I Sot Theory

- 1. Language of math & our focus Practical set theory Eg, induction frecursion - size of sets (cardinality)
- 2. Foundations of Mathematics - axiomatic approach

Model Theory \coprod

Study of (arbitrary) mathematical structures in a fixed Sormal language "the theory of mathematical theories"

<u>Chapter 1: Ordinals</u>

 $N = \{0,1,2,3,...\}$ counting: - enumerating, ordering > ordinals
- measuring (finite) sets > coordinals

set theoretic subtleties Develop an infinitary analogue, Start with some basics of posets.

- 1 Dell A partially ordered set (or poset) is a set E with a binary relation RSE2 satisfying:

 1 Reflexivity: (a,a) ER for all a E;

 2. Antisymmetric: (a,b) ER and (b,a) ER then a=b;

 - 3. Transitivity: if (a,b) ER and (b,c) ER then (a,c) ER. A poset is linear (or total) if for all a,b e E, either aRb or bRa. (Notation we often write aRb instead of (a,b) ER.)

- ex A any set, E = P(A) = set of all subsets of A

 R = containement

 As long as A has more than one element, (E,R) is a

 poset that is not total.
- ex (Z, ≤) is a total ordering
- 2 Dell A strict poset is like a poset except reflexivity is replaced by autirestexivity. $\neg (aRa)$ for all $a \in E$ A strict posed is linear (or total) if for all $a : b \in E$, either aRb or bRa or a = b.

 ex. (Z, <)

Remark: If (E,R) is a poset, define aR+b to mean aRb and a+b. Then (E,R+) is a strict poset.

If (E,R) is a strict poset, define oR=b to mean aRb or a=b.

Then (E,R=) is a poset

Notation: We will usually denote a poset by (E, ≤) and its corresponding strict poset by (E, <).

Dell A linearly ordered set (E, \leq) is well-ordered if every non-empty subset of E has a least element. That is, if $D \subseteq E$ and $D \neq \emptyset$ then there exists a E such that $\alpha \leq b$ for all $b \in D$.

- ex (M, \leq) is well-ordered, (Z, \leq) is not, neither is (R, \leq)
- 3 Lemma: Well-orderings are rigid. That is, if (E, \leq) is a well-ordering then the only automorphism of (E, \leq) is the identity. An automorphism of a poset (E, \leq) is a bijection f: E > E such that $a \leq b \mathrel{(=)} f(a) \leq f(b)$ for all $a,b \in E$.

Proof: Suppose I is an automorphism of (E, E). Consider D= face; f(a) +a}

Suppose towards a contradiction, that fit id. Then D= Ø. Let acD be least.

Case 1: fla) < a. Then fla) & D so f(fla) = fla). Hence fla) = a as f is injective. Contradiction.
Case 2: a < f(a). Then f'(a) < a as f' is an automorphism.

So f'(a) &D, so f'(a)=a, so a=f(a). Contradiction.

By totallity, these are the only possibilities. So find.

4 Corollary: If (E, ≤) and (F, ≤) are isomorphic well-orderings then there is a unique isomorphism f: E= F!

An isomorphism of posets (E, \leq) and (F, \leq) is a bijection f: E > F such that $\alpha \leq b \iff F(\alpha) \leq F(b)$ for all $\alpha, b \in E$.

5 Lemma: A distinct well-ordering is not isomorphic to any initial segment of itself

For a linear order (E,<) and beE, let Exb := {xeE; xxb}. Then (Ea, <) is the initial segment of (E,<) determined by b

ex $((0,1),<) \xrightarrow{\cong} ((0,1/2),<)$ via $\times \mapsto \times/2$

Proof: Suppose f: (E,<) -> (Exo,<) is an isomorphism, for be E Let

D= {xeE, f(x) +x}

Note beD since f(b) EExb and beExb. So D + Ø. Let a ED be least. Break into two cases again (exercise).

We are interested in (IN, <); it is the prototype of a strict well-ordered set. We want sets and membership to be basic and derive (construct all other math objects in terms of these.

Define

0:=
$$\emptyset$$

1:= $\{0\}$ = $\{\emptyset\}$
2:= $\{\emptyset, \}$ = $\{\emptyset, \{\emptyset\}\}$
3:= $\{\emptyset, \}, \{\emptyset\}, \{\emptyset\}, \{\emptyset\}\}$
 $\{\emptyset, \{\emptyset\}\}\}$
S(n):= $\{\emptyset, \{\emptyset\}\}$
1 successor function

Note: The usual ordering is induced by E.

