NOVEMBER TUESDAY * features : in charge of essential I/O functionalities Task dispatching Cault recovery embedded with RDD8 (Ruilient distributed dataset) - handles partitioning data across all nodes in clustile. - It hadds them in memory pool of the cluster as a single unit 2. Spark SQL 3 works to acress structured & simi-structured into. & enables powerful, interactive, analytical app. across both streaming & historical data. > Features & Cast based optimizes Mid query failt taxlerent Full compatibility with existing nive data 3. Spark streaming - Add on to core Jepark API which allows scalable, high - throughput, fault-tolerant stream processing of live data streams. - Accese data from Sources like kapka, Aume, kinesis or TEP cocket A - It was Nicho- batching for heat time etterming

• • • 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 •

DAY 326-039 WEEK 48

2022

(Resilient Distributed Sataset) K RDD It is a fundamental data structure of Spark. & It is an immutable collection of objects which computes on the different node of the duster

every dataset in RDB is logically partitioned

across services to be computed on diff. node.

- Resilient, i.e fault tolerant with the help of RDD lineage graph (DAG), to becompute missing or damaged partitions due to node

jailures. - Distributed, data residing on multiple nodes

2 bataset, represents records of data you work with.

· features

- In-memory computation - stones everything in RAM

- Lazy Evaluations - spark computes transformation when an

action requires a result for driver program. fault Tolkance.

→ Because of lineage - Immutability

- cannot be changed or updated

- Partitioning + fundamental unit of paralleliem in spark Rob - (an create a partition the some transformation

MANTESSMINT (SSMINTESSMINTESSMINTESS DEC - Kell 3 18 15 Ce 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 · 2022 > Users can state which RDDs they will reuse & choose a storage strategy for them.

- Every app, contain it executor
NOV MIW DESSMIWIESSMIWIESSMIWIESS - A unit of work that will be sent to one executar.

the application & creates spark context of

& maintains are of the states of spark dustre

1t must inreface with duster manager

in order to actually get physical resource

2. Spark Context - coordinate the spark applications, running as independent sets of priocesses on a ducker - It acquires executors on nodes in cluster.

is then, it sends your app, code to the executor - At last, the spark context sinds task to me executors to run.

3 ander Manager -> Autratio resources accross app.

4. Worker Nede I slave node. sole is to run the approde in dustre.

5. Executor , sinds rights to driver.

- An executor is a process launched for an application on worker node - It huns tasks & keeps data in memory or

diek storage across them. It read & write data to external source

2022

NOVEMBER

Limitations No inbuilt optimization engine.

Handling structured data - need to specify the schema.

Performance limitation - Al it involves the overhead of Garbage collection & Java sirialization which are

expensive when data grows. Storage limitation - spill over, if size of RDD is langur than R'AM then it uses disk for remaining John RDD size.

· Ways to Create RDDs in Spark.

1. Parallelized collection (parallelizing) - by taking an existing collection in the program 2 passing it to spark context's parallised method !

> used in initial stage, as it creates RDD quikly

2. External batacit (Referencing a dataset). - we can use textfile method

> It takes URL of file & read it as a collection of line. + loading CON JSON, text file.

3. Creating RDD from existing RDD.

Transformation is the way to weath an NOV MINIESSMINTESSMINTESS 2022 RDD3 from outrady 16 existing 24 RODD, 50 · · ·

. RDD Persistence & Caching

- Persistance is an optimization technique in which saves the ributt of RDD evaluation. a using this we save the intermediate result

so that we can use it it required 4 It reduces the computation breakhead

> when we use the cache () method we can store all the RDD in-memory

- we can persist the RBD in memory & use it efficiently across parallel operations.

- The diff. b/w cache () & pursist () is that using Cache () the default storage level is nemory only while using persuit () we can use various storage bevels

> When we persist RDD each node stores any partition of it that it computes in memory & makes it rebeable for future use. This process speed up the computation. - when the RDD is computed for first time, it

is kept in the memory on the stade. The cache

memory of the spark is fault tolerant so whenever any partition of DRAD is lost, it can be reconciled by thankformation operation that originally created it

· · · 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31 ·

5. Siek Jonly

Storage well of pursuit () 1: Memory only (spill over in disk) 2. Memory and disk (serialized Java do 3. memory only lex 4. Memory and disk sek

·" Unpercet RDD - Spark automatically drops out the old

data partition in LAU (least recently up) fashion. > 'we can remove manually using

· 2 R bb Operations

13. Transformation - produces new RDD from existing RDD. - Applying transformation built and RDD linear with the entire parent RDDs of the final 5 RDD dependency graph. - Are laxy in nature i.e they get executed when we call an action.

