

Switchboard EVM

Audit

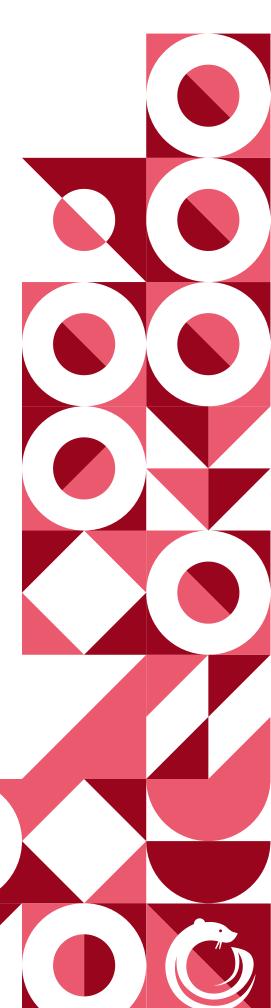


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01 | Executive Summary

Overview

Switchboard engaged OtterSec to perform an assessment of the switchboard-evm program. This assessment was conducted between June 16th and June 30th, 2023. For more information on our auditing methodology, see Appendix B.

Key Findings

Over the course of this audit engagement, we produced 11 findings in total.

In particular, we discovered vulnerabilities that allow an unlimited amount of funds to be stolen (OS-SWB-ADV-00, OS-SWB-ADV-01) alongside a vulnerability that allows for a bounded amount of funds to be stolen under an adversarial block builder assumption (OS-SWB-ADV-03). We also discovered an incorrect garbage collection implementation (OS-SWB-ADV-02). These vulnerabilities may compromise critical protocol assumptions.

Additionally, we advised removing a potentially unsafe delegate call which may escalate into a security threat with future code additions (OS-SWB-SUG-00). We also made recommendations around making a constant configurable (OS-SWB-SUG-01), removing unnecessary modifiers (??), and gas optimization (OS-SWB-SUG-02).

02 | **Scope**

The source code was delivered to us in a git repository at github.com/switchboard-xyz/switchboard-evm. This audit was performed against commit 8c4f247.

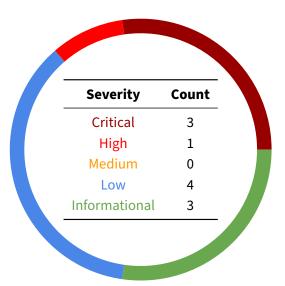
A brief description of the programs is as follows.

Name	Description
switchboard-evm	A smart contract that supports a trust-minimized, decentralized oracle network. Its core components are: Oracle: Posts off-chain updates in a round-robin fashion, which periodically provides a heartbeat to signal readiness.
	 AttestationQueue: Verifies attestation quotes on-chain, granting or revoking oracle permissions based on enclave contents.

03 | Findings

Overall, we reported 11 findings.

We split the findings into **vulnerabilities** and **general findings**. Vulnerabilities have an immediate impact and should be remediated as soon as possible. General findings do not have an immediate impact but will help mitigate future vulnerabilities.



04 | Vulnerabilities

Here, we present a technical analysis of the vulnerabilities we identified during our audit. These vulnerabilities have *immediate* security implications, and we recommend remediation as soon as possible.

Rating criteria can be found in Appendix A.

ID	Severity	Status	Description
OS-SWB-ADV-00	Critical	Resolved	An attacker may steal funds by increasing rewards after quote updates.
OS-SWB-ADV-01	Critical	Resolved	SbFunction.forward has no checks on msg.value, allowing unlimited theft of funds.
OS-SWB-ADV-02	Critical	Resolved	Garbage collection occurs on valid oracles.
OS-SWB-ADV-03	High	Resolved	An attacker may profit by utilizing unbounded gas rebates.
OS-SWB-ADV-04	Low	Resolved	In openInterval, only the interval balance gets updated while the aggregator balance remains unmodified.
OS-SWB-ADV-05	Low	Resolved	The condition for reverting in runTransactions and forward is incorrect.
OS-SWB-ADV-06	Low	Resolved	The guarded modifier implemented in Receiver allows reentrant calls.
OS-SWB-ADV-07	Low	Resolved	removeMrEnclaveToOracleQueue performs removals even when the specified mrEnclave is not present in the array.