We have already implicitly been using some axioms of set theory

Emplysed Axiom: There exists a set that has no members, devoted by Q.

Next we need to produce {x} from x. In fact, we introduce:

Pairset Axiom: Given sets x,y there exists a set whose elements are precisely x,y, denoted {x,y}.

Note if x=y then {x,y}= {x}

Extensionality: Two sets are equal is and only if they have the same members.

To get Sn) from n we need:

(Notation: Given sets x,y we say x is a subset of y, x = y, if every element of x is on element of y.)

Unionsel Axiom: Gluen a set of theme exists a set whose members 2015 01 071 are the members of the members of x, denoted by Ux. ie yeUx if and only it yez for some zex

So XUY:= Ufx,y3. In particular, S(n)= Ufn, in 3.

While this gives us a way to produce each natural number it does not prove the existence of the "set of natural numbers". Why not add an axiom soujing: There exists a set whose members are precisely the natural numbers. ie yew (=> (y=0) or (y=1) or---

Problem: We want an axiom to assert the existence of a set having certain properties, where the properties are definite condition.

6 Det) xey and x=y are definite conditions, where x,y are either sets or indeterminates standing for sets. IF P.Q are definite conditions then so are:

- not P, -P - Pand Q, Pro

- Por Q, PVQ

(Note that P=> 0) is definite since it is equivalent to -PVQ.)

- for all x, P, VXP

- there exists x, P, JxP A condition is definite if it arises in this way in finitely many steps.

Note:

- Emptyset Axiom: There exists a set x satisfying - By (yex) ??

- Poirsel Axiom Given Xiy, there exists a set p satisfying (Xep) \((yep) \lambda (zep =) (z : x) \((z : y) \).

An operation H on sets is definite if the condition y = H(x)iz a destrite condition.

ex The successor S(x) = XU {x} is definite since y=S(x) if and only if (zey) (> ((zex) V(z=x)))

We will obtain the set of natural numbers as the smallest set containing O and closed under S.

Infinity Axiom: There exists a set I contained 0 and closed under the successor function:

(DEI) ^ (YX (XEI => S(X) EI))

exercise: if H(x) is a definite operation
then H(x)ey is a definite condition

We want to pick out the smallest such I. Suppose I is as given by the Infinity Axiom. $\bigcap \{ J \subseteq I ; O \in J \text{ and } if \chi \in J \text{ then } S(\chi) \in J \}$

We need more axioms to do this.

Powerset Axiom: Given a set x there exists a set whose elements are the subsets of x, denoted P(x). $\forall z(z \in P(x) \leftrightarrow \forall t (t \in z \rightarrow t \in x))$

Separation Axiom: Given a set x and a definite condition P, there exists a set whose elements are precisely the elements of x which satisfy P, denoted by {yex; P(y)}.

Important that:

1. Pis definite

2. we are inside a (lixed) set x

Now, in

NITEP(I); OET, and XET -> S(X) ET3

as before:

· I is as given by the infinity axiom

· PCI) is given by the powersel axiom

· the set is oldined by the axiom of separation

· we can prove intersections exist by the axioms given so for (exercise)

7. Def The set of natural numbers is the set w:= N{JeP(I);0€3, and x€3=5(x)€3}.

Where I is any set given by the infinity axions.

Exercise: Prove that the above is independent of choice of I.

A useful and immediate consequence of this construction of wis:

8 Induction Principle: If Jew, OEJ, and S(X) & T whenever XeJ, then Jew.
Proof: We have Jew by assumption, and we get we J by definition of w.

One proves a ld about w this way.

I Lemma: Every element of wis a subset of w. Acof: Let

J-{new; new?

Then OEJ since O= \$\phi\$ is a subset of any set vacuously. Now suppose neJ. We would S(n) \in \text{Well S(n) = nunn } \in \text{We would S(n) \in \text{We.}}
If teS(n) then t=n or ten. If t=n then tew, and it ten than tew as new. Hence S(n) \in \text{We.} Thus S(n) \in \text{J}. Therefore \text{J} = \text{W.}

Similarly arguments yield:

10 Roposition: (w, e) is a strict well-ordering. Roof: see notes, 1.13.

This motivates:

" Dell An ordinal is a set & such that

- is every element of a is a subsect of or;
- ii) (0,8) is a strict well-ordering.

