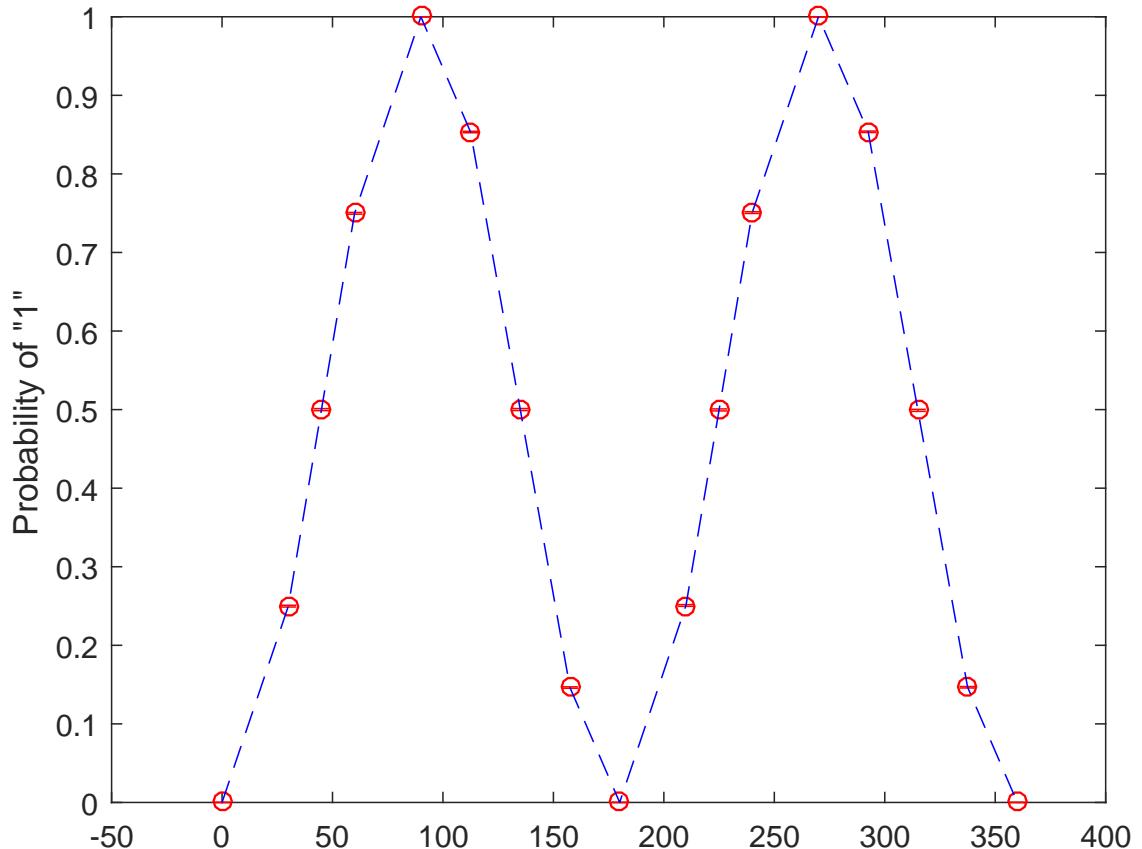


Figura 6.16: Probability of outputs a number "1" depending on the polarization angle.

6.4.3 Experimental Analysis

In order to have a real experimental quantum random number generator, a setup shown in figure 6.17 was built in the lab. To simulate a single photon source we have a CW-Pump laser with 1550 nm wavelength followed by an interferometer Mach-Zehnder in order to have a pulsed beam. The interferometer has an input signal given by a Pulse Pattern Generator. This device also gives a clock signal for the Single Photon Detector (APD-Avalanche Photodiode) which sets the time during which the window of the detector is open. After the MZM there is a Variable Optical Attenuator (VOA) which reduces the amplitude of each pulse until the probability of one photon per pulse is achieved. Next, there is a polarizer controller followed by a Linear Polarized, which is set at 45° , then a Polarization Beam Splitter (PBS) and finally, one detector at the end of each output of the PBS. The output

signals from the detector will be received by a Processing Unit. Regarding to acquired the output of the detectors, there is an oscilloscope capable of record 1×10^6 samples.

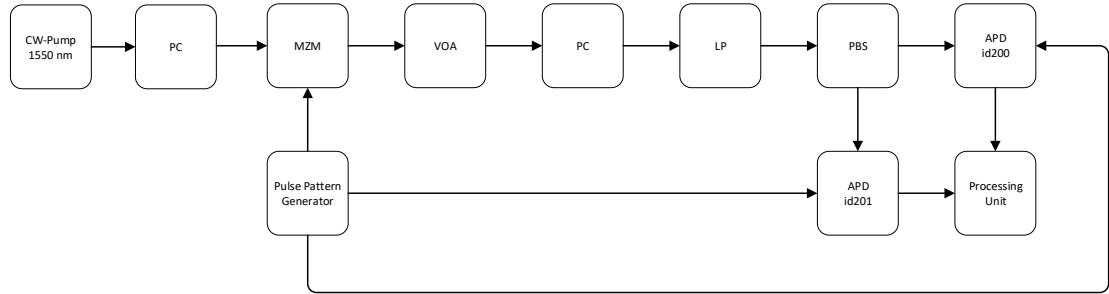


Figura 6.17: Experimental setup to implement a quantum random number generator.

6.4.4 Open Issues

References

- [1] Thorlabs. *Thorlabs Balance Amplified Photodetectors: PDB4xx Series Operation Manual*, 2014.

6.5 BB84 with Discrete Variables

Students Name	:	Mariana Ramos and Kevin Filipe
Starting Date	:	November 7, 2017
Goal	:	BB84 implementation with discrete variables.

BB84 is a key distribution protocol which involves three parties, Alice, Bob and Eve. Alice and Bob exchange information between each other by using a quantum channel and a classical channel. The main goal is continuously build keys only known by Alice and Bob, and guarantee that eavesdropper, Eve, does not gain any information about the keys.

6.5.1 Protocol Analysis

Students Name	:	Kevin Filipe (7/11/2017)
Goal	:	BB84 - Protocol Description

BB84 protocol was created by Charles Bennett and Gilles Brassard in 1984 [?]. It involves two parties, Alice and Bob, sharing keys through a quantum channel in which could be accessed by a eavesdropper, Eve. A basic model is depicted in figure 6.18.

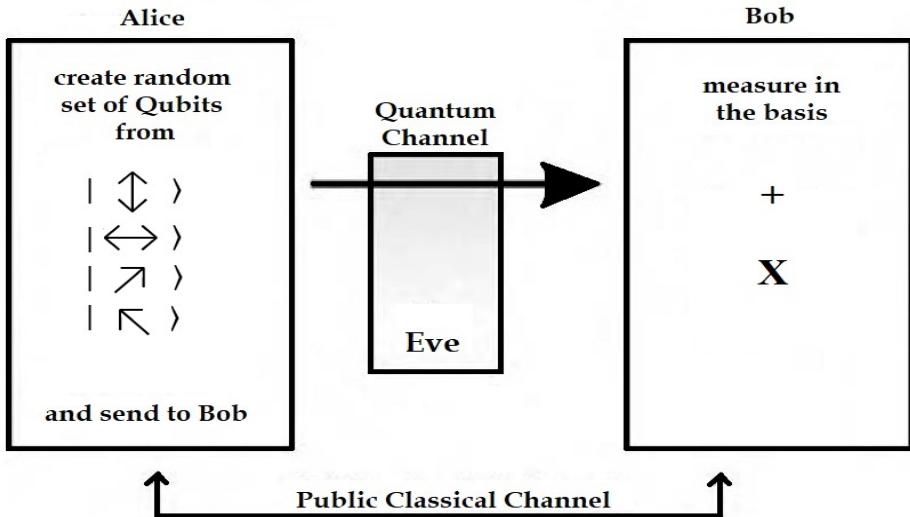


Figura 6.18: Basic QKD Model. Alice and Bob are connected by 2 communication channels, quantum and classical, with an eavesdropper, Eve, in the quantum communication channel (figure adapted from [?]).

BB84 protocol uses bit encoding into photon state polarization. Two non-orthogonal basis are used to encode the information, the rectilinear and diagonal basis, + and x respectively. The following table shows this bit encoding.

	<i>Rectilinear Basis, +</i>	<i>Diagonal Basis, x</i>
0	0	-45
1	90	45

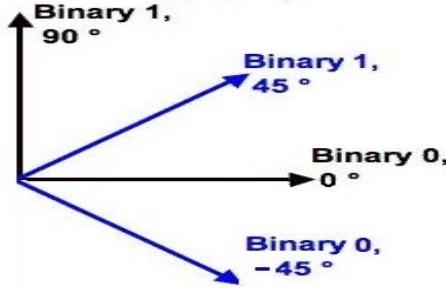


Figura 6.19: Simple representation of the bit encoding using the corresponding bases. [?]

The protocol is implemented with the following steps:

1. Alice generates two random bit strings. The random string, R_{A1} , corresponds to the data to be encoded into photon state polarization. R_{A2} is a random string in which 0 and 1 corresponds to the rectilinear, $+$, and diagonal, \times , basis of B_A , respectively.

$$R_{A1} = \{0, 1, 1, 0, 1, 0, 0, 1, 1, 0, 1, 1, 1, 0, 0, 1, 0, 0, 0, 1\}$$

$$R_{A2} = \{0, 0, 1, 0, 1, 1, 1, 0, 1, 1, 1, 0, 1, 0, 0, 0, 1, 0, 1, 0\}$$

=

$$\{+, +, \times, +, \times, \times, +, \times, \times, \times, +, \times, +, +, +, \times, +, \times, +, +\}$$

2. Alice transmits a train of photons, S_{AB} , obtained by encoding the bits, R_{A1} with the respective photon polarization state R_{A2} .

$$S_{AB} = \{\rightarrow, \uparrow, \nwarrow, \rightarrow, \nwarrow, \nearrow, \nearrow, \uparrow, \nwarrow, \nearrow, \nwarrow, \uparrow, \nwarrow, \rightarrow, \rightarrow, \rightarrow, \uparrow, \nearrow, \rightarrow, \nearrow, \uparrow\}.$$

3. Bob generates a random string, R_B , to receive the photon trains with the correspondent basis.

$$R_B = \{0, 1, 1, 1, 0, 1, 0, 0, 1, 1, 0, 0, 1, 1, 0, 0, 1, 1, 0, 0\}$$

=

$$\{+, \times, \times, \times, +, \times, +, +, \times, \times, +, +, \times, \times, +, +, \times, +, +\}.$$

4. Bob performs the incoming photon states measurement, M_B with its generated random basis, R_B . If the two photon detectors don't click, means the bit was lost during transference and so there is attenuation. If both photon detectors click, a false positive was detected. In the measurements, M_B , the attenuation is represented by a -1 and the false positives to -2. The measurements done in rectilinear or diagonal basis are represented by 0 and 1, respectively.

$$M_B = \{0, 1, 1, 1, -1, 1, 0, 0, -2, 1, 0, 0, -2, 1, 0, 0, 1, -1, 0, 0\}$$

The second phase, uses the classical communication channel:

1. After the measurement, Bob sends to Alice the values of M_B .
2. Alice remove the bits from R_{A2} corresponding to -1 and -2, performs a negated XOR and mixes the bits using a known algorithm by Alice and Bob.

R_{A2}	0	0	1	0	1	1	0	1	1	0	0	0	0	1	1	0
R_B	0	1	1	1	1	0	0	1	0	0	1	0	0	1	0	0
BAB	1	0	1	0	1	0	1	1	0	1	0	1	1	1	0	1

3. Using also the negated XOR, Bob also mixes the bits with the same algorithm used by Alice and obtains the same key, K_{AB}

$$K_{AB} = \{0, 1, 0, 1, 0, 1, 0, 1, 0, 1\}.$$

To determine the Quantum Bit Error Rate (QBER), Bob will reveal a bit sequence from the deduced key to Alice. Alice then returns to Bob the estimated QBER value, mQBER, with a confidence interval, [qLB, qUB]. To verify if the channel is reliable or not the flowchart presented in figure 6.20 must be followed and a acceptable QBER limit, qLim, must be imposed. Alice performs the QBER and confidence interval calculation by using the equations in the Bit Error Rate section, but applied to this protocol. By following the presented flowchart if the qLimit is inside the confidence interval, Bob should use more bits from key until qLim is bigger than qLB and qUB. If this condition is achieved, the QBER is successfully estimated and a reliable communication can be achieved. Otherwise the link is declared as unreliable and no further communication is made.

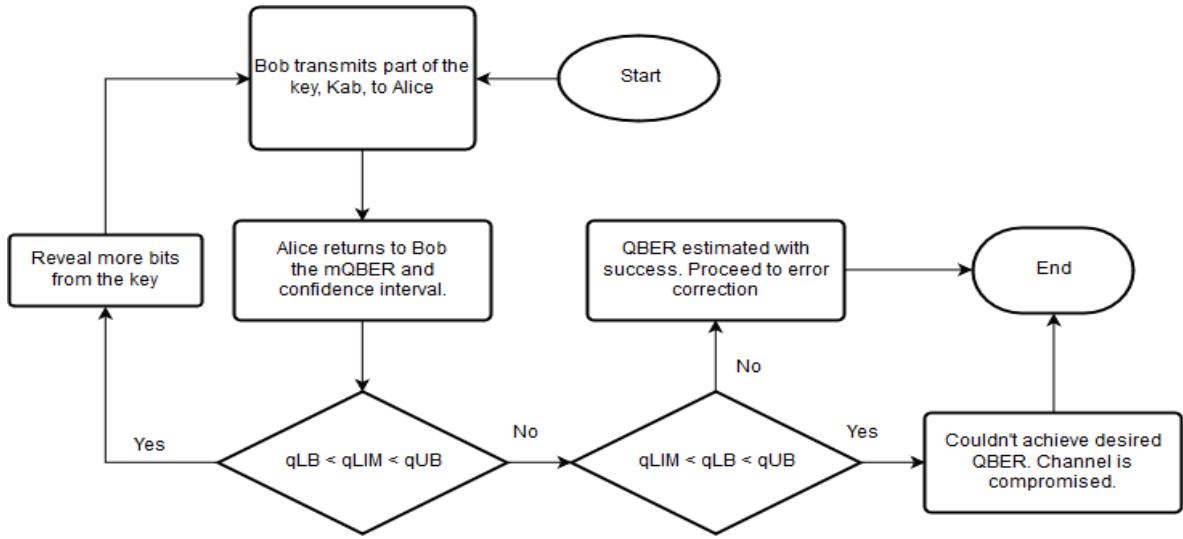


Figura 6.20: Flowchart to determine if the channel is reliable or not.

The presence of a eavesdropper will carry the risk of changing the bits. This will produce disagreement between Bob and Alice in the bits they should agree. When Eve measures and retransmits a photon she can deduce correctly with a probability of 50%. So by learning the correct polarization of half of the photons, the induced error is 25%. Alice and Bob can detect Eve presence by sacrifice the secrecy of some bits in order to test.

References

- [1] Bennett, C. H. and Brassard, G. Quantum Cryptography: Public key distribution and coin tossing. International Conference on Computers, Systems and Signal Processing, Bangalore, India, 10-12 December 1984, pp. 175-179.
- [2] Mart Haitjema, A Survey of the Prominent Quantum Key Distribution Protocols
- [3] Christopher Gerry, Peter Knight, "Introductory Quantum Optics"Cambridge University Press, 2005

Basis	
0	+
1	×

	Basis "+"
0	$\rightarrow (0^\circ)$
1	$\uparrow (90^\circ)$

	Basis "x"
0	$\searrow (-45^\circ)$
1	$\nearrow (45^\circ)$

1. Alice randomly generate a bit sequence with length ks being, in this case, $k = 2$ and $s = 4$ as it was defined at the beginning. Therefore, she must define two sets randomly: S_{A1} which contains the basis values; and S_{A2} , which contains the key values.

In that case, lets assume she gets the following sets S_{A1} and S_{A2} :

$$S_{A1} = \{0, 1, 1, 0, 0, 1, 0, 1\},$$

$$S_{A2} = \{1, 1, 0, 0, 0, 1, 0, 0\}.$$

2. Next, Alice sends to Bob throughout a quantum channel ks photons encrypted using the basis defined in S_{A1} and according to the keys defined in S_{A2} .

In the current example, Alice sends the photons, throughout a quantum channel, according to the following,

$$S_{AB} = \{\uparrow, \nearrow, \searrow, \rightarrow, \rightarrow, \nearrow, \rightarrow, \searrow\}.$$

$$S_{AB} = \{90^\circ, 45^\circ, -45^\circ, 0^\circ, 0^\circ, 45^\circ, 0^\circ, -45^\circ\}.$$

3. Bob also randomly generates ks bits, which are going to define his measurement basis, S_{B1} . He will measure the photons sent by Alice. Lets assume:

$$S_{B1} = \{0, 1, 0, 1, 0, 1, 1, 1\}.$$

When Bob receives photons from Alice, he measures them using the basis defined in S_{B1} . In the current example, S_{B1} corresponds to the following set:

$$\{+, \times, +, \times, +, \times, \times, \times\}.$$

Bob will get ks results:

$$S_{B1'} = \{1, 1, 0, 1, 0, 1, 1, 0\}.$$

4. Bob will send a *Hash Function* result HASH1 to Alice. This value will do Bob's commitment with the measurements done. In this case, this *Hash Function* is calculated from *SHA-256* algorithm for each pair (Basis from S_{B1} and measured value from $S_{B1'}$), i.e Bob sends to Alice sk pairs as his commitment. In this case, Bob sends eight pairs encoded using a *Hash Function* which is also send to Alice. From that moment on Bob cannot change his commitment neither the basis which he uses to measure the photons sent by Alice.
5. Once Alice has received the confirmation of measurement from Bob, she sends throughout a classical channel the basis which she has used to codify the photons, which in this case we assumed $S_{A1} = \{0, 1, 1, 0, 0, 1, 0, 1\}$.
6. In order to know which photons were measured correctly, Bob does the operation $S_{B2} = S_{B1} \oplus S_{A1}$. In the current example the operation will be:

$$\begin{array}{c|cccccccc} S_{B1} & 0 & 1 & 0 & 1 & 0 & 1 & 1 & 1 \\ S_{A1} & 0 & 1 & 1 & 0 & 0 & 1 & 0 & 1 \\ \hline \oplus & 1 & 1 & 0 & 0 & 1 & 1 & 0 & 1 \end{array}$$

In this way, Bob gets

$$S_{B2} = \{1, 1, 0, 0, 1, 1, 0, 1\}.$$

When Bob uses the right basis he gets the values correctly, when he uses the wrong basis he just guess the value. The values "1" correspond to the values he measured correctly and "0" to the values he just guessed.

Next, Bob sends to Alice, through a classical channel, information about the minimum number between "ones" and "zeros", i.e

$$n = \min(\#0, \#1) = 3,$$

where $\#0$ represents the number of zeros in S_{B2} and $\#1$ the number of ones in S_{B2} . At this time, Alice must be able to know if Bob is being honest or not. Therefore, she will open Bob's commitment from *step 4* and she verify if the number n sent by Bob is according with the commitment values sent by him. In other words, she opens a number of pairs committed by Bob which is known from the beginning.