OS-SWB-ADV-00 [crit] | Reward Manipulation

Description

An attacker may steal funds from the diamond by increasing rewards after updating an enclave before using it for validation. The reward is taken in updateEnclave, which posts the enclave's cid.

```
function updateEnclave(
   address enclaveId, // enclave id
   bytes calldata cid
) external payable guarded(1) {
    /* ... */

   // we pay the verifier queue's reward to verify the enclave
   uint256 verifierReward = queue.reward;
   if (msg.value < verifierReward) {
      revert ErrorLib.InsufficientBalance();
   }

   // Add balance to enclave
   EnclaveLib.fundEnclave(enclaveId, msg.value);

   /* ... */
}</pre>
```

The reward provided in updateEnclave is distributed upon failing or verifying an enclave. For instance, in failEnclave, a verifier enclave marks the subject enclave as failed and gives rewards to the verifier enclave's owner.

```
contracts/src/enclave/Enclave.sol

function failEnclave(
   address verifierId, // caller's enclave address
   address enclaveId, // enclave address to be verified
   uint256 enclaveIdx
) external guarded(1) {
   /* ... */

   // write enclave failure
   EnclaveLib.setEnclaveFailure(enclaveId);

   // do payout
   if (queue.reward > 0) {
      emit EnclavePayoutEvent(verifierId, enclaveId, queue.reward);
      payable(msg.sender).transfer(queue.reward);
   }
}
```

The vulnerability arises from the possibility of increasing the reward amount between reward provision and reward distribution.

Proof of Concept

For the proof of concept, consider a scenario where an attacker executes the following steps:

- 1. The attacker creates an AttestationQueue with zero rewards.
- 2. The attacker creates an enclave and invokes updateEnclave and forceOverrideVerify, making it available as a verifier enclave. Because the attestation queue reward is zero, no funds are required to call updateEnclave.
- 3. The attacker increases the reward amount by calling setAttestationQueueConfig. The new reward amount is equal to the Ether balance of the diamond.
- 4. The attacker utilizes the verifier enclave to call failEnclave on itself. The attacker, who owns the verifier enclave, will be able to withdraw the queue.reward amount from the diamond, equal to its entire Ether balance.

This method allows the attacker to extract Ether from the contract.

Remediation

Prevent reward payouts in failEnclave and validateEnclave if the reward amount is larger than the enclave balance.

Patch

Fixed in a58f237 and a937a28.

OS-SWB-ADV-01 [crit] | Unchecked Fund Provision

Description

SbFunction. forward allows external calls with user-provided msg.values without checking if sufficient funds are provided. An attacker may steal all Ether from the contract by forging a transaction with a value equal to the diamond balance.

```
function forward(
    TransactionLib.Transaction[] memory transactions,
    bytes[] memory signatures
) external guarded(257) {
    // Just run the functions and verify the signatures
    for (uint256 i = 0; i < transactions.length; i++) {
        TransactionLib.Transaction memory transaction = transactions[i];

    /* ... */

    // EIP2771: call the function with the encoded data and the from address
        (bool success, bytes memory returnData) = transaction.to.call{
            value: transaction.value
            }(abi.encodePacked(transaction.data, transaction.from));
      }
}</pre>
```

Remediation

Disallow transfer of funds through forward by setting the call value to zero.

Patch

Fixed in e833e48.

OS-SWB-ADV-02 [crit] | Incorrect Oracle Garbage Collection

Description

oracleHeartbeat performs garbage collection on a valid oracle instead of an expired one. Each oracle has a field named numRows, which indicates the number of queues associated with it. An oracle may only be added to a queue if its numRows value is zero. Therefore, setting numRows of an oracle back to zero should always be accompanied by queue removal to prevent the oracle from existing in multiple queues simultaneously.

However, oracleHeartbeat erroneously performs setNumRows on oracleId instead of gcOracleId during garbage collection.

