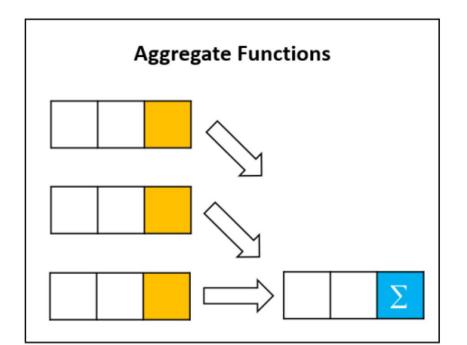
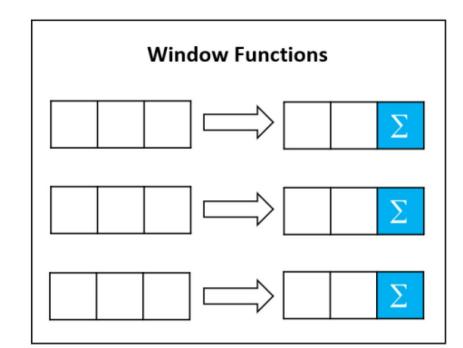
## Window Functions In SQL



- Window functions: These are special SQL functions that perform a calculation across a set of related rows.
- How it works: Instead of operating on individual rows, a window function operates on a group or 'window' of rows that are somehow related to the current row. This allows for complex calculations based on these related rows.
- **Window definition**: The **'window'** in window functions refers to a **set of rows**. The window can be defined using different criteria depending on the requirements of your operation.
- **Partitions**: By using the **PARTITION BY** clause, you can divide your data into smaller sets or **'partitions'**. The window function will then be applied individually to each partition.
- Order of rows: You can specify the order of rows in each partition using the ORDER BY clause. This order influences how some window functions calculate their result.
- **Frames**: The **ROWS/RANGE** clause lets you further narrow down the window by defining a **'frame'** or subset of rows within each partition.
- Comparison with Aggregate Functions: Unlike aggregate functions that return a single result per group, window functions return a single result for each row of the table based on the group of rows defined in the window.
- Advantage: Window functions allow for more complex operations that need to take into account not just the current row, but also its 'neighbours' in some way.









# **Example**

|    | Results 📋 | Messages   | _              |           |              |             |
|----|-----------|------------|----------------|-----------|--------------|-------------|
|    | order_id  | order_date | customer_name  | city      | order_amount | grand_total |
| 1  | 1002      | 2017-04-02 | David Jones    | Arington  | 20000.00     | 37000.00    |
| 2  | 1007      | 2017-04-10 | Andrew Smith   | Arington  | 15000.00     | 37000.00    |
| 3  | 1008      | 2017-04-11 | David Brown    | Arington  | 2000.00      | 37000.00    |
| 4  | 1001      | 2017-04-01 | David Smith    | GuildFord | 10000.00     | 50500.00    |
| 5  | 1006      | 2017-04-06 | Paum Smith     | GuildFord | 25000.00     | 50500.00    |
| 6  | 1004      | 2017-04-04 | Michael Smith  | GuildFord | 15000.00     | 50500.00    |
| 7  | 1010      | 2017-04-25 | Peter Smith    | GuildFord | 500.00       | 50500.00    |
| 8  | 1005      | 2017-04-05 | David Williams | Shalford  | 7000.00      | 13000.00    |
| 9  | 1003      | 2017-04-03 | John Smith     | Shalford  | 5000.00      | 13000.00    |
| 10 | 1009      | 2017-04-20 | Robert Smith   | Shalford  | 1000.00      | 13000.00    |



# **Window Function Syntax**

```
function_name (column) OVER (
    [PARTITION BY column_name_1, ..., column_name_n]
    [ORDER BY column_name_1 [ASC | DESC], ..., column_name_n [ASC | DESC]]
)
```