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The Zermelo-Frankel axioms of set theory are:

- Extensionality
- Emphyset
- -Paliset
- Unionset
- -Insinity
- Powerset
- Separation
- Replacement Axiom

Replacement Axion: If H is a definite operation on sets and x is a set than there exists a set y which is the image of Hlx. ie There exists a set y such that $z \in y \in Z$ -H(t) for some tex

Aside: Contersian products and functions

Def) Given sets x,y, the ordered pair is $(x,y) = 12x3, \{x,y\}3.$

Exercise : (x,y) = (x,'y') (=) x=x' and y=y'

Given X,Y sets, xeX, yeY, (x,y) & P(P(XUY)).

The set of all ordered pairs exists, and is denoted by X×Y=P(P(XuY)):
Pairs (X,Y) = { peP(X,Y); p={x,y}, xeX, yeY}
Singletons(X) = { SEP(X); S={x} Some xeX}

**XX—Pairs (Singletons (X), Pairs (X,Y))

ce correction 20150114

Cartesian products exist.

Defl Given sets X,Y, a function f: X-> Y is a subset T's X×Y such that few all xex, there is a unique element of Y, which we denote by f(x), such that (x,f(x)) eT.

Remark We are identifying a function with its graph. Sometimes in practice we trad functions as operations.

Exercise. A function is precisely and operation restricted to a set.

Recall the definition of an ordinal.

ex Each natural number, and w, are ordinals.

12 Lemma: Suppose of B are ordinals. If orch then ar B. Proof Let D= Bld. Note D + Q. Let de O be the boost element, with respect to (β, ϵ) . We show $\alpha = d$. Note deBimplies deB. So we claim x=d as subsets of B. Claim: dea. Verification If not, let xed a. Then xcBla=D. SoxED and x<d. Contradicts chaire of d. Claim Red. Verification: Let xea. As x, deB, either x=d or x<d or d<x. We cannot have x=d as d &a. We cannot have dex, (Wrong grood: dex, x < a => d < a => dea => *) as dex implies dex, and red implies x sox los or Ban ordinal), and so ded. Contradiction. Therefore xxd, so xxd, as desired. Thus a=de B. 123 13 Proposition.

(a) Every member of an ordinal is an ordinal.

(b) No ordinal is a member of itself.

(c) If a is an ordinal then so is its successor S(a) := aufaz.

(d) The intersection of two ordinals is an ordinal.

Proof!

well-ordering. Let yex. We want yex. Let zey. As x,y,zeB and z<y<x, we get z<x. So zex. Hence yex.

(b) Suppose aca for an ordinal of. As a is a shirld poset, OKO. This

is a contradiction as a : a.

(c) Easy enough

(d) Exercise.

1

14 Roposition:

(a) If a, B are ordinals then either NEB or BEA or a= B.

(rea) . (der)

(b) Any set of ordinals is a strict well-ordering under E. Proof:

(a) Note CARB is an ordinal with GABE G and GABEB. If both containements are proper, then GABE GAB by bemma 12. So GABEA (Wlog), so GEB => GEB or GCB (=> GEB by lem 10). It I be a set of ordinals. Antireflexivity follows From 13(b). Linearity follows from 14(a). We show that EB well-cordered. Suppose ASE with A+D. Let VEA.

Case 1: anA +cp. Let A: anA, a non-empty subset of A. Let a ∈ A be least of A'. Assume, for a contradiction, that there is a be A with loca. Then bea, so be anA=A', contradicting that a was least in A'. Thus a is least in A.

Cose 2: an A=\$. Then VBEA, B&a, so either B=d of arB.
Thus a is beast in A.

Remark: We only use the set E because the collection of all ordinals isn't a set.

Correction:

XxY = {pe P(P(XUY)); Zun | uep nuep n \$ Vz(zep -> z-uvz=v) A Jay (xeu ~ Yz(ze V-> z=x) A revayev A As(Sen-) S=XNS=A) VyeXVdel)]

11 Proposition (cont.)

(c) If E is a set of ordinals then sup E = UE is amordinal.

(d) There is no set consisting of all ordinals.

Boot.

(c) sup[is a set of ordinals by 13(a). 14(b) -> (sup E, E) is a

strict well-ordering. Let acsupE. So are y for some yeE. So as y is an ordinal. If KED then XEYS SUPE. So OF SUPE. Therefore SUPE is an ardial.
(d) Suppose such a set E exists. Then 14(b) -s (E,E) is a shirt Well-ordering. And to the And if are and xear then by 13(1), x is an ordinal so xet. Honce are. Therefore E is an ordinal. Therefore E & E. Contradicting 13(6). 10

Convention: Given ordinals of and B we write of B for deB.