7. If $n < s$, being s the message's size, Alice and Bob will repeat the steps from 1 to 7. In this case, $n = 3$ which is smaller than $s = 4$. Therefore, Alice and Bob repeat the steps from 1 to 7 in order to enlarge Bob's sets S_{B1} and S_{B2} as well as Alice's sets S_{A1} and S_{A2} .
8. Lets assume :

$$S_{B1} = \{1, 1, 0, 0, 0, 1, 0, 0, 1, 0, 0, 0, 0, 1, 1\}.$$

At Alice's side the new sets S_{A1} , which contains the basis values, and S_{A2} , which contains the key values, will be the following:

$$S_{A1} = \{0, 1, 1, 0, 0, 1, 0, 1, 1, 1, 0, 0, 1, 1, 1, 0\},$$

$$S_{A2} = \{1, 1, 0, 0, 0, 1, 0, 0, 1, 0, 1, 0, 0, 0, 1, 1\}.$$

Finally, for $S_{B2} = S_{B1} \oplus S_{A1}$ Bob gets the following sequence:

$$S_{B2} = \{1, 1, 0, 0, 1, 1, 0, 1, 0, 1, 0, 0, 1, 1, 0, 1\}.$$

Note that the sets were enlarge in the second iteration.

9. At this time, Bob sends again to Alice, through a classical channel, the minimum number between "ones" and "zeros", $n = \min(\#0, \#1)$. In this case, n is equal to 7 which is the number of zeros.
10. Alice checks if $n > s$ and acknowledge to Bob that she already knows that $n > s$. In this case, $n = 7$ and $s = 4$ being $n > s$ a valid condition.
11. Next, Bob defines two new sub-sets, I_0 and I_1 . I_0 is a set of values with photons array positions which Bob just guessed the measurement since he did not measure them with the same basis as Alice, I_1 is a set of values with photons array positions which Bob measured correctly since he used the same basis as Alice used to encoded them.

In this example, Bob defines two sub-sets with size $s = 4$:

$$I_0 = \{3, 4, 7, 11\},$$

and

$$I_1 = \{2, 5, 6, 13\},$$

where I_0 is the sequence of positions in which Bob was wrong about basis measurement and I_1 is the sequence of positions in which Bob was right about basis measurement. Bob sends to Alice the set S_b

Thus, if Bob wants to know m_0 he must send to Alice throughout a classical channel the set $S_0 = \{I_1, I_0\}$, otherwise if he wants to know m_1 he must send to Alice throughout a classical channel the set $S_1 = \{I_0, I_1\}$.

12. With both the received set S_b and the hash function value HASH1, Alice must be able to prove that Bob has being honest.
13. Lets assume Bob sent $S_0 = \{I_1, I_0\}$. Alice defines two encryption keys K_0 and K_1 using the values in positions defined by Bob in the set sent by him. In this example, lets assume:

$$K_0 = \{1, 0, 1, 0\}$$

$$K_1 = \{0, 0, 0, 1\}.$$

Alice does the following operations:

$$m = \{m_0 \oplus K_0, m_1 \oplus K_1\}.$$

$$\begin{array}{c|cccc} m_0 & 0 & 0 & 1 & 1 \\ K_0 & 1 & 0 & 1 & 0 \\ \hline \oplus & 1 & 0 & 0 & 1 \end{array}$$

$$\begin{array}{c|cccc} m_1 & 0 & 0 & 0 & 1 \\ K_1 & 0 & 0 & 0 & 1 \\ \hline \oplus & 0 & 0 & 0 & 0 \end{array}$$

Adding the two results, m will be:

$$m = \{1, 0, 0, 1, 0, 0, 0, 0\}.$$

After that, Alice sends to Bob the encrypted message m through a classical channel.

14. When Bob receives the message m , in the same way as Alice, Bob uses S_{B1} , values of positions given by I_1 and I_0 and does the decrypted operation. In this case, he does following operation:

$$\begin{array}{c|cccccccc} m & 1 & 0 & 0 & 1 & 0 & 0 & 0 & 0 \\ & 1 & 0 & 1 & 0 & 0 & 1 & 1 & 0 \\ \hline \oplus & 0 & 0 & 1 & 1 & 0 & 1 & 1 & 0 \end{array}$$

The first four bits corresponds to message 1 and he received $\{0, 0, 1, 1\}$, which is the right message m_0 and $\{0, 1, 1, 0\}$ which is a wrong message for m_1 .

6.5.2 Simulation Setup

6.5.3 Simulation Analysis

Students Name : Mariana Ramos
Starting Date : November 7, 2017
Goal : Perform a simulation of the setup presented bellow in order to implement BB84 communication protocol.

In this sub section the simulation setup implementation will be described in order to implement the BB84 protocol. In figure 6.21 a top level diagram is presented. Then it will be presented the block diagram of the transmitter block (Alice) in figure 6.22, the receiver block (Bob) in figure 6.23 and finally the eavesdropper block (Eve) in figure 6.24.

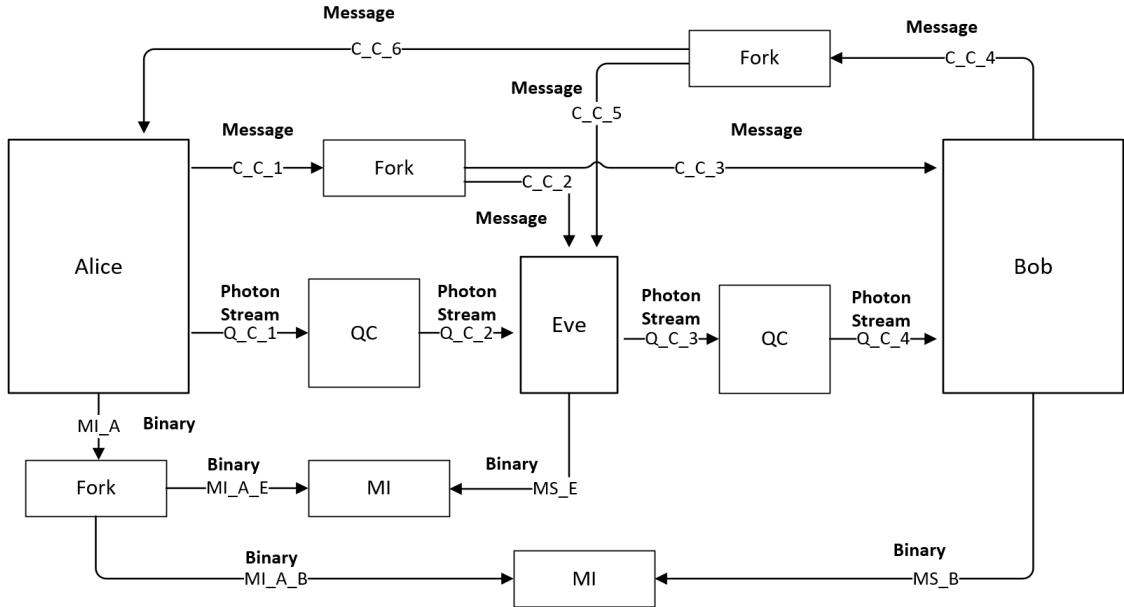


Figura 6.21: Simulation diagram at a top level

Figure 6.21 presents the top level diagram of our simulation. The setup contains three parties, Alice, Eve and Bob where the communication between them is done throughout two classical and one quantum channel. In the middle of the classical channel there is a Fork's diagram which has one input and two outputs. In the case of the classical channel C_C_4 which has the information sent by Bob, the fork's block enables Alice and Eve have access to it. In the quantum communication, the information sent by Alice can be intercepted by Eve and changed by her, or can follow directly to Bob as we can see later in figure 6.24. Furthermore, for mutual information calculation there must be two blocks MI, one to calculate the mutual information between Alice and Eve, and other to calculate the mutual information between Alice and Bob.

In figure 6.22 one can observe a block diagram of the simulation at Alice's side. As it is shown in the figure, Alice must have one block for random number generation which is responsible for basis generation to polarize the photons, and for key random generation in order to have a random state to encode each photon. Furthermore, she has a Processor block for all logical operations: array analysis, random number generation requests, and others. This block also receives the information from Bob after it has passed through a fork's block. In addition, it is responsible for set the initial length l of the first array of photons which will send to Bob. This block also must be responsible for send classical information to Bob. Finally, Processor block will also send a real continuous time signal to single photon generator, in order to generate photons according to this signal, and finally this block also sends to the polarizer a real discrete signal in order to inform the polarizer which basis it should use. Therefore, she has two more blocks for quantum tasks: the single photon generator and the polarizer block which is responsible to encode the photons generated

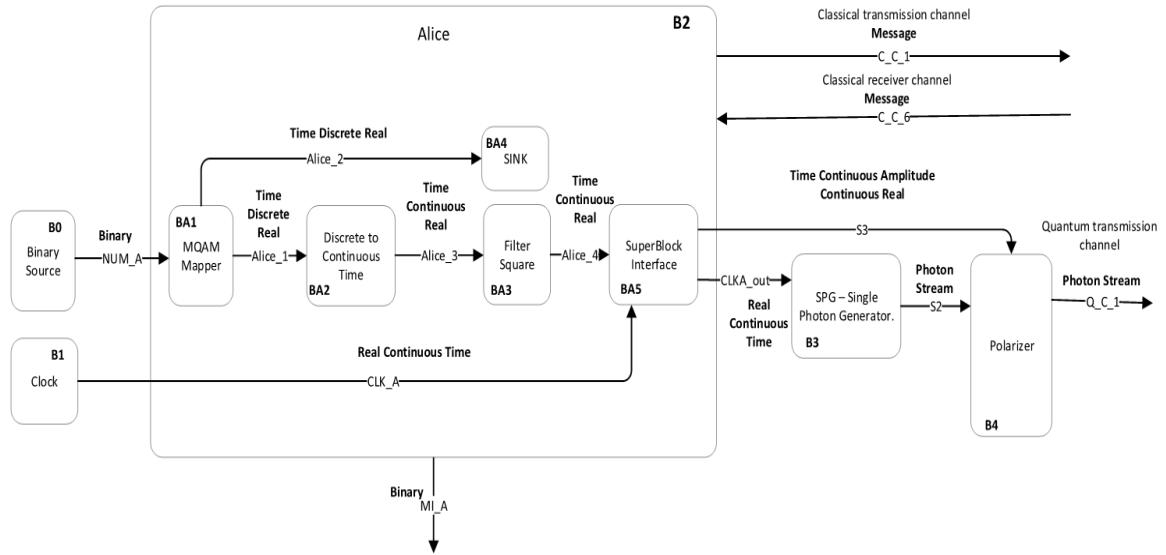


Figura 6.22: Simulation diagram at Alice's side

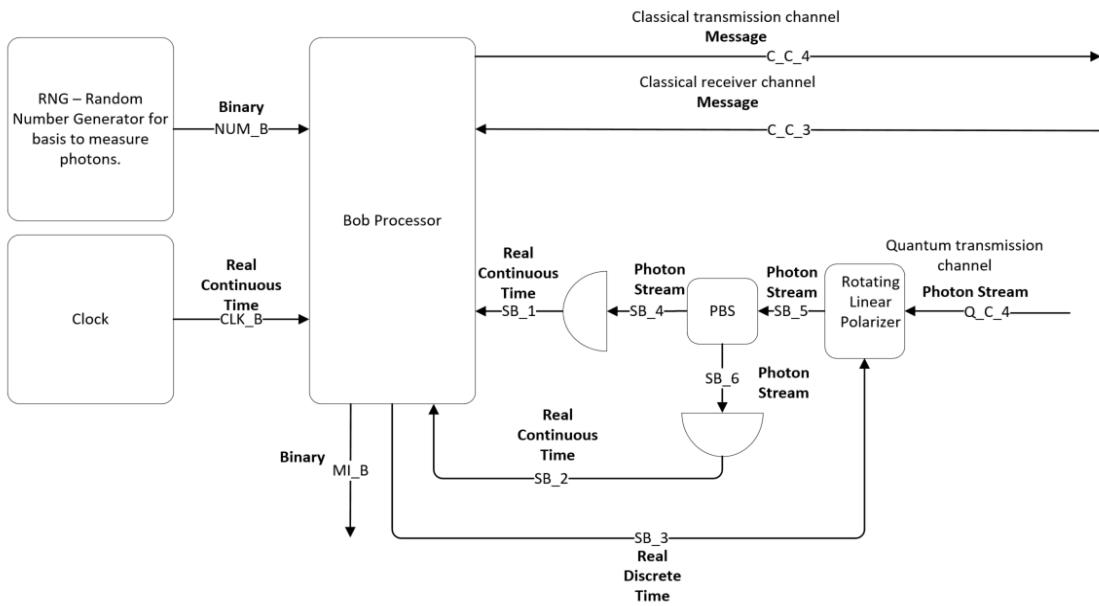


Figura 6.23: Simulation diagram at Bob's side

from the previous block and send them throughout a quantum channel from Alice to Bob.

Finally, Alice's processor has an output to Mutual Information top level block, Ms_A .

In figure 6.23 one can observe a block diagram of the simulation at Bob's side. From this

side, Bob has one block for Random Number Generation which is responsible for randomly generate basis values which Bob will use to measure the photons sent by Alice throughout the quantum channel. Like Alice, Bob has a Processor block responsible for all logical tasks, analysing functions, requests for random number generator block, etc. Additionally, it receives information from Alice throughout a classical channel after passed through a fork's block and a quantum channel. However, Bob only sends information to Alice throughout a classical channel. Furthermore, Bob has one more block for single photon detection which receives from processor block a real discrete time signal, in order to obtain the basis it should use to measure the photons.

Finally, Bob's processor has an output to Mutual Information top level block, Ms_B .

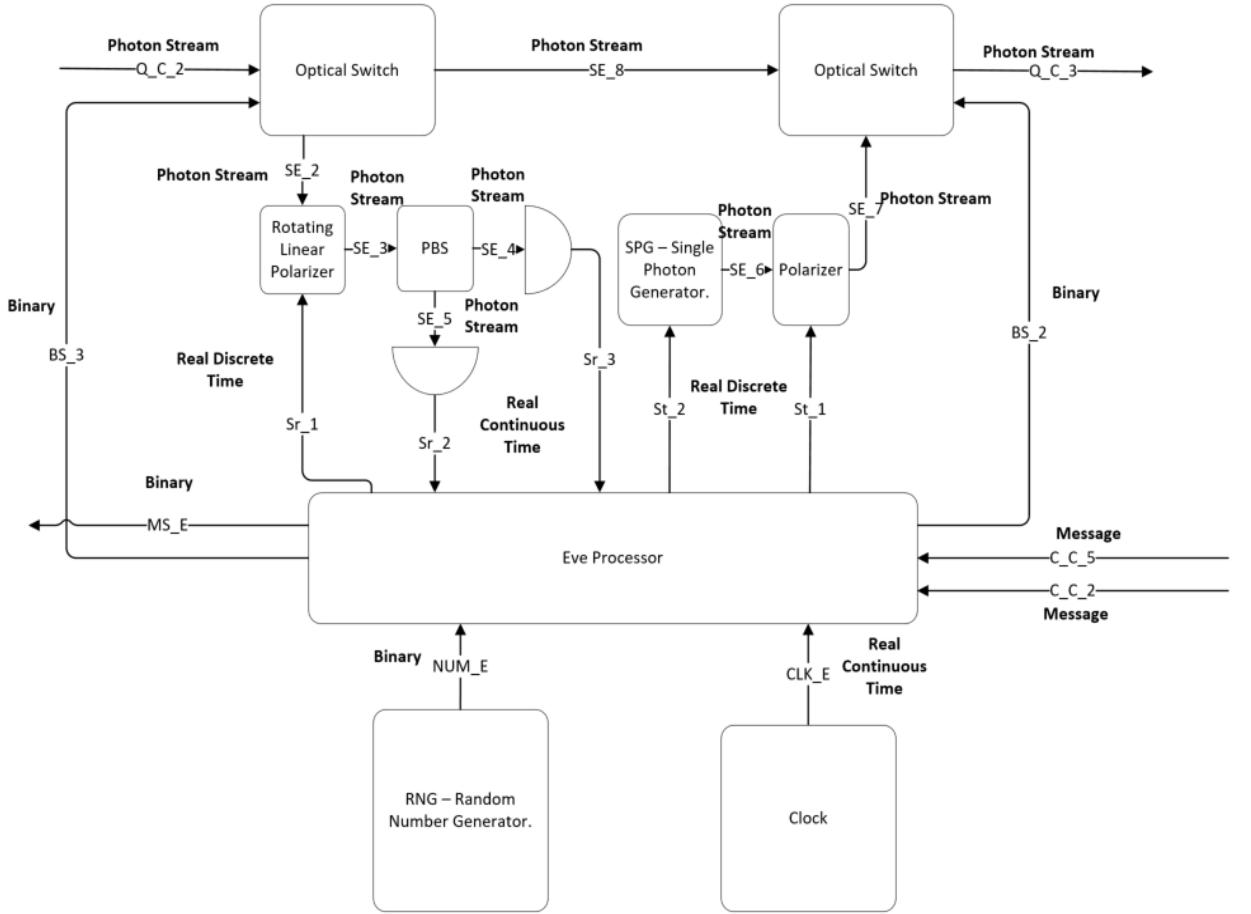


Figura 6.24: Simulation diagram at Eve's side

Figure 6.24 presents the Eve's side diagram. Eve's processor has two receiver classical signals, one from Alice (C_C_2) and other from Bob (C_C_5). About quantum channel, Eve received a quantum message from Alice through the channel Q_C_1 and depends on her

decision the photon can follows directly to Bob or the photon's state can be changed by her. In this case, the photon is received by a block similar to Bob's diagram 6.23 and this block sends a message to Eve's processor in order to reveal the measurement result. After that, Eve's processor sends a message to Alice's diagram similar to figure 6.22 and this block is responsible for encode the photon in a new state. Now, the changed photon is sent to Bob.

In addition, Eve's diagram has one more output Ms_E which is a message sent to the mutual information block as an input parameter.

Tabela 6.8: System Signals

Signal name	Signal type	Status
C_C_1 ... C_C_6	Message	
Q_C_1 .. Q_C_4	Photon Stream	
Ms_A, Ms_B, Ms_E	Binary	
NUM_A , NUM_B, NUM_E	Binary	
CLK_A, CLK_B, CLK_E	Real continuous time	
SB_1, SB_2, Sr_1, Sr_2	Real continuous time	
SA_1, SA_2, St_1, St_2	Real discrete time	
SA_3	Photon Stream	
S_2, S_3, S_4, S_5	Photon Stream	
BS_1, BS_2	Binary	

Table 6.11 presents the system signals as well as them type.

Tabela 6.9: System Input Parameters

Parameter	Default Value	Description
SymbolRate	100K	
NumberOfBits	Number of photons that Alice sends to Bob	

Tabela 6.10: Header Files

File name	Description	Status
binary_source.h		
single_photon_source.h		
single_photon_detector.h		
optical_switch.h		Missing
polarization_beam_splitter.h		
mutual_information.h		Missing
bit_error_rate.h		
clock.h		
fiber.h		
qrng_decision_circuit.h		
message_to_send.h		Missing
message_to_receive.h		Missing
netxpto.h		

Tabela 6.11: Source Files

File name	Description	Status
binary_source.cpp		
single_photon_sourcer.cpp		
single_photon_detector.cpp		
optical_switch.cpp		Missing
polarization_beam_splitter.cpp		
mutual_information.cpp		Missing
bit_error_rate.cpp		
clock.cpp		
fiber.cpp		
qrng_decision_circuit.cpp		
message_to_send.cpp		Missing
message_to_receive.cpp		Missing
netxpto.cpp		
bb84_sdf.cpp		

Capítulo 7

Library

7.1 Add

Header File	:	add.h
Source File	:	add.cpp
Version	:	20180118

Input Parameters

This block takes no parameters.

Functional Description

This block accepts two signals and outputs one signal built from a sum of the two inputs. The input and output signals must be of the same type, if this is not the case the block returns an error.