- function\_name: This is the window function you want to use. Examples include ROW\_NUMBER(), RANK(), DENSE\_RANK(), SUM(), AVG(), and many others.
- (column): This is the column that the window function will operate on. For some functions like SUM(salary)
- **OVER** (): This is where you define the window. The parentheses after OVER contain the specifications for the window.
- PARTITION BY column\_name\_1, ..., column\_name\_n: This clause divides the result set into partitions upon which the window function will operate independently. For example, if you have PARTITION BY salesperson\_id, the window function will calculate a result for each salesperson independently.
- ORDER BY column\_name\_1 [ASC | DESC], ..., column\_name\_n [ASC | DESC]: This clause specifies the order of
  the rows in each partition. The window function operates on these rows in the order specified. For example, ORDER
  BY sales\_date DESC will make the window function operate on rows with more recent dates first.



# **Different Types of Window Functions**

There are three main categories of window functions in SQL: **Ranking functions**, **Value functions**, and **Aggregate functions**. Here's a brief description and example for each:

### **Ranking Functions:**

• ROW\_NUMBER(): Assigns a unique row number to each row, ranking start from 1 and keep increasing till the end of last row

SELECT Studentname,
Subject,
Marks,
ROW\_NUMBER() OVER(ORDER BY Marks desc)
RowNumber
FROM ExamResult;

|   | Studentname | Subject | Marks | RowNumber |
|---|-------------|---------|-------|-----------|
| 1 | Isabella    | english | 90    | 1         |
| 2 | Olivia      | english | 89    | 2         |
| 3 | Lily        | Science | 80    | 3         |
| 4 | Lily        | english | 70    | 4         |
| 5 | Isabella    | Science | 70    | 5         |
| 6 | Lily        | Maths   | 65    | 6         |
| 7 | Olivia      | Science | 60    | 7         |
| 8 | Olivia      | Maths   | 55    | 8         |
| 9 | Isabella    | Maths   | 50    | 9         |



• **RANK()**: Assigns a rank to each row. Rows with equal values receive the same rank, with the next row receiving a rank which skips the duplicate rankings.

SELECT Studentname,
Subject,
Marks,
RANK() OVER(ORDER BY Marks DESC) Rank
FROM ExamResult
ORDER BY Rank;

|   | Studentname | Subject | Marks | Rank |
|---|-------------|---------|-------|------|
| 1 | Isabella    | english | 90    | 1    |
| 2 | Olivia      | english | 89    | 2    |
| 3 | Lily        | Science | 80    | 3    |
| 4 | Lily        | english | 70    | 4    |
| 5 | Isabella    | Science | 70    | 4    |
| 6 | Lily        | Maths   | 65    | 6    |
| 7 | Olivia      | Science | 60    | 7    |
| 8 | Olivia      | Maths   | 55    | 8    |
| 9 | Isabella    | Maths   | 50    | 9    |



• **DENSE\_RANK()**: Similar to RANK(), but does not skip rankings if there are duplicates.

SELECT Studentname,
Subject,
Marks,
DENSE\_RANK() OVER(ORDER BY Marks DESC) Rank
FROM ExamResult
ORDER BY Rank;

|   | Studentname | Subject | Marks | Rank |           |
|---|-------------|---------|-------|------|-----------|
| 1 | Isabella    | english | 90    | 1    |           |
| 2 | Olivia      | english | 89    | 2    |           |
| 3 | Lily        | Science | 80    | 3    |           |
| 4 | Lily        | english | 70    | 4    |           |
| 5 | Isabella    | Science | 70    | 4    |           |
| 6 | Lily        | Maths   | 65    | 5    | * Similar |
| 7 | Olivia      | Science | 60    | 6    | Rank      |
| 8 | Olivia      | Maths   | 55    | 7    |           |
| 9 | Isabella    | Maths   | 50    | 8    |           |