By exploiting this vulnerability, an attacker may create an oracle that exists in two queues simultaneously, violating crucial assumptions made by other parts of the code. For example, saveResults assumes that an oracle exists in only one queue and that an oracle's effective heartbeat indicates it is not expired. However, if an oracle is present in multiple queues, it may continue to perform heartbeats even when it is already expired.

Remediation

Change setNumRows(oracleId, 0) to setNumRows(gcOracleId, 0).

Patch

Fixed in e833e48.

OS-SWB-ADV-03 [high] Unbounded Gas Rebate

Description

Block builders may exploit gas rebates, as they do not pay for gas costs. This enables them to set an unusually high gas price, potentially resulting in the theft of Ether from the diamond.

saveResult provides gas rebates equal to the gas spent in the function multiplied by tx.gasprice. Regular users gain no benefits from setting a high tx.gasprice as they are responsible for paying the increased gas cost. However, this does not apply to block builders. An oracle authority may collude with a block builder to set an abnormally high gas price to receive inflated rebates.

```
function saveResult(
   address aggregatorId,
   address oracleId,
   int256 result
) internal {
   // snapshot gas left
   uint256 remainingGas = gasleft();

   /* computation intensive work */

   // handle payment for oracle results - don't err out for now, just don't pay
   uint256 fullReward = reward + (remainingGas - gasleft()) * tx.gasprice;

   // remove funds from interval balance if we can (and if reward is not 0)
   if (aggregator.balanceLeftForInterval >= fullReward && reward != 0) {
        /* ... */

        // pay oracle for gas spend + reward - handing off trust to the oracle
        operator address
            payable(msg.sender).transfer(fullReward);
    }
}
```

The severity of this issue is mitigated by the fact that balanceLeftForInterval bounds the amount of rewards extractable.

Remediation

Place a constant boundary on spent gas.

Patch

Fixed in ee13d95 and 168ec96.

OS-SWB-ADV-04 [low] | Neglected Balance Update

Description

openInterval does not update the aggregator. balance field. Consequently, subsequent withdrawals may fail due to an inaccurate aggregator. balance value, even if there are sufficient funds available within the aggregator.

```
contracts/src/Switchboard/aggregator/Aggregator.sol

function openInterval(address aggregatorId) external payable guarded(4) {
    AggregatorLib.Aggregator storage aggregator = AggregatorLib.aggregators(
        aggregatorId
    );
    uint256 reward = OracleQueueLib.oracleQueues(aggregator.queueId).reward;

if (msg.value < reward * (aggregator.config.minOracleResults + 1)) {
        revert ErrorLib.InsufficientBalance();
    }

AggregatorLib.refreshInterval(
        aggregatorId,
        reward * (aggregator.config.minOracleResults + 1)
    );
}</pre>
```

For instance, the vulnerability may result in an integer overflow during the execution of withdrawIntervalAndBalance, potentially resulting in failures of legitimate oracle result post transactions.

```
contracts/src/Switchboard/aggregator/AggregatorLib.sol

function withdrawIntervalAndBalance(
   address aggregatorId,
   uint256 amount
) internal {
   Aggregator storage aggregator = diamondStorage().aggregators[
        aggregatorId
   ];
   aggregator.balanceLeftForInterval -= amount;
   aggregator.balance -= amount;
}
```

Remediation

Update aggregator.balance in openInterval.

Patch

Fixed in fdaed57.

OS-SWB-ADV-05 [low] | Incorrect Transaction Expiration Checks

Description

The implementation of the deadline expiration checks in runTransactions and forward is incorrect, resulting in these function invocations reverting unconditionally.

```
function runTransactions(
   address functionId,
   uint256 reward,
   TransactionLib.Transaction[] memory transactions,
   bytes[] memory signatures
) internal {
   FunctionLib.SbFunction storage fn = FunctionLib.funcs(functionId);
   for (uint256 i = 0; i < transactions.length; i++) {
        /* ... */

        if (block.timestamp < transaction.expirationTimeSeconds) {
            revert ErrorLib.TransactionExpired();
        }

        /* ... */
   }
}</pre>
```

The condition for expiration checks if the block timestamp is greater than the expiration deadline. Therefore, the direction of the inequality should be inverted to accurately reflect this condition.