Input Signals

Number: 2

Type: Real, Complex or Complex_XY signal (ContinuousTimeContinuousAmplitude)

Output Signals

Number: 1

Type: Real, Complex or Complex_XY signal (ContinuousTimeContinuousAmplitude)

7.2 Balanced Beam Splitter

Header File	:	balanced_beam_splitter.h
Source File	:	balanced_beam_splitter.cpp
Version	:	20180124

Input Parameters

Name	Type	Default Value
Matrix	array <t_complex, 4>	$\left\{ \left\{ \frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{1}{\sqrt{2}}, \frac{-1}{\sqrt{2}} \right\} \right\}$
Mode	double	0

Functional Description

The structure of the beam splitter can be controlled with the parameter mode.

When **Mode = 0** the beam splitter will have one input port and two output ports - **1x2 Beam Splitter**. If Mode has a value different than 0, the splitter will have two input ports and two output ports - **2x2 Beam Splitter**.

Considering the first case, the matrix representing a 2x2 Beam Splitter can be summarized in the following way,

$$M_{BS} = \frac{\sqrt{2}}{2} \begin{bmatrix} 1 & 1 \\ 1 & -1 \end{bmatrix} \quad (7.1)$$

The relation between the values of the input ports and the values of the output ports can be established in the following way

$$\begin{bmatrix} A' \\ B' \end{bmatrix} = M_{BS} \begin{bmatrix} A \\ B \end{bmatrix} \quad (7.2)$$

Where, A and B represent the inputs and A' and B' represent the outputs of the Beam Splitter.

Input Signals

Number: 1 or 2

Type: Complex

Output Signals

Number: 2

Type: Complex

References

- [1] Thorlabs. *Thorlabs Balance Amplified Photodetectors: PDB4xx Series Operation Manual*, 2014.

7.3 Binary Source

Header File	:	binary_source.h
Source File	:	binary_source.cpp

This block generates a sequence of binary values (1 or 0) and it can work in four different modes:

- | | |
|-----------------|-----------------------------|
| 1. Random | 3. DeterministicCyclic |
| 2. PseudoRandom | 4. DeterministicAppendZeros |

This blocks doesn't accept any input signal. It produces any number of output signals.

Input Parameters

Parameter	Type	Values	Default
mode	string	Random, PseudoRandom, DeterministicCyclic, DeterministicAppendZeros	PseudoRandom
probabilityOfZero	real	$\in [0,1]$	0.5
patternLength	int	Any natural number	7
bitStream	string	sequence of 0's and 1's	0100011101010101
numberOfBits	long int	any	-1
bitPeriod	double	any	1.0/100e9

Tabela 7.1: Binary source input parameters

Methods

```
BinarySource(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig, OutputSig){};
```

```
void initialize(void);
```

```
bool runBlock(void);
```

```
void setMode(BinarySourceMode m) BinarySourceMode const getMode(void)
```

```
void setProbabilityOfZero(double pZero)
```

```
double const getProbabilityOfZero(void)
```

```
void setBitStream(string bStream)
```

```

string const getBitStream(void)

void setNumberOfBits(long int nOfBits)

long int const getNumberOfBits(void)

void setPatternLength(int pLength)

int const getPatternLength(void)

void setBitPeriod(double bPeriod)

double const getBitPeriod(void)

```

Functional description

The *mode* parameter allows the user to select between one of the four operation modes of the binary source.

Random Mode Generates a 0 with probability *probabilityOfZero* and a 1 with probability $1 - \text{probabilityOfZero}$.

Pseudorandom Mode Generates a pseudorandom sequence with period $2^{patternLength} - 1$.

DeterministicCyclic Mode Generates the sequence of 0's and 1's specified by *bitStream* and then repeats it.

DeterministicAppendZeros Mode Generates the sequence of 0's and 1's specified by *bitStream* and then it fills the rest of the buffer space with zeros.

Input Signals

Number: 0

Type: Binary (DiscreteTimeDiscreteAmplitude)

Output Signals

Number: 1 or more

Type: Binary (DiscreteTimeDiscreteAmplitude)

Examples

Random Mode

PseudoRandom Mode As an example consider a pseudorandom sequence with *patternLength*=3 which contains a total of 7 ($2^3 - 1$) bits. In this sequence it is possible to find every combination of 0's and 1's that compose a 3 bit long subsequence with the exception of 000. For this example the possible subsequences are 010, 110, 101, 100, 111, 001 and 100 (they appear in figure 7.1 numbered in this order). Some of these require wrap.

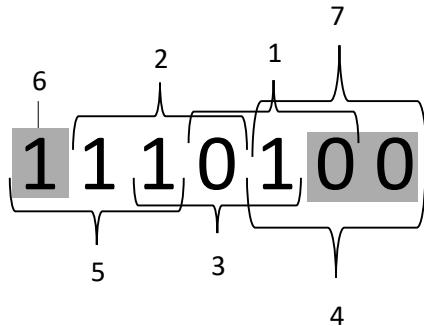


Figura 7.1: Example of a pseudorandom sequence with a pattern length equal to 3.

DeterministicCyclic Mode As an example take the *bit stream* '0100011101010101'. The generated binary signal is displayed in.

DeterministicAppendZeros Mode Take as an example the *bit stream* '0100011101010101'. The generated binary signal is displayed in 7.2.

Sugestions for future improvement

Implement an input signal that can work as trigger.

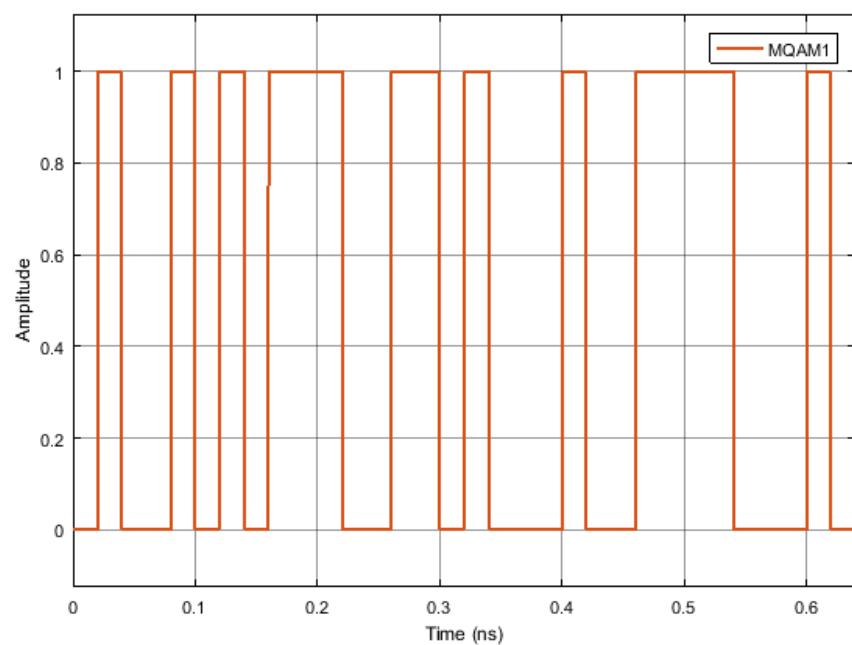


Figura 7.2: Binary signal generated by the block operating in the *Deterministic Append Zeros* mode with a binary sequence 01000...

7.4 Bit Error Rate

Header File	:	bit_error_rate.h
Source File	:	bit_error_rate.cpp
Version	:	20171810 (Responsible: Daniel Pereira)

Input Parameters

Name	Type	Default Value
Confidence	double	0.95
MidReportSize	integer	0
LowestMinorant	double	1×10^{-10}

Input Signals

Number: 2

Type: Binary (DiscreteTimeDiscreteAmplitude)

Output Signals

Number: 1

Type: Binary (DiscreteTimeDiscreteAmplitude)

Functional Description

This block accepts two binary strings and outputs a binary string, outputting a 1 if the two input samples are equal to each other and 0 if not. This block also outputs *.txt* files with a report of the estimated Bit Error Rate (BER), $\widehat{\text{BER}}$ as well as the estimated confidence bounds for a given probability α .

The block allows for mid-reports to be generated, the number of bits between reports is customizable, if it is set to 0 then the block will only output the final report.

Theoretical Description

The $\widehat{\text{BER}}$ is obtained by counting both the total number received bits, N_T , and the number of coincidences, K , and calculating their relative ratio:

$$\widehat{\text{BER}} = 1 - \frac{K}{N_T}. \quad (7.3)$$

The upper and lower bounds, BER_{UB} and BER_{LB} respectively, are calculated using the Clopper-Pearson confidence interval, which returns the following simplified expression for

$N_T > 40$ [?]:

$$\text{BER}_{\text{UB}} = \widehat{\text{BER}} + \frac{1}{\sqrt{N_T}} z_{\alpha/2} \sqrt{\widehat{\text{BER}}(1 - \widehat{\text{BER}})} + \frac{1}{3N_T} \left[2 \left(\frac{1}{2} - \widehat{\text{BER}} \right) z_{\alpha/2}^2 + (2 - \widehat{\text{BER}}) \right] \quad (7.4)$$

$$\text{BER}_{\text{LB}} = \widehat{\text{BER}} - \frac{1}{\sqrt{N_T}} z_{\alpha/2} \sqrt{\widehat{\text{BER}}(1 - \widehat{\text{BER}})} + \frac{1}{3N_T} \left[2 \left(\frac{1}{2} - \widehat{\text{BER}} \right) z_{\alpha/2}^2 - (1 + \widehat{\text{BER}}) \right], \quad (7.5)$$

where $z_{\alpha/2}$ is the $100(1 - \frac{\alpha}{2})$ th percentile of a standard normal distribution.

References

- [1] Thorlabs. *Thorlabs Balance Amplified Photodetectors: PDB4xx Series Operation Manual*, 2014.

7.5 Bit Decider

Header File	:	bit_decider.h
Source File	:	bit_decider.cpp
Version	:	20170818

Input Parameters

Name	Type	Default Value
decisionLevel	double	0.0

Functional Description

This block accepts one real discrete signal and outputs a binary string, outputting a 1 if the input sample is greater than the decision level and 0 if it is less or equal to the decision level.

Input Signals

Number: 1

Type: Real signal (DiscreteTimeContinuousAmplitude)

Output Signals

Number: 1

Type: Binary (DiscreteTimeDiscreteAmplitude)

7.6 Clock

Header File	:	clock.h
Source File	:	clock.cpp

This block doesn't accept any input signal. It outputs one signal that corresponds to a sequence of Dirac's delta functions with a user defined *period*.

Input Parameters

Parameter	Type	Values	Default
period	double	any	0.0
samplingPeriod	double	any	0.0

Tabela 7.2: Binary source input parameters

Methods

Clock()

```
Clock(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig,
OutputSig)
```

```
void initialize(void)
```

```
bool runBlock(void)
```

```
void setClockPeriod(double per)
```

```
void setSamplingPeriod(double sPeriod)
```

Functional description

Input Signals**Number:** 0**Output Signals****Number:** 1

Type: Sequence of Dirac's delta functions.
(TimeContinuousAmplitudeContinuousReal)

Examples**Sugestions for future improvement**

7.7 Clock_20171219

This block doesn't accept any input signal. It outputs one signal that corresponds to a sequence of Dirac's delta functions with a user defined *period*, *phase* and *sampling period*.

Input Parameters

Parameter: period{ 0.0 };

Parameter: samplingPeriod{ 0.0 };

Parameter: phase {0.0};

Methods

Clock()

Clock(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig, OutputSig)

void initialize(void)

bool runBlock(void)

void setClockPeriod(double per) double getClockPeriod()

void setClockPhase(double pha) double getClockPhase()

void setSamplingPeriod(double sPeriod) double getSamplingPeriod()

Functional description

Input Signals

Number: 0

Output Signals

Number: 1

Type: Sequence of Dirac's delta functions.
(TimeContinuousAmplitudeContinuousReal)

Examples

Sugestions for future improvement

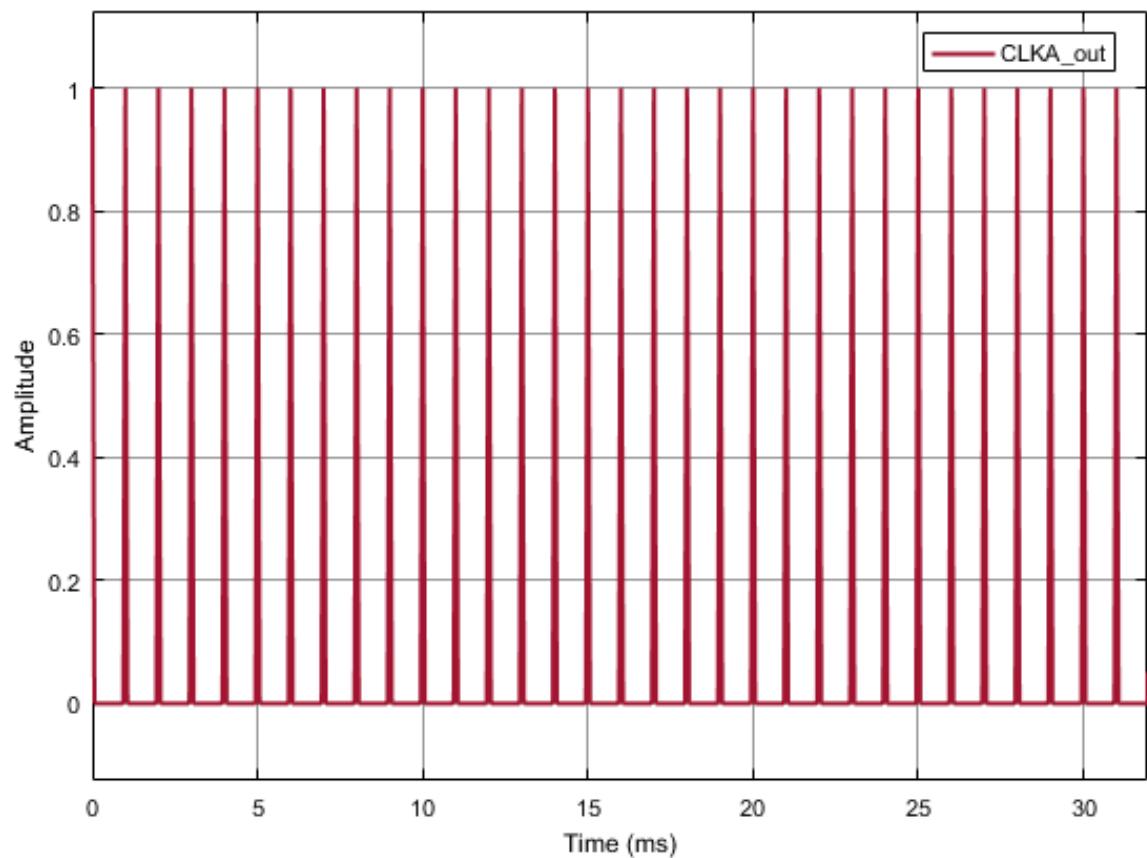


Figura 7.3: Example of the output signal of the clock without phase shift.

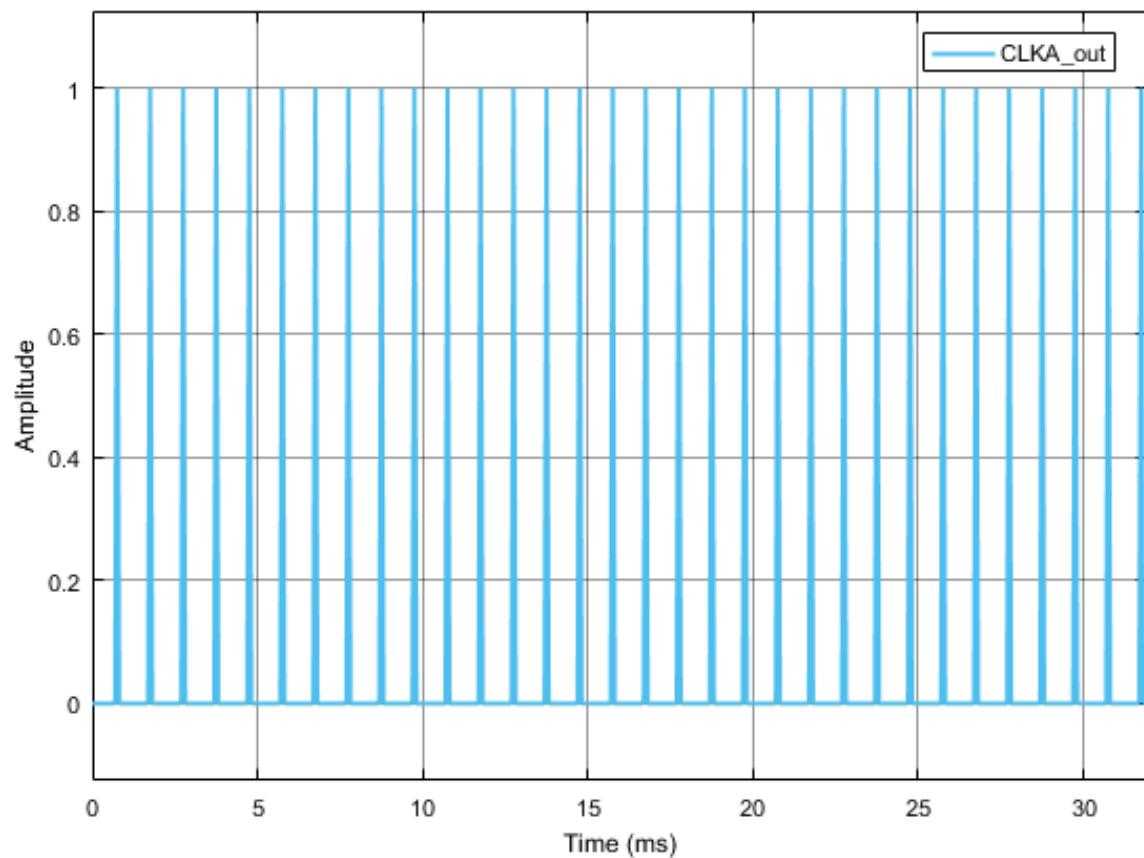


Figura 7.4: Example of the output signal of the clock with phase shift.

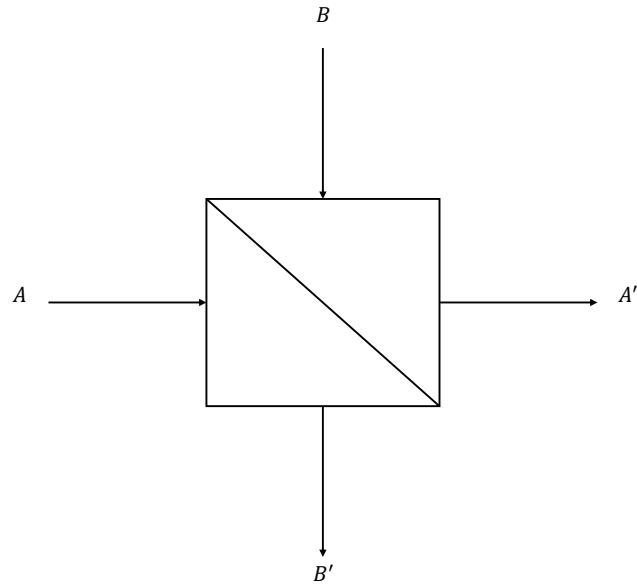


Figura 7.5: 2x2 coupler

7.8 Coupler 2 by 2

In general, the matrix representing 2x2 coupler can be summarized in the following way,

$$\begin{bmatrix} A' \\ B' \end{bmatrix} = \begin{bmatrix} T & iR \\ iR & T \end{bmatrix} \cdot \begin{bmatrix} A \\ B \end{bmatrix} \quad (7.6)$$

Where, A and B represent inputs to the 2x2 coupler and A' and B' represent output of the 2x2 coupler. Parameters T and R represent transmitted and reflected part respectively which can be quantified in the following form,

$$T = \sqrt{1 - \eta_R} \quad (7.7)$$

$$R = \sqrt{\eta_R} \quad (7.8)$$

Where, value of the $\sqrt{\eta_R}$ lies in the range of $0 \leq \sqrt{\eta_R} \leq 1$.