**Value Functions:** These functions perform calculations on the values of the window rows.

• **FIRST\_VALUE():** Returns the first value in the window.

```
SELECT
employee_name,
department,
hours,
FIRST_VALUE(employee_name) OVER (
PARTITION BY department
ORDER BY hours
) least_over_time
FROM
overtime;
```

|   | employee_name     | department | hours | least_over_time |
|---|-------------------|------------|-------|-----------------|
| • | Diane Murphy      | Accounting | 37    | Diane Murphy    |
|   | Jeff Firrelli     | Accounting | 40    | Diane Murphy    |
|   | Mary Patterson    | Accounting | 74    | Diane Murphy    |
|   | Gerard Bondur 💙   | Finance    | 47    | Gerard Bondur   |
|   | William Patterson | Finance    | 58    | Gerard Bondur   |
|   | Anthony Bow       | Finance    | 66    | Gerard Bondur   |
|   | Leslie Thompson   | IT         | 88    | Leslie Thompson |
| _ | Leslie Jennings   | IT         | 90    | Leslie Thompson |
|   | Loui Bondur 🗸     | Marketing  | 49    | Loui Bondur     |
|   | Gerard Hernandez  | Marketing  | 66    | Loui Bondur     |
|   | George Vanauf     | Marketing  | 89    | Loui Bondur     |
|   | Steve Patterson   | Sales      | 29    | Steve Patterson |
|   | Foon Yue Tseng    | Sales      | 65    | Steve Patterson |
|   | Julie Firrelli    | Sales      | 81    | Steve Patterson |
|   | Barry Jones       | SCM        | 65    | Barry Jones     |
|   | Pamela Castillo   | SCM        | 96    | Barry Jones     |
|   | Larry Bott        | SCM        | 100   | Barry Jones     |



• **LAST\_VALUE():** Returns the last value in the window.

SELECT employee\_name, department,salary, LAST\_VALUE(employee\_name) OVER ( PARTITION BY department ORDER BY salary ) as max\_salary FROM Employee;

| employee_name | department | salary | max_salary |
|---------------|------------|--------|------------|
| Vishal        | Accounting | 40000  | Ravi       |
| Ravi          | Accounting | 60000  | Ravi       |
| Nilesh        | Finance    | 55000  | Abdul      |
| Sushant       | Finance    | 65000  | Abdul      |
| Abdul         | Finance    | 68000  | Abdul      |
| Jai           | IT         | 45000  | Mohit      |
| Aman          | IT         | 60000  | Mohit      |
| Mohit         | IT         | 70000  | Mohit      |



LAG(): Returns the value of the previous row.

SELECT
Year,
Quarter,
Sales,
LAG(Sales, 1, 0) OVER(
PARTITION BY Year
ORDER BY Year,Quarter ASC)
AS NextQuarterSales
FROM ProductSales;

|    | Year | Quarter | Sales     | NextQuarterSales |
|----|------|---------|-----------|------------------|
| 1  | 2017 | 1       | 55000.00  | 0.00             |
| 2  | 2017 | 2       | 78000.00  | 55000.00         |
| 3  | 2017 | 3       | 49000.00  | 78000.00         |
| 4  | 2017 | 4       | 32000.00  | 49000.00         |
| 5  | 2018 | 1       | 41000.00  | 0.00             |
| 6  | 2018 | 2       | 8965.00 \ | 41000.00         |
| 7  | 2018 | 3       | 69874.00  | 8965.00          |
| 8  | 2018 | 4       | 32562.00  | 69874.00         |
| 9  | 2019 | 1       | 87456.00  | 0.00             |
| 10 | 2019 | 2       | 75000.00  | 87456.00         |
| 11 | 2019 | 3       | 96500.00  | 75000.00         |
| 12 | 2019 | 4       | 85236.00  | 96500.00         |



• **LEAD():** Returns the value of the next row.