13 Lemma:

(a) a ordinal then S(a) > a and there is nothing in between.
(b) E + \$\phi\$ set of ordinals then supE is the beast upper bound of E.

(1) E set of ordinals, there exists a beat ordinal ad in E.

(a) S(a)= auza's so a eS(a) so a < S(a). If x < S(a) then X=Q OV REQ => X < Q.

(M) First show upper bound. are E, supE < a => supE & UE = supE, a contradiction.

then acys, yet -> 9 < y => a and an upper bound.

(C) Let Q= 35(S(M supE)). We claim E & X, XEE => X = SUPE = X<S(supt) = 5 X<S(S(supt)). So E = Q.

If E= a => sup E= sup(S(S(sup E))) = S(sup E) *

(Exercise: Why B=S(supE) won't do?)
Let μ be least in $\alpha \in Suppose x \notin E$ ordinal.

If $x < \alpha$ then $x \ge \mu$ by charge of the ruise $x \ge \alpha \ge \mu$. Therefore μ is the least ordinal not in E

1.6 Transfinite induction/recursion

16 Theorem: Suppose P(X) is a definite condition satisfying:

(*) — If α is an ordinal and P(B) for all Pxα, then P(α).

Thun P is true of all prelimate.

Proof (P(b)) holds xiocuouslyi) Suppose P(α) is false. Let

D = PB ≤ α, P(B) is false 3 ≠ Φ.

Let α pe D be least. If B < α, then B ≠ D so P(B) is frue.

By (*), P(α) is true. Contractiction.

Det) A successor ordinal is an ordinal of the form Sa) for some or. A limit ordinal is an ordinal thad is not a successor.

A reformulation:

- 17 Corollary: (Second form of transfinhe induction)
 P(x) definhe condition such that:
 - 1. P(0) is true;
 - 2. If P(B) then P(S(B));
 - 3. If 0 >0 is a limit ordinal and P(B) for all BKO then P(O). Then P is true of all ordinals.

Ironstinite Recordion

Induction is used to prove definite statements about all ordinals. Recursion is used to construct definite operations an ordinals.

A partial function on ordinals is a function whose domain is a set of ordinals.

Note: If F is a definite operation on ordinals and α is an ordinal then F by is a partial function on ordinals.

Theorem: Suppose G is a definite operation on partial functions on ordinals. Then there exists a unique definite operation on ordinals, F, satisfying

F(a) = G(Fla)

Proof (sketch):

Second Form of transfinite recursion

typo in

19 Corollary: Suppose G. is a set, G2 is a definite operation on sets, and Gs is a definite operation on partial functions on ordinals. Then there is a definite operation of on ordinals satisfying:

Orderal Arithmetic

20 Del (Ordinal addition)

Fix an ordinal B, and define B+a for all ordinals a by transmite rewriton (2nd Sorn):

Note: In terms of translivite recursion:

$$G_2 = S$$

 $G_3(f) = \sup\{im(f)\}$

Remarks:

(b) not communative: I+w=supflen; n<w3= w+ 5(w)=w+1

(c) ordinal arithmetic restricted to finite ordinals is usual arithmetic

21 Del (Ordinal Roduct)

Remark: Not commutative

w.5 = m.1 + m = (m.0 + m) + m = (0 + m) + m

= (sup {O+n; n<w}) + w = w+w = sup {w+n; n<w} > w+1>w.

22 Proposition: Q, B, S ordinals.

(a) 0<B (=> S+0<S+B

(b) S+a= S+B => A=B

(c) (Q+B)+8 = Q+(B+8)

(d) for all 8+0, 0<B<=> 5.0< S.B

(e) for all 8 +0, S. a = S.B => A = B.

(f) (a,B); - a,(B;s)

23 Dell (ordinal exponentiation)

Ora limit Ba := enb{Bx, x ca}

24 Theorem Every strict well-ordering is isomorphic to an ordinal. Moreover, both the ordinal and the isomorphism are unique. Proof: Suppose (E, L) is a strict well-ordering. Uniquenous of the isomorphism dollows from corollary 4 (rigidity). Next we show uniqueness of the ordinal. Note for ordinals 9, B, a< to it and only it a is an initial segment of B. Now suppose (E, L) was isomorphic to both a and B. Then either A<B or a=B or B< a. If a<B then since fit- B is an isomorphism, and a is an initial segment of B, we got an initial segment $E_{S'(G)}$ of $(E_{1}\times)$ So $E_{B'(G)}\cong \alpha\cong E$.