It is worth to mention that if we put $\eta_R = 1/2$ then it leads to a special case of "Balanced Beam splitter" which equally distribute the input power into both output ports.

7.9 Decoder

Header File	:	decoder.h
Source File	:	decoder.cpp

This block accepts a complex electrical signal and outputs a sequence of binary values (0's and 1's). Each point of the input signal corresponds to a pair of bits.

Input Parameters

Parameter	Type	Values	Default
m	int	≥ 4	4
iqAmplitudes	vector<t_complex>	—	{ { 1.0, 1.0 }, { -1.0, 1.0 }, { -1.0, -1.0 }, { 1.0, -1.0 } }

Tabela 7.3: Binary source input parameters

Methods

Decoder()

```
Decoder(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig,
OutputSig)
```

```
void initialize(void)
```

```
bool runBlock(void)
```

```
void setM(int mValue)
```

```
void getM()
```

```
void setIqAmplitudes(vector<t_iqValues> iqAmplitudesValues)
```

```
vector<t_iqValues> getIqAmplitudes()
```

Functional description

This block makes the correspondence between a complex electrical signal and pair of binary values using a predetermined constellation.

To do so it computes the distance in the complex plane between each value of the input signal and each value of the *iqAmplitudes* vector selecting only the shortest one. It then converts the point in the IQ plane to a pair of bits making the correspondence between the input signal and a pair of bits.

Input Signals

Number: 1

Type: Electrical complex (TimeContinuousAmplitudeContinuousReal)

Output Signals

Number: 1

Type: Binary

Examples

As an example take an input signal with positive real and imaginary parts. It would correspond to the first point of the *iqAmplitudes* vector and therefore it would be associated to the pair of bits 00.

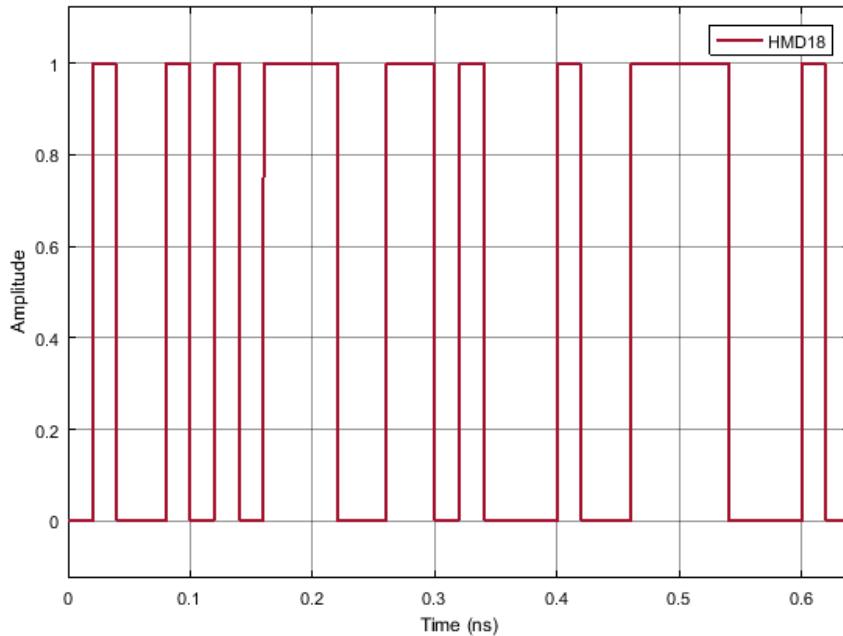


Figura 7.6: Example of the output signal of the decoder for a binary sequence 01. As expected it reproduces the initial bit stream

Sugestions for future improvement

7.10 Discrete To Continuous Time

Header File	:	discrete_to_continuous_time.h
Source File	:	discrete_to_continuous_time.cpp

This block converts a signal discrete in time to a signal continuous in time. It accepts one input signal that is a sequence of 1's and -1's and it produces one output signal that is a sequence of Dirac delta functions.

Input Parameters

Parameter	Type	Values	Default
numberOfSamplesPerSymbol	int	any	8

Tabela 7.4: Binary source input parameters

Methods

```
DiscreteToContinuousTime(vector<Signal * > &inputSignals, vector<Signal * > &outputSignals) :Block(inputSignals, outputSignals){};

void initialize(void);

bool runBlock(void);

void setNumberOfSamplesPerSymbol(int nSamplesPerSymbol)

int const getNumberOfSamplesPerSymbol(void)
```

Functional Description

This block reads the input signal buffer value, puts it in the output signal buffer and it fills the rest of the space available for that symbol with zeros. The space available in the buffer for each symbol is given by the parameter *numberOfSamplesPerSymbol*.

Input Signals

Number : 1

Type : Sequence of 1's and -1's. (DiscreteTimeDiscreteAmplitude)

Output Signals

Number : 1

Type : Sequence of Dirac delta functions (ContinuousTimeDiscreteAmplitude)

Example

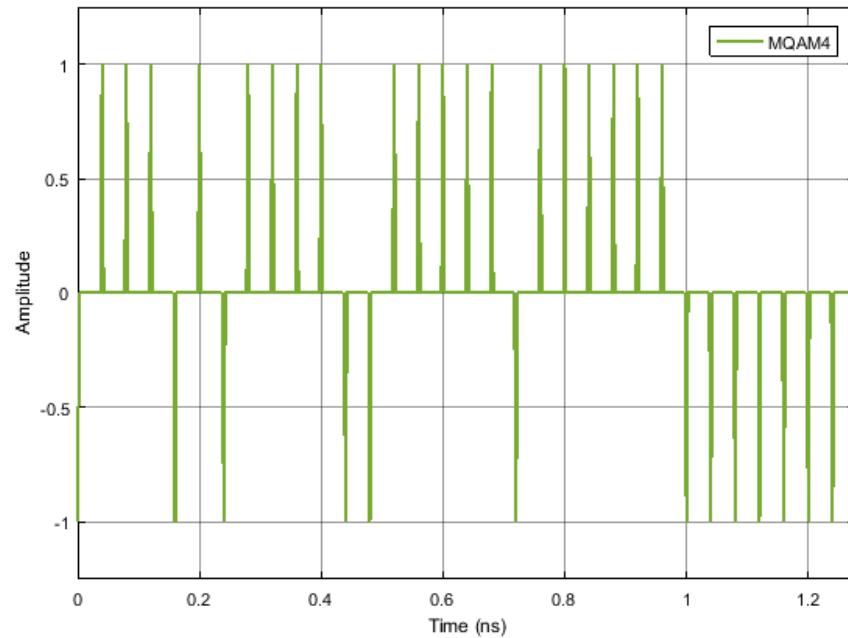


Figura 7.7: Example of the type of signal generated by this block for a binary sequence 0100...

7.11 Fork

Header File	:	fork_20171119.h
Source File	:	fork_20171119.cpp
Version	:	20171119 (Student Name: Romil Patel)

Input Parameters

— NA —

Input Signals

Number: 1

Type: Any type (BinaryValue, IntegerValue, RealValue, ComplexValue, ComplexValueXY, PhotonValue, PhotonValueMP, Message)

Output Signals

Number: 2

Type: Same as applied to the input.

Number: 3

Type: Same as applied to the input.

Functional Description

This block accepts any type signal and outputs two replicas of the input signal.

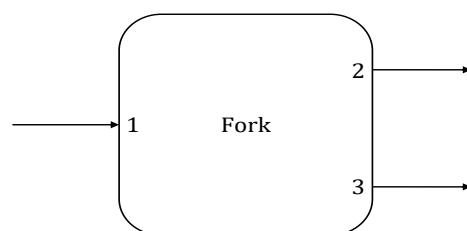


Figura 7.8: Fork

7.12 MQAM Receiver

Header File	:	<code>m_qam_receiver.h</code>
Source File	:	<code>m_qam_receiver.cpp</code>

Warning: `homodyne_receiver` is not recommended. Use `m_qam_homodyne_receiver` instead.

This block of code simulates the reception and demodulation of an optical signal (which is the input signal of the system) outputing a binary signal. A simplified schematic representation of this block is shown in figure 7.9.

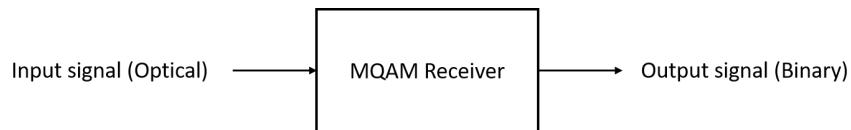


Figura 7.9: Basic configuration of the MQAM receiver

Functional description

This block accepts one optical input signal and outputs one binary signal that corresponds to the M-QAM demodulation of the input signal. It is a complex block (as it can be seen from figure 7.10) of code made up of several simpler blocks whose description can be found in the `lib` repository.

It can also be seen from figure 7.10 that there's an extra internal (generated inside the homodyne receiver block) input signal generated by the `Clock`. This block is used to provide the sampling frequency to the `Sampler`.

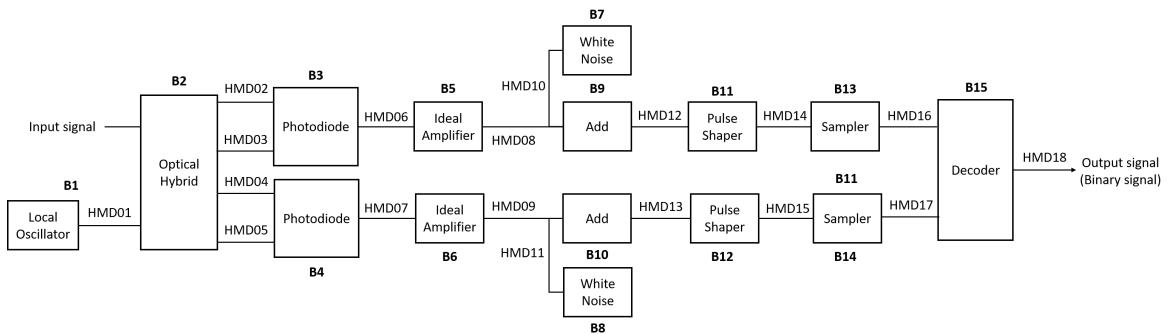


Figura 7.10: Schematic representation of the block homodyne receiver.

Input parameters

This block has some input parameters that can be manipulated by the user in order to change the basic configuration of the receiver. Each parameter has associated a function that allows for its change. In the following table (table 7.9) the input parameters and corresponding functions are summarized.

Input parameters	Function	Type	Accepted values
IQ amplitudes	setIqAmplitudes	Vector of coordinate points in the I-Q plane	Example for a 4-QAM mapping: { { 1.0, 1.0 }, { -1.0, 1.0 }, { -1.0, -1.0 }, { 1.0, -1.0 } }
Local oscillator power (in dBm)	setLocalOscillatorOpticalPower_dBm	double(t_real)	Any double greater than zero
Local oscillator phase	setLocalOscillatorPhase	double(t_real)	Any double greater than zero
Responsivity of the photodiodes	setResponsivity	double(t_real)	$\in [0,1]$
Amplification (of the TI amplifier)	setAmplification	double(t_real)	Positive real number
Noise amplitude (introduced by the TI amplifier)	setNoiseAmplitude	double(t_real)	Real number greater than zero
Samples to skip	setSamplesToSkip	int(t_integer)	
Save internal signals	setSaveInternalSignals	bool	True or False
Sampling period	setSamplingPeriod	double	Given by $symbolPeriod / samplesPerSymbol$

Tabela 7.5: List of input parameters of the block MQAM receiver

Methods

HomodyneReceiver(vector<Signal *> &inputSignal, vector<Signal *> &outputSignal)
(constructor)

```
void setIqAmplitudes(vector<t_iqValues> iqAmplitudesValues)
vector<t_iqValues> const getIqAmplitudes(void)
void setLocalOscillatorSamplingPeriod(double sPeriod)
void setLocalOscillatorOpticalPower(double opticalPower)
void setLocalOscillatorOpticalPower_dBm(double opticalPower_dBm)
void setLocalOscillatorPhase(double lOscillatorPhase)
void setLocalOscillatorOpticalWavelength(double lOscillatorWavelength)
void setSamplingPeriod(double sPeriod)
void setResponsivity(t_real Responsivity)
void setAmplification(t_real Amplification)
void setNoiseAmplitude(t_real NoiseAmplitude)
void setImpulseResponseTimeLength(int impResponseTimeLength)
void setFilterType(PulseShaperFilter fType)
void setRollOffFactor(double rOffFactor)
void setClockPeriod(double per)
void setSamplesToSkip(int sToSkip)
```

Input Signals

Number: 1

Type: Optical signal

Output Signals

Number: 1

Type: Binary signal

Example

Sugestions for future improvement

7.13 IQ Modulator

Header File	:	iq_modulator.h
Source File	:	iq_modulator.cpp

This blocks accepts one input signal continuous in both time and amplitude and it can produce either one or two output signals. It generates an optical signal and it can also generate a binary signal.

Input Parameters

Parameter	Type	Values	Default
outputOpticalPower	double	any	$1e - 3$
outputOpticalWavelength	double	any	$1550e - 9$
outputOpticalFrequency	double	any	speed_of_light/outputOpticalWavelength

Tabela 7.6: Binary source input parameters

Methods

```
IqModulator(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig, OutputSig){};
```

```
void initialize(void);
bool runBlock(void);
void setOutputOpticalPower(double outOpticalPower)
void setOutputOpticalPower_dBm(double outOpticalPower_dBm)
void setOutputOpticalWavelength(double outOpticalWavelength)
void setOutputOpticalFrequency(double outOpticalFrequency)
```

Functional Description

This block takes the two parts of the signal: in phase and in amplitude and it combines them to produce a complex signal that contains information about the amplitude and the phase.

This complex signal is multiplied by $\frac{1}{2}\sqrt{outputOpticalPower}$ in order to reintroduce the information about the energy (or power) of the signal. This signal corresponds to an optical signal and it can be a scalar or have two polarizations along perpendicular axis. It is the signal that is transmitted to the receptor.

The binary signal is sent to the Bit Error Rate (BER) measurement block.

Input Signals

Number : 2

Type : Sequence of impulses modulated by the filter
(ContinuousTimeContiousAmplitude)

Output Signals

Number : 1 or 2

Type : Complex signal (optical) (ContinuousTimeContinuousAmplitude) and binary signal (DiscreteTimeDiscreteAmplitude)

Example

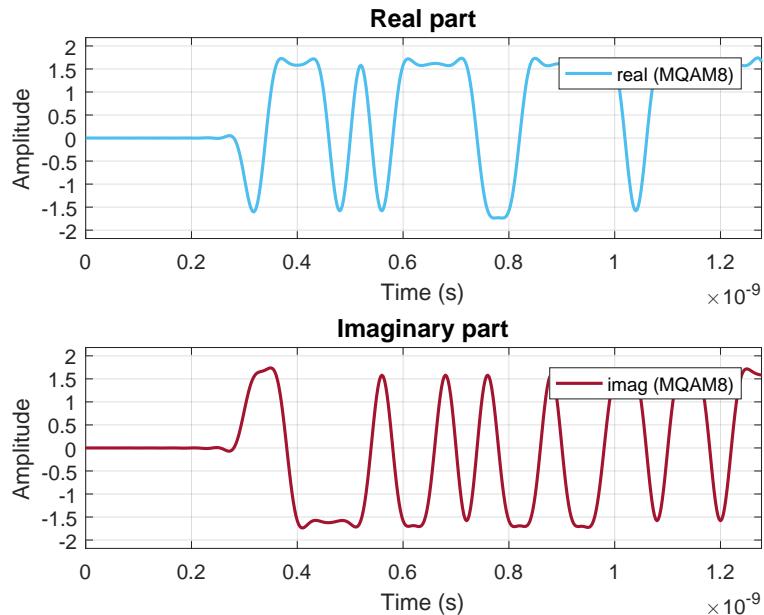


Figura 7.11: Example of a signal generated by this block for the initial binary signal 0100...

7.14 Local Oscillator

Header File	:	local_oscillator.h
Source File	:	local_oscillator.cpp
Version	:	20180130

This block simulates a local oscillator with constant power and initial phase. It produces one output complex signal and it doesn't accept input signals.

Input Parameters

Parameter	Type	Values	Default
opticalPower	double	any	1e - 3
outputOpticalWavelength	double	any	1550e - 9
outputOpticalFrequency	double	any	SPEED_OF_LIGHT / outputOpticalWavelength
phase	double	$\in [0, \frac{\pi}{2}]$	0
samplingPeriod	double	any	0.0
symbolPeriod	double	any	0.0
signaltoNoiseRatio	double	any	0.0
laserLineWidth	double	any	0.0
laserRIN	double	any	0.0

Tabela 7.7: Binary source input parameters

Methods

LocalOscillator()

```

LocalOscillator(vector<Signal * > &InputSig, vector<Signal * > &OutputSig)
:Block(InputSig, OutputSig){};

void initialize(void);

bool runBlock(void);

void setSamplingPeriod(double sPeriod);

void setSymbolPeriod(double sPeriod);

void setOpticalPower(double oPower);

void setOpticalPower_dBm(double oPower_dBm);

void setWavelength(double wlength);

```

```
void setFrequency(double freq);  
void setPhase(double lOscillatorPhase);  
void setSignaltoNoiseRatio(double sNoiseRatio);  
void setLaserLinewidth(double laserLinewidth);  
void setLaserRIN(double laserRIN);
```

Functional description

This block generates a complex signal with a specified phase given by the input parameter *phase*.

Input Signals

Number: 0

Output Signals

Number: 1

Type: Optical signal

7.15 MQAM Mapper

Header File	:	m_qam_mapper.h
Source File	:	m_qam_mapper.cpp

This block does the mapping of the binary signal using a m -QAM modulation. It accepts one input signal of the binary type and it produces two output signals which are a sequence of 1's and -1's.

Input Parameters

Parameter	Type	Values	Default
m	int	2^n with n integer	4
iqAmplitudes	vector<t_complex>	—	{ { 1.0, 1.0 }, { -1.0, 1.0 }, { -1.0, -1.0 }, { 1.0, -1.0 } }

Tabela 7.8: Binary source input parameters

Methods

```
MQamMapper(vector<Signal * > &InputSig, vector<Signal * > &OutputSig)
:Block(InputSig, OutputSig) {}

void initialize(void);

bool runBlock(void);

void setM(int mValue);

void setIqAmplitudes(vector<t_iqValues> iqAmplitudesValues);
```

Functional Description

In the case of $m=4$ this block attributes to each pair of bits a point in the I-Q space. The constellation used is defined by the *iqAmplitudes* vector. The constellation used in this case is illustrated in figure 7.12.