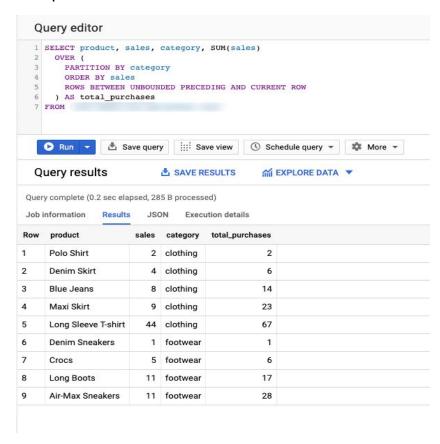
SELECT Year,
Quarter,
Sales,
LEAD(Sales, 1, 0) OVER(
PARTITION BY Year
ORDER BY Year,Quarter ASC)
AS NextQuarterSales
FROM ProductSales;

|    | Year | Quarter | Sales    | NextQuarterSales |                              |
|----|------|---------|----------|------------------|------------------------------|
| 1  | 2017 | 1       | 55000.00 | 78000.00         |                              |
| 2  | 2017 | 2       | 78000.00 | 49000.00         |                              |
| 3  | 2017 | 3       | 49000.00 | 32000.00         |                              |
| 4  | 2017 | 4       | 32000.00 | 0.00             | Lead function on             |
| 5  | 2018 | 1       | 41000.00 | <b>8965.00</b>   | PARTITION for<br>Year column |
| 6  | 2018 | 2       | 8965.00  | 69874.00         |                              |
| 7  | 2018 | 3       | 69874.00 | 32562.00         |                              |
| 8  | 2018 | 4       | 32562.00 | 0.00             |                              |
| 9  | 2019 | 1       | 87456.00 | 75000.00         |                              |
| 10 | 2019 | 2       | 75000.00 | 96500.00         |                              |
| 11 | 2019 | 3       | 96500.00 | 85236.00         |                              |
| 12 | 2019 | 4       | 85236.00 | 0.00             |                              |



### **Aggregation Functions:** These functions perform calculations on the values of the window rows.

- SUM()
- MIN()
- MAX()
- AVG()





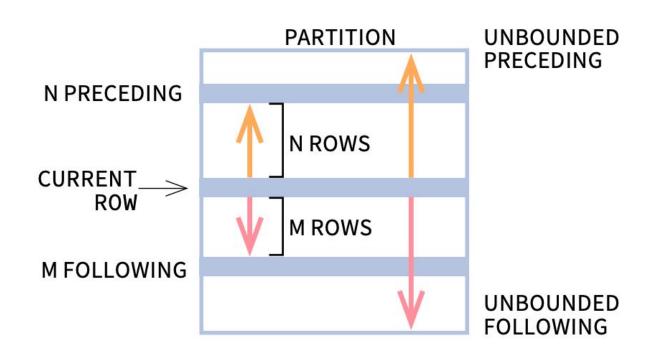
#### Frame Clause in Window Functions

- The frame clause in window functions defines the subset of rows ('frame') used for calculating the result of the function for the current row.
- It's specified within the OVER() clause after PARTITION BY and ORDER BY.
- The frame is defined by two parts: a start and an end, each relative to the current row.
- Generic syntax for a window function with a frame clause:

```
function_name (expression) OVER (
    [PARTITION BY column_name_1, ..., column_name_n]
    [ORDER BY column_name_1 [ASC | DESC], ..., column_name_n [ASC | DESC]]
    [ROWS|RANGE frame_start TO frame_end]
)
```

- The frame start can be:
  - UNBOUNDED PRECEDING (starts at the first row of the partition)
  - N PRECEDING (starts N rows before the current row)
  - CURRENT ROW (starts at the current row)
- The frame end can be:
  - UNBOUNDED FOLLOWING (ends at the last row of the partition)
  - N FOLLOWING (ends N rows after the current row)
  - CURRENT ROW (ends at the current row)
- For **ROWS**, the frame consists of N rows coming before or after the current row.
- For **RANGE**, the frame consists of rows within a certain value range relative to the value in the current row.