Remediation

Invert the inequality.

Patch

Fixed in a58f237 and 4e1d94c.

OS-SWB-ADV-06 [low] | Incorrect Reentrancy Guard Implementation

Description

The implementation of the guard modifier is incorrect and may be bypassed through specific call paths.

```
contracts/src/Switchboard/util/Recipient.sol

function hasEntryCode(uint256 code) internal view returns (bool) {
    return diamondStorage().code & code != 0;
}

modifier guarded(uint256 code) {
    uint256 startCode = UtilLib.getCode();

    // if this is the first call to a guarded function, set the code (permissions if (startCode == 0) {
        UtilLib.setCode(code);
}

    // if the code is the same as the code - this represents a re-entry, gotta revert if (startCode != 0 && startCode == code) {
        revert if (startCode != 0 && !UtilLib.hasEntry();
    }

    // if the code is not a subset of the code - entry is not allowed if (startCode != 0 && !UtilLib.hasEntryCode(code)) {
        revert ErrorLib.InvalidEntry();
    }

    -;

    // reset the code to the previous value - will end up being 0
    UtilLib.setCode(startCode);
}
```

The guard modifier allows entry if bits intersect between the callee's code and the caller's code. However, this implementation allows for a nested call path such as $guarded(258) \rightarrow guarded(257) \rightarrow guarded(1)$ to bypass the intended restriction. This is a violation as 258 & 1 = 0, but calling guarded(257) in the middle enables this bypass.

Remediation

Re-implement reentrancy guards to prevent reentering unless self-forwarding behavior is intended.

Patch

Fixed in d57d785 and d9ab8a0.

OS-SWB-ADV-07 [low] | Incorrect Array Removal Logic

Description

removeMrEnclaveToOracleQueue fails to revert when an element is missing, leading to the unintended removal of an entry. The function utilizes a linear search on the mrEnclaves array to locate the index intended for removal. However, it mistakenly proceeds with the removal process regardless of the presence of mrEnclave, removing the default index.

```
contracts/src/Switchboard/util/Recipient.sol

function removeMrEnclaveFromOracleQueue(
   address queueId,
   bytes32 mrEnclave
) internal {
   DiamondStorage storage ds = diamondStorage();
   AttestationConfig storage config = ds.queueAttestationConfigs[queueId];

   uint256 idx = 0;
   for (uint256 i = 0; i < config.mrEnclaves.length; i++) {
      if (config.mrEnclaves[i] == mrEnclave) {
       idx = i;
       break;
      }
   }
   uint256 lastIdx = config.mrEnclaves.length - 1;
   config.mrEnclaves[idx] = config.mrEnclaves[lastIdx];
   config.mrEnclaves.pop();
}</pre>
```

Remediation

Perform element removal only when the element is present.

Patch

Fixed in 7dc4922.

05 | General Findings

Here, we present a discussion of general findings during our audit. While these findings do not present an immediate security impact, they represent anti-patterns and may lead to security issues in the future.

ID	Description
OS-SWB-SUG-00	saveResults uses a delegate call for invoking oracleHeartbeat, which is susceptible to msgSender forgery.
OS-SWB-SUG-01	verifyEnclave has a fixed expiration threshold of 20 seconds, which may necessitate modifications depending on chain properties or conditions.
OS-SWB-SUG-02	saveResults calculates a function selector in a gas-costly manner by invoking keccak256 on the function signature string.

OS-SWB-SUG-00 | Potentially Dangerous Self Call Pattern

Using a delegate call in saveResults for invoking oracleHeartbeat raises concerns due to the potential interaction with EIP2771. Introducing EIP2771 alters the behavior of determining the sender of the message in a self-call context, where the sender is determined by inspecting the call data rather than relying on the msg.sender value, as implemented in getMsgSender.

```
function getMsgSender() internal view returns (address payable signer) {
   signer = payable(msg.sender);
   if (msg.data.length >= 20 && signer == address(this)) {
      assembly {
       signer := shr(96, calldataload(sub(calldatasize(), 20)))
      }
   }
}
```

Thus, if saveResults is called in a self-call context, getMsgSender will return oracleId. By setting oracleId to the address of an oracle queue authority, an attacker may bypass authority checks.