Input Signals

Number : 1

Type : Binary (DiscreteTimeDiscreteAmplitude)

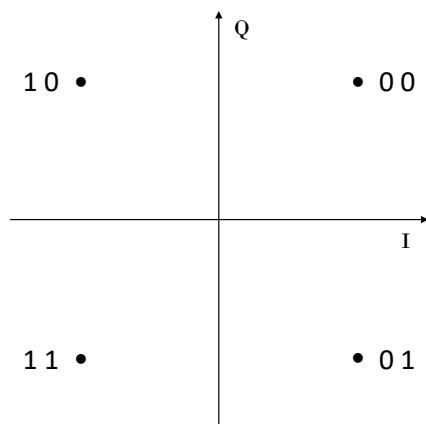


Figura 7.12: Constellation used to map the signal for $m=4$

Output Signals

Number : 2

Type : Sequence of 1's and -1's (DiscreteTimeDiscreteAmplitude)

Example

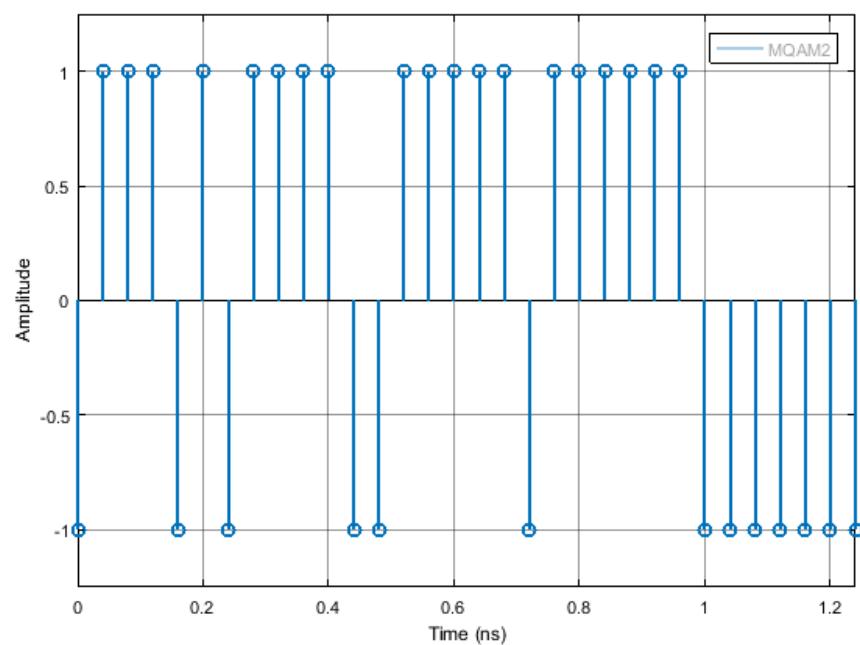


Figura 7.13: Example of the type of signal generated by this block for the initial binary signal 0100...

7.16 MQAM Transmitter

Header File	:	m_qam_transmitter.h
Source File	:	m_qam_transmitter.cpp

This block generates a MQAM optical signal. It can also output the binary sequence. A schematic representation of this block is shown in figure 7.14.

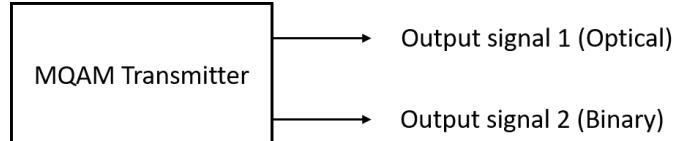


Figura 7.14: Basic configuration of the MQAM transmitter

Functional description

This block generates an optical signal (output signal 1 in figure 7.15). The binary signal generated in the internal block Binary Source (block B1 in figure 7.15) can be used to perform a Bit Error Rate (BER) measurement and in that sense it works as an extra output signal (output signal 2 in figure 7.15).

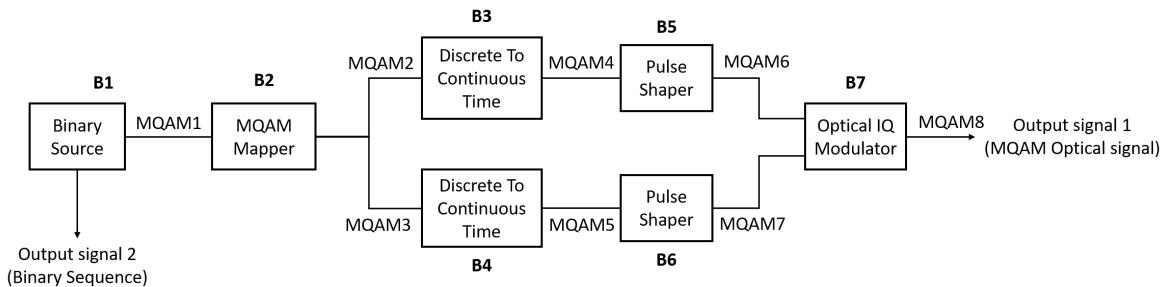


Figura 7.15: Schematic representation of the block MQAM transmitter.

Input parameters

This block has a special set of functions that allow the user to change the basic configuration of the transmitter. The list of input parameters, functions used to change them and the values that each one can take are summarized in table 7.9.

Input parameters	Function	Type	Accepted values
Mode	setMode()	string	PseudoRandom Random DeterministicAppendZeros DeterministicCyclic
Number of bits generated	setNumberOfBits()	int	Any integer
Pattern length	setPatternLength()	int	Real number greater than zero
Number of bits	setNumberOfBits()	long	Integer number greater than zero
Number of samples per symbol	setNumberOfSamplesPerSymbol()	int	Integer number of the type 2^n with n also integer
Roll off factor	setRollOffFactor()	double	$\in [0,1]$
IQ amplitudes	setIqAmplitudes()	Vector of coordinate points in the I-Q plane	Example for a 4-qam mapping: { { 1.0, 1.0 }, { -1.0, 1.0 }, { -1.0, -1.0 }, { 1.0, -1.0 } }
Output optical power	setOutputOpticalPower()	int	Real number greater than zero
Save internal signals	setSaveInternalSignals()	bool	True or False

Tabela 7.9: List of input parameters of the block MQAM transmitter

Methods

MQamTransmitter(vector<Signal *> &inputSignal, vector<Signal *> &outputSignal);
(constructor)

void set(int opt);

void setMode(BinarySourceMode m)

BinarySourceMode const getMode(void)

void setProbabilityOfZero(double pZero)

double const getProbabilityOfZero(void)

void setBitStream(string bStream)

```
string const getBitStream(void)

void setNumberOfBits(long int nOfBits)

long int const getNumberOfBits(void)

void setPatternLength(int pLength)

int const getPatternLength(void)

void setBitPeriod(double bPeriod)

double const getBitPeriod(void)

void setM(int mValue) int const getM(void)

void setIqAmplitudes(vector<t_iqValues> iqAmplitudesValues)

vector<t_iqValues> const getIqAmplitudes(void)

void setNumberOfSamplesPerSymbol(int n)

int const getNumberOfSamplesPerSymbol(void)

void setRollOffFactor(double rOffFactor)

double const getRollOffFactor(void)

void setSeeBeginningOfImpulseResponse(bool sBeginningOfImpulseResponse)

double const getSeeBeginningOfImpulseResponse(void)

void setOutputOpticalPower(t_real outOpticalPower)

t_real const getOutputOpticalPower(void)

void setOutputOpticalPower_dBm(t_real outOpticalPower_dBm)

t_real const getOutputOpticalPower_dBm(void)
```

Output Signals

Number: 1 optical and 1 binary (optional)

Type: Optical signal

Example

Sugestions for future improvement

Add to the system another block similar to this one in order to generate two optical signals with perpendicular polarizations. This would allow to combine the two optical signals and generate an optical signal with any type of polarization.

7.17 Alice QKD

This block is the processor for Alice does all tasks that she needs. This block accepts binary, messages, and real continuous time signals. It produces messages, binary and real discrete time signals.

Input Parameters

Parameter: double RateOfPhotons{1e3}

Parameter: int StringPhotonsLength{ 12 }

Methods

```
AliceQKD (vector <Signal*> &inputSignals, vector <Signal*> &outputSignals) :  
Block(inputSignals, outputSignals) {};  
void initialize(void);  
bool runBlock(void);  
void setRateOfPhotons(double RPhotons) { RateOfPhotons = RPhotons; }; double const  
getRateOfPhotons(void) { return RateOfPhotons; };  
void setStringPhotonsLength(int pLength) { StringPhotonsLength = pLength; }; int const  
getStringPhotonsLength(void) { return StringPhotonsLength; };
```

Functional description

This block receives a sequence of binary numbers (1's or 0's) and a clock signal which will set the rate of the signals produced to generate single polarized photons. The real discrete time signal **SA_1** is generated based on the clock signal and the real discrete time signal **SA_2** is generated based on the random sequence of bits received through the signal **NUM_A**. This last sequence is analysed by the polarizer in pairs of bits in which each pair has a bit for basis choice and other for direction choice.

This block also produces classical messages signals to send to Bob as well as binary messages to the mutual information block with information about the photons it sent.

Input Signals

Number : 3

Type : Binary, Real Continuous Time and Messages signals.

Output Signals

Number : 3

Type : Binary, Real Discrete Time and Messages signals.

Examples

Sugestions for future improvement

7.18 Polarizer

This block is responsible of changing the polarization of the input photon stream signal by using the information from the other real time discrete input signal. This way, this block accepts two input signals: one photon stream and other real discrete time signal. The real discrete time input signal must be a signal discrete in time in which the amplitude can be 0 or 1. The block will analyse the pairs of values by interpreting them as basis and polarization direction.

Input Parameters

Parameter: m{4}

Parameter: Amplitudes { {1,1}, {-1,1}, {-1,-1}, { 1,-1} }

Methods

```
Polarizer (vector <Signal*> &inputSignals, vector <Signal*>&outputSignals) :  
Block(inputSignals, outputSignals) {};  
void initialize(void);  
bool runBlock(void);  
void setM(int mValue);  
void setAmplitudes(vector <t_iqValues> AmplitudeValues);
```

Functional description

Considering m=4, this block attributes for each pair of bits a point in space. In this case, it is be considered four possible polarization states: 0° , 45° , 90° and 135° .

Input Signals

Number : 2

Type : Photon Stream and a Sequence of 0's and '1s (DiscreteTimeDiscreteAmplitude).

Output Signals

Number : 1

Type : Photon Stream

Examples

Sugestions for future improvement

7.19 Bob QKD

This block is the processor for Bob does all tasks that she needs. This block accepts and produces:

1.

2.

Input Parameters

Parameter:

Parameter:

Methods

Functional description

Input Signals

Examples

Sugestions for future improvement

7.20 Eve QKD

This block is the processor for Eve does all tasks that she needs. This block accepts and produces:

1.

2.

Input Parameters

Parameter:

Parameter:

Methods

Functional description

Input Signals

Examples

Sugestions for future improvement

7.21 Rotator Linear Polarizer

This block accepts a Photon Stream signal and a Real discrete time signal. It produces a photon stream by rotating the polarization axis of the linearly polarized input photon stream by an angle of choice.

Input Parameters

Parameter: m[2]

Parameter: axis { {1,0}, { $\frac{\sqrt{2}}{2}$, $\frac{\sqrt{2}}{2}$ } }

Methods

```
RotatorLinearPolarizer(vector <Signal*> &inputSignals, vector <Signal*> &outputSignals) :
Block(inputSignals, outputSignals) {};
void initialize(void);
bool runBlock(void);
void setM(int mValue);
void setAxis(vector <t_iqValues> AxisValues);
```

Functional description

This block accepts the input parameter m, which defines the number of possible rotations. In this case m=2, the block accepts the rectilinear basis, defined by the first position of the second input parameter axis, and the diagonal basis, defined by the second position of the second input parameter axis. This block rotates the polarization axis of the linearly polarized input photon stream to the basis defined by the other input signal. If the discrete value of this signal is 0, the rotator is set to rotate the input photon stream by 0°, otherwise, if the value is 1, the rotator is set to rotate the input photon stream by an angle of 45°.

Input Signals

Number : 2

Type : Photon Stream and a Sequence of 0's and '1s (DiscreteTimeDiscreteAmplitude)

Output Signals

Number : 1

Type : Photon Stream

Examples

Sugestions for future improvement

7.22 Optical Switch

This block has one input signal and two input signals. Furthermore, it accepts an additional input binary input signal which is used to decide which of the two outputs is activated.

Input Parameters

No input parameters.

Methods

```
OpticalSwitch(vector <Signal*> &inputSignals, vector <Signal*> &outputSignals) :  
Block(inputSignals, outputSignals) {};  
void initialize(void);  
bool runBlock(void);
```

Functional description

This block receives an input photon stream signal and it decides which path the signal must follow. In order to make this decision it receives a binary signal (0's and 1's) and it switch the output path according with this signal.

Input Signals

Number : 1

Type : Photon Stream

Output Signals

Number : 2

Type : Photon Stream

Examples

Sugestions for future improvement

7.23 Optical Hybrid

Header File	:	optical_hybrid.h
Source File	:	optical_hybrid.cpp

This block simulates an optical hybrid. It accepts two input signals corresponding to the signal and to the local oscillator. It generates four output complex signals separated by 90° in the complex plane. Figure 7.16 shows a schematic representation of this block.

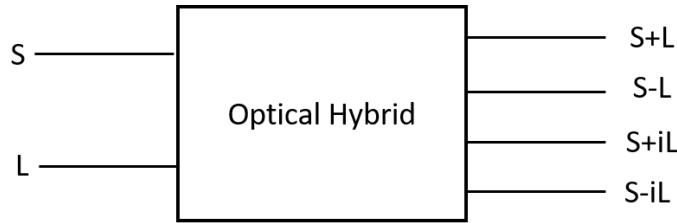


Figura 7.16: Schematic representation of an optical hybrid.

Input Parameters

Parameter	Type	Values	Default
outputOpticalPower	double	any	$1e - 3$
outputOpticalWavelength	double	any	$1550e - 9$
outputOpticalFrequency	double	any	$SPEED_OF_LIGHT / outputOpticalWavelength$
powerFactor	double	≤ 1	0.5

Tabela 7.10: Optical hybrid input parameters

Methods

OpticalHybrid()

```
OpticalHybrid(vector<Signal * > &InputSig, vector<Signal * > &OutputSig)
:Block(InputSig, OutputSig)
```

```
void initialize(void)
```

```
bool runBlock(void)
```

```
void setOutputOpticalPower(double outOpticalPower)
```

```
void setOutputOpticalPower_dBm(double outOpticalPower_dBm)
```

```
void setOutputOpticalWavelength(double outOpticalWavelength)  
void setOutputOpticalFrequency(double outOpticalFrequency)  
void setPowerFactor(double pFactor)
```

Functional description

This block accepts two input signals corresponding to the signal to be demodulated (S) and to the local oscillator (L). It generates four output optical signals given by $powerFactor \times (S + L)$, $powerFactor \times (S - L)$, $powerFactor \times (S + iL)$, $powerFactor \times (S - iL)$. The input parameter $powerFactor$ assures the conservation of optical power.

Input Signals

Number: 2

Type: Optical (OpticalSignal)

Output Signals

Number: 4

Type: Optical (OpticalSignal)

Examples

Sugestions for future improvement

7.24 Photodiode pair

Header File	:	photodiode_old.h
Source File	:	photodiode_old.cpp

This block simulates a block of two photodiodes assembled like in figure 7.17. It accepts two optical input signals and outputs one electrical signal. Each photodiode converts an optical signal to an electrical signal. The two electrical signals are then subtracted and the resulting signals corresponds to the output signal of the block.

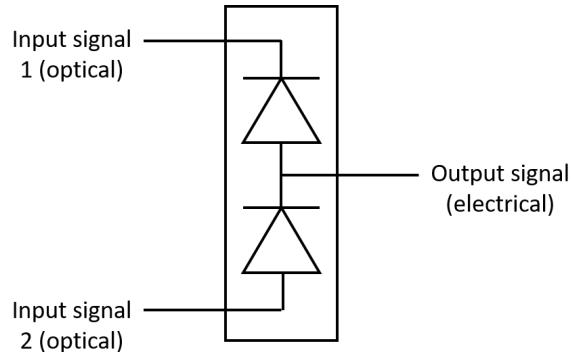


Figura 7.17: Schematic representation of the physical equivalent of the photodiode code block.

Input Parameters

Parameter: responsivity{1}

Parameter: outputOpticalWavelength{ 1550e-9 }

Parameter: outputOpticalFrequency{ SPEED_OF_LIGHT / wavelength }

Methods

Photodiode()

```
Photodiode(vector<Signal * > &InputSig, vector<Signal * > &OutputSig)
:Block(InputSig, OutputSig)
```

void initialize(void)

bool runBlock(void)

void setResponsivity(t_real Responsivity)

Functional description

This block accepts two input optical signals. It computes the optical power of the signal (given by the absolute value squared of the input signal) and multiplies it by the *responsivity* of the photodiode. This product corresponds to the current generated by the photodiode. This is done for each of the input signals. The two currents are then subtracted producing a single output current, that corresponds to the output electrical signal of the block.

Input Signals

Number: 2

Type: Optical (OpticalSignal)

Output Signals

Number: 1

Type: Electrical (TimeContinuousAmplitudeContinuousReal)

Examples

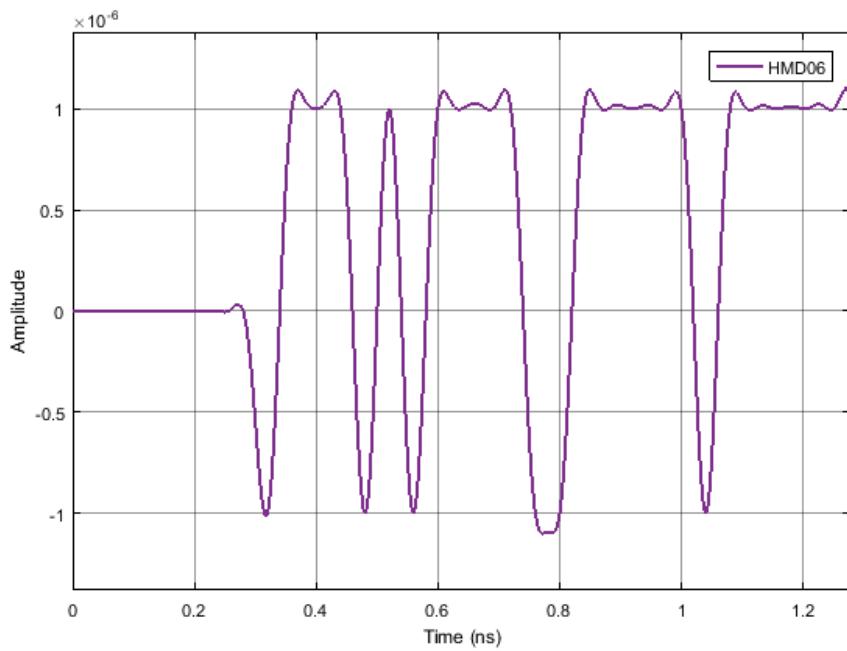


Figura 7.18: Example of the output singal of the photodiode block for a bunary sequence 01

Sugestions for future improvement

7.25 Pulse Shaper

Header File	:	pulse_shaper.h
Source File	:	pulse_shaper.cpp

This block applies an electrical filter to the signal. It accepts one input signal that is a sequence of Dirac delta functions and it produces one output signal continuous in time and in amplitude.

Input Parameters

Parameter	Type	Values	Default
filterType	string	RaisedCosine, Gaussian	RaisedCosine
impulseResponseTimeLength	int	any	16
rollOffFactor	real	$\in [0, 1]$	0.9

Tabela 7.11: Pulse shaper input parameters

Methods

```

PulseShaper(vector<Signal * > &InputSig, vector<Signal * > OutputSig)
:FIR_Filter(InputSig, OutputSig){};

void initialize(void);

void setImpulseResponseTimeLength(int impResponseTimeLength)

int const getImpulseResponseTimeLength(void)

void setFilterType(PulseShaperFilter fType)

PulseShaperFilter const getFilterType(void)

void setRollOffFactor(double rOffFactor)

double const getRollOffFactor()

```

Functional Description

The type of filter applied to the signal can be selected through the input parameter *filterType*. Currently the only available filter is a raised cosine.