## **ROWS BETWEEN Example**



SELECT date, revenue,
SUM(revenue) OVER (
ORDER BY date
ROWS BETWEEN UNBOUNDED PRECEDING AND CURRENT ROW) running\_total
FROM sales
ORDER BY date;

## Input Table

# Output Table

| sales     |            |         | date       | revenue | running_total |
|-----------|------------|---------|------------|---------|---------------|
| record_id | date       | revenue | 2021-09-01 | 1515.45 | 1515.45       |
| 1         | 2021-09-01 | 1515.45 | 2021-09-02 | 2345.35 | 3860.80       |
| 2         | 2021-09-02 | 2345.35 |            |         |               |
| 3         | 2021-09-03 | 903.99  | 2021-09-03 | 903.99  | 4764.79       |
| 4         | 2021-09-04 | 2158.55 | 2021-09-04 | 2158.55 | 6923.34       |
| 5         | 2021-09-05 | 1819.80 | 2021-09-05 | 1819.80 | 8743.14       |



### **RANGE BETWEEN Example**

```
SELECT
shop,
date,
revenue_amount,
MAX(revenue_amount) OVER (
ORDER BY DATE
RANGE BETWEEN INTERVAL '3' DAY PRECEDING
AND INTERVAL '1' DAY FOLLOWING
) AS max_revenue
FROM revenue_per_shop;
```

| shop   | date       | revenue_amount | max_revenue |
|--------|------------|----------------|-------------|
| Shop 1 | 2021-05-01 | 12,573.25      | 18,847.54   |
| Shop 2 | 2021-05-01 | 11,348.22      | 18,847.54   |
| Shop 1 | 2021-05-02 | 14,388.14      | 18,847.54   |
| Shop 2 | 2021-05-02 | 18,847.54      | 18,847.54   |
| Shop 1 | 2021-05-03 | 9,845.29       | 18,847.54   |
| Shop 2 | 2021-05-03 | 14,574.56      | 18,847.54   |
| Shop 1 | 2021-05-04 | 11,500.63      | 18,847.54   |
| Shop 2 | 2021-05-04 | 16,897.21      | 18,847.54   |
| Shop 1 | 2021-05-05 | 9,634.56       | 21,489.22   |
| Shop 2 | 2021-05-05 | 14,255.87      | 21,489.22   |
| Shop 1 | 2021-05-06 | 11,248.33      | 21,489.22   |
| Shop 2 | 2021-05-06 | 21,489.22      | 21,489.22   |
| Shop 2 | 2021-05-07 | 15,517.22      | 21,489.22   |
| Shop 1 | 2021-05-07 | 14,448.65      | 21,489.22   |

### **Output Table**

# **Common Table Expression**



A Common Table Expression (CTE) in SQL is a named temporary result set that exists only within the execution scope of a single SQL statement. Here are some important points to note about CTEs:

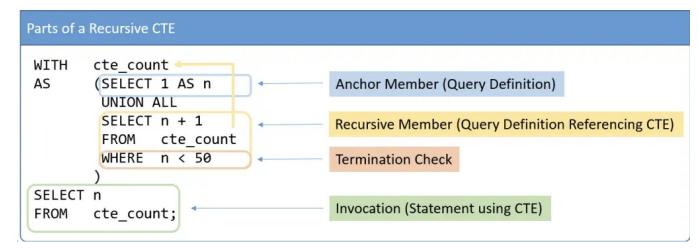
- CTEs can be thought of as alternatives to derived tables, inline views, or subqueries.
- They can be used in SELECT, INSERT, UPDATE, or DELETE statements.
- CTEs help to simplify complex queries, particularly those involving multiple subqueries or recursive queries.
- They make your guery more readable and easier to maintain.
- A CTE is defined using the WITH keyword, followed by the CTE name and a query. The CTE can then be
  referred to by its name elsewhere in the query.