- Two types of transformation: 1. Novedus transformation -> All elements that are seq, to compute the records in single partition live in the single

partition of parent RDD.

map(), filter(), flat Map(), maplating
sample(), union(). NOV MTWTFSSMTWTFSSMTWTFSSMTWTFSSMTW

2. Wide transformation - All elements that meg, to compute the necords

in the single partition may live in many partitions of parent RDO! - eg. groupbykey'(), reduceby key(), Jain() intersection (), distinct (), contacion (), supartition (), coalisce ()

22. Action works with actual data when action is perform bous not form new RDD. - The values of action are stored to topeach

drivers or to the external storage system > Action is one of the way of sending data from Executor to the driver > count() - no. of elements is returned

collect () - returns the entire RDDs content to driver program. take (1) - returns n no. by elements from RAD.

top () - extract top elethents from 'RDD. count By Value () - returns, occurrence of each element - reduce () - operation like addition, takes two fold () - Take zero value as input, Timput, diff. - reduce throws an exception for empty

collection, but fold is defined for empty " collection aggregate(), foreach ()

· · · 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29

2022 • 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30

& distributed collection of data organized in → bata organised into named columns. → (an eay that it is a relational table with Processes large amount of structured data

-> Contains schema (illustration of the structure -> Immutability, in-memory, resilient, distribut > sources data from structured data file, tables in rive, external dbs on existing ROD.

- Available in scala, Java, Python & R. > Lata frame over RDD 1 frounds memory management - data is stored in off hear memory in

2022 **IUNE**

DAY 153-212 WEEK 23 **THURSDAY**

@ Optimized Execution plan - query optimizer, where execution plan is heated for the execution of a quiny

· Limitation of RDD " Does not have any built-in optimization engine.

2 No provision to handle structured data. · features of DataFrame

2 named column, equivalent to table in ROBM. > Deals with structured Esemistructured data o formate. eg. Auro, CSV, elastic search, cassandro. - beals with storage eystoms-HDRS, HWS, MYSGIER s catalyst supports optimization, 4 phases: 1. Analyze logical plan to solve references 2. Logical plan optimization 3. Physical planning

· 4. cade generation to ampile part of query to java byterade · Limitation - Does not provide provision for compile time type safety! so, we need to make a structure in order to manipulate.

M T W T F S S M T W T F S S M T W T F S S M T W T F S S M T W T F S S • • • 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 31

computing capability

* batafram.

al data)

23 Similar to table in ROBNS.

good optimization technique.

 JUN
 M T W T F S S M T W T F S S M T W T F S S M T W T F S S M T W T F S S M T W T F S S M T W T F S S

 2022
 • 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19 20 21 22 23 24 25 26 27 28 29 30 • • •

	FRIDAY					JUNE
* batas	et					
- Stron - Repres - Provid	ly type	1 & 1	map	to rel	ation	u schem
- Repre	s'ether	Ducti	p' bere	leine	with	evicogn
2 Provid	us both	type	eafu	ty & o	gect-	outhle
progr	amming	intie	face.	J		
10			V			
·11 Featur	. ó ÷	· · · · · · · · · · · · · · · · · · ·			1	, :
- Ostini		00.00.	Costali	D	45.41.0-	Limien
- Optimi - 2 Analysi - at com	Lat com	Dita 7	ima	chart	of of	A A and
at iom	all tim	mark to	MICE, -	Chuck	Synto	x & ana
- Pennist	ent stop	10.01	ervio	182061	Po	مامام سمامام
- Persist - Faster	Comput	TIN L	1 Hoan	(110	ه مو م	wayan
2 Less m	emory	Consum	Ation	- etru	cture	Ol data
2 less m	ataret	is kno	wn.	Silv		of data
3						
4	18			1 4 4 5		
5						
6						
		ν,				į.
7						
	. 7	,			3	2 1
		,				
-	e c					