The filter's transfer function is defined by the vector *impulseResponse*. The parameter *rollOffFactor* is a characteristic of the filter and is used to define its transfer function.

Input Signals

Number : 1

Type : Sequence of Dirac Delta functions (ContinuousTimeDiscreteAmplitude)

Output Signals

Number : 1

Type : Sequence of impulses modulated by the filter (ContinuousTimeContinuousAmplitude)

Example

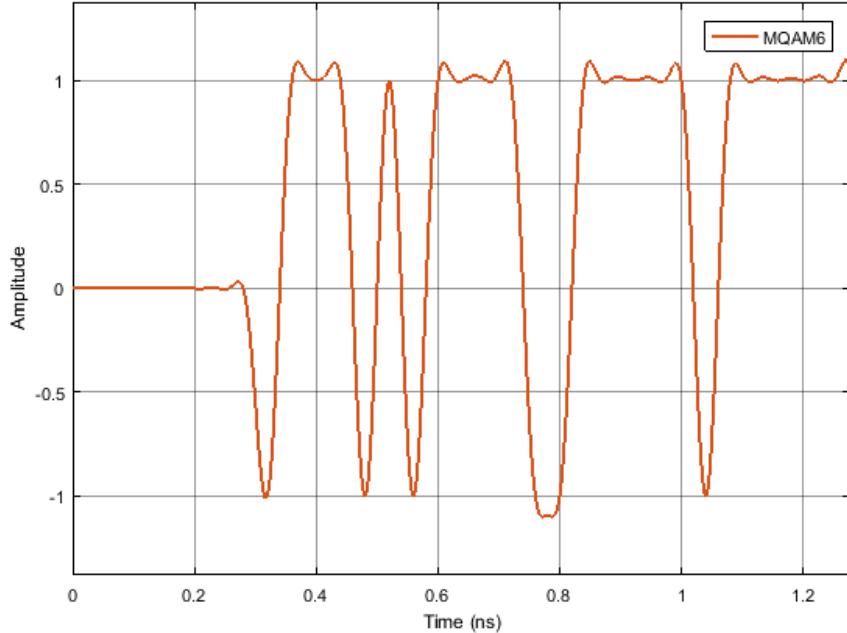


Figura 7.19: Example of a signal generated by this block for the initial binary signal 0100...

Sugestions for future improvement

Include other types of filters.

7.26 Sampler

Header File	:	sampler.h
Source File	:	sampler_20171119.cpp

This block can work in two configurations: with an external clock or without it. In the latter it accepts two input signals one being the clock and the other the signal to be demodulated. In the other configuration there's only one input signal which is the signal.

The output signal is obtained by sampling the input signal with a predetermined sampling rate provided either internally or by the clock.

Input Parameters

Parameter	Type	Values	Default
samplesToSkip	int	any (smaller than the number of samples generated)	0

Tabela 7.12: Sampler input parameters

Methods

Sampler()

Sampler(vector<Signal *> &InputSig, vector<Signal *> &OutputSig) :Block(InputSig, OutputSig)

void initialize(void)

bool runBlock(void)

void setSamplesToSkip(t_integer sToSkip)

Functional description

This block can work with an external clock or without it.

In the case of having an external clock it accepts two input signals. The signal to be demodulate which is complex and a clock signal that is a sequence of Dirac delta functions with a predetermined period that corresponds to the sampling period. The signal and the clock signal are scanned and when the clock has the value of 1.0 the correspondent complex value of the signal is placed in the buffer corresponding to the output signal.

There's a detail worth noting. The electrical filter has an impulse response time length of 16 (in units of symbol period). This means that when modulating a bit the spike in the signal corresponding to that bit will appear 8 units of symbol period later. For this reason there's

the need to skip the earlier samples of the signal when demodulating it. That's the purpose of the *samplesToSkip* parameter.

Between the binary source and the current block the signal is filtered twice which means that this effect has to be taken into account twice. Therefore the parameter *samplesToSkip* is given by $2 * 8 * \text{samplesPerSymbol}$.

Input Signals

Number: 1

Type: Electrical real (TimeContinuousAmplitudeContinuousReal)

Output Signals

Number: 1

Type: Electrical real (TimeDiscreteAmplitudeContinuousReal)

Examples

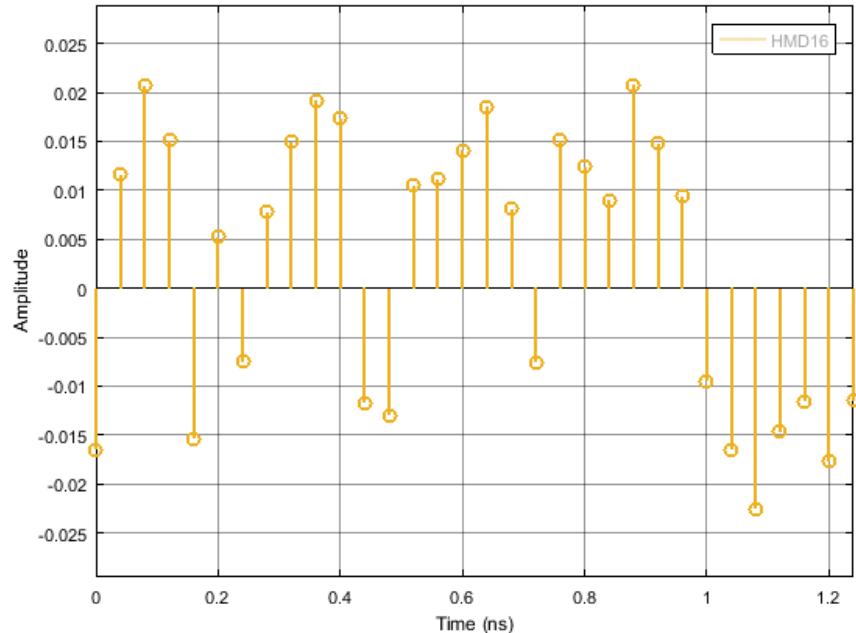


Figura 7.20: Example of the output signal of the sampler

Sugestions for future improvement

7.27 Sink

Header File	:	sink.h
Source File	:	sink.cpp

This block accepts one input signal and it does not produce output signals. It takes samples out of the buffer until the buffer is empty. It has the option of displaying the number of samples still available.

Input Parameters

Parameter	Type	Values	Default
numberOfSamples	long int	any	-1

Tabela 7.13: Sampler input parameters

Methods

`Sink(vector<Signal *> &InputSig, vector<Signal *> &OutputSig)`

`bool runBlock(void)`

`void setNumberOfSamples(long int nOfSamples)`

`void setDisplayNumberOfSamples(bool opt)`

Functional Description

Capítulo 8

Mathlab Tools

8.1 Generation of AWG Compatible Signals

Student Name	:	Francisco Marques dos Santos
Starting Date	:	September 1, 2017
Goal	:	Convert simulation signals into waveform files compatible with the laboratory's Arbitrary Waveform Generator
Directory	:	mtools

This section shows how to convert a simulation signal into an AWG compatible waveform file through the use of a matlab function called `sgnToWfm`. This allows the application of simulated signals into real world systems.

8.1.1 `sgnToWfm_20171121`

Structure of a function

```
[dataDecimate, data, symbolPeriod, samplingPeriod,
type, numberOfRowsInSection, samplingRate, samplingRateDecimate] = sgnToWfm_20171121
(fname_sgn, nReadr, fname_wfm)
```

Inputs

fname_sgn: Input filename of the signal (*.sgn) you want to convert. It must be a real signal (Type: TimeContinuousAmplitudeContinuousReal).

nReadr: Number of symbols you want to extract from the signal.

fname_wfm: Name that will be given to the waveform file.

Outputs

A waveform file will be created in the Matlab current folder. It will also return six variables in the workspace which are:

dataDecimate: A vector which contains decimated signal data by an appropriate decimation factor to make it compatible with the AWG.

data: A vector with the signal data.

symbolPeriod: Equal to the symbol period of the corresponding signal.

samplingPeriod: Sampling period of the signal.

type: A string with the name of the signal type.

numberOfSymbols: Number of symbols retrieved from the signal.

samplingRate: Sampling rate of the signal.

samplingRateDecimate: Reduced sampling rate which is compatible with AWG. (i.e. less than 16 GSa/s).

Functional Description

This matlab function generates a *.wfm file given an input signal file (*.sgn). The waveform file is compatible with the laboratory's Arbitrary Waveform Generator (Tektronix AWG70002A). In order to recreate it appropriately, the signal must be real, not exceed $8 * 10^9$ samples and have a sampling rate equal or bellow 16 GS/s.

This function can be called with one, two or three arguments:

Using one argument:

```
[dataDecimate, data, symbolPeriod, samplingPeriod, type, numberOfSymbols, samplingRate, samplingRateDecimate] = sgnToWfm('S6.sgn');
```

This creates a waveform file with the same name as the *.sgn file and uses all of the samples it contains.

Using two arguments:

```
[dataDecimate, data, symbolPeriod, samplingPeriod, type, numberOfSymbols, samplingRate, samplingRateDecimate] = sgnToWfm('S6.sgn',256);
```

This creates a waveform file with the same name as the signal file name and the number of samples used equals nReadr x samplesPerSymbol. The samplesPerSymbol constant is defined in the *.sgn file.

Using three arguments:

```
[dataDecimate, data, symbolPeriod, samplingPeriod, type, numberOfSymbols, samplingRate, samplingRateDecimate] = sgnToWfm('S6.sgn',256,'myWaveform.wfm');
```

This creates a waveform file with the name "myWaveform" and the number of samples used equals nReadr x samplesPerSymbol. The samplesPerSymbol constant is defined in the *.sgn file.

8.1.2 Loading a signal to the Tektronix AWG70002A

The AWG we will be using is the Tektronix AWG70002A which has the following key specifications:

Sampling rate up to 16 GS/s: This is the most important characteristic because it determines the maximum sampling rate that your signal can have. It must not be over 16 GS/s or else the AWG will not be able to recreate it appropriately.

8 GSample waveform memory: This determines how many data points your signal can have.

After making sure this specifications are respected you can create your waveform using the function. When you load your waveform, the AWG will output it and repeat it constantly until you stop playing it.

1. Using the function sgnToWfm: Start up Matlab and change your current folder to mtools and add the signals folder that you want to convert to the Matlab search path. Use the function accordingly, putting as the input parameter the signal file name you want to convert.

2. AWG sampling rate: After calling the function there should be waveform file in the mtools folder, as well as a variable called samplingRate in the Matlab workspace. Make sure this is equal or bellow the maximum sampling frequency of the AWG (16 GS/s), or else the waveform can not be equal to the original signal. If it is higher you have to adjust the parameters in the simulation in order to decrease the sampling frequency of the signal(i.e. decreasing the bit period or reducing the samples per symbol).

3. Loading the waveform file to the AWG: Copy the waveform file to your pen drive and connect it to the AWG. With the software of the awg open, go to browse for waveform on the channel you want to use, and select the waveform file you created (Figure 7.1).

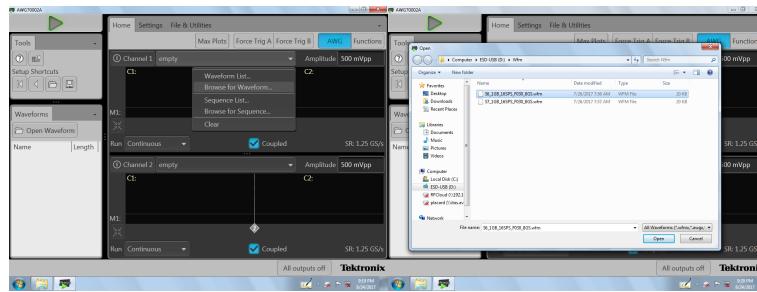


Figura 8.1: Selecting your waveform in the AWG

Now you should have the waveform displayed on the screen. Although it has the same shape, the waveform might not match the signal timing wise due to an incorrect sampling rate configured in the AWG. In this example (Figure 7.2), the original signal has a sample

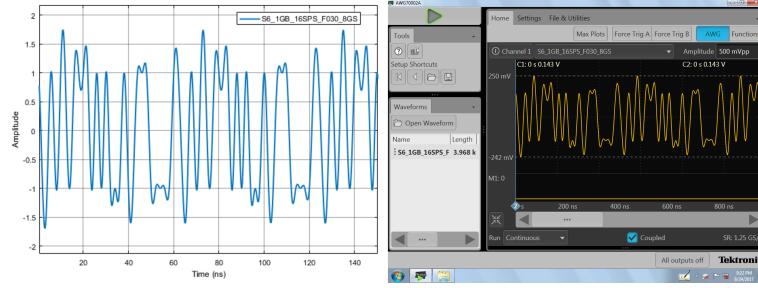


Figura 8.2: Comparison between the waveform in the AWG and the original signal before configuring the sampling rate

rate of 8 GS/s and the AWG is configured to 1.25 GS/s. Therefore it must be changed to the correct value. To do this go to the settings tab, clock settings, and change the sampling rate to be equal to the one of the original signal, 8 GS/s (Figure 7.3). Compare the waveform in

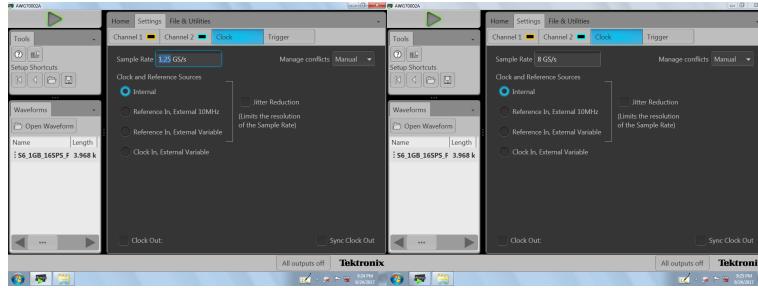


Figura 8.3: Configuring the right sampling rate

the AWG with the original signal, they should be identical (Figure 7.4).

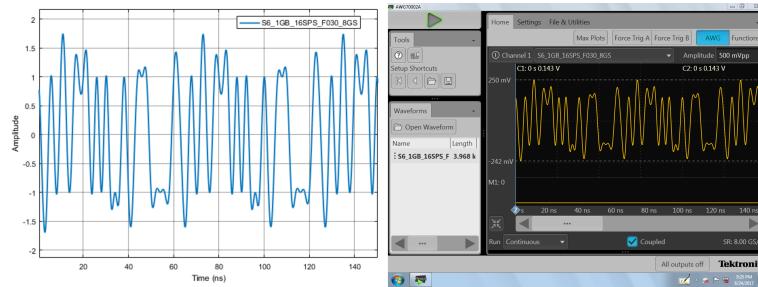


Figura 8.4: Comparison between the waveform in the AWG and the original signal after configuring the sampling rate

4. Generate the signal: Output the wave by enabling the channel you want and clicking on the play button.

9.1 Overlap-Save Method

Linear filtering can be easily implemented in time-domain resorting to the use of finite impulse response (FIR) digital filters and convolution property as,

$$y(n) = \sum_{k=0}^{R-1} x(n-k)h(k), \quad (9.1)$$

where $x(n)$ is the input signal, $h(k)$ is the FIR filter coefficients, R is the length of FIR filter coefficients and $y(n)$ represents the filtered output signal. Analysing this equation we can note that, for a block signal of length R , the required number of operations for the direct implementation of equation (9.2) evolves with R^2 , $\mathcal{O}(R^2)$. This limitation imposed the emergence of algorithm, where the linear convolution is calculated faster than the directly implementing of (9.2). In this sense, it is used the computation of linear filtering in frequency domain resorting to the use of fast Fourier transform (FFT) and inverse fast Fourier transform (IFFT) algorithms as However, for long input sequence, the direct implementation of frequency domain filtering in real-time is limited by the limited memory of the digital processors. Hence, the filtering in frequency-domain is implemented by sectioning or block the input signal, such that the practical implementation of FFT and IFFT is feasible. In order to implement the non-cyclic convolution with the finite-length of cyclic convolution that the FFT provides, overlap-save and overlap-add method are use, enabling that the complexity evolves in log scale $\mathcal{O}(N \log_2 N)$. The general method is to split the input signal into manageable blocks, then apply the FFT to to perform the linear convolution and at the end recombine the output blocks such that it is avoided the wrap-around errors due to the circular convolution imposed by FFT.

The overlap-save method can be computed in the following steps:

1. Step 1:

Parameter: Determine the length R of impulse response, $h(n)$;

2. Step 2:

Parameter: Define the size of FFT and IFFT operation, N ;

3. Step 3:

Parameter: Determine the length of block L to section the input sequence $x(n)$, considering that $N = L + R - 1$;

4. Step 4:

Parameter: Pad $L - 1$ zeros to the impulse response $h(n)$ to obtain the length N ;

5. Step 5:

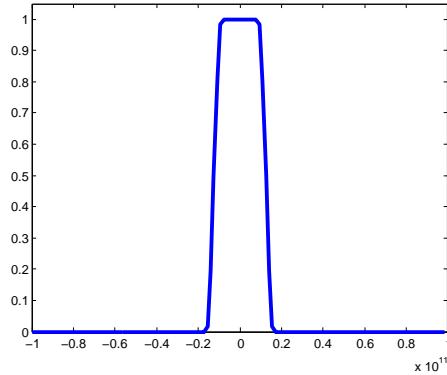


Figura 9.1: Frequency response of raised-cosine filter.

Parameter: Make the segments of the input sequences of length L , $x_i(n)$, where index i correspond to the i^{th} block. Overlap $R - 1$ samples of the previous block at the beginning of the segmented block to obtain a block of length N . In the first block, it is added $R - 1$ null samples;

6. Step 6:

Parameter: Compute the circular convolution of segmented input sequence $x_i(n)$ and $h(n)$ described as,

$$y_i(n) = x_i(n) \circledast h(n). \quad (9.2)$$

This is obtained in the following steps:

Description: Compute the FFT of x_i and h both with length N ;

Description: Compute the multiplication of $X_i(f)$ and the transfer function $H(f)$;

Description: Compute the IFFT of the multiplication result to obtain the time-domain block signal, y_i ;

7. Step 7:

Parameter: Discarded $R - 1$ initial samples from the y_i , and save only the error-free $N - R - 1$ samples in the output record.

In the Figure 9.4 it is illustrated an example of overlap-save method.

9.1.1 Frequency Response of Filter

The frequency response of filter can be directly defined by using the frequency-domain formula, or it can be equivalently calculated from the FFT of impulse response of the filter. In this sense, we present an example of FIR filter (*raised-cosine filter*) to illustrate these two cases.

Frequency-domain Formula

The frequency-domain description of raised-cosine filter can be given as,

$$H(f) = \begin{cases} 1, & |f| \leq \frac{1-\beta}{2T} \\ \frac{1}{2} \left[1 + \cos \left(\frac{\pi T}{\beta} \left[|f| - \frac{1-\beta}{2T} \right] \right) \right], & \frac{1-\beta}{2T} < |f| \leq \frac{1+\beta}{2T} \\ 0, & \text{otherwise} \end{cases}$$

where, f is the frequency, $0 \leq \beta \leq 1$ corresponds to the roll-off factor and T is the reciprocal of the symbol-rate. A plot of direct implementation of frequency response of raised-cosine filter is presented in Figure 9.1, for a roll-off factor of 0.3. The frequency response, $H(f)$, calculated for N frequency bins, which can be defined as,

$$f = -\frac{f_{windowTF}}{2} : \frac{f_{windowTF}}{N} : \left(\frac{f_{windowTF}}{2} - \frac{f_{windowTF}}{N} \right), \quad (9.3)$$

where $f_{windowTF}$ is the sampling frequency. This, imposes that the length of $H(f)$ is N as expected for the N -point FFT and the transfer function multiplication.