Here's a basic example of a CTE:

```
WITH sales_cte AS (
    SELECT sales_person, SUM(sales_amount) as total_sales
    FROM sales_table
    GROUP BY sales_person
)
SELECT sales_person, total_sales
FROM sales_cte
WHERE total_sales > 1000;
```



Recursive CTE: This is a CTE that references itself. In other words, the CTE query definition refers back to the CTE name, creating a loop that ends when a certain condition is met. Recursive CTEs are useful for working with hierarchical or tree-structured data.





### Subqueries in SQL

 IN: The IN operator allows you to specify multiple values in a WHERE clause. It returns true if a value matches any value in a list.

SELECT \* FROM Orders WHERE ProductName IN ('Apple', 'Banana');

• **NOT IN:** The NOT IN operator excludes the values in the list. It returns true if a value does not match any value in the list.

SELECT \* FROM Orders WHERE ProductName NOT IN ('Apple', 'Banana');

- ANY: The ANY operator returns true if any subquery value meets the condition.
- **ALL:** The ALL operator returns true if all subquery value meets the condition.
- EXISTS: The EXISTS operator returns true if the subquery returns one or more records.
- NOT EXISTS: The NOT EXISTS operator returns true if the subquery returns no records.

### **Views**



A view in SQL is a virtual table based on the result-set of an SQL statement. It contains rows and columns, just like a real table. The fields in a view are fields from one or more real tables in the database.

Here are some key points about views:

- You can add SQL functions, WHERE, and JOIN statements to a view and display the data as
  if the data were coming from one single table.
- A view always shows up-to-date data. The database engine recreates the data every time a
  user queries a view.
- Views can be used to encapsulate complex queries, presenting users with a simpler interface to the data.
- They can be used to restrict access to sensitive data in the underlying tables, presenting only non-sensitive data to users.





CREATE VIEW View\_Products AS SELECT ProductName, Price FROM Products WHERE Price > 30;

# **Employee**

| EmployeeID | Ename | DeptID | Salary |
|------------|-------|--------|--------|
| 1001       | John  | 2      | 4000   |
| 1002       | Anna  | 1      | 3500   |
| 1003       | James | 1      | 2500   |
| 1004       | David | 2      | 5000   |
| 1005       | Mark  | 2      | 3000   |
| 1006       | Steve | 3      | 4500   |
| 1007       | Alice | 3      | 3500   |

CREATE VIEW emp\_view AS SELECT DeptID, AVG(Salary) FROM Employee GROUP BY DeptID;

Create View of grouped records on Employee table

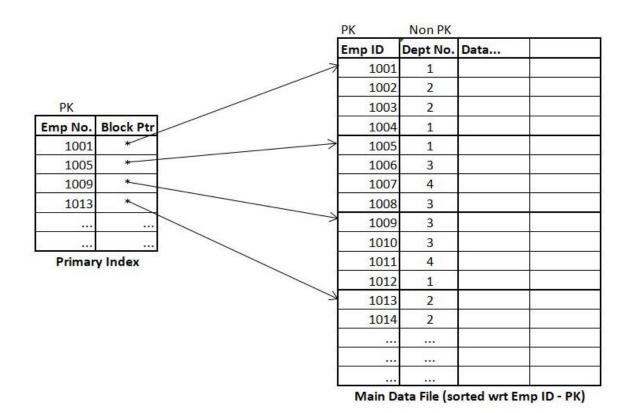
emp\_view

| DeptID | AVG(Salary) |
|--------|-------------|
| 1      | 3000.00     |
| 2      | 4000.00     |
| 3      | 4250.00     |

# Indexing



Indexing in databases involves creating a data structure that improves the speed of data retrieval operations on a database table. Indexes are used to quickly locate data without having to search every row in a table each time a database table is accessed.