This contractions lemma 5, as no well-ordering is isomorphic to an initial segment of itself. Similarly we convol have 1200, so or B. Tirally we show existence. Let

A: {XEE; (Ex, X) is isomorphic to an orderal} Note if E = & then a = 0 works. So assure E + \$P. Let xe E be loost wit L. Then Ex=\$ =0 50 xcA. Hence A+0.

Now define f on A by letting f(a) be the ordinal which is isomorphic to (E_{α}, λ) . Then im(f) is a set of ordinals, so let α be the beost ordinal not im(f).

We claim:

I) $\alpha = im(f)$, so $f: A \rightarrow d$ is surjective;

2) $f: A \rightarrow d$ is injective;

3) $f: (A, \lambda) = (\alpha, \epsilon)$;

Y) E = A.

Exercise (or see notes). Thus $(E, \lambda) = (A, \lambda) \geq (\alpha, \epsilon)$.

Ordinals one a good translinite extension of a for enumeration but not for measuring size: w+1>w have the same size.

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Del Two sets A, B are equinumerous, denoted IAI=1BI, if there is a bijection I:A-B.

Note: We have not yet defined IAI.

Proposition (Schröder-Bernstein)

[Al=181 H and only it there exist injections F:A-B and g:B-A.

Lemma: For any inhabe ordinal α , $|\alpha| = |\alpha+1|$. Proof: Define $g: \alpha+1 \rightarrow \alpha$ by $g(x) = \begin{cases} x+1 & \text{if } x < \omega \\ x & \text{otherwise} \end{cases}$

Dell A cordinal is an ordinal that is not equinomerous to any (shirtly) lesser ordinal.

Examples:

- · new, n is a condinal
- · w is a cordinal

Note: Every infinite cardinal is a limite ordinal. (By the lemma.) The converse is lake. Ex lw+wl=lwl.

Defl A set is countable if it is equinumerous with w or some naw.

Proposition. For every set E there is a unique cardinal, denoted h(E), which is the beast ordinal that is not equinumerous to any subset of E.

Note: h(w) is an uncountable cardinal.

Proof. Suppose such an ordinal exists, h(E), then h(E) is cardinal: Indeed, if B<h(E) then 1B1=1A1, simple ASE. Hence h(E)+B. Consider

W= {(A, X); A = E, X well-ordering on B}

Wisa set by separation. If $(A, X) \in W$, theorem 1.36 says that $(A, X) \cong (\alpha, \epsilon)$. Write $\alpha : f(A, X)$. We have a definite operation. By replacement, Im(f) is a set. We claim that $Im(f) = \{\alpha; \alpha = |A|, \text{ some } A \subseteq E\}$.

(skeleh, see notes)

So h(f) least ordinal not in Tun(f) works. Uniqueness is any a

We have lots of uncountable cardinals. We really want every set to be equinumerous with a cardinal. This would imply every set admits a strict well-orderings. This cannot be proved from ZF. We need a 9th axiom.

Dell Suppose F is a set (of sets). A choice function on F is a dunction C: F -> UF such that C(F) & F for all F& F.

Axiom of choice: Every set of non-empty sets has a choice function

Remarks:

(1) This too is a set existence axiom:

There exists a subset TSFXUF such that for all FEF

there is a unique deUF such that (F,d)EF, and such
that deF."

(2) It is not the case that whenever you choose elements from sets, you are necessarily using the axiom of choice.

ex Suppose Fraset of non-empty subsets of w. Than C.F.-W.F. given by ((1) being the beast of F is a choice Sunction. We didn't use choice.

ex Suppose F : {A} So, A+D. Let xeA. Then C: F -> UF given by c(F)-x is a choice function. We did not use choice.

Theorem: The following are equivalent:

1) Axiom of choice; (AC)

2) Well Ordering Principle; (WOP) (every set admits a strict well-ordering)
3) Zom's Lemma. (If (E.R) is a strict exect in which every totally

3) Zom's Lemma. (If (E, P) is a shield poset in which every totally ordered subset (chair!) has an upper bound, then

(E,R) has a maximal devicent)

Proof: (1)=>(2): Let A be a set. Let L:P(A)EA)->UP(A)=A be a choice

function. Define Frecursively on ordinals to A by

Function. Define Frecureively on ordinals to A by $F(\alpha) = \begin{cases} U(A) Im(F(\alpha)) & \text{if } A | Im(F(\alpha) + \phi), \\ 0 & \text{otherwise.} \end{cases}$

where Θ is an ardinal not in A.