FFT of Impulse Response

Alternatively to the frequency-domain formula, we can obtain the frequency response of filter by calculating the FFT of its impulse response. In the case of raised-cosine filter, the impulse response is given as,

$$h(t) = \frac{\sin(\pi t/T)}{\pi t/T} \frac{\cos(\pi t\beta/T)}{1 - (2\beta t/T)^2}, \quad (9.4)$$

where t is the time. Figure 9.2 shows a plot of impulse response of raised-cosine filter for a roll-off factor of 0.3. Before calculating the FFT of the impulse response we must zero-pad the

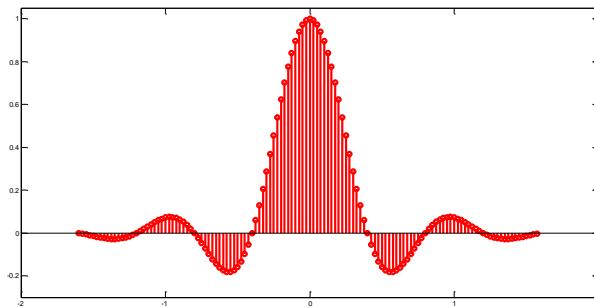


Figura 9.2: Impulse response of raised-cosine filter.

impulse response, which has the length R , to the length N . In this case N is the FFT length, which is efficiently defined as power of 2. The zero-padding process can be performed by appending $L - 1$ zeros at the end of impulse response, as shown in the Figure 9.3.

In both cases, the frequency response of the filter will be limited to the frequency interval $f_{windowTF}, [-\frac{f_{windowTF}}{2}, \frac{f_{windowTF}}{2}]$, and this range show us the N frequency components

Figura 9.3: Zero-padding of impulse response $h(n)$.

obtained from FFT. The minimum frequency bin is $-\frac{f_{windowTF}}{2}$ and the maximum bin is $\frac{f_{windowTF}}{2}$, in which $f_{windowTF}$ corresponds to the sampling frequency. We can note that the spectral width of $H(f)$ is $f_{windowTF}$, which is the inverse of sampling period, dt . It is also important to note that, for a given sampling frequency, the frequency resolution, Δf , of $H(f)$ depends on the parameter N and it increases with N , $\Delta f = \frac{f_{windowTF}}{N}$.

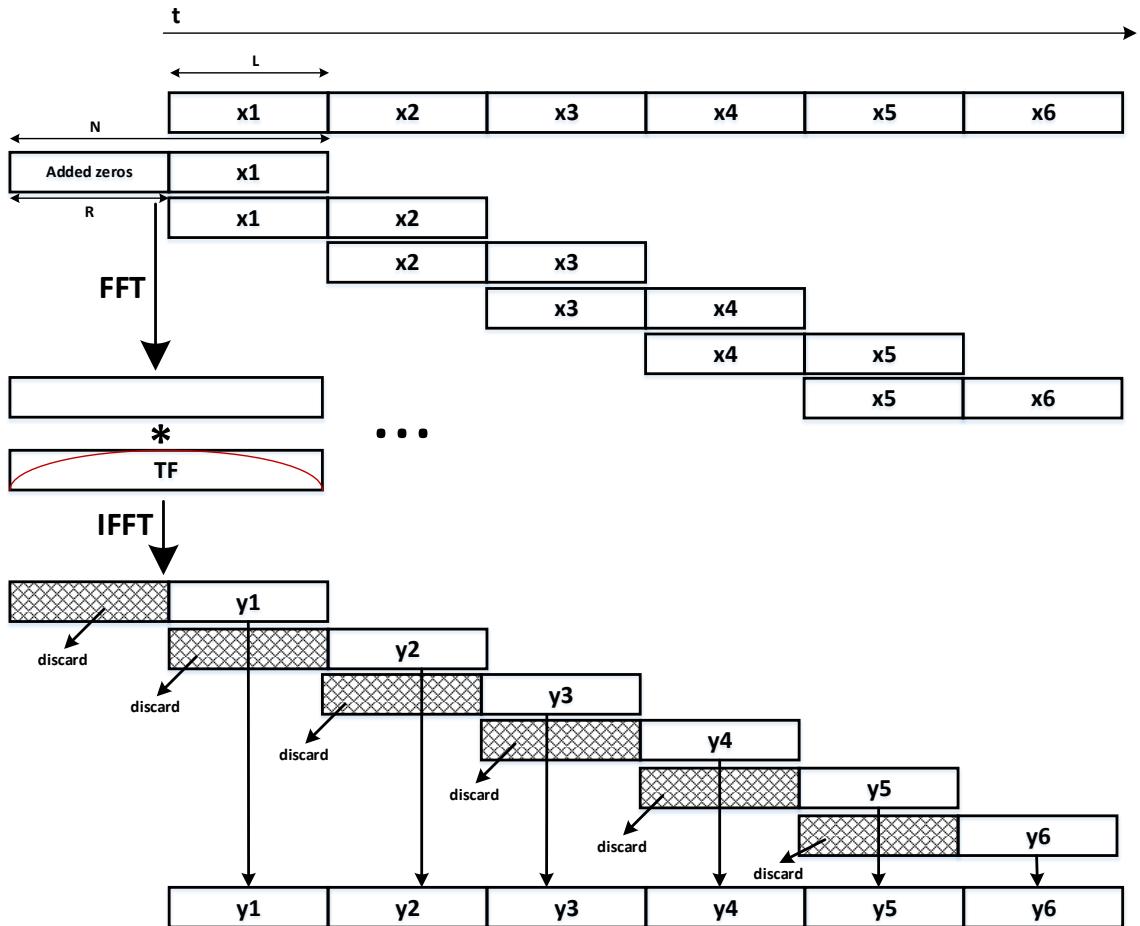


Figura 9.4: Illustration of Overlap-save method.

References

- [1] Blahut, R.E. *Fast Algorithms for Digital Signal Processing*, Addison-Wesley, Reading, MA, 1985.
- [2] Steven W. Smith. *The Scientist and Engineer's Guide to Digital Signal Processing*. California Technical Publishing, San Diego, CA, USA, 1997.

9.2 FFT

Algorithm

An alternate representation of the Fourier transform can give a flexibility to utilize the same code for the FFT and IFFT with appropriate input arguments. The algorithm for the FFT will follow the following formula,

$$X_k = \frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} x_n \cdot e^{-i2\pi kn/N} \quad 0 \leq k \leq N - 1 \quad (9.5)$$

Similarly, for IFFT,

$$x_n = \frac{1}{\sqrt{N}} \sum_{k=0}^{N-1} X_k \cdot e^{i2\pi kn/N} \quad 0 \leq k \leq N - 1 \quad (9.6)$$

From the above manipulation discussed in equations 9.5 and 9.6, we can write only one script for the implementations Fourier algorithm and manipulate its functionality as a FFT or IFFT by applying an appropriate input arguments.

The generalized form for the algorithm can be given as,

$$OUT = \frac{1}{\sqrt{N}} \sum_{n=0}^{N-1} IN \cdot e^{s \cdot i2\pi kn/N} \quad 0 \leq k \leq N - 1 \quad (9.7)$$

where,

IN → Input complex signal

OUT → Output complex signal & *s* → '-1' for FFT and '1' for IFFT

Function description

To perform FFT operation, the input argument to the function can be given as follows,

$$OUT = transform(IN, -1)$$

in a similar way, IFFT can be manipulated as,

$$OUT = transform(IN, 1)$$

Flowchart

This algorithm converts time domain signal to frequency domain for both real and complex signal type. In case of real signal, we have to manipulate imaginary part to be zero for the execution of algorithm. The figure 9.5 displays top level architecture of the FFT algorithm. If the length of the input signal is 2^N , then it'll execute Radix-2 algorithm otherwise it'll execute Bluestein algorithm.

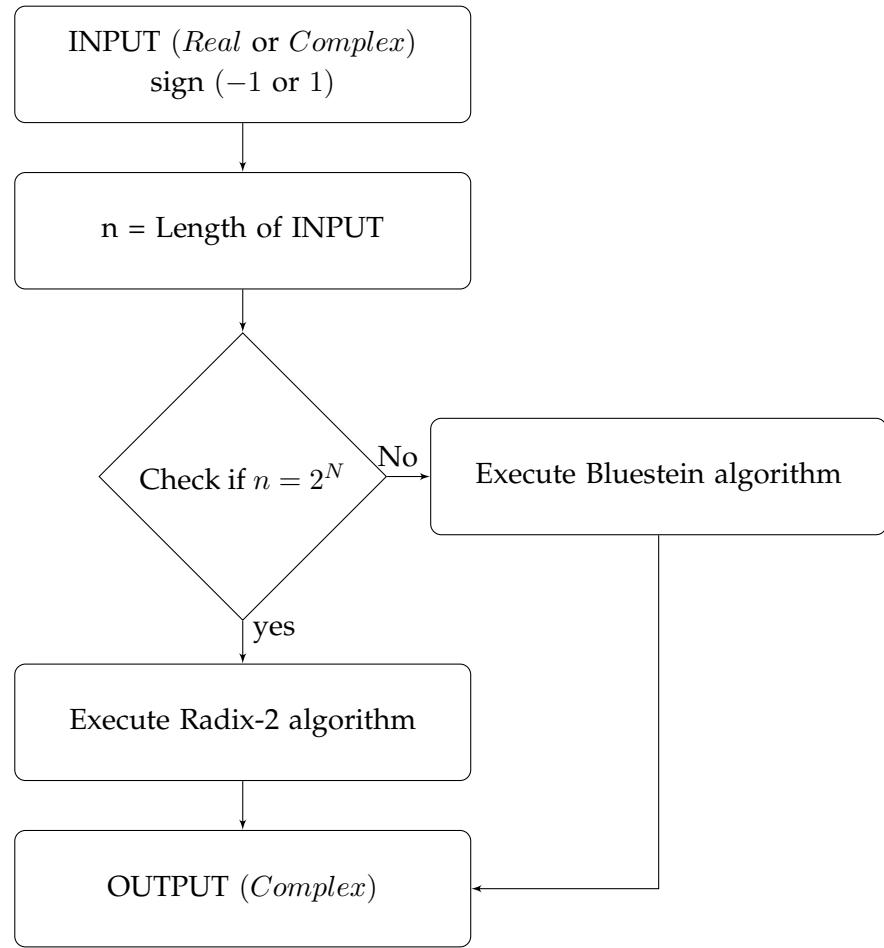


Figura 9.5: Top level architecture of the FFT algorithm

Radix-2 algorithm

The architecture of the algorithm is to accept time domain complex signal and generate frequency domain complex output signal. In the case of real input signal, first we have to convert it into the complex form by adding zero to the imaginary part.

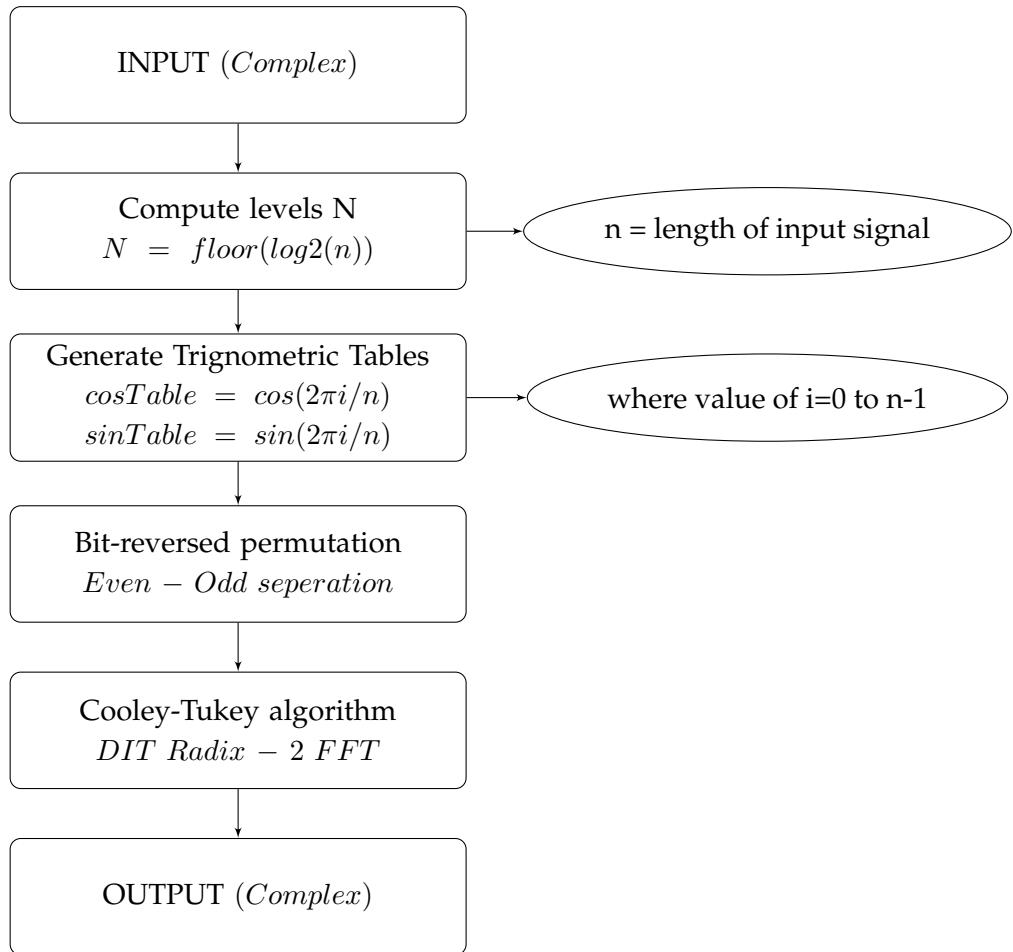


Figura 9.6: Flowchart of Cooley-Tukey DIT Radix-2 algorithm

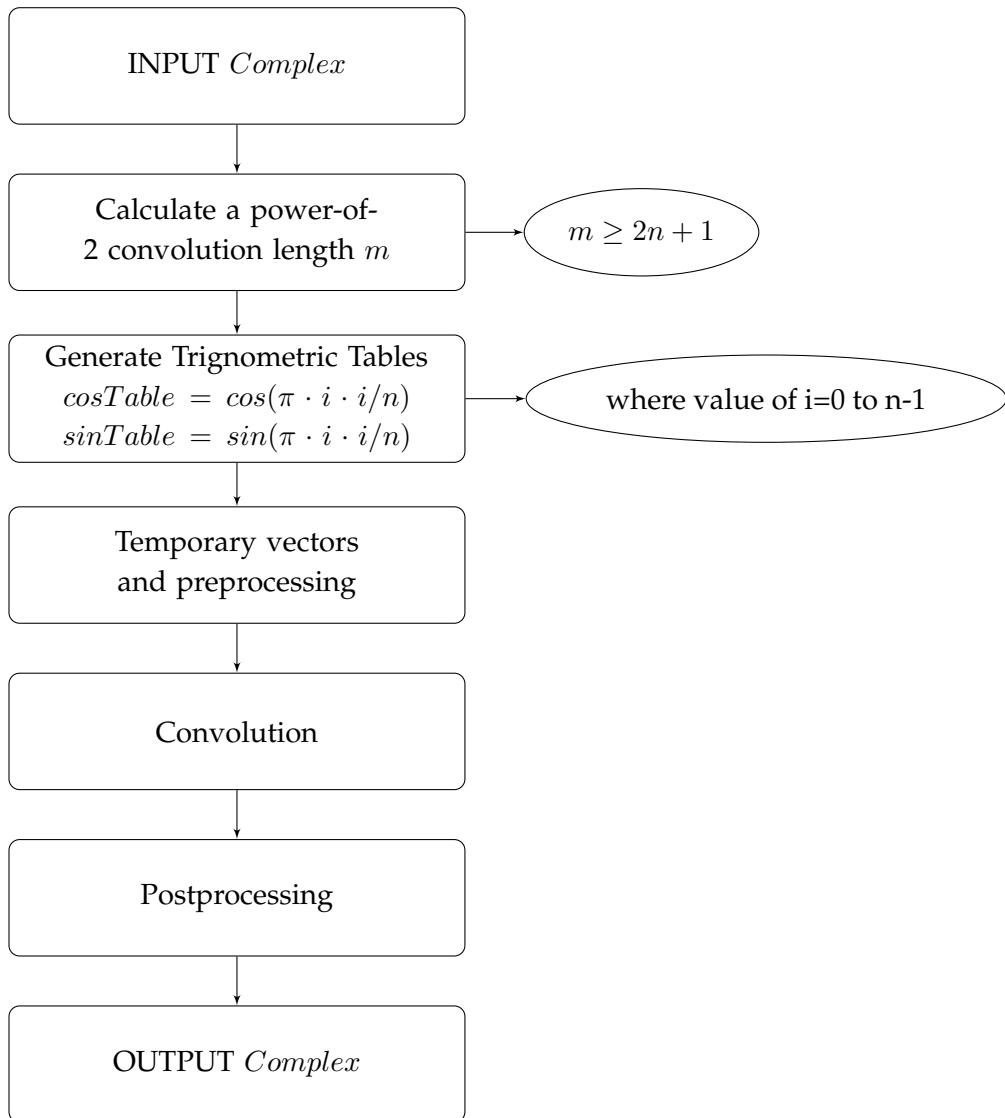
Bluestein algorithm

Figura 9.7: Flowchart of Bluestein algorithm

Capítulo 10

Multi-Party Computation

10.1 Introduction

Multi-Party Computation (MPC), also known as Secure Function Evaluation, allows two or more parties to correctly compute a function of their private inputs without exposure, i.e., without the input of one party being revealed to the other parties. X

A generic function f receives as input a set $A = \{a_1, a_2, \dots, a_n\}$ of arguments, where a_i is the input of the i -th party, and $1 \leq i \leq n$, and outputs a value z , which represents the result of the joint computation of f , as shown in Figure 10.1. X

The output of f is given by the following expression,

$$z = f(a_1, a_2, \dots, a_n) \quad (10.1)$$

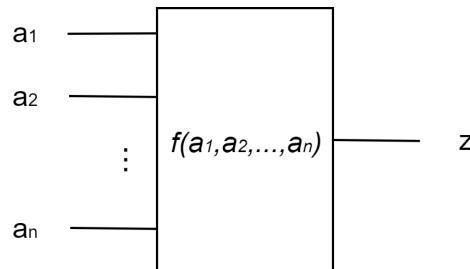


Figure 10.1: Multi-Party Computation Diagram

10.2 Two-Party Computation

Two-Party Computation (2PC) is a specific case of MPC, where a generic function f receives as input a set ~~A~~ $\{a, b\}$ of arguments, where a is the input from the first party and b is the input from the second, and outputs a value c , as shown in Figure 10.2. ~~X~~

The output of f is given by the following expression,

$$c = f(a, b) \quad (10.2)$$

An example of a 2PC problem, is **Yao's Millionaires' Problem**, in which 2 millionaires, Alice and Bob, are interested in knowing which one of them is richer without revealing their wealth. The goal of this problem is to solve the inequality $a \geq b$, where a is Alice's input and b is Bob's input, without revealing the values of a and b .

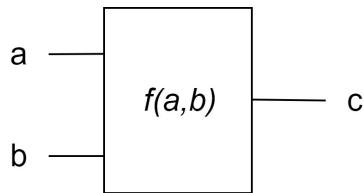


Figure 10.2: Two-Party Computation Diagram

10.3 MPC Example

Let's consider 2 parties, Alice and Bob. Alice has a set $M = \{m_0, m_1\}$ of messages and Bob has the index, i , of the message he wants to receive from Alice, m_i .