### Why is Indexing Important?



Indexes are crucial for enhancing the performance of a database by:

- 1. **Speeding up Query Execution:** Indexes reduce the amount of data that needs to be scanned for a query, significantly speeding up data retrieval operations.
- 2. **Optimizing Search Operations:** Indexes help in efficiently searching for records based on the indexed columns.
- 3. *Improving Sorting and Filtering*: Indexes assist in sorting and filtering operations by providing a structured way to access data.
- 4. **Enhancing Join Performance**: Indexes on join columns improve the performance of join operations between tables.

## **Advantages of Indexing**

- 1. Faster Data Retrieval: Indexes make search queries faster by providing a quick way to locate rows in a table.
- 2. **Efficient Use of Resources**: Reduced query execution time translates to more efficient use of CPU and memory resources.
- 3. **Improved Performance for Large Tables**: Indexes are particularly beneficial for large tables where full table scans would be time-consuming.
- 4. **Better Sorting and Filtering**: Indexes can improve the performance of ORDER BY, GROUP BY, and WHERE clauses.

### **How to Choose the Right Indexing Column**

- 1. **Primary Key and Unique Constraints**: Always index columns that are primary keys or have unique constraints, as they uniquely identify rows.
- 2. Frequently Used Columns in WHERE Clauses: Index columns that are frequently used in WHERE clauses to filter data.
- 3. **Columns Used in Joins**: Index columns that are used in join conditions to speed up join operations.
- 4. **Columns Used in ORDER BY and GROUP BY:** Index columns that are used in ORDER BY and GROUP BY clauses for faster sorting and grouping.
- 5. **Selectivity of the Column**: Choose columns with high selectivity (columns with many unique values) to maximize the performance benefits of the index.



- Use Column Names Instead of \* in a SELECT Statement
- Avoid including a HAVING clause in SELECT statements

The HAVING clause is used to filter the rows after all the rows are selected and it is used like a filter. It is quite useless in a SELECT statement. It works by going through the final result table of the query parsing out the rows that don't meet the HAVING condition.

#### Example:

Original query:

SELECT s.cust\_id,count(s.cust\_id)

FROM SH.sales s

GROUP BY s.cust\_id

HAVING s.cust\_id != '1660' AND s.cust\_id != '2';

#### Improved query:

SELECT s.cust\_id,count(cust\_id)

FROM SH.sales s

WHERE s.cust id != '1660'

AND s.cust id !='2'

GROUP BY s.cust\_id;



Eliminate Unnecessary DISTINCT Conditions

Considering the case of the following example, the DISTINCT keyword in the original query is unnecessary because the table\_name contains the primary key p.ID, which is part of the result set.

### Example:

Original query:
SELECT DISTINCT \* FROM SH.sales s
JOIN SH.customers c
ON s.cust\_id= c.cust\_id
WHERE c.cust\_marital\_status = 'single';

Improved query:
SELECT \* FROM SH.sales s JOIN
SH.customers c
ON s.cust\_id = c.cust\_id
WHERE c.cust\_marital\_status='single';



Consider using an IN predicate when querying an indexed column

The IN-list predicate can be exploited for indexed retrieval and also, the optimizer can sort the IN-list to match the sort sequence of the index, leading to more efficient retrieval.

Example:
Original query:
SELECT s.\*
FROM SH.sales s

WHERE s.prod\_id = 14

OR s.prod\_id = 17;

Improved query:

SELECT s.\*

FROM SH.sales s

WHERE s.prod\_id IN (14, 17);



Try to use UNION ALL in place of UNION

The UNION ALL statement is faster than UNION, because UNION ALL statement does not consider duplicate s, and UNION statement does look for duplicates in a table while selection of rows, whether or not they exist.

#### Example:

Original query:

SELECT cust\_id

FROM SH.sales

UNION

SELECT cust id

FROM customers;

### Improved query:

SELECT cust\_id

FROM SH.sales

**UNION ALL** 

SELECT cust\_id

FROM customers;