We claim that $F(\alpha) = \theta$ for some ordinal $\alpha < h(A)$.

Assume, for a contradiction, that F doesn't half

Note that if $F(\alpha) + \theta$ $\forall \alpha \in E$ for some set E of ordinals, then

FIE: E-A is injective. So if the claim Pails, then FTh(A): h(A) - A is an injection Contradiction

So let a be the least ordinal such that F(x) = O. So we have

Fra. a - A. As F(a) = O, Im (Fra) = A. Thus Fra is a bijection. We can now use this bijection and the well-ordering on a to get a well-ordering on A.

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(2)=>(3): Suppose (E,R) is a strict poset satisfying

every chain has an upper bound.
Assume, for a contradiction, that (E,R) has no maximal element.
Let < be a strict well-ordering on E. Let h(E) be the least ordinal not equinumerous with any subset of E. Build an embedding of h(E) lubo E, by recursion:

F:h(E) >E

F(0) = e (some fixed element of E

any orch(E): F(0+1) = Eleast element XEE such that F(a) RX (possible as F(a) 15 nd maximal in E)

OKAKHEY limb. F(B) = <-least XEE st F(X)RX YXKB

(possible since Im(FIB) is a chain in E) (by (4))

Thus F is injective.

(3)=>(1). Let I be a sel of non-empty sets, Want a choice Sunction on I. Let

1= { all partial choice functions }

Torn's lewing implies that there is a maximal f: G-OG choice function, G=F. If GFF, bet FEFIG, XeF. Then

Further, G=F. If GFF, bet FEFIG, XeF. Then

FUS(F, X)?

is a partial choice function strictly extending f, contradicting maximality. So G= F and f is a choice function or F.

We work in ZFC from now on. (9 axions)

Lemma: Every set is equinumerous with a unique cardinal. Proof: Uniquees is clear: distinct cardinals are not equinumerous. Existence: Let X be a set. Let < be a strict well-ordering on X. Then $1.36 \Rightarrow (X, <) \cong (\alpha, \epsilon)$ for some ordinal α . In particular, 1X = (a). Let

S= {B= 03 |B|- 1X13. Let BES be least. Then BIS a carclinal if of < B and BI-IJI then IXI=IXI SU YES, contradicting minimality of F. Then BB as desired.

Denote the unique cardinal equinumerous with X by IXI, and call it the <u>cardinality</u> of X.

Lemma: Given A,B, we have IAI < 181 if and only if there is an injection AUB. Proof IF IAIS 181 then A C = 1A | = 1B | C >> B and Here Gore A C>B.

Conversely, suppose ALB. Then
[ALBA-BL->1B] and so IA/ IBI. If IBI < IA! then IBI - IAI. So S-B implies IAI=1BI, This controdicts that IAI is a cardinal. Thus

Corollary: Given A,B, there exists an embedding A - B or an IF(A) ! (A) embedding B - A.

Lemma: Suppose f: A-B & a function. Then | Im(s) | < |A| Proof: Let

F. {f'(b); b \in Im(f) \is. f'(b) nears \(\(\lambda \) \(\lambda \) Then F exists by replacement as bi-of (6) to definite, and is a self of non-comply subsets of A. By AC, there is a choice functional or J. Let g. Im(f) - A be given by

g(b) = c(f.(b)).

Then g is injective.

1A1 < 1B1.

Recall that X is countable means |X| < w.

Lemma: A counteble union of countable sets is countable.

Proof: Suppose (A) < w and for each a.e.h. la! < w. For yeA, let Sx = fach; yeas. Note Sx + Ø, and x -> Sx is definite.

So let S = fox; yeUA's. We can get a choice function

C: S -> A (since & US - A).

Tor each a.e.A, there is un embedding a.c., w. As before, using replacement and choice, we get for each a.e.A an inj.

fa: a -> w. We also have g: A -> w. Now define F:

UA - wxw by F(y) = (g(c(x)), fam (x)). Then

F is injective (check), so IVA ! < |wxw!.

It remains to show that luxus! < |w!. But we know how to do this (exercise).

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Ordinal-emmeration of cardinals

Remark: k cardinal, h(k) & the least ordinal not equinumerous with a subset of k, is a cardinal.

N(K)>K since h(K) = lh(K) |> |K| + K

If I cardinal, A>R then IXI= A>K-IRI, so I is not equinumerous with any subset of R. Hence I > h(R).

