# **Segmentation and Paging**

### • Segmentation:

- Divide process memory into segments based on their functionality.
- Segment Table holds base (segment start address) and limit (segment size) registers.
- Logical memory request has format < segment, offset>
  - Request successful if *offset* < *limit*, otherwise *segmentation fault*.

#### Paging:

- Divide main memory into fixed-size Frames (generally powers of 2).
- Divide processes' memory into frame-sized chunks called Pages (generally ~4KB).
- Logical memory request has format < page, offset>
- Process Page Tables stored by a process track in which Frames processes' Pages are stored.
- Process Table stored in main memory containing [page, frame] entries.
  - Page Tables are very long; require optimizations.

# Shared Pages:

- Multiple processes of same program often run together.
- Possible if program code non-reentrant (i.e. stateless because of no static or global variable usage).
- Processes storing common-code can share *Pages*.

### Page Table Structure

#### Hierarchical:

- Table contains *outer-page* entries, which contains *inner-page* table
- If p bits used for page address, then p-k used to outer-page, k used for inner-page

# ■ Hashed (with Linked-Lists):

- page address used as hash-value
- Page is searched for in the Linked-list using page address

#### Inverted:

- Replace *page table* with *frame table*.
  - frame table maps frame number to page number (and Process ID)
- Required to search *entire* frame table (*very*, *very* slow)
  - Operation can be *optimized* using hardware optimizations

# • Hardware Support using Translation Lookaside Buffer:

- Simply a *cache* for *page table* containing recently used entries
- Cache hit means [page, frame] (key-value pair) entry is in cache
- *Cache miss* means entry not in cache; slower main-memory access required to *page table*.
  - If *page* is not in main-memory, *page-fault* occurrs
  - page-fault is a trap; process will initiate reading from hard-disk.