Currently, reentering saveResults is infeasible due to the absence of external address calls in functions marked with guarded (4). However, we advise addressing this issue nonetheless due to future code additions.

Remediation

Handle EIP2771 in the self-call to oracleHeartbeat.

Patch

Fixed in d9ab8a0.

OS-SWB-SUG-01 | Constant Expiry Threshold

Description

verifyEnclave relies on a fixed expiration threshold of 20 seconds, which may prove inadequate in scenarios where the chain is congested or if the contract is deployed on an EVM sidechain with a different block generation time. The function currently reverts if the time reported by the verifier enclave authority and the execution time (block.timestamp) differ by more than 20 seconds. While this measure is commendable, it is advisable to make the threshold configurable as this contract may be deployed on chains with varying properties or under changing chain conditions.

Remediation

Implement setter functions to set this threshold.

Patch

Fixed in ca211e0.

OS-SWB-SUG-02 | Gas Optimization

Description

The invocation of keccak256 for selector computations in saveResults may be optimized. Using keccak256 for selector calculations involves a computationally expensive operation, resulting in higher gas consumption. An alternative approach to improve efficiency is to replace the invocation of keccak256 with the .selector syntax available in Solidity.

Remediation

Replace the call with Oracle.oracleHeartbeat.selector or validate the oracle by observing state variables.

Patch

Fixed in 55ac77b.

$\land\mid$ Vulnerability Rating Scale

We rated our findings according to the following scale. Vulnerabilities have immediate security implications. Informational findings can be found in the General Findings section.

Critical

Vulnerabilities that immediately lead to loss of user funds with minimal preconditions

Examples:

- Misconfigured authority or access control validation
- · Improperly designed economic incentives leading to loss of funds

High

Vulnerabilities that could lead to loss of user funds but are potentially difficult to exploit.

Examples:

- Loss of funds requiring specific victim interactions
- Exploitation involving high capital requirement with respect to payout

Medium

Vulnerabilities that could lead to denial of service scenarios or degraded usability.

Examples:

- · Malicious input that causes computational limit exhaustion
- Forced exceptions in normal user flow

Low

Low probability vulnerabilities which could still be exploitable but require extenuating circumstances or undue risk.

Examples:

Oracle manipulation with large capital requirements and multiple transactions

Informational

Best practices to mitigate future security risks. These are classified as general findings.

Examples:

- Explicit assertion of critical internal invariants
- · Improved input validation

B | Procedure

As part of our standard auditing procedure, we split our analysis into two main sections: design and implementation.

When auditing the design of a program, we aim to ensure that the overall economic architecture is sound in the context of an on-chain program. In other words, there is no way to steal funds or deny service, ignoring any chain-specific quirks. This usually requires a deep understanding of the program's internal interactions, potential game theory implications, and general on-chain execution primitives.

One example of a design vulnerability would be an on-chain oracle that could be manipulated by flash loans or large deposits. Such a design would generally be unsound regardless of which chain the oracle is deployed on.

On the other hand, auditing the implementation of the program requires a deep understanding of the chain's execution model. While this varies from chain to chain, some common implementation vulnerabilities include reentrancy, account ownership issues, arithmetic overflows, and rounding bugs.

As a general rule of sum, implementation vulnerabilities tend to be more "checklist" style. In contrast, design vulnerabilities require a strong understanding of the underlying system and the various interactions: both with the user and cross-program.

As we approach any new target, we strive to get a comprehensive understanding of the program first. In our audits, we always approach targets with a team of auditors. This allows us to share thoughts and collaborate, picking up on details that the other missed.

While sometimes the line between design and implementation can be blurry, we hope this gives some insight into our auditing procedure and thought process.