So h(K) is the least ordinal greater than K, we denote it by Kt

Note K'+ K+1

next condition that addition

We construct recursively the following sequence of whinite auditals:

More w More wat Boo limit No = sup { No ; } < B}

Lemma: For all of, No is an infinite cardinal. Proof: By induction. Limit stage: suppose 1300 is a limit ordinal

Suppose a< No ordinal Heree a< Ny for some y< B. By induction, it is an infinite coordinal. Hence Idl< 1815 181. Thus MB 13 a cardinal.

Lemma IP a<B then No < No. Proof Exercise: induction.

So there is a strictly increasing ordinal-enumeration of audinals.

Lemma For all ordinals a, a < Na. Proof: Induction. Limit stage: B>O limit ordinal. For all geB, of Ka< NB, so B= SUBB = NB

Note: Equality is possible in this lemma, though only for limit alaribro

Proposition: Every infinite cardinal is of the form No, for some ordinal &.

Proof: Suppose Kis amade cardinal. K < MK < MKHI

So it suffices to show by induction on ordinals, For every ordinal B, and every infinite cardinal K< NB, (*) there is an ordinal ox& such that K= Na.

B+1: Suppose & K< NB+1 = NB. So either K- NB Vor K< NB / by induction Blimit: K< NB = SUP ENFIYEBS, & K< Ny Soi some J<B by induction

Dell A successor cordinal is one of the form Non for some ordinal a. A limit cardinal is one of the form No for some Unit ordinal B or O.

Note: All (instite) cordinals are limit ordinals. And K is a successor cardinal if and only if K= 1+ for some cardinal 1.

Theorem (Canhoi's Diagonalization Theorem)

For any set E, le/ IP(E)1.

Proof E - P(E), em les implies le l'é [P(E)].

Suppose |E|= |P(E)|, say Eifn(P(E). Let D= {eeE; me&f(e)} & P(E).

So Defler for some eff. Then ef A fond only it ed A. Contradiction.

The Continuum Hypothesis says that N, = (P(N.)). (Equivalently, every subset of IR is either countable or asize ITE)

The generalized Continuum Hypothesis says that K+= IP(K) for all infinite coordinals K.

These are independent of ZFC. We usually do not assume either. vor their negations.

Cordinal Alikhmetic

Det K. K. ordinals. The cardinal sum 12,+ K2 = |X,UX21 where X1/1 Xz= \$\, |X1/= K1, |X2/= K2.

Note: This does not depend on the choice of X, Xz.

The cardinal product R.K. = |X1xX2|

where Kil=K, Kil=Kz.

Again, this does not depend on the choice of X, X2.

Note: Courdinal sum and product are not the same as ordinal sum and product.

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Properties:
(a) condited simproduct are commutative and associative
(6) K; (Kz+Kz) = K,·K2 = K,·K3
(C) O+K=K, O.K=O YK
Proof: Exercise.
Theorem: Suppose k is an infinite coordinal. Then K.K.K.
Proof: K-KxK, XI->(X.X), SO K-IKI = IKXKI-K.K.
For the converse, for any ordinal a, define on Nax Na an ordering
 < as follows:
       (x,y,) < (x2,y2) if max (x,y,) < max (x2,y2)
                         or may (xily)= max (xzilyz) and xxxz
                                                      and x1: Xz and y1<4
Claim: (Nar Na, <) is a strict well-ordering (see notes).
Claim: Every initial segment of (No x No, <) is of size less
than Va.
Verilication We proceed by induction on or. Easy if 0-0.
Suppose a>0 and 5 to the withol segment for (xo.yo).
 Let z=max{xo,yo}+1, so that S= ZxZ. Now 12/=2< No.
 50 Z: KB For some BKO. Hence
        151 = 12x21 = 121/121 = NB. NB= NB< Na
by inductive assumption.
Now suppose K- Ha and Ha < No. Ha By the first claims
(Nax No, <) Ba strict well-ordering, and so is isomorphic to
(x, e) for some owned x. Now