For this example we will be assuming that $i = 1$ and $M = \{11, 10\}$. Both messages are 2-bit binary numbers.

1. Bob generates a 4-bit key, $K_b = 0111$ where only the bits of indexes $2i$ and $2i + 1$ are known (C) to him. Since $i=1$, Bob only knows the last 2 bits of K_b . The remaining 2 bits of K_b are random (R), thus unknown to Bob.

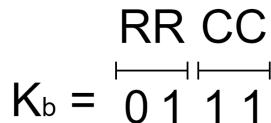


Figure 10.3: Known (C) and Unknown (R) bits from Bob's perspective

2. Bob sends to Alice the set $I = \{2, 3\}$ corresponding to the indexes of the bits known to him.
3. Alice generates a 4-bit key, $K_a = 1011$ where the last 2 bits of K_a must be the same as the last 2 of K_b . All 4 bits from the key are unknown to Alice. She doesn't know which

2 bits from K_a are equal to K_b (only Bob has that information), i.e., from her perspective all bits are random.

$$K_a = \begin{matrix} RR & RR \\ \hline 0 & 0 & 1 & 1 \end{matrix}$$

Figure 10.4: Known (C) and Unknown (R) bits from Alice's perspective

4. Alice encrypts the messages with K_a and sends the encrypted information to Bob. In this example, 1101 will be sent to Bob.

$$\begin{array}{r} 0011 \\ \oplus \quad 1110 \\ \hline 1101 \end{array} \begin{array}{l} \xleftarrow{K_a} \\ \xleftarrow{\text{concat}(m_0, m_1)} \\ \xleftarrow{\text{encrypted information}} \end{array}$$

Figure 10.5: Alice encrypts the messages with K_a

5. Bob decrypts the information received from Alice with K_b and retrieves the message m_1 , which correspond to bits of indexes 2 and 3 (since $i = 1$).
The first two bits of the decrypted information do not represent m_0 .

$$\begin{array}{r} 1011 \\ \oplus \quad 1101 \\ \hline 0110 \end{array} \begin{array}{l} \xleftarrow{K_b} \\ \xleftarrow{\text{encrypted information}} \\ \begin{array}{c} \hline \text{meaningless} & \text{m1} \\ \hline \text{value} & \end{array} \end{array}$$

Figure 10.6: Bob decrypts the information with K_b

10.4 Oblivious Transfer

Oblivious Transfer (OT) is a protocol that allows the sender to transfer one piece of information to a receiver, but remains oblivious as to what piece has been transferred.

For example, consider two parties, Alice and Bob, where Alice has a set $M = \{m_0, m_1\}$ of messages and Bob wants to obtain one of those messages, but won't reveal which one. By OT, Bob is able to receive the message he wanted from Alice without Alice knowing which one she sent to Bob.

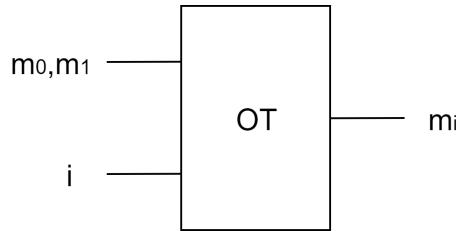


Figure 10.7: Oblivious Transfer Diagram

10.5 Garbled Circuit Protocol

Introduced in 1986 by Andrew Yao, the Garbled Circuit protocol (GC) addresses the case of Two-Party Computation (2PC), without the presence of a trusted third party.

GC allows a secure evaluation of a function given as a Boolean circuit that is represented as a series of logic gates. The circuit is known to both parties.

10.6 Hardware Description Languages

Contrary to Programming Languages such as C or C++, which are used to specify a set of instructions to a computer, Hardware Description Languages (HDL) are computer languages used to describe the structure and behavior of digital logic circuits. They allow for the synthesis of HDL description code into a netlist (specification of physical electronic components, such as AND gates or NOT gates, and how they are connected together).

10.7 TinyGarble

TinyGarble is a GC framework that takes advantage of powerful logic synthesis techniques, provided by both HDL synthesis tools and TinyGarble's custom libraries, in order to improve the overall efficiency of the GC protocol.

It is possible to describe the circuit using High-Level Programming Languages (HLPL) such as C, although High-Level Synthesis (HLS) is required. HLS is performed by High-Level Synthesis tools, such as SPARK for the C language.

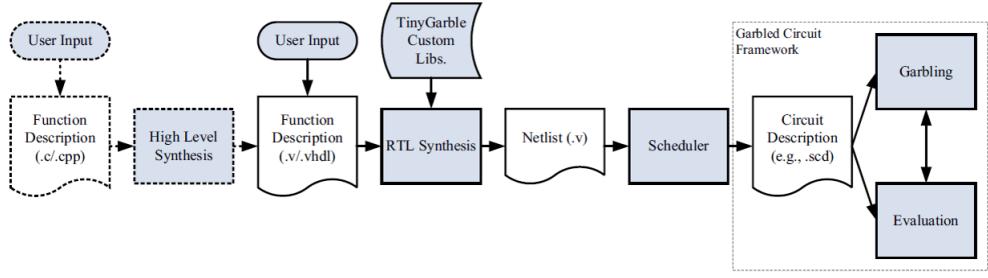


Figure 10.8: TinyGarble Flow Diagram

10.8 ARM2GC

Although circuit description in HLPL is possible, it is not very efficient when compared to HDL circuit description. ARM2GC addresses this problem, significantly improving the performance of garbled circuits described in HLPL.

In the case of 2PC, ARM2GC's approach to GC is based on the ARM processor architecture and consists in providing a public parameter to function f , so that its output would be given by the following expression,

$$z = f(a, b, p) \quad (10.3)$$

, where p represents a public parameter, known to both parties.

In ARM2GC the Boolean circuit required to perform GC is that of a processor to which the compiled binary of the function is given as a public input ($p = \text{compiled binary of the function}$). This optimization is performed by the SkipGate algorithm.

References

- [1] Thorlabs. *Thorlabs Balance Amplified Photodetectors: PDB4xx Series Operation Manual*, 2014.

Capítulo 11

Building C++ Projects Without Visual Studio

This is a guide on how to build C++ projects without having Microsoft Visual Studio installed. All the necessary files will be available in the `\msbuild\` folder on this repository.

11.1 Installing Microsoft Visual C++ Build Tools

Run the file `visualcppbuildtools_full.exe` and follow all the setup instructions;

11.2 Adding Path To System Variables

Please follow this step-by-step tutorial carefully.

1. Open the **Control Panel**.
2. Select the option **System and Security**.
3. Select the option **System**.
4. Select **Advanced System Settings** on the menu on the left side of the window.
5. This should have opened another window. Click on **Environment variables**.
6. Check if there is a variable called **Path** in the **System Variables** (bottom list).
7. If it doesn't exist, create a new variable by pressing **New** in **System Variables** (bottom list). Insert the name **Path** as the name of the variable and enter the following value `C:\Windows\Microsoft.Net\Framework\v4.0.30319`. Jump to step 10.
8. If it exists, click on the variable **Path** and press **Edit**. This should open another window;
9. Click on **New** to add another value to this variable. Enter the following value: `C:\Windows\Microsoft.Net\Framework\v4.0.30319`.
10. Press **Ok** and you're done.

11.3 How To Use MSBuild To Build Your Projects

You are now able to build (compile and link) your C++ projects without having Visual Studio installed on your machine. To do this, please follow the instructions below:

1. Open the **Command Line** and navigate to your project folder (where the .vcxproj file is located).
2. Enter the command:
`msbuild <filename> /tv:14.0 /p:PlatformToolset=v140,TargetPlatformVersion=8.1,OutDir=".\"`, where <filename> is your .vcxproj file.

After building the project, the .exe file should be automatically generated to the current folder.

The signals will be generated into the sub-directory `\signals\`, which must already exist.

11.4 Known Issues

11.4.1 Missing ucrtbased.dll

In order to solve this issue, please follow the instructions below:

1. Navigate to `C:\Program Files (x86)\Windows Kits\10\bin\x86\ucrt\`
2. Copy the following file: `ucrtbased.dll`
3. Paste this file in the following folder: `C:\Windows\System32\`
4. Paste this file in the following folder: `C:\Windows\SysWOW64\`

Attention:you need to paste the file in BOTH of the folders above.

Git creates and maintains a database that store all saved versions of a repository, i.e. all committed states of a folder. To create this database for a specific folder the Git application must be installed in the computer. After Git installation, open the Git console and go to the specific folder and execute the following command:

```
git init
```

This command creates a Git database in the folder *.git* in the root of your repository.

12.1 Data Model

To understand Git is fundamental to understand the Git data model. Git stores objects. Git stores the following objects:

- commits - text files that store a description of the repository;
- trees - text files that store a description of a folder;
- blobs - the files that exist in your repository added with a header;
- tags - is a friendly named pointer to a commit;

The objects are stored in the folder *.git/objects*. Each store object is identified by its SHA1 hash message, i.e. 20 bytes which identifies unequivocally the object. Note that 20 bytes can be represented by a 40 characters hexadecimal string. The ID of each object is the 40 characters hexadecimal string. Note that Git creates a new object every time a file or a folder is changed. In that was the Git store all committed versions of a file.

A commit object like any other object is identified by a SHA1 hash message, and has the following information: a pointer for the root tree of your repository, a pointer for the previous commit, the identification of the author of the commit with a time stamp, the committer with a time stamp and a commit message. See below an example of a commit file contend:

```
tree 2c04e4bad1e2bcc0223239e65c0e6e822bba4f16
parent bd3c8f6fed39a601c29c5d101789aaa1dab0f3cd
author NetXPTO <netxpto@gmail.com> 1514997058 +0000
```

committer NetXPTO <netxpto@gmail.com> 1514997058 +0000

2018-01-03, after meeting Daniel

A tree object is identified by a SHA1 hash message, and has a list of blobs and trees that are inside that tree. See below an example of a tree file contend:

100644	<i>blob</i>	<i>bdb0cabc87cf50106df6e15097dff816c8c3eb34</i>	<i>.gitattributes</i>
100644	<i>blob</i>	<i>50492188dc6e12112a42de3e691246dafdad645b</i>	<i>.gitignore</i>
100644	<i>blob</i>	<i>8f564c4b3e95add1a43e839de8adbfd1ceccf811</i>	<i>bfg-1.12.16.jar</i>
040000	<i>tree</i>	<i>de44b36d96548240d98cb946298f94901b5f5a05</i>	<i>doc</i>
040000	<i>tree</i>	<i>8b7147dbfdc026c78fee129d9075b0f6b17893be</i>	<i>garbage</i>
040000	<i>tree</i>	<i>bdfcd8ef2786ee5f0f188fc04d9b2c24d00d2e92</i>	<i>include</i>
040000	<i>tree</i>	<i>040373bd71b8fe2fe08c3a154cada841b3e411fb</i>	<i>lib</i>
040000	<i>tree</i>	<i>7a5fce17545e55d2faa3fc3ab36e75ed47d7bc02</i>	<i>msbuild</i>
040000	<i>tree</i>	<i>b86efba0767e0fac1a23373aa95884a47c495c5</i>	<i>mtools</i>
040000	<i>tree</i>	<i>1f981ea3a52bccf1cb00d7cb6dfdc687f33242ea</i>	<i>references</i>
040000	<i>tree</i>	<i>86d462afd7485038cc916b62d7cbfc2a41e8cf47</i>	<i>sdf</i>
040000	<i>tree</i>	<i>13bfce10b78764b24c1e3dfbd0b10bc6c35f2f7b</i>	<i>things_to_do</i>
040000	<i>tree</i>	<i>232612b8a5338ea71ab6a583d477d41f17ebae32</i>	<i>visualizerXPTO</i>
040000	<i>tree</i>	<i>1e5ee96669358032a4a960513d5f5635c7a23a90</i>	<i>work_in_progress</i>

A blob is identified by a SHA1 hash message, and has a header plus the compressed contend of a file. The header is very simple is as only information of the type of the object, a blob in this case, and the object contend size. The Git objects are compressed using the zlib library.

12.2 Database Folders and Files

12.2.1 Objects Folder

Git stores the database and the associated information in a set of folders and files inside the the folder *.git* in the root of your repository.

The folder *.git/objects* stores information about all objects (commits, trees and blobs). The objects are stored in files inside folders. The name of the folders are the 2 first characters of the SHA1 40 characters hexadecimal string. The name of the files are the other 38 hexadecimal characters of the SHA1. The information is compressed to save same space but it can be access using low-level Git commands, also named in Git terminology plumbing commands as opposed to porcelain commands, i.e. high-level Git commands.

12.2.2 Refs Folder

The *refs* folder has the following folders *refs/heads*, *refs/remotes*, and *refs/tags*.

Heads Folder

The *refs/heads* folder has information about all branches. Each branch is associated with a text file, the name of the file is the name of the branch. The file has inside the hash of the commit pointed by that branch.

Remotes Folder

The *refs/remotes* folder has information about all remotes and theirs branches. Each remote has a subfolder inside *refs/remotes* with the name of the remote. Inside each folder are files associated with each branches in that remote. Each branch is associated with a text file, the name of the file is the name of the branch. The file has inside the hash of the commit pointed by that branch.

12.2.3 HEAD File

The HEAD file has information about the current branch. HEAD is a pointer to a branch.

Detached HEAD State

When your HEAD doesn't point to branch, points to a commit you are in a detached HEAD state.

12.2.4 Index File

The index is also known in Git as the stage or the cache area. The stage is what is your next commit. It is almost everything that is in your repository, with the exception of new files or folders that were not add to the index yet. To see the files that are in the index at a given time use *git ls-files -stage*. Use *git add <file_name>* to add files to the index, and *git rm -cached <file_name>* to remove files from the index.

12.3 Merging

Merging his about creating a new commit joining the work of two branches that have diverged. The normal merging strategy is a three-way-merging. First a common state is identified. Second are identified files that are not changed in both branches, that files are incorporated in the merged-commit. Files that change in just one branch are also directly incorporated in the merge-commit. The only problem are files that change in both commits. For this files Git try to incorporate both changes comparing line by line. If a line is not changed in both files that line goes for the new merged-file. If it find different lines Git has to decide if that lines are new added lines or if are lines that were changed in both files. To do that Git has to compare one line with the below lines from the other file, if it finds a match it is going to say that the above lines are new added lines in that case the new added lines

are incorporated in the merged-file. If it doesn't find a match it is going to say that both lines were changed and a conflict is generated.

12.3.1 Solving Conflicts

Conflicts are originated by merges that Git cannot handle. In this case Git leave both versions in the merged-file and signaled it so that the user can go their and solve it. The simplest way conflicts are signaled in the following way:

```
«««< HEAD
This is changed in my branch.
=====
This is changed in their branch.
»»»> theirs
```

12.3.2 Fast-Forwards

Happens when there exist a linear history between my branch and the branch that I am going to merge. Git only has to update my pointer (HEAD) to theirs pointer.

12.4 Commands

12.4.1 git add

```
git add <file_name>
```

Computes the SHA1 hash message for the file and add it to the index.

12.4.2 git branch

```
git branch -set-upstream-to=<remote>/<branch> Develop.Nelson
git checkout -b <branch_name> SHA1
```

12.4.3 git cat-file

```
git cat-file -t <hash>
```

Shows the type of the object stored in the Git database associated with the <hash>.

```
git cat-file -p <hash>
```

Shows the content of the object stored in the Git database associated with the <hash>.

12.4.4 git clone**12.4.5 git checkout****12.4.6 git commit****12.4.7 git diff**

show the changes with relation to the previous commit.

12.4.8 git fetch -all**12.4.9 git hash-object**

git hash-object <file>

Calculates the SHA1 hash message of the blob object associated the <file>.

git hash-object -w <file>

Calculates the SHA1 hash message of the blob object associated the <file> and stores the object in the Git database of the repository.

12.4.10 git init**12.4.11 git log**

shows a list of commits from reverse time order (goes back in time), i.e. shows the history of your repository. The history of a repository can be represented by a directed acyclic graph (dag for short), pointing in the forward direction in time.

options

-graph

shows a graphical representation of the repository commits history.

12.4.12 git merge

git merge -abort

12.4.13 git rebase

git rebase -abort

12.4.14 git reflog

Show the list of past commits. Can be used to recover a previous commit by doing the following:

```
git reflog
```

After finding the commit create a branch in that commit with:

```
git checkout -b <branch_name> <commit>
```

12.4.15 git reset

git show

Shows what is new in the last commit.

12.4.16 git rm

```
git rm -cached <file_name>
```

Removes the file from the index but leaves the files in the disk as untracked file.

```
git rm -f <file_name>
```

Removes the file from the stage area and from the disk.

git status

12.4.17 git -version

12.5 The Configuration Files

There is a config file for each repository that is stored in the `.git/` folder with the name `config`.

There is a config file for each user that is stored in the `c:/users/<user name>/` folder with the name `.gitconfig`.

To open the `c:/users/<user name>/.gitconfig` file type:

```
git config -global -e
```

12.6 Alias

12.7 Applications

12.7.1 Meld

12.7.2 GitKraken

12.8 Error Messages

12.8.1 Large files detected

Clean the repository with the [BFG Repo-Cleaner](#).

Run the Java program:

```
java -jar bfg-1.12.16.jar --strip-blobs-bigger-than 100M
```

This program is going to remote from your repository all files larger than 100MBytes. After do:

```
git push --force.
```

Capítulo 13

Simulating VHDL programs with GHDL

This guide will help you simulate VHDL programs with the open-source simulator GHDL.

13.1 Adding Path To System Variables

Please follow this step-by-step tutorial:

1. Open the **Control Panel**.
2. Select the option **System and Security**.
3. Select the option **System**.
4. Select **Advanced System Settings** on the menu on the left side of the window.
5. This should have opened another window. Click on **Environment variables**.
6. Check if there is a variable called **Path** in the **System Variables** (bottom list).
7. **If it doesn't exist**, create a new variable by pressing **New** in **System Variables** (bottom list). Insert the name **Path** as the name of the variable and enter your absolute path to the folder `\LinkPlanner\vhdl_simulation\ghdl\bin`.
Example: `C:\repos\LinkPlanner\vhdl_simulation\ghdl\bin`.
Jump to step 10.
8. **If it exists**, click on the variable **Path** and press **Edit**. This should open another window;
9. Click on **New** to add another value to this variable. Enter your absolute path to the folder `\LinkPlanner\vhdl_simulation\ghdl\bin`.
Example: `C:\repos\LinkPlanner\vhdl_simulation\ghdl\bin`.
10. Press **Ok** and you're done.

13.2 Using GHDL To Simulate VHDL Programs

This guide will only cover the simulation of the VHDL module in this repository. This simulation will take an .sgn file and output its binary information, removing the header. There are two ways to simulate this module.

13.2.1 Requirements

Place a .sgn file in the directory `\vhdl_simulation\input_files\` and rename it to `SIGNAL.sgn`

13.2.2 Option 1

Execute the batch file `simulation.bat`, located in the directory `\vhdl_simulation\` in this repository.

13.2.3 Option 2

Open the **Command Line** and navigate to your project folder (where the .vhd file is located). Execute the following commands:

```
ghdl -a -std=08 signal_processing.vhd
ghdl -a -std=08 vhdl_simulation.vhd
ghdl -e -std=08 vhdl_simulation
ghdl -r -std=08 vhdl_simulation
```

Additional information: The first two commands are used to compile the program and will generate .cf files. Do not remove these file until the simulation is complete.

The third command is used to elaborate the simulation.

The last command is used to run the simulation. If you want to simulate the same program again, you will just need to execute this command (as long as you don't delete the .cf files).

13.2.4 Simulation Output

The simulation will output the file `SIGNAL.sgn`. This file will contain all the processed binary information of the input file, with the header.