(x, e) for some owned x. Now
so Na Ban milial segment of Y. So f'(No) Ban intial
segment of (Kax Ka, <). The second claim says If (Na)/ Na.
But I is a Wijection. Contradiction Thus No 2 No. W.
                                                                2015 02 1
 Conlary: Ki, Kz non-zero cardinals, not both finite. Then
```

Ritke= Ki-Ke= max { Ki, Ke's

Proof: Suppose K, < Kz, so Kz is infinite. Since Kz La Kxxx

X La (0,x)

we have Ros Ki. Hz. But Ki. Hz & Kz. Kz so Ki. Kz & Kz. Kz. So by theorem, all are equality, so that Ki. Kz = Kz = max {Ki, Kz}.

Next, K= SK+ H= SK+ R= SK2, 50 Equality holds and KiTH= K= MON

 $K_2 \leq K_1 + K_2$ (\downarrow) $\leq K_2 + K_2$ ($k \leq \lambda = 3 \text{ key} \leq \lambda = \gamma$) $= 2 \cdot K_2$ (bij. $2 \times K_2 \leftarrow 3 \times \gamma$, $\times n\gamma = \emptyset$, $|\chi| = |\gamma| = K_2$)

= Max (2, Kz) (above) = Kz (kz namile)

So equality holds and Hi+ Az= Kz= max fHi, Hzs.

Def) Given sets I and X, am I-sequence in X is simply a function f: I-sX, often denoted by (ai; ie I), where ai= f(i).

Suppose (Xi; ie I) is an I-sequence of sets (in some X). Then the Cartesian Roduct of (Xi; ie I) is

X Vi

the set of all I-sequences (airieI) in UX such that air Xi.

Note II I and each Xi are non-empty then by AC, is Xi + 8.

Dell Suppose (Ki; icI) is a sequence of cardinals. Then

Z K: = VX

where (Yi; ie]) is a sequence of sets such that IY: I = Ki and Yin Yj= & I i +i, and

TTRE = X YELL

where (Yi, iEI) is a sequence of sets with Wil= ki.

Well-defined: Suppose (Yi; ieI), [Yi]= Ki. Using AC, we get a sequence of bijections (f.; ieI), fi: Yi-Yi. Now

Infinite sums also trivialize:

Proposition. If I is an infinite condinal and (ki; i<1) is a sequence of non-zero condinals then

Proof: Let

Conversely,

$$\lambda = \sum_{i \in \Lambda} \{(as \text{ obove with } Kei, \text{ plus } \lambda \cdot 1 = \lambda)\}$$

$$\leq \sum_{i \in \Lambda} \{k_i, \{(ae)\}\} \{(ae)\} \{k\}\}$$

$$= \sum_{i \in \Lambda} \{k_i, \{(ae)\}\} \{(ae)\} \{(ae)\} \{(ae)\}$$

$$\leq \sum_{i \in \Lambda} \{k_i, \{(ae)\}\} \{(ae)\} \{(ae)\} \{(ae)\}$$

$$\leq \sum_{i \in \Lambda} \{k_i, \{(ae)\}\} \{(ae)\} \{$$

maxfixt= h. K s (K K;) = Ki.

What about infinite products?

Example:
$$\lambda$$
 the infinite coordinal, $K_i = 2 \ \forall i \in \lambda$.

$$\begin{array}{c|c}
\hline
1 & X & X & = |P(\lambda)| > \lambda, 2 \\
\hline
(a_i; i \in \lambda) \mapsto \{i \in \lambda; a_i = 1\}
\end{array}$$

Defl K. X cordinals

$$K^{\lambda} := |\{f_{\text{functions}}, f_{\text{rom}}, \chi_{\text{to}}, \kappa_{\text{t}}\}|$$

$$= \prod_{i \in \lambda} K.$$

Theorem (König's Theorem):
(Ki; icI), (hi; icI) sequences of cardinals, Kichi VicI
Then

Proof: (KijieI) pairwise disjoint, IXII-Ki.
(YijieI), IYII- Xi

Si Xie Yi injective (not surjective)
Construct

| Charles UX; Fx Y: choose crey (f(Xi)) iet Y: fix F(x):= fi(x) xe Xi f(x):= cg xe Xi Charles C xi xe Xi |
|---|
| Check f is injective. So [Ki < T] Ai. |
| Suppose toward a contradiction that there is a surjective map |
| For each if I. Coordinate projection |
| is surjective by AC, & (Yitd as hi> Ki=0). Hence the think is surjective. Let |
| Consider Xi Wishling Vie C |
| which is not surjective as kink him. Choose the Etin Kin. Let ae X Yi |
| such that a io = C. Let so Xe YXi |
| - such that TT:o(h(x))= a For each jEI, his |
| Note his not surjective as $K_j < \lambda_j$. Let $t \in X_i$ |

such that

Ge Yolhi (Xi)

View. Let

Ne V Xi

such that h(x)=C. Let j be such that $x \in X_i$. Then $h_j(x)=c_j d h_j(X_j)$, a contradiction

(2)